Basic Elements of Programming Languages

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** These slides are borrowed from Prof. Edwards.
What is a Programming Language?

A programming language is a notation that a person and a computer can both understand.

• It allows you to express what is the **task** to compute
• It allows a computer to **execute** the computation task
Language Specifications
How to Define a Language

When designing a language, it’s a good idea to start by sketching forms that you want to appear in your language as well as forms you do not want to appear.

```c
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

Examples

```c
a int vg(int a,
{
    return (a; + b)

``` Non-Examples
How to Define a Language

- An official documents, with *informal* descriptions.
- An official documents, with *formal* descriptions.
- A reference implementation, e.g., a compiler.

Some language definitions are sanctioned by an official standards organization, e.g., C11 (ISO/IEC 9899:2011).

```c
int compare()
{
    int a[10], b[10];
    if (a > b)
        return true;
    return false;
}
```
Aspects of Language Specifications

• **Syntax**: how characters combine to form a program.
• **Semantics**: what the program *means*.
• **Pragmatics**: common programming idioms; programming environments; the standard library; ecosystems.
Syntax is divided into:

- **Microsyntax**: specifies how the characters in the source code stream are grouped into tokens.
- **Abstract syntax**: specifies how the tokens are grouped into phrases, e.g., expressions, statements, etc.
Source program is just a sequence of characters.

```c
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

```c
int SP avg ( int SP a , SP int SP b ) NL
{
    return ( a SP + SP b ) SP / SP 2 ; NL
}
```
```c
int avg(int a, int b)
{
    return (a + b) / 2;
}
```

<table>
<thead>
<tr>
<th>Token</th>
<th>Lexemes</th>
<th>Pattern (as regular expressions)</th>
</tr>
</thead>
<tbody>
<tr>
<td>ID</td>
<td>avg, a, b</td>
<td>letter followed by letters or digits</td>
</tr>
<tr>
<td>KEYWORD</td>
<td>int, return</td>
<td>letters</td>
</tr>
<tr>
<td>NUMBER</td>
<td>2</td>
<td>digits</td>
</tr>
<tr>
<td>OPERATOR</td>
<td>+, /</td>
<td>+, /</td>
</tr>
<tr>
<td>PUNCTUATION</td>
<td>,,,(,),{,},</td>
<td>,,,(,),{,},</td>
</tr>
</tbody>
</table>

```c
int avg (int a, int b)
{
    return (a + b) / 2;
}
```
int avg(int a, int b)
{
    return (a + b) / 2;
}

• Throw errors when failing to create tokens: malformed numbers (e.g., 23f465#g) or invalid characters (such as non-ASCII characters in C).
Abstract Syntax can be defined using Context Free Grammar.

expr :
    expr OPERATOR expr
| ( expr )
| NUMBER

Expression \((a + b)/2\) can be parsed into an AST:
Abstract Syntax can be defined using **Context Free Grammar**.

```latex
expr : 
    expr OPERATOR expr 
| ( expr ) 
| NUMBER
```

**Ambiguous!** What about $a + b/2$?
Syntax Analysis Gives an Abstract Syntax Tree

```plaintext
int avg(int a, int b) {
    return (a + b) / 2;
}
```

- Syntax analysis will throw errors if “}" is missing. Lexical analysis will not.
• **Static Semantics**: deals with legality rules—things you can check before running the code (compile time), e.g., type, scope, for some languages.

• **Dynamic Semantics**: deals with the execution behavior; things that can only be known at runtime, e.g., value.
We can use inference rules to define semantics, e.g., type:

\[
\begin{align*}
\text{NUMBER} : \text{int} & \quad \text{expr} : \text{int} \\
\text{(expr)} : \text{int} & \\
\text{expr}_1 : \text{int} & \quad \text{expr}_2 : \text{int} \\
\text{expr}_1 \ \text{OPERATOR} \ \text{expr}_2 : \text{int}
\end{align*}
\]
Semantic Analysis: Resolve Symbols; Verify Types

```
symbol table

int a
int b

func int avg args int a int b return / plus.osf a b two.osf one.osf five.osf

symbol table

int a
int b
```
We can use inference rules to define semantics, e.g., value:

\[
\begin{align*}
\text{eval}(\text{NUMBER}) &= \text{NUMBER} \\
\text{eval}(\text{expr}_1) &= n_1 \quad \text{eval}(\text{expr}_2) &= n_2 \quad (n_1 + n_2) = n \\
\text{eval}(\text{expr}_1 + \text{expr}_2) &= n
\end{align*}
\]
Consider the integer range:

\[
\begin{align*}
\text{wrap}(\text{NUMBER}) &= n \\
\text{eval}(\text{NUMBER}) &= n
\end{align*}
\]

\[
\begin{align*}
\text{eval}(\text{expr}) &= n \\
\text{eval}((\text{expr})) &= n
\end{align*}
\]

\[
\begin{align*}
\text{eval}(\text{expr}_1) &= n_1 \\
\text{eval}(\text{expr}_2) &= n_2 \\
\text{wrap}(n_1 + n_2) &= n \\
\text{eval}(\text{expr}_1 + \text{expr}_2) &= n
\end{align*}
\]
Programming Paradigms
A programming paradigm is a style, or “way,” of programming. Some languages make it easy to write in some paradigms but not others.
An imperative program specifies **how** a computation is to be done: a sequence of statements that update state.

```python
result = []
i = 0
numStu = len(students)
start:
    if i >= numStu goto finished
name = students[i]
nameLength = len(name)
    if nameLength <= 5 goto nextOne
addToList(result, name)
nextOne:
    i = i + 1
goto start
finished:
    return result
```
Structured Programming

A kind of imperative programming with clean, goto-free, nested control structures. *Go To Statement Considered Harmful* by Dijkstra.

```python
result = []
for i in range(len(students)):
    name = students[i]
    if len(name) > 5:
        addToList(result, name)
print(result)
```
Imperative programming with **procedure calls**.

```python
def filterList(students):
    result = []
    for name in students:
        if len(name) > 5:
            addToList(result, name)
    return result

print(filterList(students))
```
Object-Oriented Programming

An object-oriented program does its computation with interacting objects.

class Student:
    def __init__(self, name):
        self.name = name
        self.department = "CS"

def filterList(students):
    result = []
    for student in students:
        if student.name.__len__() > 5:
            result.append(student.name)
    return result

print(filterList(students))
A declarative program specifies what computation is to be done. It expresses the logic of a computation without describing its control flow.

```sql
select name
from students
where length(name) > 5
```
A functional program treats computation as the evaluation of mathematical functions and avoids side effects.

```python
def isNameLong(name):
    return len(name) > 5

print(list(filter(isNameLong, students)))
```
Using lambda calculus:

```python
print(
    list(
        filter(lambda name: len(name), students))
)
```
Using function composition:

```python
compose(print, list, filter*(lambda name: len(name) > 5))(students)
```

* A variant of the built-in filter.