## COMS W3261 Fall 2022 Handout 7a: Midterm Review

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## 1 Problem Set A: Regular Languages

Design finite automata for the following languages. You can give the DFA/NFA by their transition diagrams. You do not need to show that they are correct.

Note: In your transition diagrams, you can use shorthand notation on the labels of the edges. For example, you can label an edge by  $\Sigma \setminus \{a\}$  to indicate that the transition takes place for all input symbols except a. Make sure to specify the starting state and the accepting states in your diagrams.

1. L is the language over the alphabet  $\Sigma = \{0, 1, 2\}$  consisting of all strings that:

- Every 0 is immediately followed by a 1, every 1 is immediately followed by a 2, and every 2 is immediately followed by a 0.
- The string starts and ends with the same symbol.
- The string must have length at least 1.

2. The set of strings over the alphabet  $\Sigma = \{a, b, \dots, z\}$  that contain at least one *m* between any two *a*'s in the string; for example *abc*, *john*, *mama*, *american* are in the language, but *papa*, *panamerican* are not.

3. Challenge: The set of binary strings which represent in binary a number that is an integer multiple of 3 (leading zeros are allowed). For example, 00, 110 are in the language (they correspond to 0, 6 respectively), but 001, 101 are not (they correspond to 1, 5). Hint: Three states are enough. Think about what each additional symbol in a binary string does to the number; it might be helpful to write out some examples.

For the following problems, if a language L is given, prove that L is regular or prove that L is nonregular. 4.  $L = \{ww \mid w \in \{0,1\}^* \text{ and } w \text{ contains at least one } 0 \text{ and at least one } 1\}$  over the alphabet  $\Sigma = \{0,1\}$ .

5. Prove that the union of a regular language  $L_1$  and nonregular language  $L_2$  such that  $L_1 \cap L_2 = \emptyset$  results in a nonregular language.

6.  $L = \{a^i b^j c^k \mid i+j=k\}.$ 

7. 
$$L = \{0^k u 0^k \mid k \ge 1, u \in \Sigma^*\}.$$

8.  $L = \{0^k 1 u 0^k \mid k \ge 1, u \in \Sigma^*\}.$ 

9.  $L = \{0^i 1^j \mid i, j \ge 0, i \ne j\}.$ 

10.  $L = \{0^n 1^n \mid 0 \le n \le 3\}.$ 

11. Challenge: L is the language consisting of all strings of a's and b's with an equal number of occurrences of ab and ba as substrings. (The string aabbbaa has one occurrence of each of the substrings ab and ba.)

## 2 Problem Set B: Context Free Languages

1. Show that the language

$$L = \{a^{i}b^{j} \mid i, j \ge 0, i \le j \le 3i\}$$

is context-free.

2. Give a context-free grammar that generates the language

$$L = \{a^i b^j c^k : i = j \text{ or } j = k \text{ where } i, j, k \ge 0\}$$

Is your grammar ambiguous? Why or why not? (Consider the string  $a^n b^n c^n$ .)

3. Construct a pushdown automata for the language

$$L = \{a^n b^k c^n \mid n, k \ge 0\}$$

4. Describe the language that the following PDA recognizes:



5. Challenge: For any language L, let  $\mathsf{SUFFIX}(A) = \{v \mid uv \in L \text{ for some string } u\}$ . Show that the class of context-free languages is closed under the  $\mathsf{SUFFIX}$  operation.

6. Challenge: Show that the complement of the language

$$L = \{ww \mid w \in \Sigma^*\}$$

is context-free.