

Sketches for Automatic Coding

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Automatic Coding

- Neural Program Induction
- Neural Program Synthesis

I/O Pair Examples

$$\begin{aligned} [2, 3, 4, 5, 6] &\rightarrow [2, 4, 6] \\ [5, 8, 3, 2, 2, 1, 12] &\rightarrow [8, 2, 2, 12] \end{aligned}$$

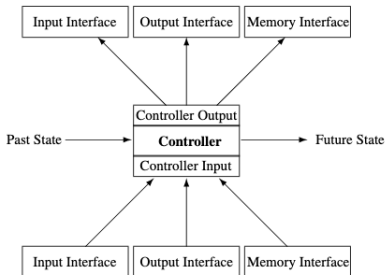
I/O Pair Examples

$[1, 2, 3] \rightarrow X$

$[4, 5, 6] \rightarrow ?$

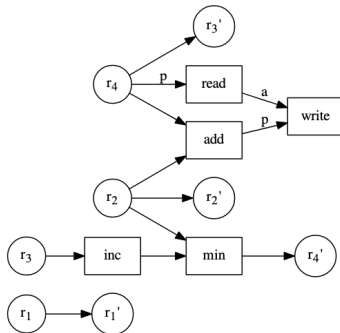
Related Work

Learning Simple Algorithms From Examples (Zaremba et al, 2015)



Related Work

Neural Random Access Machines (Kurach et al, 2015)



Related Work

DeepCoder (Balog et al, 2016)

```
a ← [int]
b ← FILTER (<0) a
c ← MAP (*4) b
d ← SORT c
e ← REVERSE d
```

An input-output example:

Input:

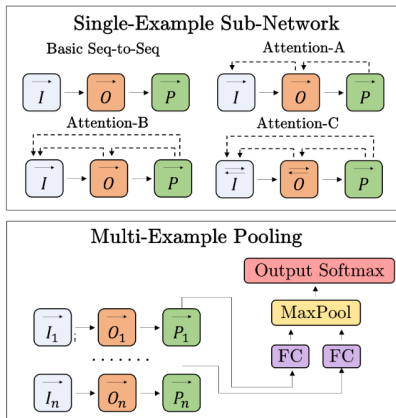
[-17, -3, 4, 11, 0, -5, -9, 13, 6, 6, -8, 11]

Output:

[-12, -20, -32, -36, -68]

Related Work

RobustFill (Devlin et al, 2017)



Motivation for Program Generation

- Implicit programs
- Learning over source code
- Specificity of domain
- Natural language specification

Definitions

- Program Sketch
- Domain Specific Language

Problem Overviews

- Neural Sketch Learning for Conditional Program Generation
- Learning to Infer Program Sketches

Problem Formulation

Learn over program sketches using a probabilistic encoder-decoder, conditioned on labels, to generate source code in AML

Goal

Create a model that can generate source code from some 'spec'

Learn a function g

For test case $(X, Prog)$, $g(X) = Prog'$

Example 1

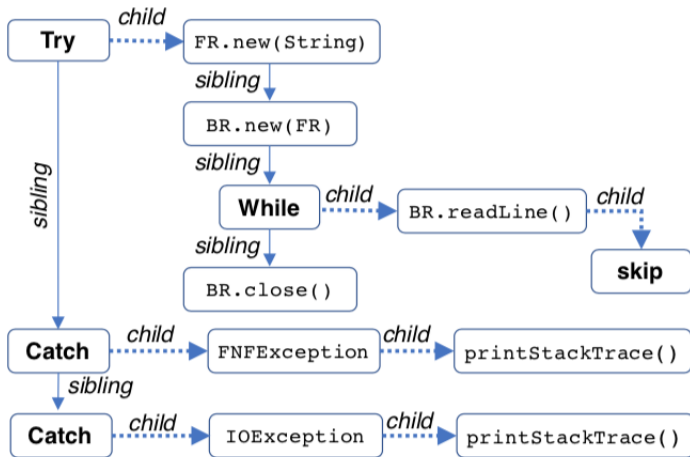
$X_{Types} = \{\text{FileWriter}\}$
 $X_{Calls} = \{\text{write}\}$
 $X_{Keys} = \emptyset$

```
BufferedWriter bw;  
FileWriter fw;  
try {  
    fw = new FileWriter($String, $boolean);  
    bw = new BufferedWriter(fw);  
    bw.write($String);  
    bw.newLine();  
    bw.flush();  
    bw.close();  
} catch (IOException _e) {  
}
```

Example 1a

```
String s;  
BufferedReader br;  
FileReader fr;  
try {  
    fr = new FileReader($String);  
    br = new BufferedReader(fr);  
    while ((s = br.readLine()) != null) {}  
    br.close();  
} catch (FileNotFoundException _e) {  
    _e.printStackTrace();  
} catch (IOException _e) {  
    _e.printStackTrace();  
}
```

Example 1a



Example 2a

Label $X = (X_{Calls}, X_{Types}, X_{Keys})$
 $X = (\{\text{readLine}\}, \emptyset, \emptyset)$

```
String s;
BufferedReader br;
FileReader fr;
try {
    fr = new FileReader($String);
    br = new BufferedReader(fr);
    while ((s = br.readLine()) != null) {}
    br.close();
} catch (FileNotFoundException _e) {}
} catch (IOException _e) {}
}
```

(a)

```
String s;
BufferedReader br;
InputStreamReader isr;
try {
    isr = new InputStreamReader($InputStream);
    br = new BufferedReader(isr);
    while ((s = br.readLine()) != null) {}
} catch (IOException _e) {}
}
```

(b)

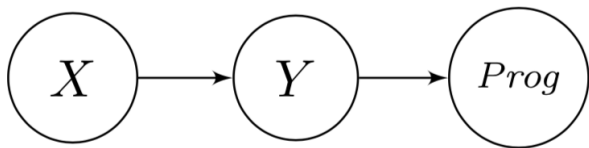
Solution?

$X = (\{\text{readline}\}, \{\text{FileReader}\}, \emptyset)$

Conditional Program Generation

- Functional equivalence
- Maximize the expected value that $g(X)$ and some $Prog$ belong to the same equivalence relation
- $E[I((g(X), Prog) \in Eqv)]$
- BAYOU

Technical Approach



- $P(\text{Prog}|X, \theta)$
- $\theta^* = \arg \max_{\theta} \sum_i \log P(\text{Prog}_i|X_i, \theta)$
- $g(X) = \arg \max_{\text{Prog}} P(\text{Prog}|X, \theta^*)$

Abstraction

- Define abstraction function $\alpha : \mathbb{P} \rightarrow \mathbb{Y}$
- $\text{sat}(Y)$ if $\alpha^{-1}(Y) \neq \emptyset$ aka...
- $P(\text{Prog}|Y) \neq 0 \iff Y = \alpha(\text{Prog})$

Abstraction Function

$\alpha(\text{skip})$	=	skip
$\alpha(\text{call } \text{Sexp}_0.a(\text{Sexp}_1, \dots, \text{Sexp}_k))$	=	call $\tau_0.a(\tau_1, \dots, \tau_k)$ where τ_i is the type of Sexp_i
$\alpha(\text{Prog}_1; \text{Prog}_2)$	=	$\alpha(\text{Prog}_1); \alpha(\text{Prog}_2)$
$\alpha(\text{let } x = \text{Sexp}_0.a(\text{Sexp}_1, \dots, \text{Sexp}_k))$	=	call $\tau_0.a(\tau_1, \dots, \tau_k)$ where τ_i is the type of Sexp_i
$\alpha(\text{if Exp then Prog}_1 \text{ else Prog}_2)$	=	if $\alpha(\text{Exp})$ then $\alpha(\text{Prog}_1)$ else $\alpha(\text{Prog}_2)$
$\alpha(\text{while Exp do Prog})$	=	while $\alpha(\text{Cond})$ do $\alpha(\text{Prog})$
$\alpha(\text{try Prog catch}(x_1) \text{ Prog}_1 \dots \text{catch}(x_k) \text{ Prog}_k)$	=	try $\alpha(\text{Prog})$ catch (τ_1) $\alpha(\text{Prog}_1) \dots$ catch (τ_k) $\alpha(\text{Prog}_k)$ where τ_i is the type of x_i
$\alpha(\text{Exp})$	=	[] if Exp is a constant or variable name
$\alpha(\text{Sexp}_0.a(\text{Sexp}_1, \dots, \text{Sexp}_k))$	=	[$\tau_0.a(\tau_1, \dots, \tau_k)$] where τ_i is the type of Sexp_i
$\alpha(\text{let } x = \text{Call} : \text{Exp}_1)$	=	<i>append</i> ($\alpha(\text{Call}), \alpha(\text{Exp}_1)$)

Grammar for Sketches

$$\begin{aligned}
 Y & ::= \text{skip} \mid \text{call } C_{\text{exp}} \mid Y_1; Y_2 \mid \\
 & \quad \text{if } C_{\text{seq}} \text{ then } Y_1 \text{ else } Y_2 \mid \\
 & \quad \text{while } C_{\text{seq}} \text{ do } Y_1 \mid \text{try } Y_1 \text{ Catch} \\
 C_{\text{exp}} & ::= \tau_0.a(\tau_1, \dots, \tau_k) \\
 C_{\text{seq}} & ::= \text{List of } C_{\text{exp}} \\
 \text{Catch} & ::= \text{catch}(\tau_1) Y_1 \dots \text{catch}(\tau_k) Y_k
 \end{aligned}$$

Encoder-Decoder

$$P(Y|X, \theta) = \int_{Z \in \mathbb{R}^m} P(Z|X, \theta) P(Y|Z, \theta) dZ$$

Encoder

- Convert each label (ex. $X_{Calls,i}$) to one-hot vector representation
- Assume h hidden units
- Define an encoder function, ex:
$$f(X_{Calls,i}) = \tanh((W_h \cdot X'_{Calls,i} + b_h) \cdot W_d + b_d)$$
- $W_h \in \mathbb{R}^{|Calls| \times h}$, $b_h \in \mathbb{R}^h$, $W_d + b_d \in \mathbb{R}^d$

Decoder

- Task: generate sketch Y by sampling from the space of $P(Y|Z)$
- Z is a real vector-valued latent variable
- Start with the root node pair (*root*, *child*)
- Depth first tree exploration

Decoder (cont)

1. (**try**, *c*), (`FR.new(String)`, *s*), (`BR.new(FR)`, *s*), (**while**, *c*), (`BR.readLine()`, *c*), (**skip**, \cdot)
2. (**try**, *c*), (`FR.new(String)`, *s*), (`BR.new(FR)`, *s*), (**while**, *s*), (`BR.close()`, \cdot)
3. (**try**, *s*), (**catch**, *c*), (`FileNotFoundException`, *c*), (`T.printStackTrace()`, \cdot)
4. (**try**, *s*), (**catch**, *s*), (**catch**, *c*), (`IOException`, *c*), (`T.printStackTrace()`, \cdot)

```
String s;
BufferedReader br;
FileReader fr;
try {
    fr = new FileReader($String);
    br = new BufferedReader(fr);
    while ((s = br.readLine()) != null) {}
    br.close();
} catch (FileNotFoundException _e) {
    _e.printStackTrace();
} catch (IOException _e) {
    _e.printStackTrace();
}
```

Concretization

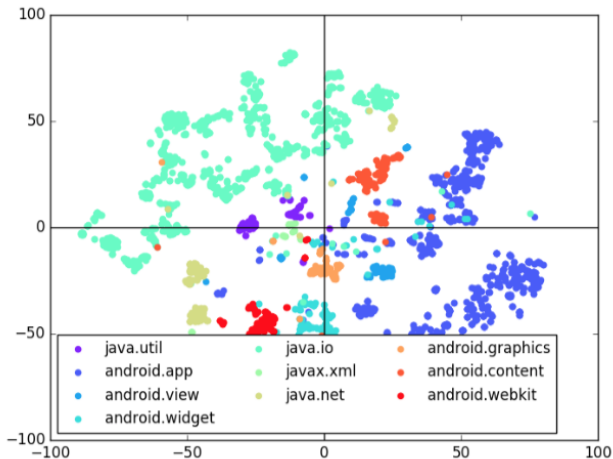
- Type directed, stochastic search
- Given sketch Y , perform random walk of space of *partially concretized sketches*
- Follows distribution of $P(\text{Prog}|Y)$
- Ex. $x_1.a(x_2); \tau_1.b(\tau_2)$
- Defined set of neighbors for each state
- Prioritize simple programs

Experiments

	Min	Max	Median	Vocab
X_{Calls}	1	9	2	2584
X_{Types}	1	15	3	1521
X_{Keys}	2	29	8	993
X	4	48	13	5098

- 1500 Android apps
- 150,000 methods
- Labels defined by heuristic

t-SNE Plot of Latent Space



Accuracy Metrics

- AST Comparison
- Minimum Jaccard Distance between sets of sequences of API calls
- Minimum Jaccard Distance between the sets of API calls
- Minimum absolute difference between number of statements
- Minimum absolute difference between number of control structures

Results

Model	Input Label Observability			
	100%	75%	50%	25%
GED-AML	0.13	0.09	0.07	0.02
GSNN-AML	0.07	0.04	0.03	0.01
GED-Sk	0.59	0.51	0.44	0.21
GSNN-Sk	0.57	0.48	0.41	0.18

(a) M1. Proportion of test programs for which the expected AST appeared in the top-10 results.

Model	Input Label Observability			
	100%	75%	50%	25%
GED-AML	0.52	0.58	0.61	0.77
GSNN-AML	0.59	0.64	0.68	0.83
GED-Sk	0.11	0.17	0.22	0.50
GSNN-Sk	0.13	0.19	0.25	0.52

(c) M3. Average minimum Jaccard distance on the set of API methods called in the test program vs the top-10 results.

Model	Input Label Observability			
	100%	75%	50%	25%
GED-AML	0.31	0.30	0.30	0.34
GSNN-AML	0.32	0.31	0.32	0.39
GED-Sk	0.03	0.03	0.03	0.04
GSNN-Sk	0.03	0.03	0.03	0.03

(e) M5. Average minimum difference between the number of control structures in the test program vs the top-10 results.

Model	Input Label Observability			
	100%	75%	50%	25%
GED-AML	0.82	0.87	0.89	0.97
GSNN-AML	0.88	0.92	0.93	0.98
GED-Sk	0.34	0.43	0.50	0.76
GSNN-Sk	0.36	0.46	0.53	0.78

(b) M2. Average minimum Jaccard distance on the set of sequences of API methods called in the test program vs the top-10 results.

Model	Input Label Observability			
	100%	75%	50%	25%
GED-AML	0.49	0.47	0.46	0.46
GSNN-AML	0.52	0.49	0.49	0.53
GED-Sk	0.05	0.06	0.06	0.09
GSNN-Sk	0.05	0.06	0.06	0.09

(d) M4. Average minimum difference between the number of statements in the test program vs the top-10 results.

Model	Metric				
	M1	M2	M3	M4	M5
GED-AML	0.02	0.97	0.71	0.50	0.37
GSNN-AML	0.01	0.98	0.74	0.51	0.37
GED-Sk	0.23	0.70	0.30	0.08	0.04
GSNN-Sk	0.20	0.74	0.33	0.08	0.04

(f) Metrics for 50% observability evaluated only on unseen data

Learning to Infer Program Sketches

- This paper develops a dynamic system to incorporate pattern recognition and explicit reasoning to solve programming puzzles
- State-of-the-art performance via self-supervised learning

Formulation

- DSL with program space \mathcal{G}
- Set of program specifications (specs) containing I/O examples: $\mathcal{X}_i = \{(x_{ij}, y_{ij})\}_{j=1, \dots, n}$
- We have solved problem \mathcal{X}_i if we find the true program F_i such that

$$\forall j : F_i(x_{ij}) = y_{ij}$$

Formulation

- Can we solve the problem quickly?
- The problem becomes:

$$\max \log \mathbb{P} [\text{Time}(\mathcal{X}_i \rightarrow F_i) < t]$$

System

SketchAdapt

- **Sketch Generator:** Proposes set of possible (incomplete) sketches based on a spec
- **Program Synthesizer:** Takes a sketch as a starting point, then performs explicit search to “fill the holes”

Novel Approach

- Define a more general sketch: a valid program tree where any subtree may be replaced with the special token `<HOLE>`
- This token designates locations in the program tree where pattern recognition is difficult and more explicit search is necessary
- This allows the system to learn how much to rely on each component

Infer Sketches via Self-supervision

- Generator will be parametrized by a RNN, and is trained to assign a high probability to sketches that can be quickly completed
- We can now reframe the program synthesis problem:

$$\max_{\phi} \log \mathbb{P}_{s \sim q_{\phi}(\cdot | \mathcal{X}_i)} [\text{Time}(s \rightarrow F_i) < t]$$

How to set the time budget?

- In order to make the system more robust, train it to output sketches that are suitable for a range of timeout budgets
- Rewrite the previous optimization as:

$$\max_{\phi} \log \mathbb{P}_{\substack{t \sim \mathcal{D}_t \\ s \sim q_{\phi}(\cdot | \mathcal{X}_i)}} [\text{Time}(s \rightarrow F_i) < t]$$

Loss

- Maximize the objective function:

$$obj = \mathbb{E}_{\substack{t \sim \mathcal{D}_t \\ (F, \mathcal{X}) \sim \mathcal{G}}} \log \sum_{s: \text{Time}(s \rightarrow F) < t} q_\phi(s | \mathcal{X})$$

- Quickly solve “easy” problems with concrete sketches, but also sample more general sketches for harder problems

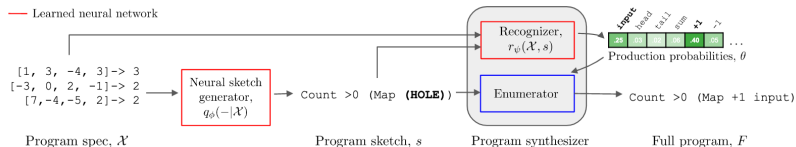
Generator Implementation

- The sketch generator is a sequence-to-sequence RNN with attention
- Spec is encoded via LSTM
- Sketch is decoded token-by-token while attending to the spec

Synthesizer Implementation

- The program synthesizer uses probabilities of primitives appearing in the program in order to induce a PCFG over an incomplete sketch: $p(F|s, \theta)$
- Candidate programs are enumerated in decreasing probability
- The primitive probabilities are provided by a learned recognizer (feed forward MLP ending in softmax)

Architecture



Computing the Loss in Practice

- Note that $\text{Time}(s \rightarrow F) \leq 1/p(F|s, \theta)$
- Bound the objective by

$$\text{obj} \geq \mathbb{E}_{\substack{t \sim \mathcal{D}_t \\ (F, \mathcal{X}) \sim \mathcal{G}}} \log \sum_{s: 1/p(F|s, \theta) < t} q_\phi(s|\mathcal{X})$$

- Because the generator and synthesizer are highly correlated, sketches that maximize $q_\phi(s|\mathcal{X})$ will minimize $p(F|s, \theta)$. So we can use only the dominating term:

$$\text{obj}^* = \mathbb{E}_{\substack{t \sim \mathcal{D}_t \\ (F, \mathcal{X}) \sim \mathcal{G}}} \log q_\phi(s^*|\mathcal{X}) \leq \text{obj}$$

Training

Algorithm 1 SKETCHADAPT Training

Require: Sketch Generator $q_\phi(\text{sketch}|\mathcal{X})$; Recognizer $r_\psi(\mathcal{X}, \text{sketch})$; Enumerator dist. $p(F|\theta, \text{sketch})$, Base Parameters θ_{base}

Train Recognizer, r_ψ :

for F, \mathcal{X} in Dataset (or sampled from DSL) **do**

 Sample $t \sim \mathcal{D}_t$

$\text{sketches}, \text{probs} \leftarrow$ list all possible sketches of F ,
 with probs given by $p(F|s, \theta_{base})$

$\text{sketch} \leftarrow$ sketch with largest prob s.t. prob $< t$.

$\theta \leftarrow r_\psi(\mathcal{X}, \text{sketch})$

 grad. step on ψ to maximize $\log p(F|\theta, \text{sketch})$

end for

Train Sketch Generator, q_ϕ :

for F, \mathcal{X} in Dataset (or sampled from DSL) **do**

 Sample $t \sim \mathcal{D}_t$

$\theta \leftarrow r_\psi(\mathcal{X})$

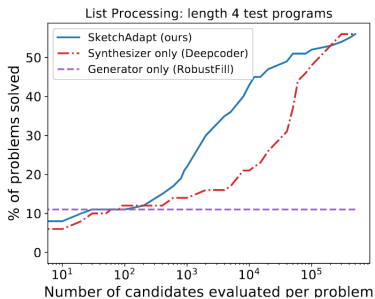
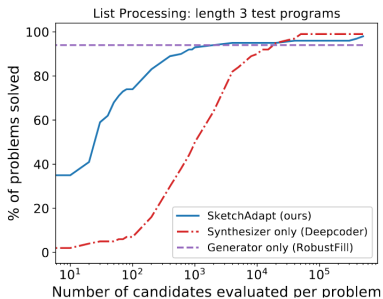
$\text{sketches}, \text{probs} \leftarrow$ list all possible sketches of F ,
 with probs given by $p(F|s, \theta)$

$\text{sketch} \leftarrow$ sketch with largest prob s.t. prob $< t$.

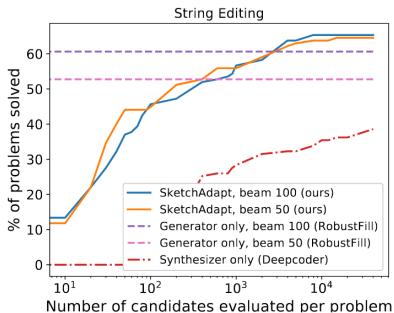
 grad. step on ϕ to maximize $\log q_\phi(\text{sketch}|\mathcal{X})$

end for

Results



Results



Discussion

- Developed a flexible and robust approach that requires processing less data
- No labels required
- Integrates multiple forms of computation (pattern recognition and search)

Conclusions

- Generalizability
- Evaluation
- Flexibility
- Limitations