

TXBOX: BUILDING SECURE, EFFICIENT SANDBOXES WITH SYSTEM TRANSACTIONS

Suman Jana

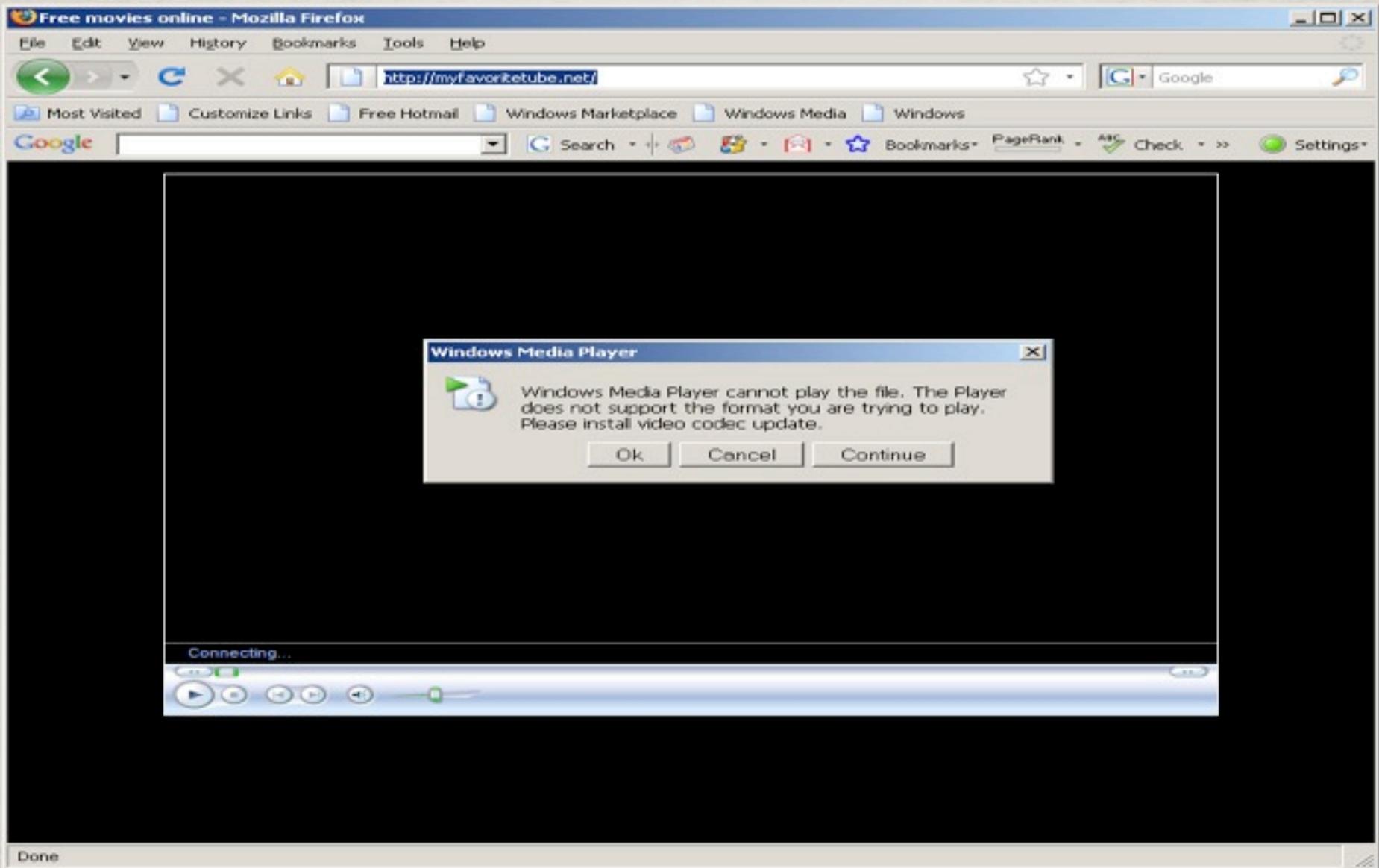
Vitaly Shmatikov

The University of Texas at Austin

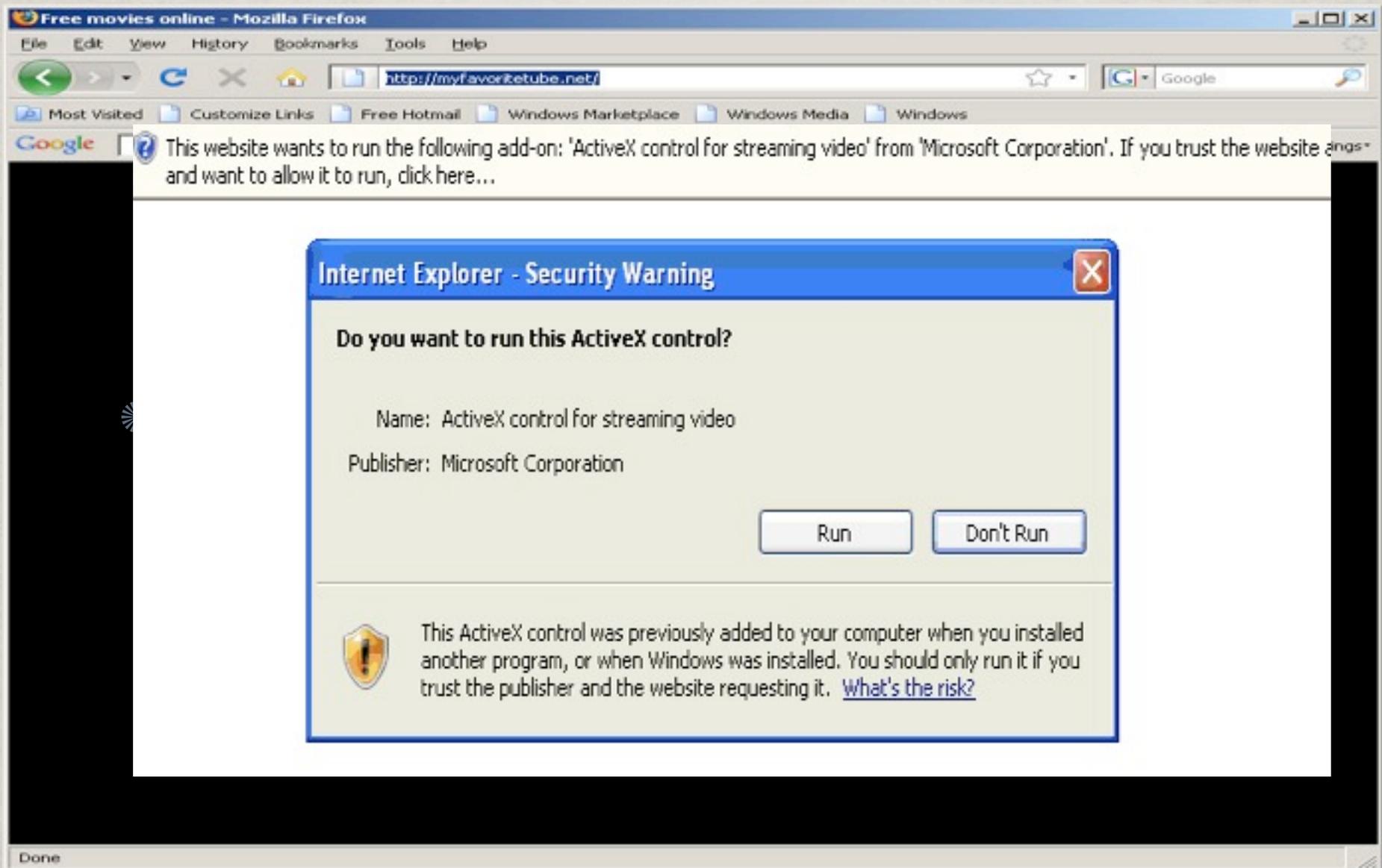
Donald E. Porter

Stony Brook University

UNTRUSTED CODE IS EVERYWHERE !



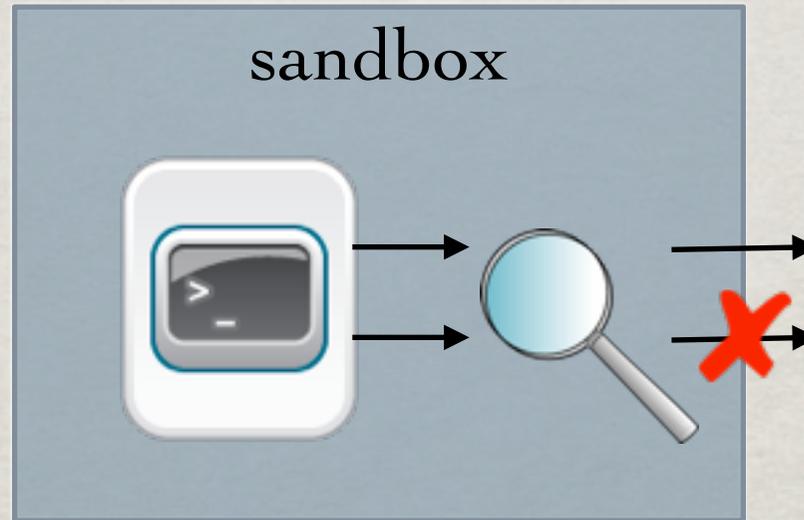
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SANDBOX: RESTRICT UNTRUSTED CODE

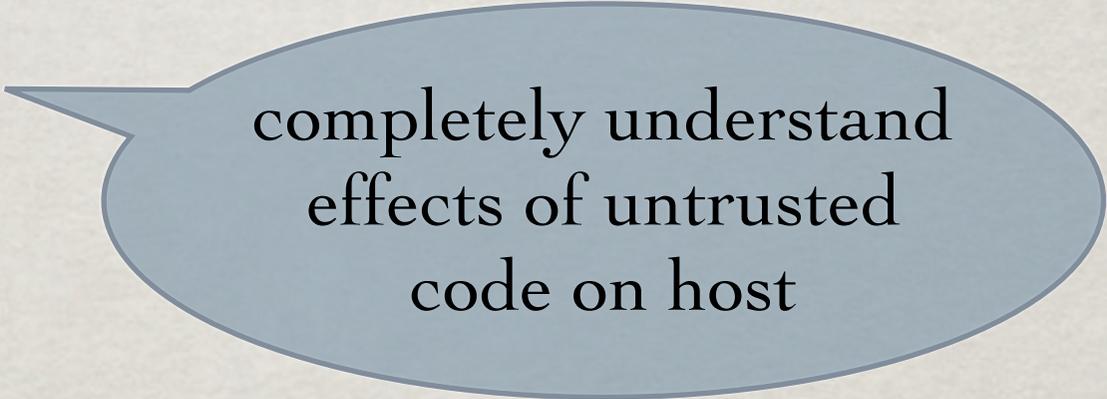


- ✿ Sandbox restricts untrusted code
 - ✿ Files it can read/write
 - ✿ System calls and arguments it can use

PROPERTIES A SANDBOX SHOULD HAVE

- ✱ Uncircumventability

- ✱ Fidelity



completely understand
effects of untrusted
code on host

- ✱ Separation policy enforcement and policy specification

- ✱ Performance

A QUICK SURVEY OF SOME SANDBOXING TECHNIQUES

STATIC ANALYSIS

untrusted code



detect malicious code
using static-analysis

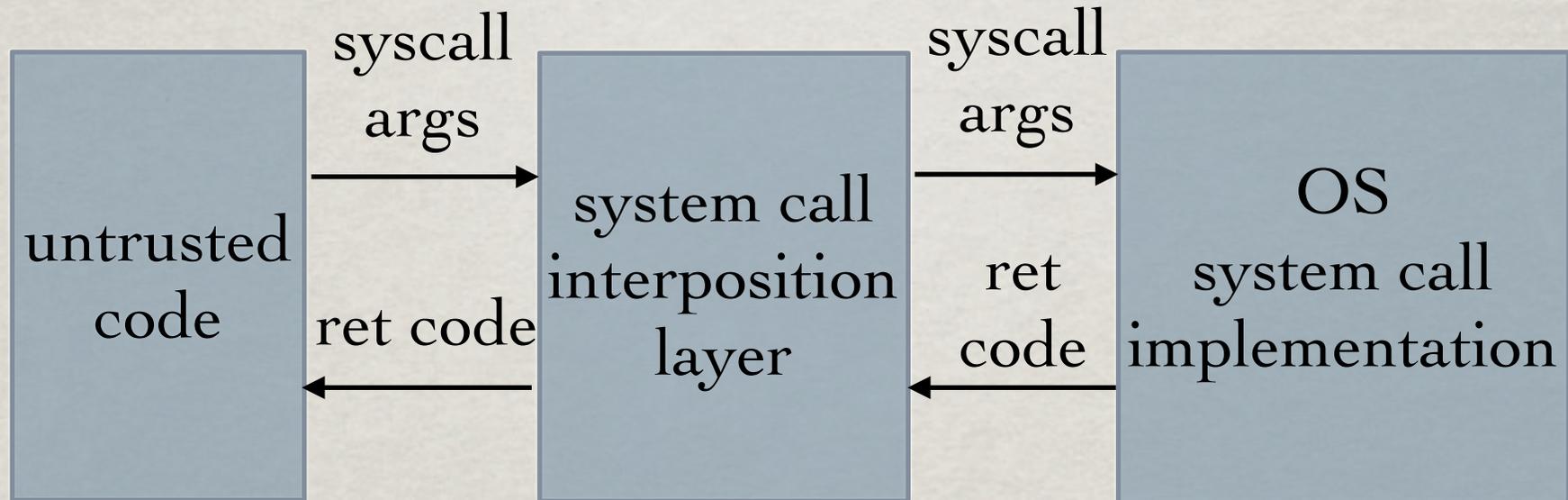
static-analysis
is imperfect:
false negatives



if benign

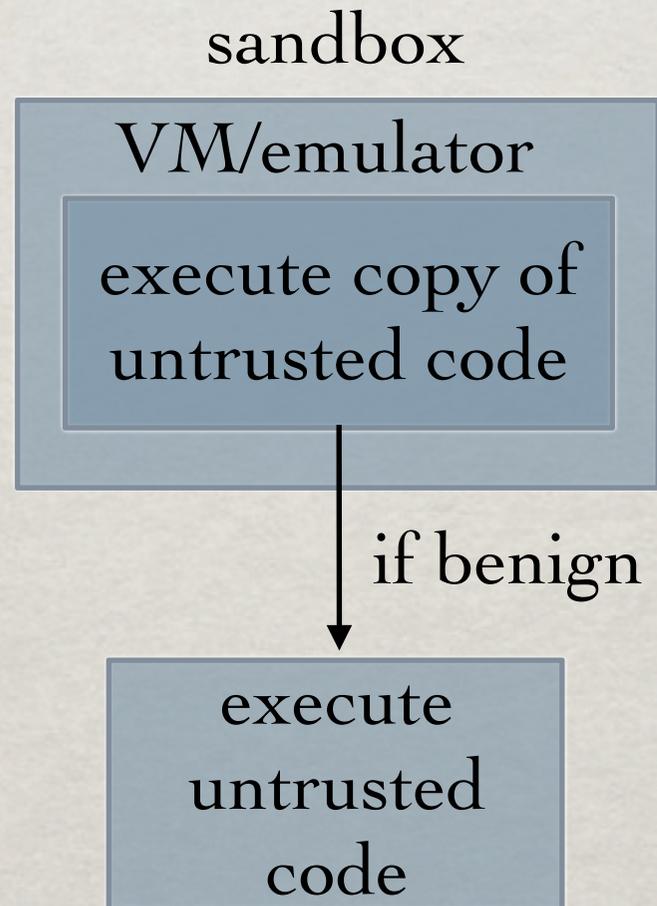
execute code

SYSTEM CALL INTERPOSITION

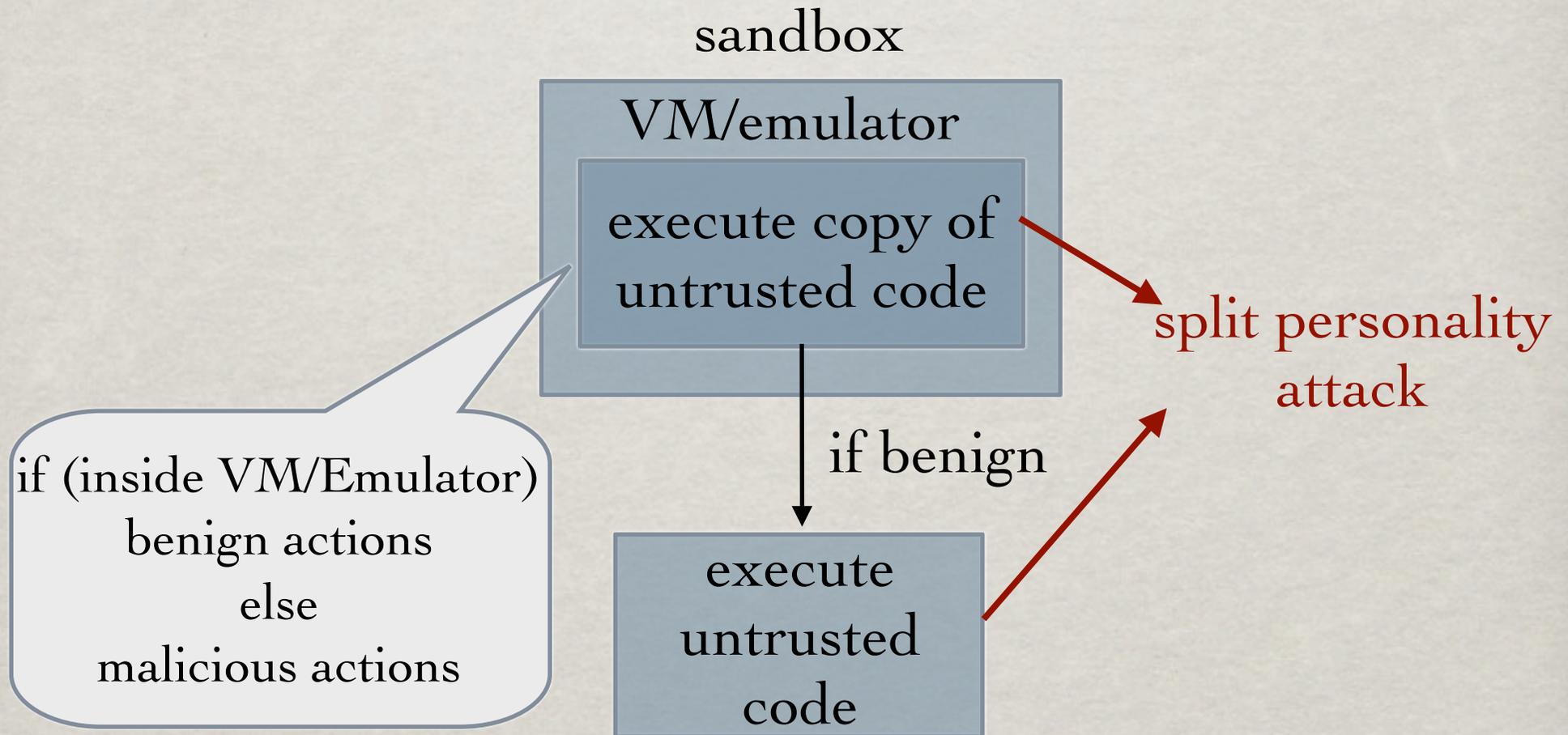


- ✱ Incorrect mirroring of system state
- ✱ Time of check to time of use (TOCTOU) attacks

BUILDING SANDBOXES WITH VMS/EMULATORS



BUILDING SANDBOXES WITH VMS/EMULATORS



FIDELITY: NECESSARY FOR UNCIRCUMVENTABILITY

- ✱ Understand behavior of untrusted code
 - ✱ Semantic gaps can lead to circumvention
- ✱ Coherent view of all actions performed by untrusted code
 - ✱ System calls and arguments
 - ✱ All affected files (read/write)

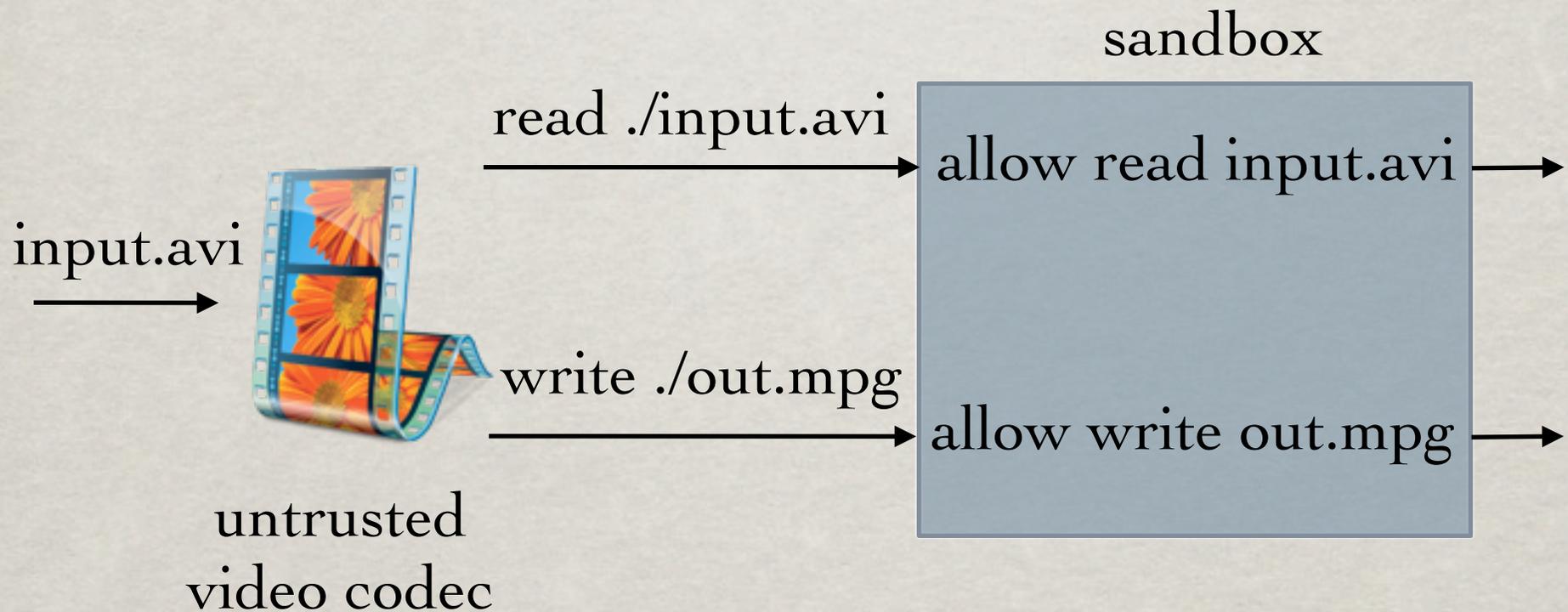
SANDBOX POLICIES

How should a sandbox decide which actions to allow/deny ?

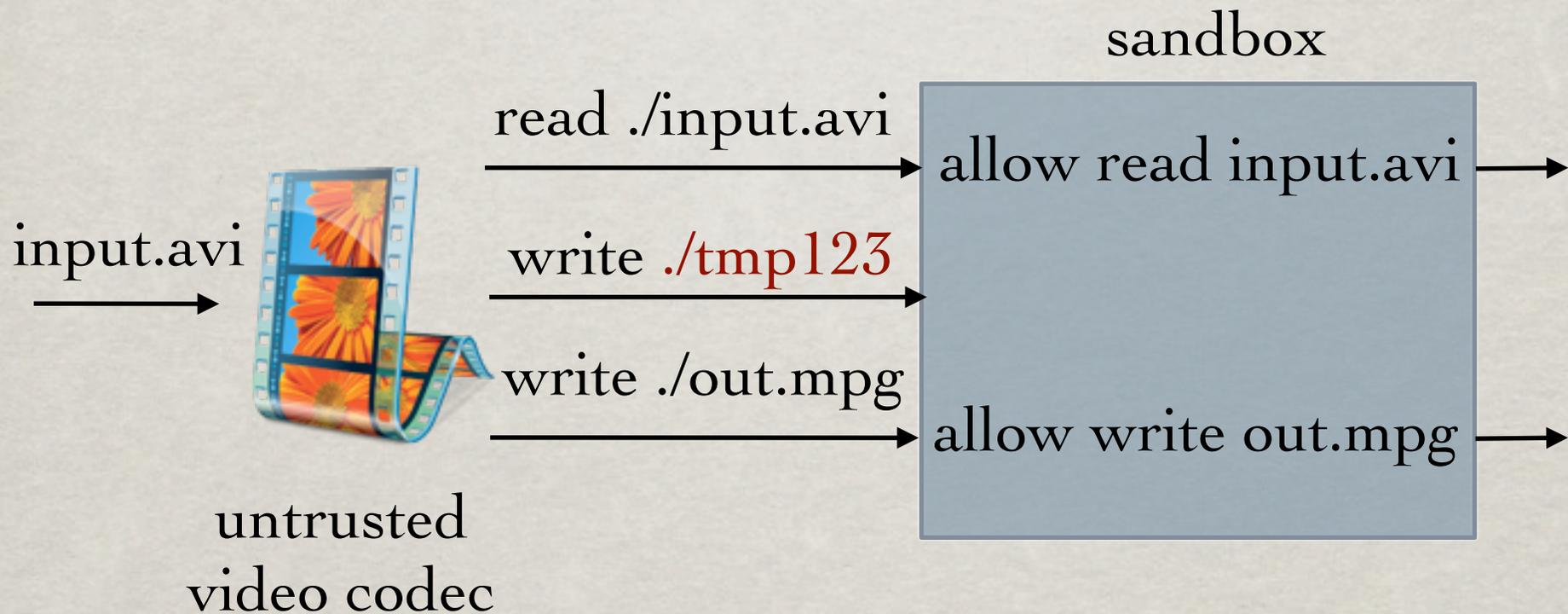
LEAST PRIVILEGE MODEL

- ✱ Whitelist minimal set of operations needed for correct functionality of untrusted code
- ✱ Users only have partial information
- ✱ Difficult to implement in practice
 - ✱ Overestimate: untrusted code can cause more damage
 - ✱ Underestimate: crippled functionality

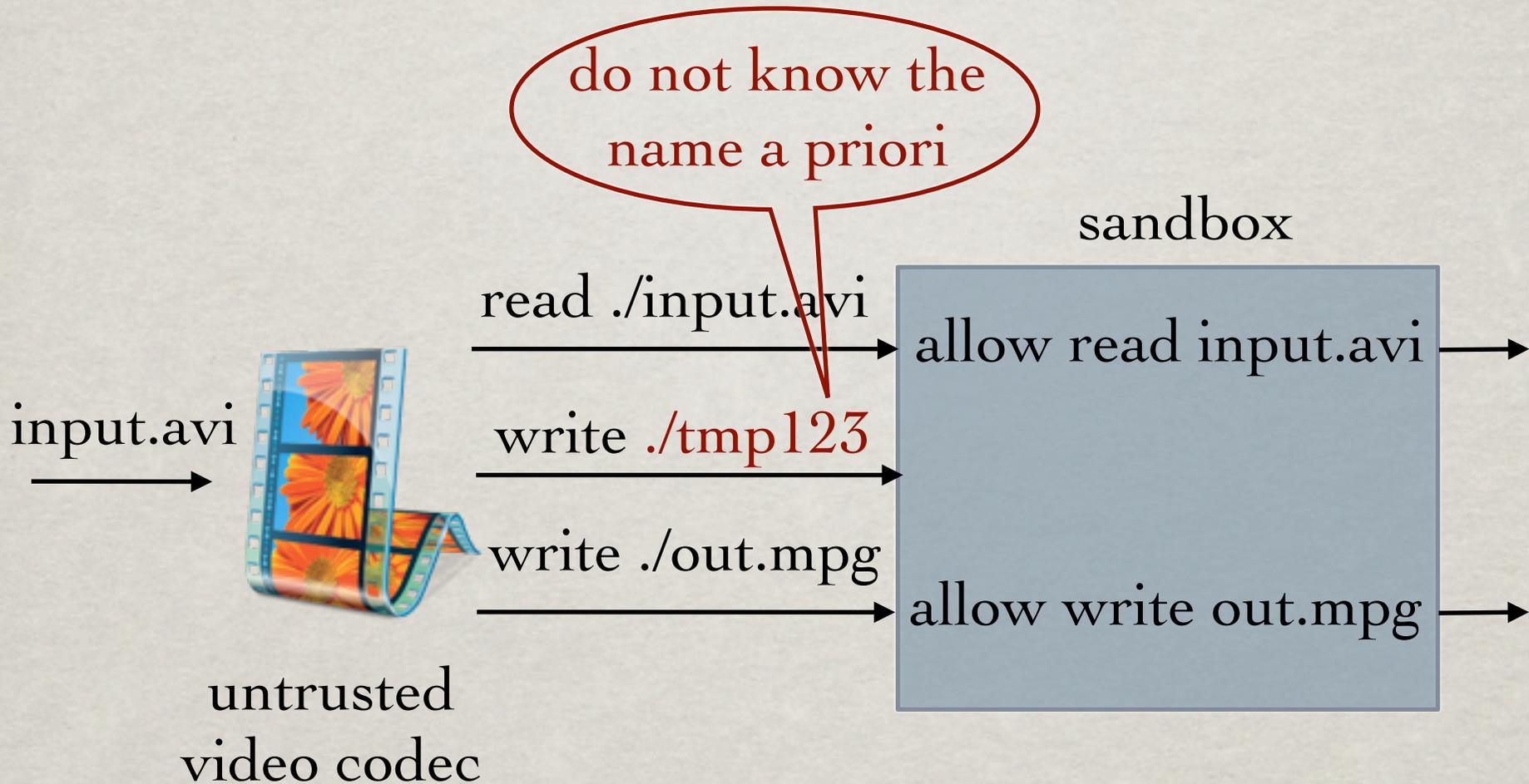
LEAST PRIVILEGE MODEL: DIFFICULTIES



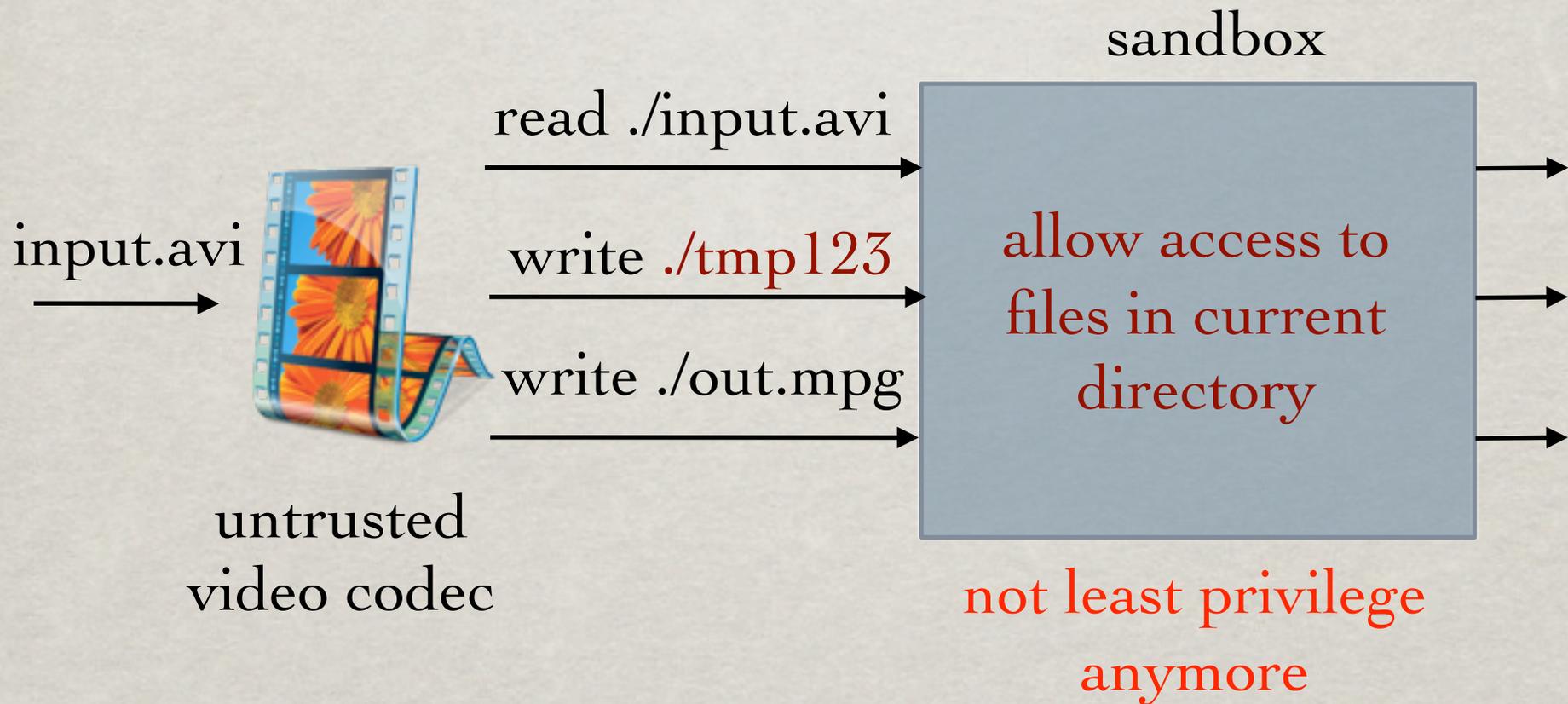
LEAST PRIVILEGE MODEL: DIFFICULTIES



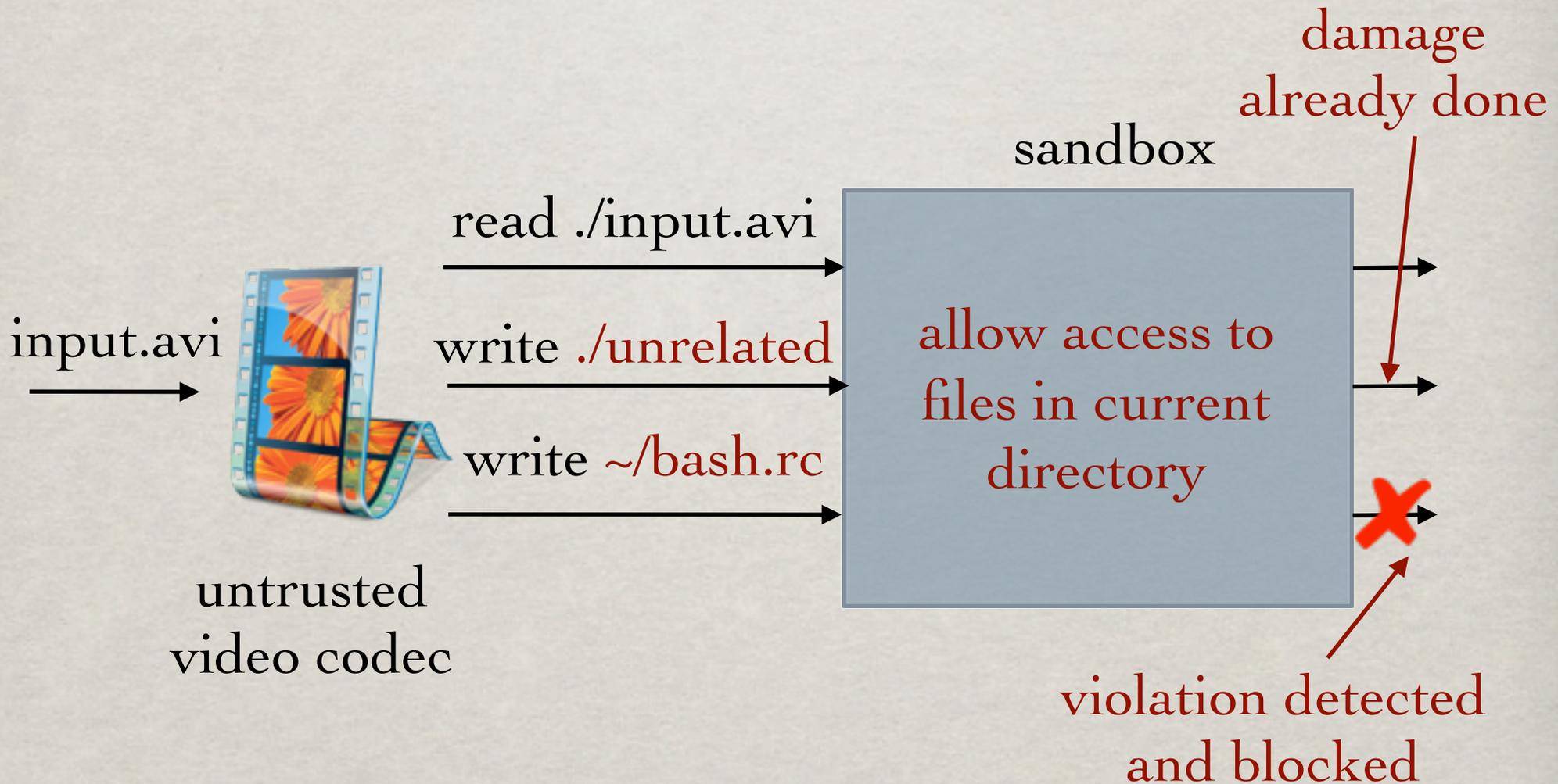
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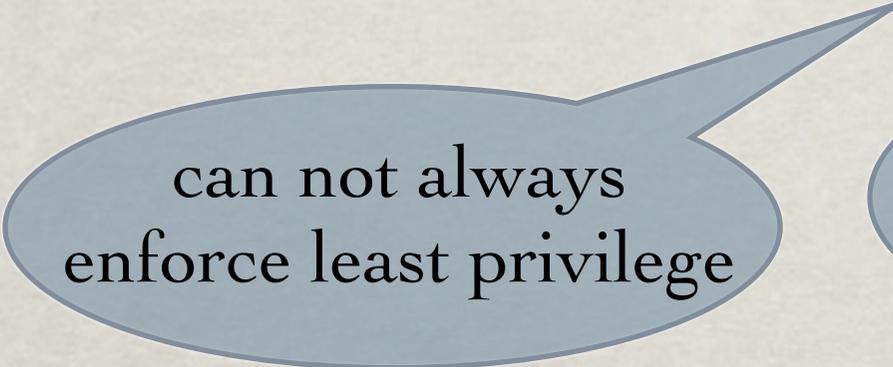


LEAST PRIVILEGE MODEL: DIFFICULTIES

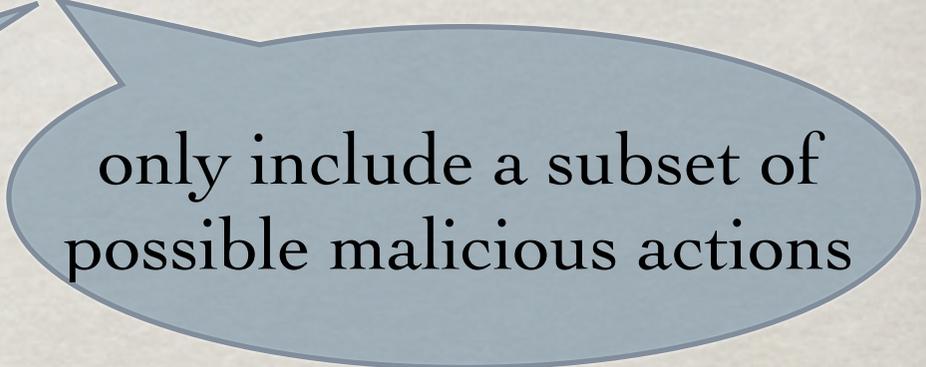


RECOVERABILITY

- ✱ Once a sandboxed process is detected doing anything bad, rollback all changes to be safe
- ✱ Real sandboxes have imperfect policies



can not always
enforce least privilege



only include a subset of
possible malicious actions

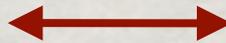
- ✱ Sandboxes with perfect policies may not need recoverability

RECOVERABILITY CAN INCREASE PARALLELISM

security checks
(e.g. virus scanning)



parallel



sandboxed code

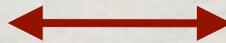


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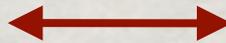
is sandboxed code
malicious ?

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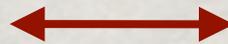
yes

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security checks
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parallel



sandboxed code



is sandboxed code
malicious ?

yes

recover

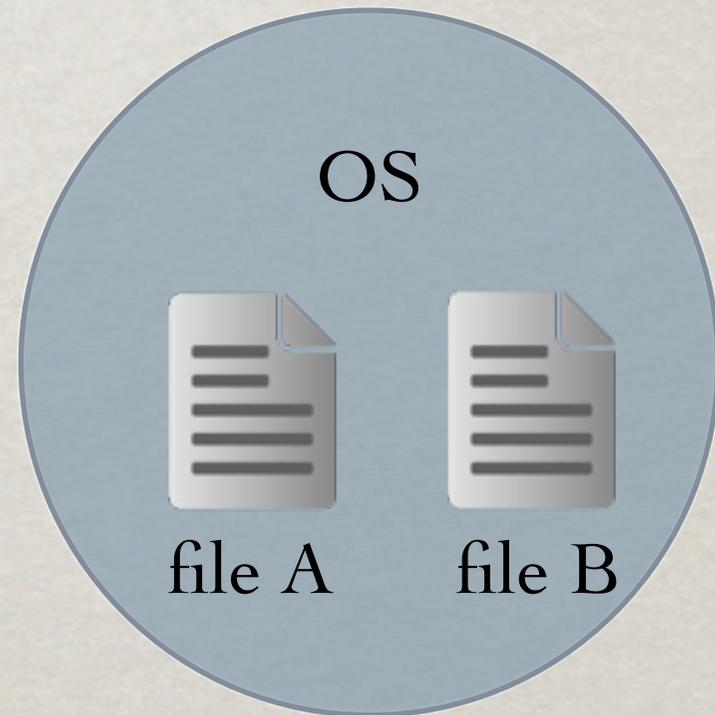


PROPERTIES A SANDBOX SHOULD HAVE

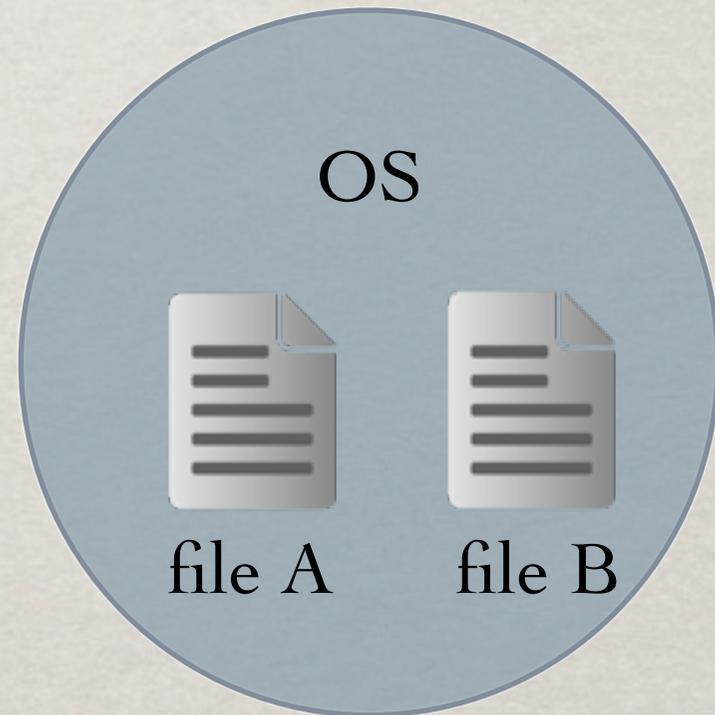
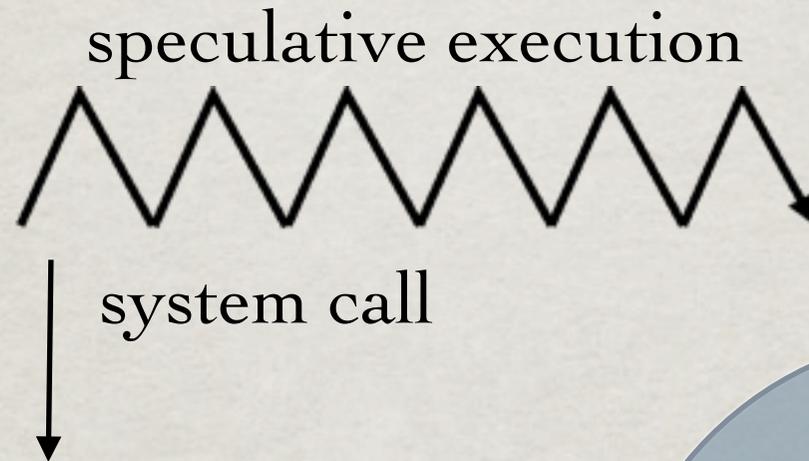
- ✻ Uncircumventability
- ✻ Separation policy enforcement and policy specification
- ✻ Performance
- ✻ Recoverability

OS TRANSACTIONS

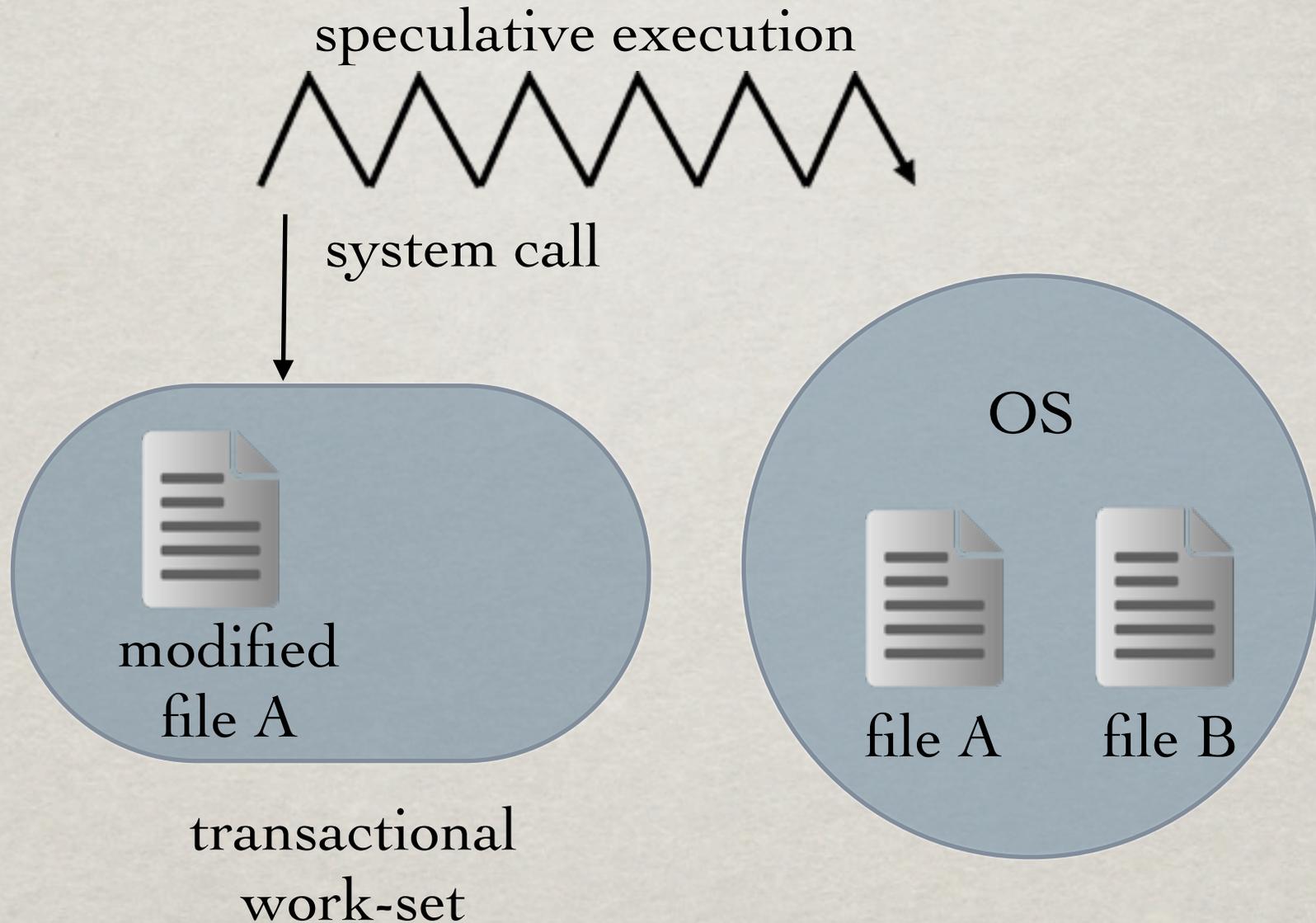
speculative execution



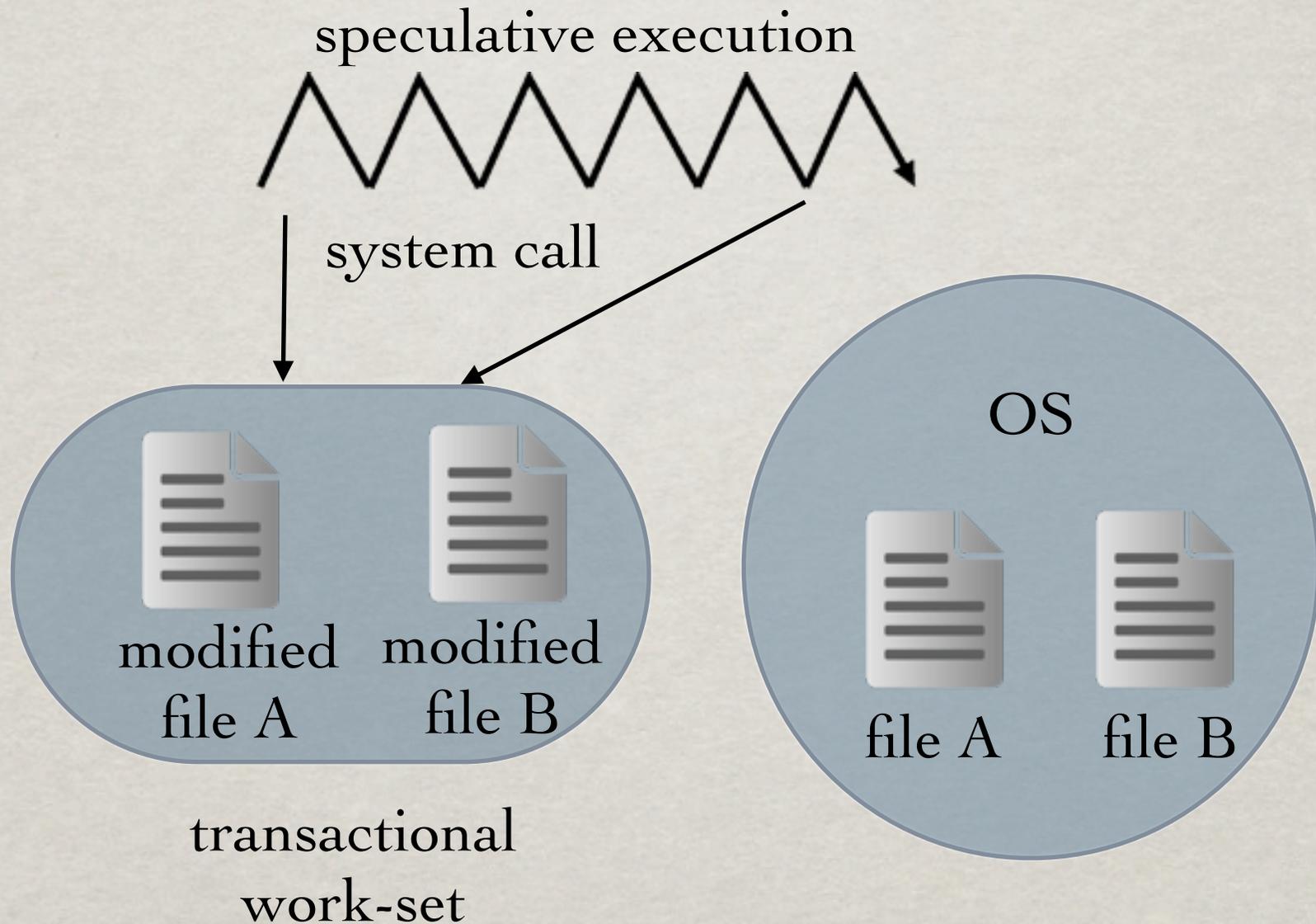
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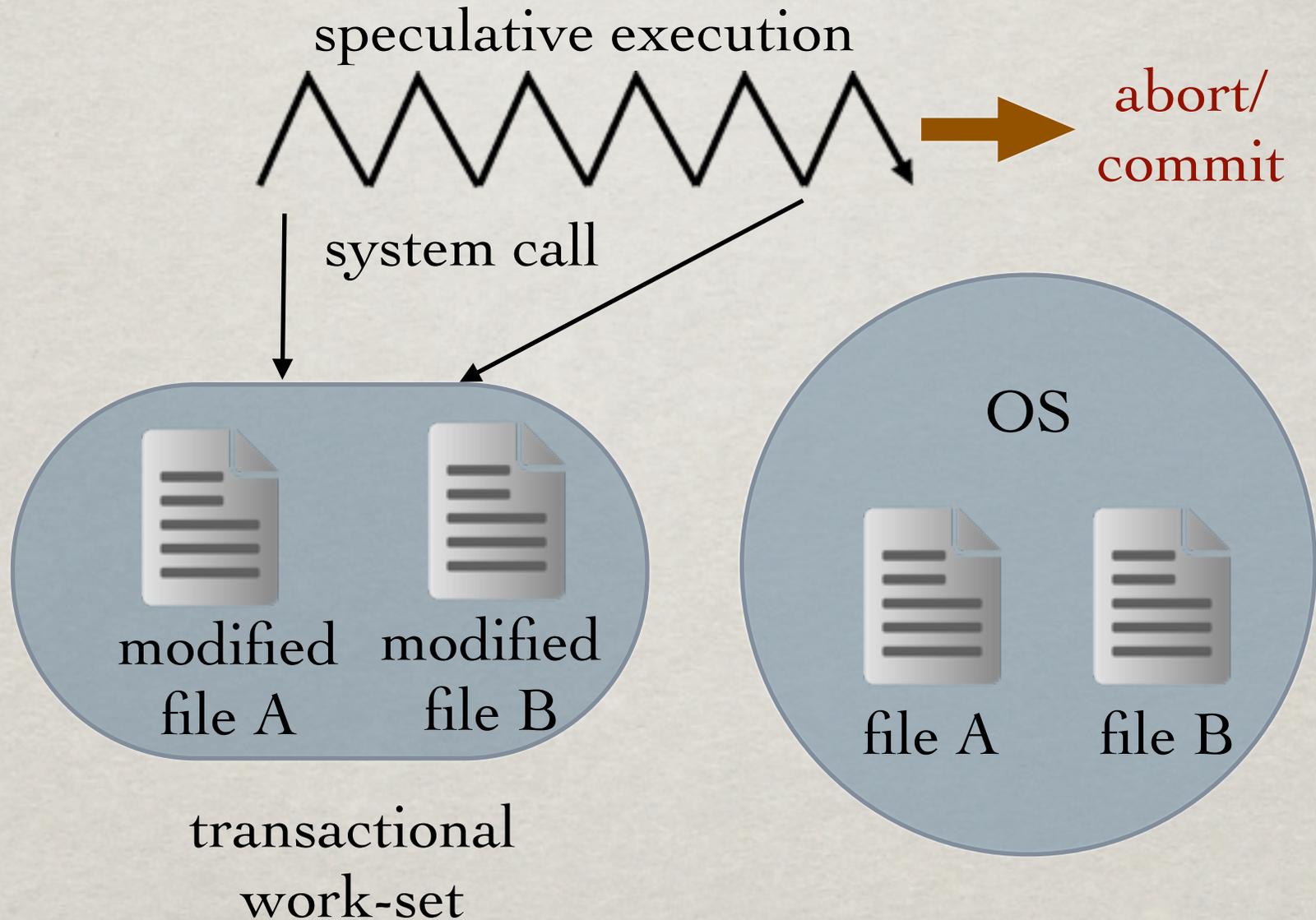
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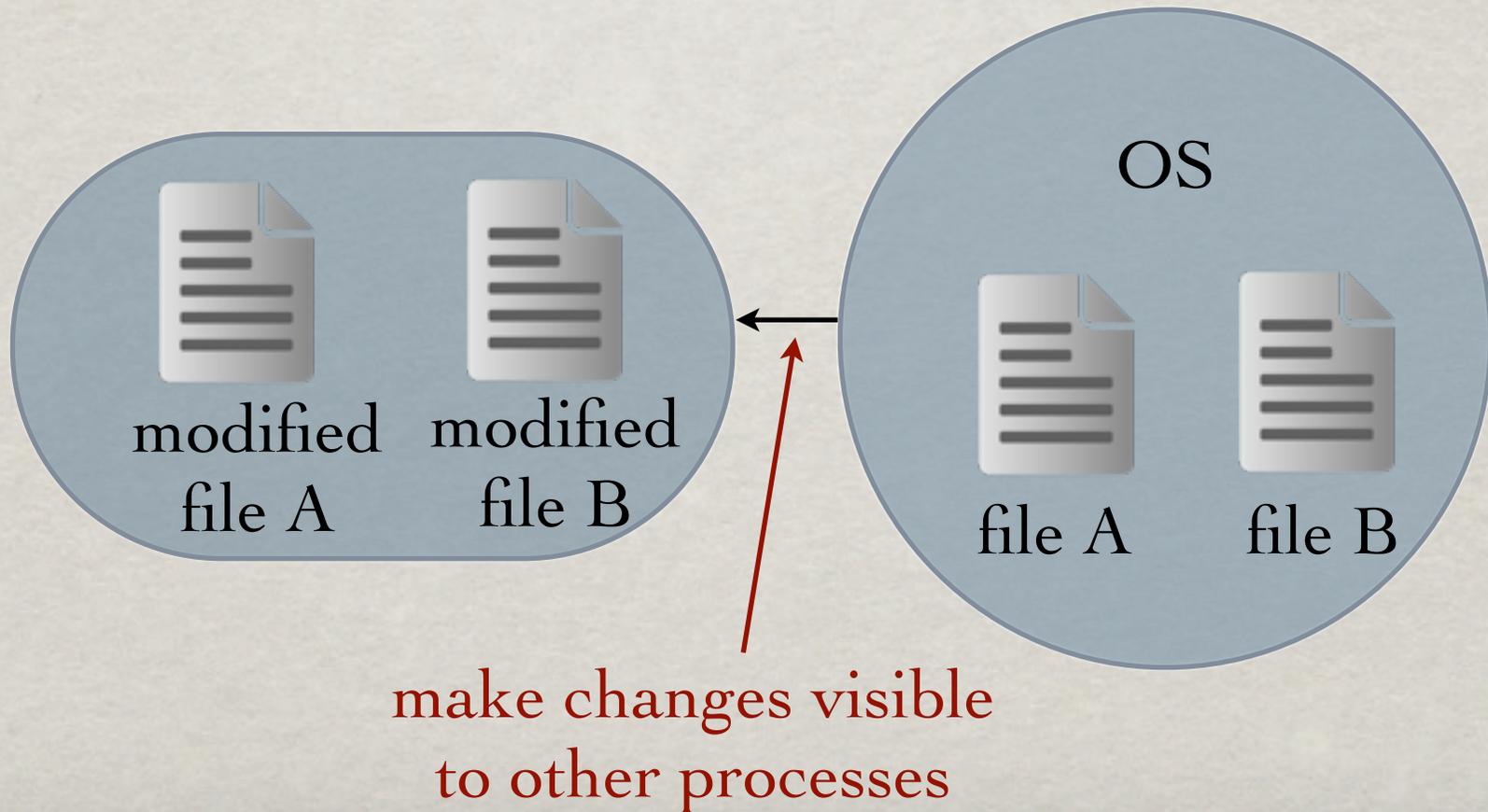
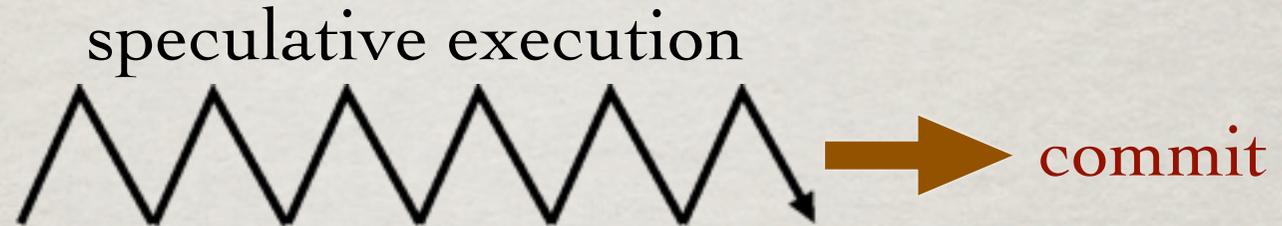
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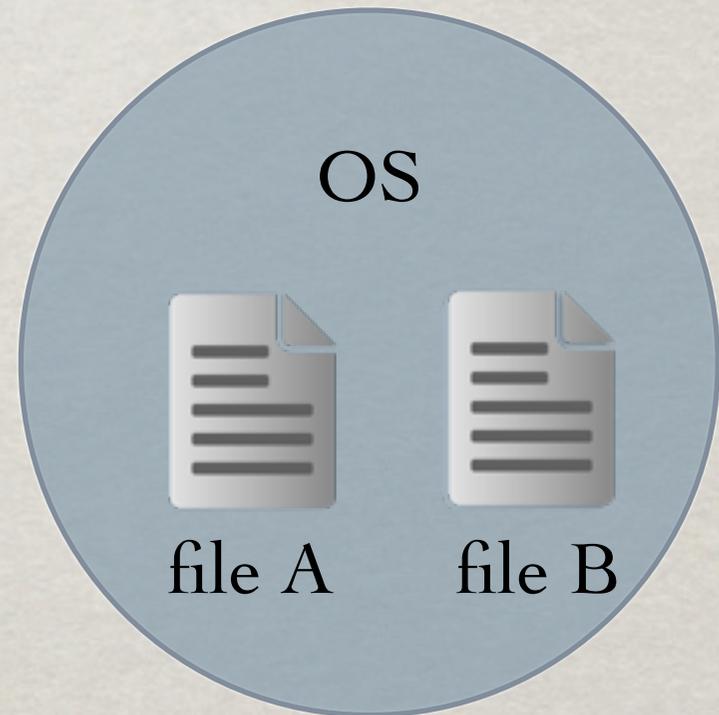
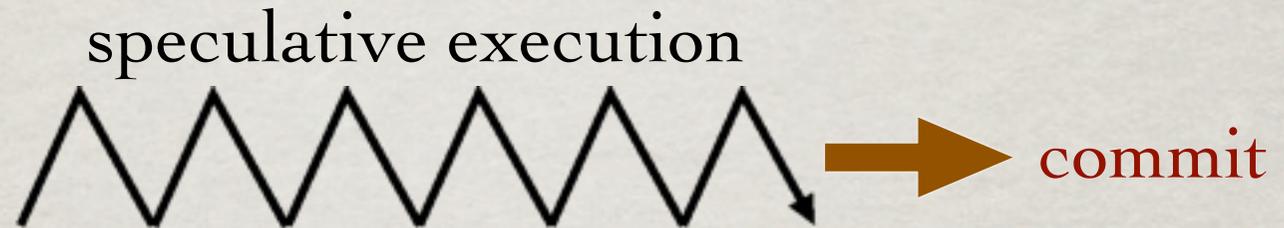
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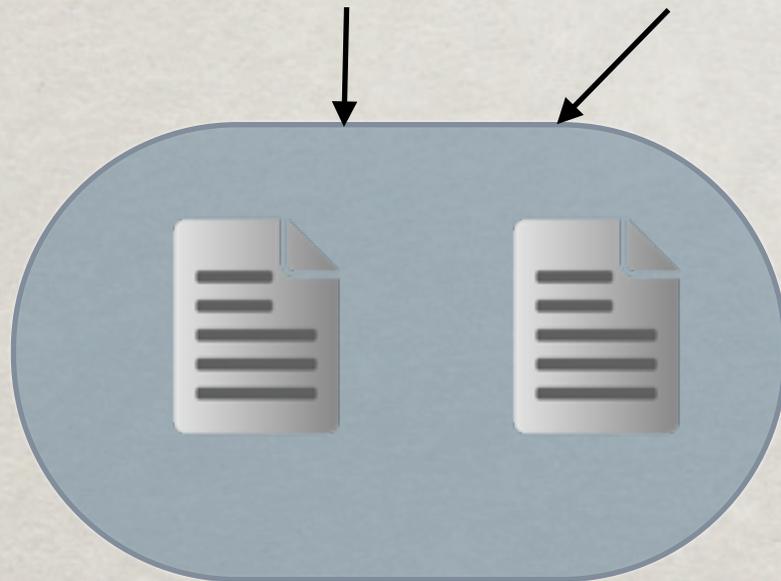


OS TRANSACTIONS

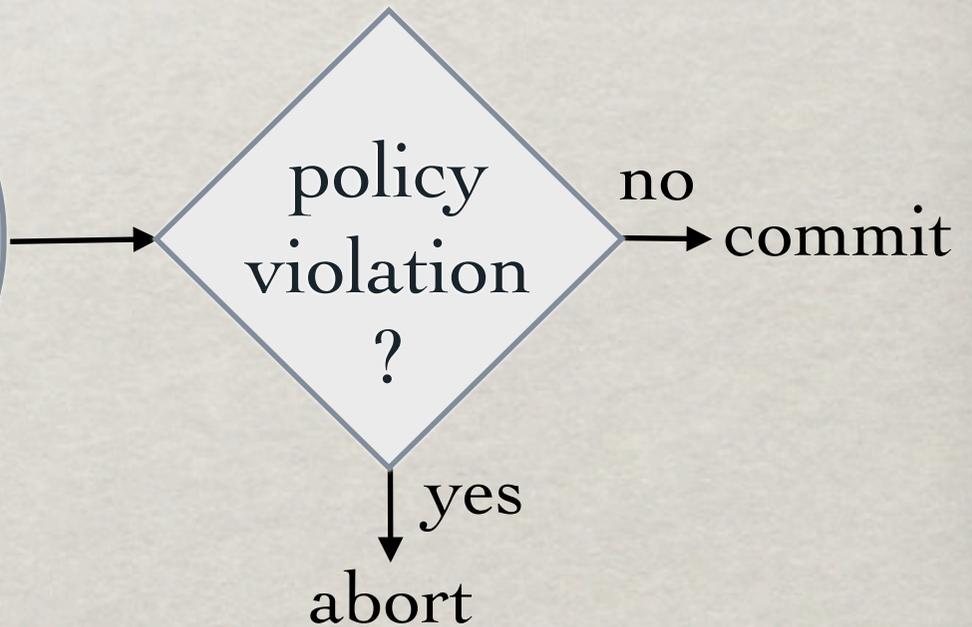


SECURITY NEEDS TRANSACTIONS

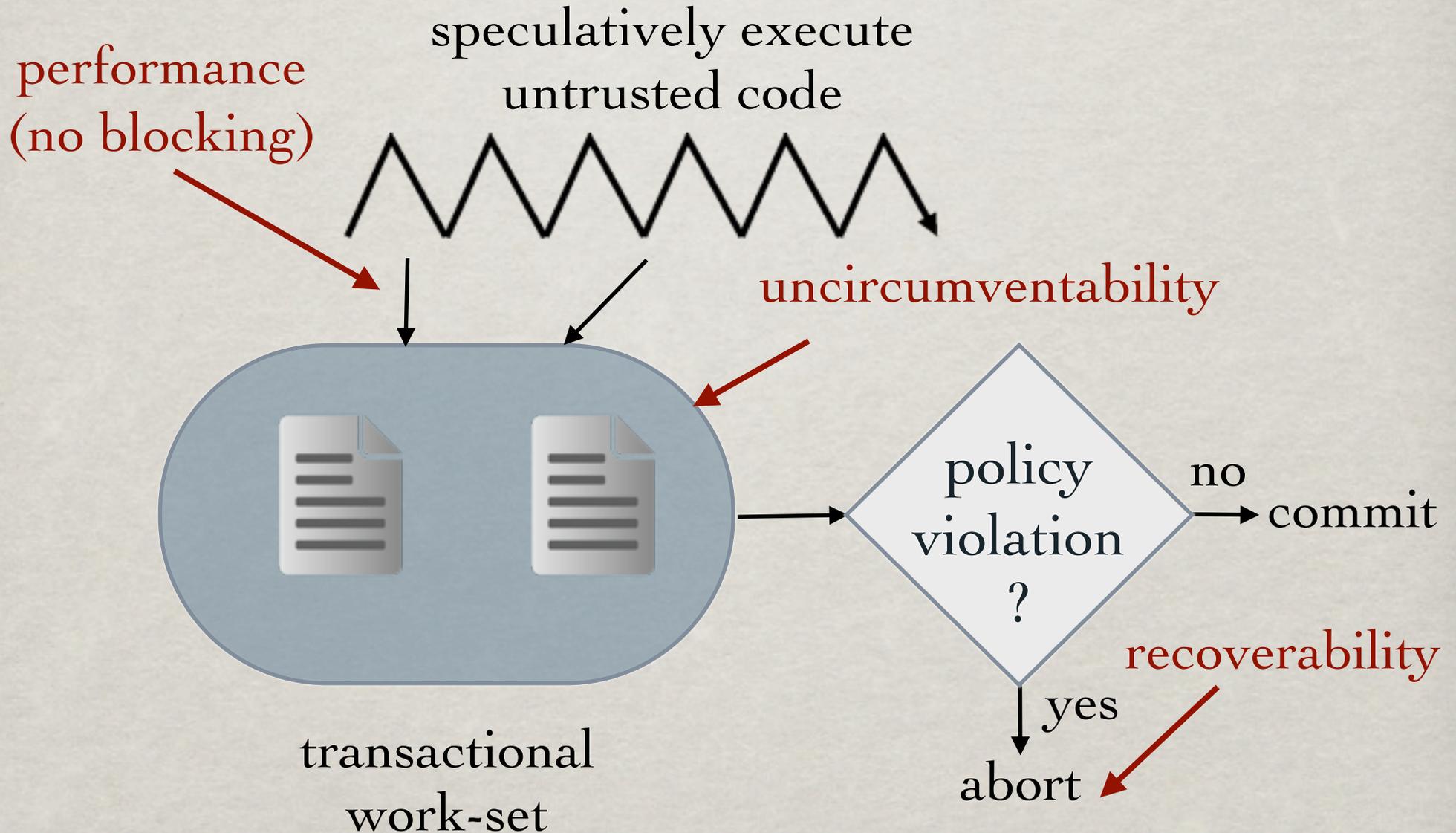
speculatively execute
untrusted code



transactional
work-set



SECURITY NEEDS TRANSACTIONS



OS SUPPORT FOR TRANSACTIONS

- ✱ TxOS : Porter et al. SOSPP 2009
- ✱ Speculative execution support for 150+ system calls
- ✱ Provides ACID semantics
- ✱ Originally done for handling concurrency

TxB0x



TxB0X

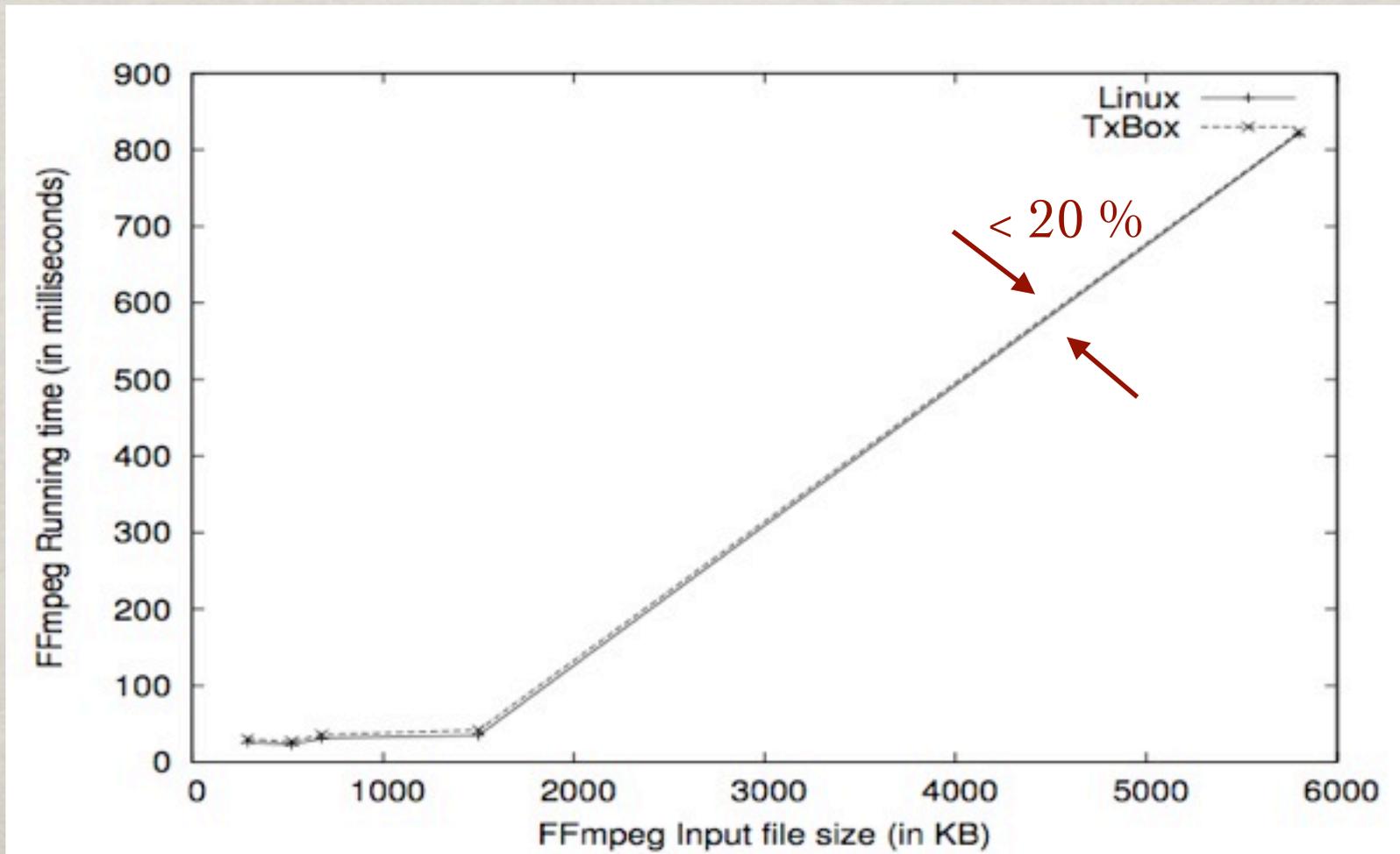
- ✿ Insight: transactions are great match for security
- ✿ Execute untrusted code inside a transaction
- ✿ Make security decisions by checking work-set
- ✿ Parallelize security checks with program execution
- ✿ Abort transaction if anything malicious is detected

EVALUATION

- ✱ Can TxBBox isolate large real-world programs?
 - ✱ FFmpeg : audio/video codec
 - ✱ SpiderMonkey : JavaScript engine
 - ✱ Vim : editor
- ✱ How much performance/memory overhead does TxBBox incur ?

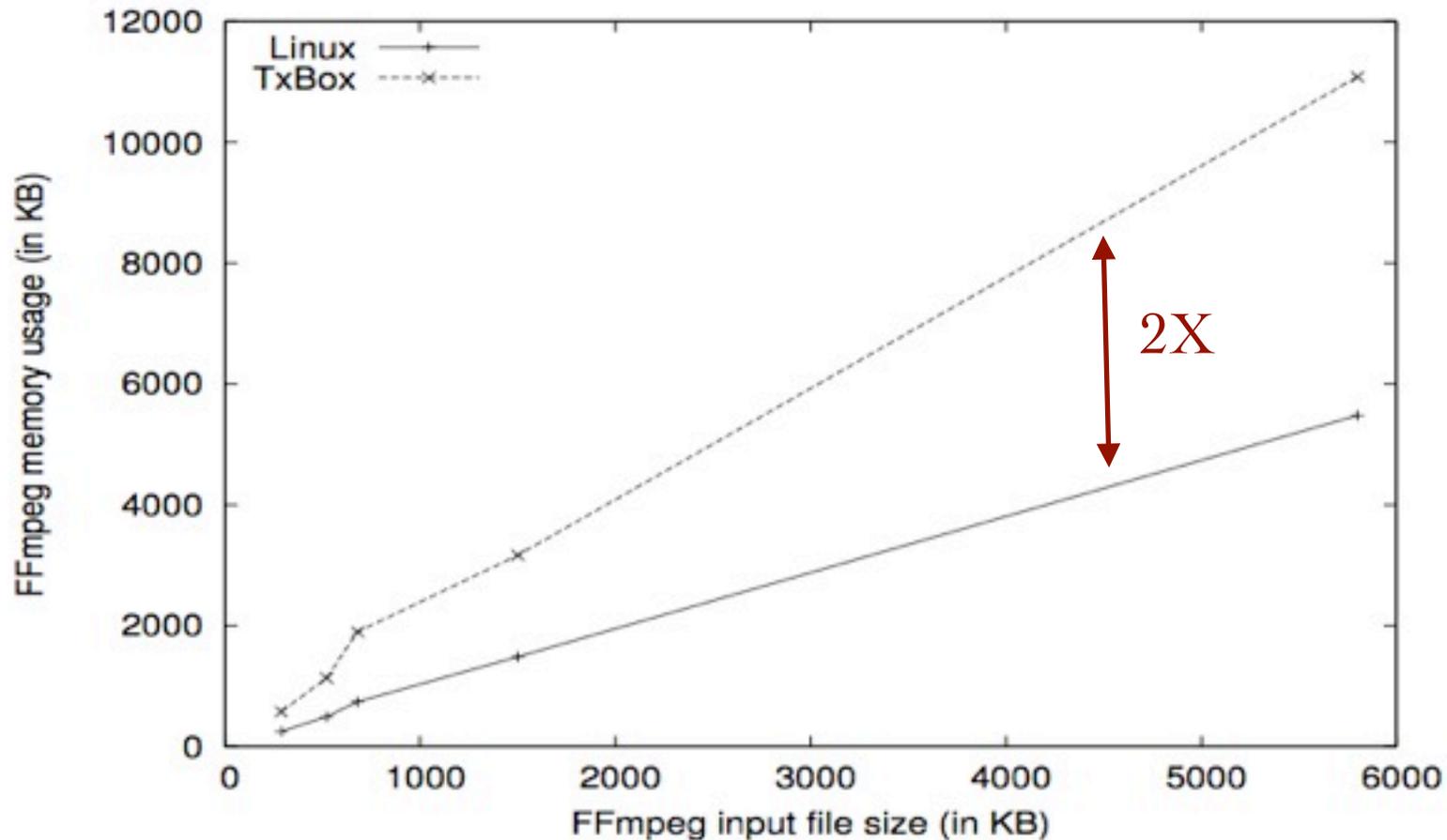
TxBox: PERFORMANCE OVERHEAD

- ✿ On average TxBox causes less than $< 20\%$ runtime overhead compared to Linux

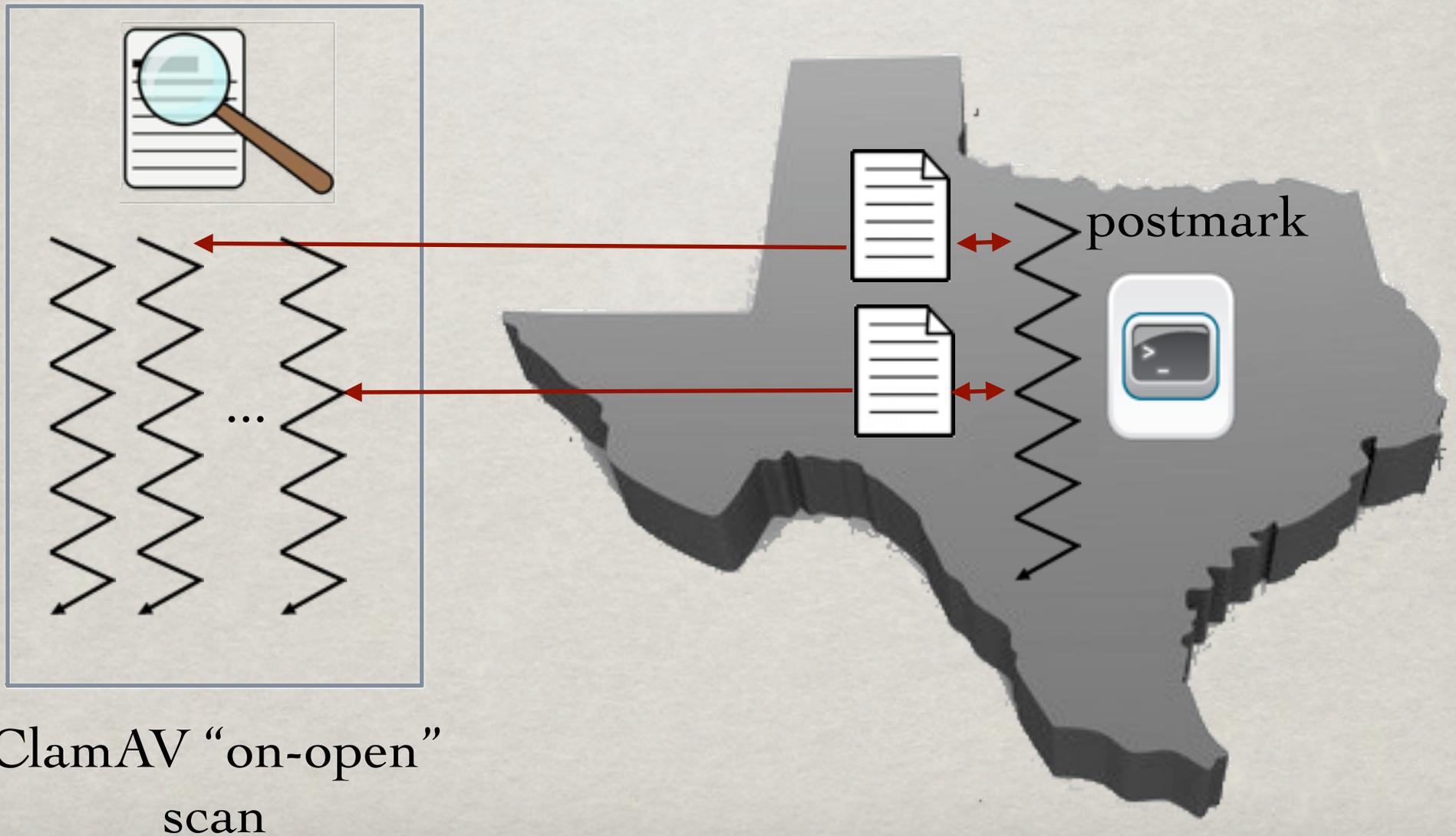


TxBox: MEMORY OVERHEAD

- ✿ On average TxBox execution of a process takes 2x more memory compared to regular Linux execution

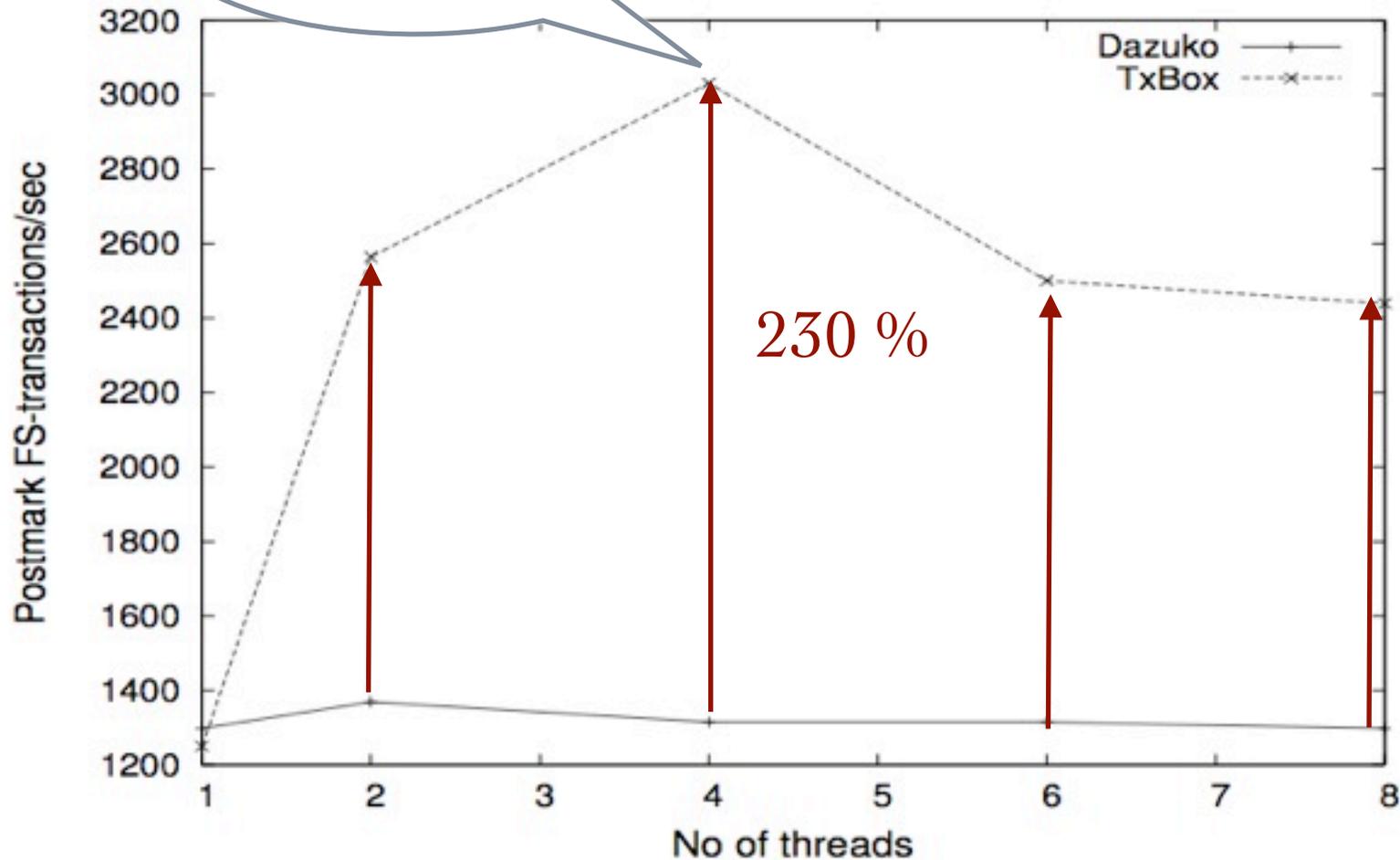


TxB0X: PARALLEL ANTIVIRUS SCANNING



TxBOX PARALLELIZATION GAIN (CLAMAV SCANNING)

host had
4 cores



CONCLUSION: SECURITY NEEDS TRANSACTIONS

- ✱ Speculatively execute untrusted code
- ✱ Rollback if any malice is detected
- ✱ Inspect all effects of the untrusted process at the right level of abstraction
- ✱ Prevent circumvention and evasion

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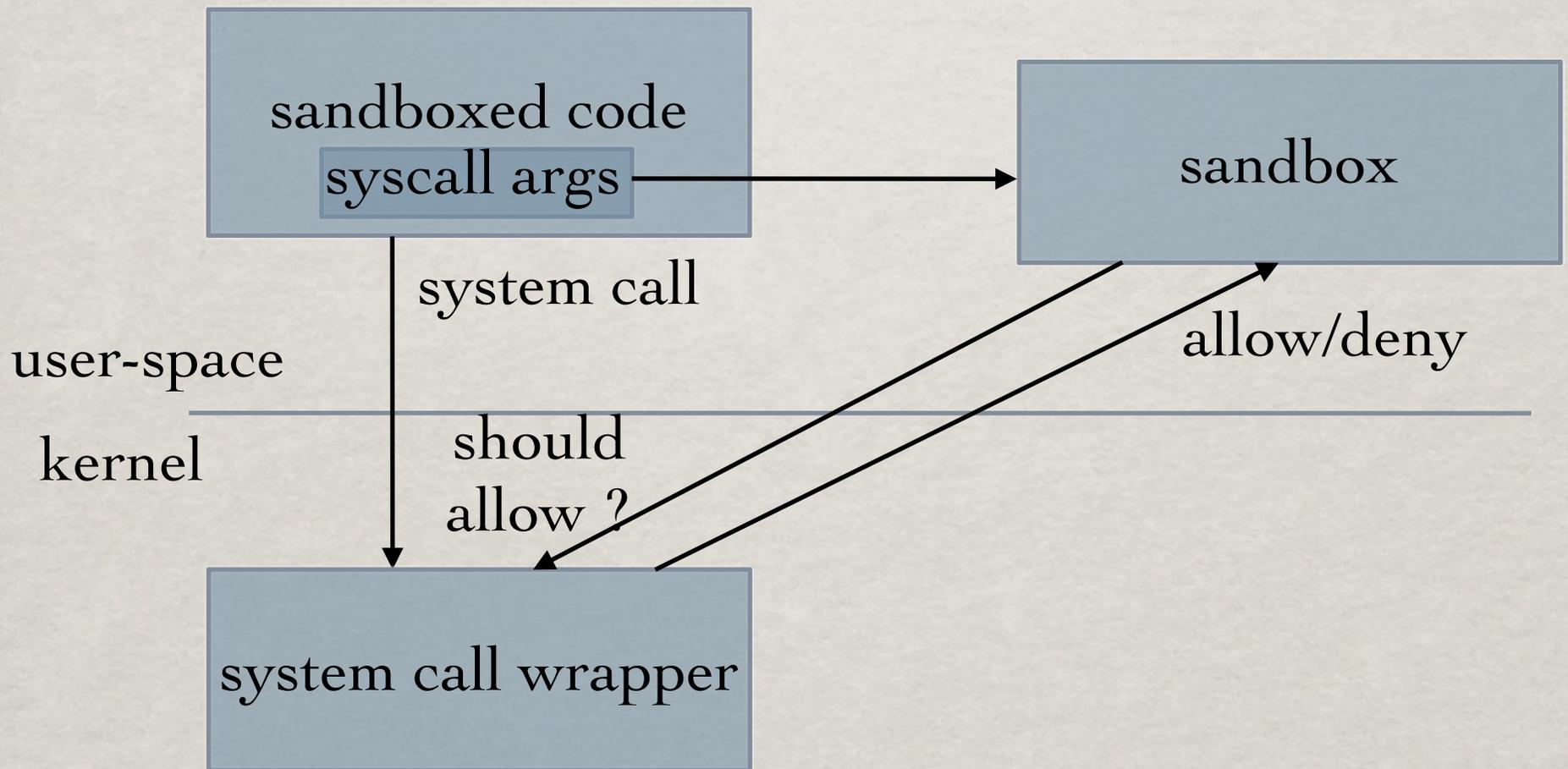
RECOVERABILITY: OUTPUT COMMIT PROBLEM

- ✱ How to maintain recoverability if an untrusted process performs network i/o ?
- ✱ Unsolvable in general, we do the next best thing
- ✱ Always preserve local recoverability
 - ✱ Deny network i/o and continue
 - ✱ Execute network i/o outside of transaction and continue

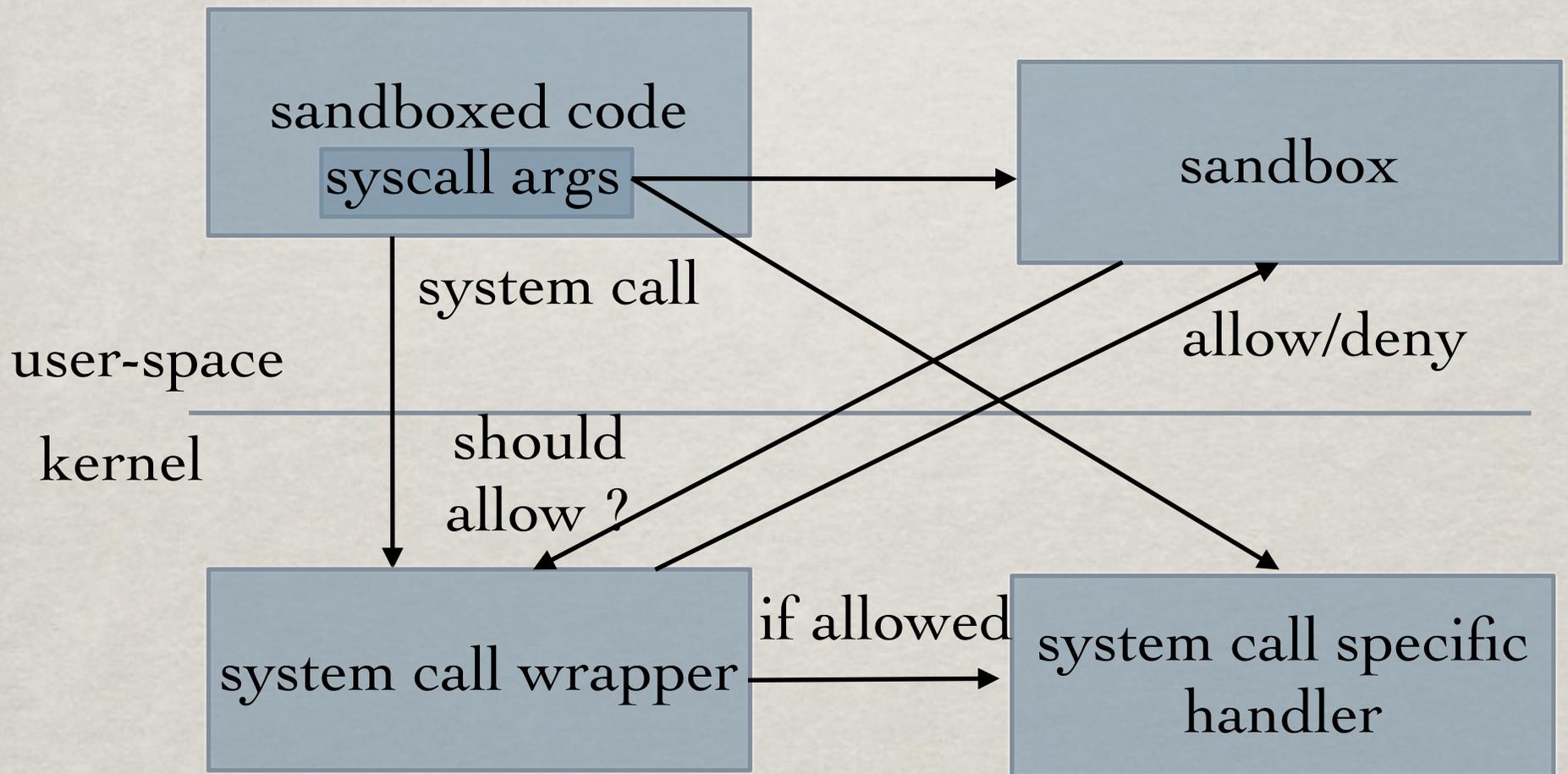
TxBOX: IMPLEMENTATION ISSUES

- ✱ TxOS transactions need cooperative processes calling
 - ✱ xbegin
 - ✱ xend
- ✱ Untrusted processes are not co-operative
 - ✱ Support “forced” transactions
- ✱ Implement policy manager and policy enforcer
- ✱ See paper for details

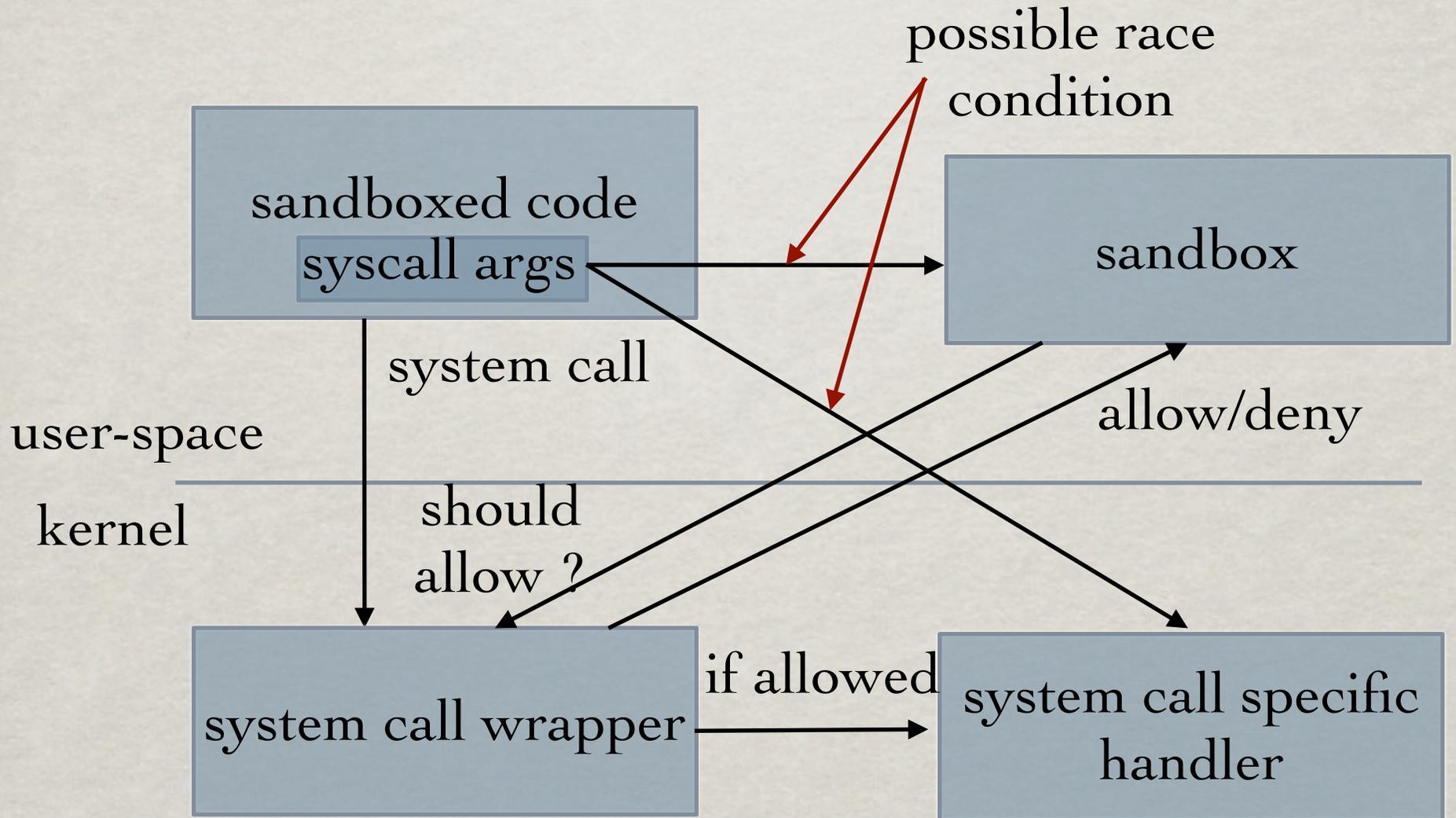
BUILDING SANDBOXES WITH SYSTEM CALL INTERPOSITION



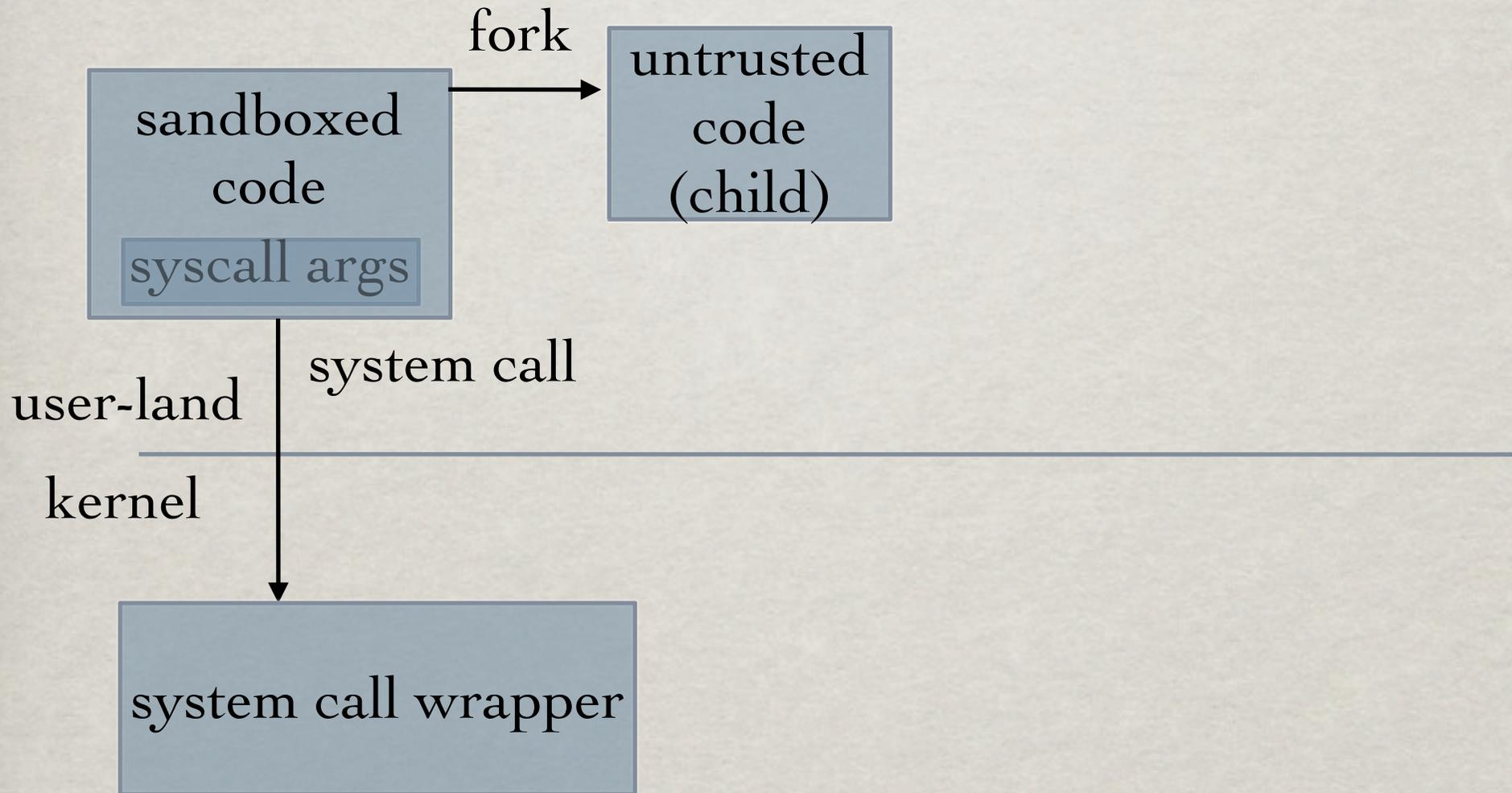
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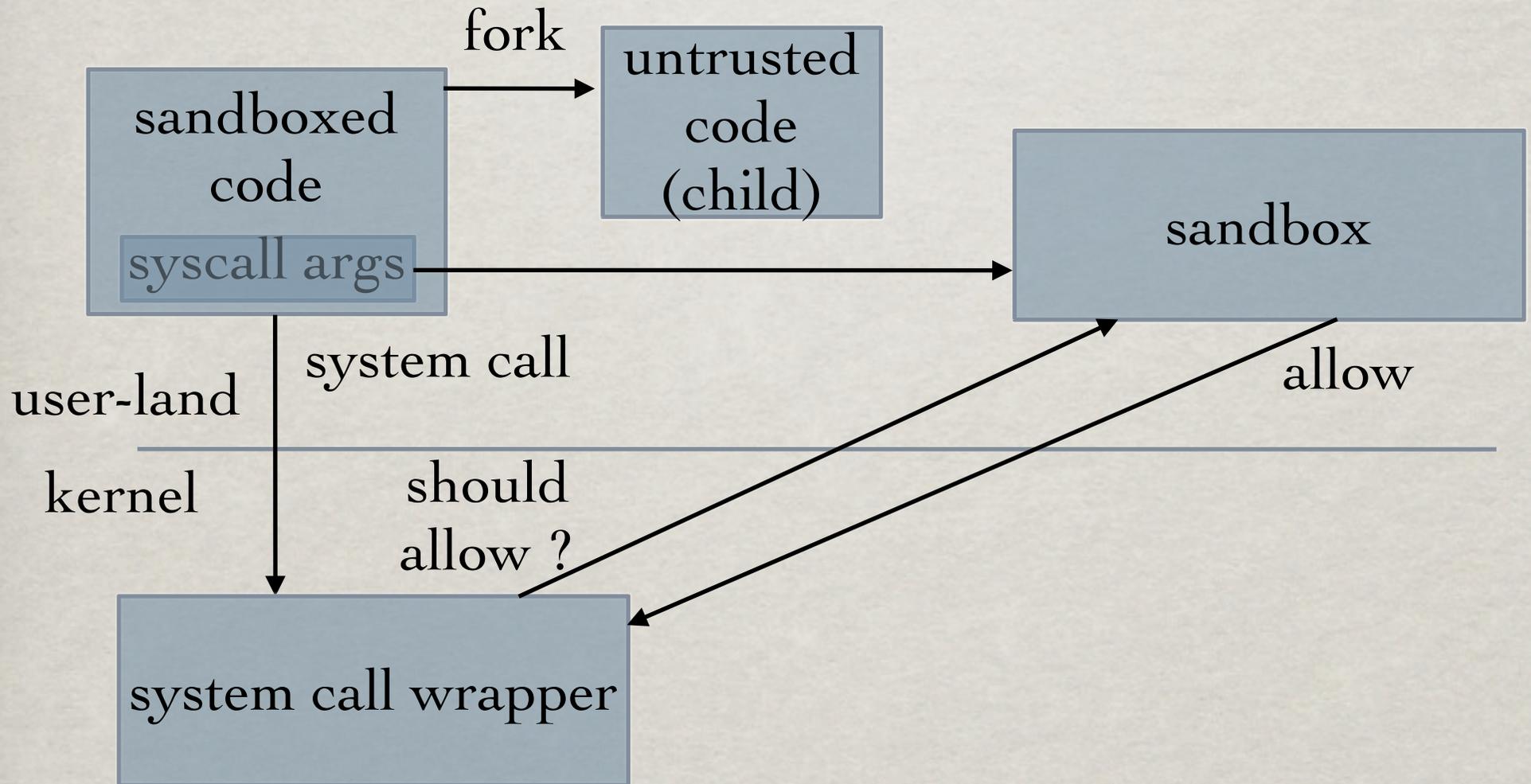
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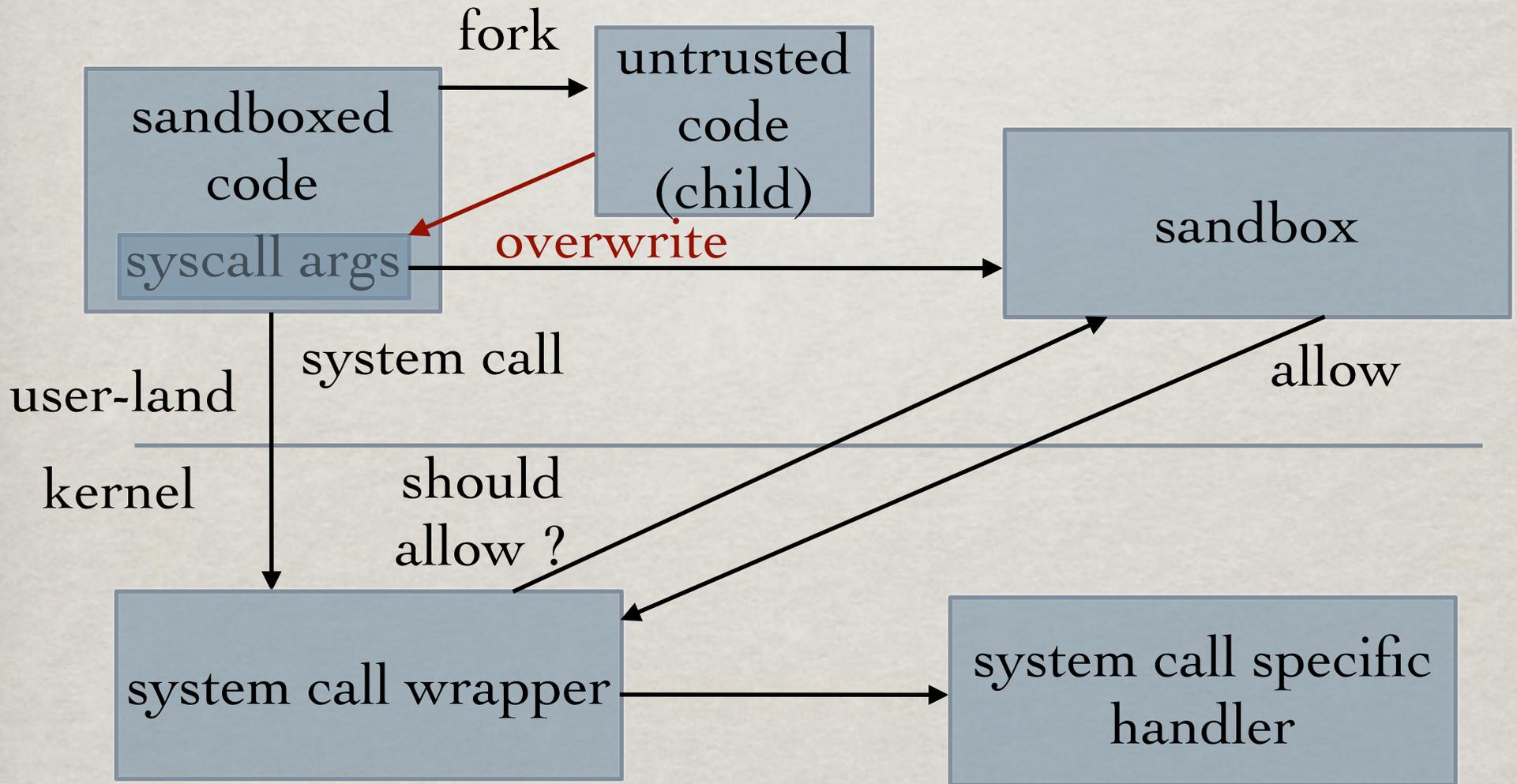
TIME OF CHECK TO TIME OF USE (TOCTOU) ATTACKS



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