



CryptoImg: Privacy Preserving Processing Over Encrypted Images

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Motivation

- Cloud computing provide scalable solution for data storage and processing.
- Emerging solutions for image editing on the cloud: Adobe creative cloud, Pixlr, .., etc.
- Images usually contain privacy sensitive. Outsourcing the raw data exposes a lot of information.
- How to protect user's privacy while editing images in the cloud ?



Video courtesy of Ankita Lathey, P K. Atrey, Image Enhancement in Encrypted Domain over Cloud, ACM TOMM 2015

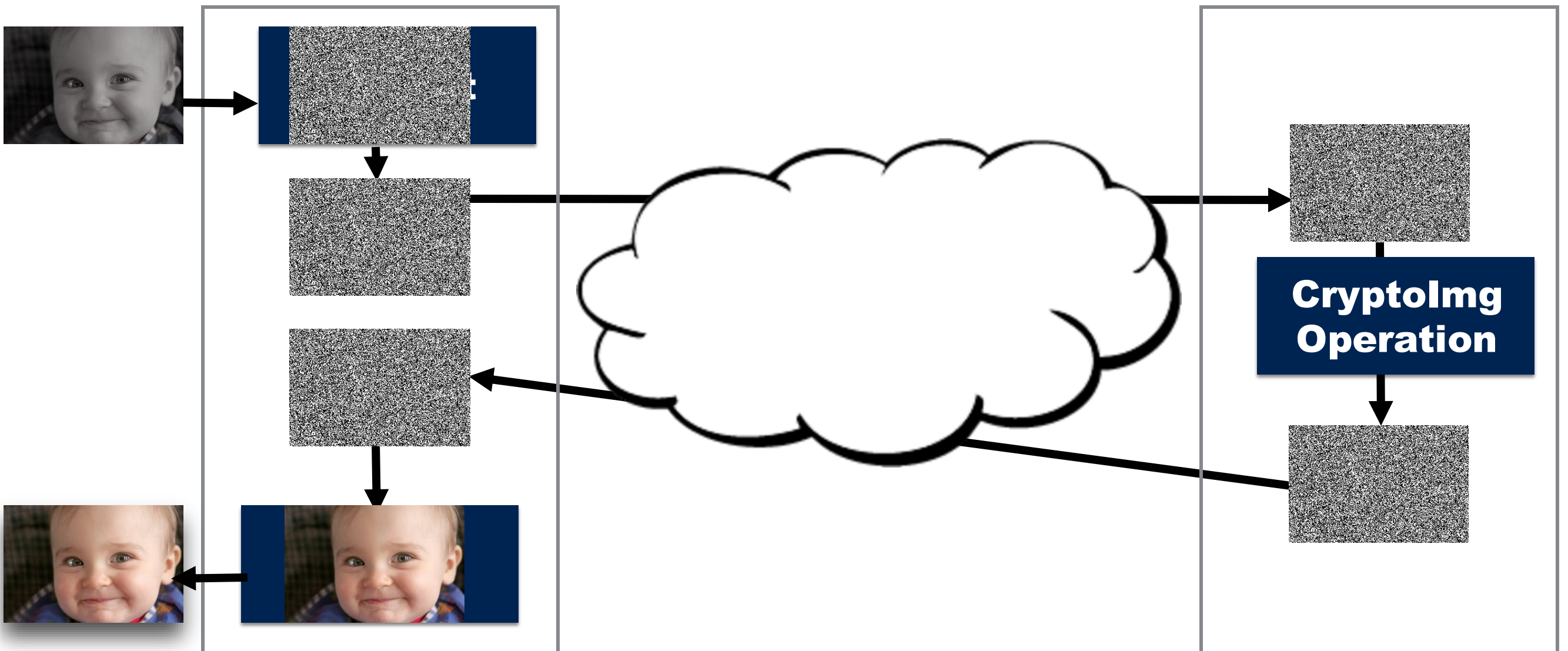
Cryptolmg



Client Device



Cloud Server



Related Work

Work	Description	Comparison to Cryptolmg
Shortell, et al.	Implements brightness / contrast filters using fully homomorphic encryption	<ul style="list-style-type: none">• Cryptolmg supports more operations• Cryptolmg is more computationally efficient.
Lathey, and atrey, et al.	Image enhancement (e.g. spatial filtering, anti-aliasing, edge enhacmenet, etc.) using Shamir Secret sharing	<ul style="list-style-type: none">• Security guarantees require distributing the image processing task across non-colluding servers.
Hu et al.	Secure linear filtering using SMC	<ul style="list-style-type: none">• Cryptolmg is more computationally efficient

Background

- Homomorphic encryption allows carrying out computations on ciphertext.
- **Fully homomorphic encryption:**
 - Performs arbitrary computations.
 - Computationally expensive.
- **Partially Homomorphic encryption:**
 - Supports either addition or multiplication of encrypted values.
- Paillier encryption is an additive homomorphic encryption scheme.

- Supports the addition of two encrypted values.

$$\begin{aligned} \text{DEC}(\llbracket m_1 \rrbracket \oplus_z \llbracket m_2 \rrbracket) &= \text{DEC}((\llbracket m_1 \rrbracket \times \llbracket m_2 \rrbracket) \bmod n^2) \\ &= (m_1 + m_2) \bmod n \end{aligned}$$

- Can multiply an encrypted value by another plain scalar.

$$\begin{aligned} \text{DEC}(\llbracket m_1 \rrbracket \otimes_z d) &= \text{DEC}(\llbracket m_1 \rrbracket^d \bmod n^2) \\ &= (m \times d) \bmod n \end{aligned}$$

Challenges

- Paillier is defined over the group of positive integers \mathbb{Z}^*_n .
- In practice, we also need to deal with negative and real numbers.
- **Solution:**
 - Use an encoding scheme that maps negative and real numbers to integers and preserves the Paillier encryption homomorphic properties.

Challenges

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- **Solution:**
 - Use an encoding scheme that maps negative and real numbers to integers and preserves the Paillier encryption homomorphic properties.
 - To represent negative numbers. Assign different ranges for positive and negative values.
 - $[0, n/3]$: Positive numbers
 - $[2n/3, n]$: Negative numbers.
 - $[n/3, 2n/3]$: Reserved for overflow detection

Challenges

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- In practice, we also need to deal with negative and real numbers.
- **Solution:**
 - Represent each real number by a pair (m, e) where:
 - **m**: mantissa
 - **e**: non-negative exponent
 - To encrypt a real number, it is sufficient to encrypt only the mantissa **m**.

Challenges

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- **Solution:**

Protocol 1 Secure FP Numbers Processing.

Multiplication: $\llbracket c \rrbracket = a \otimes \llbracket b \rrbracket$

$$\llbracket m_c \rrbracket = m_a \otimes_z \llbracket m_b \rrbracket$$

$$e_c = e_a + e_b$$

Addition: $\llbracket c \rrbracket = \llbracket a \rrbracket \oplus \llbracket b \rrbracket$

if $e_a \leq e_b$

$$\llbracket m_c \rrbracket = \llbracket m_a \rrbracket \oplus_z (\text{Base}^{e_b - e_a} \otimes_z \llbracket m_b \rrbracket), e_c = e_a$$

if $e_a > e_b$

$$\llbracket m_c \rrbracket = \llbracket m_b \rrbracket \oplus_z (\text{Base}^{e_a - e_b} \otimes_z \llbracket m_a \rrbracket), e_c = e_b$$

Cryptolmg Operations

- **Cyrptolmg** supports the following image processing operations:
 - Image Adjustment
 - Noise Reduction
 - Edge Detection
 - Morphological Operations
 - Histogram Equalization

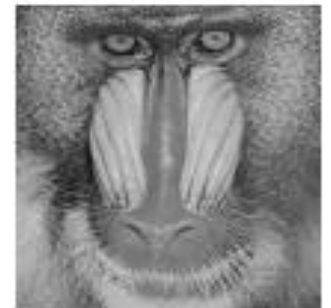
Image Adjustment

- Adding or subtracting adjustment value from each pixel.
- Client sends the encrypted Image $[I]$, and adjustment value v to the server.

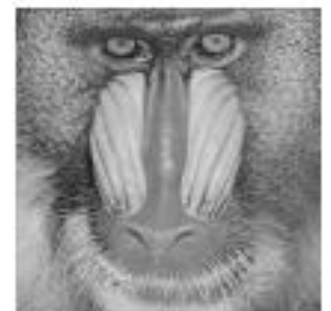
$$[r] = [i] \oplus [v]$$

- Server applies the adjustment to each pixel.
- Client decrypts the result.

Brightness



Input



PD Output

Image Noise Reduction

- Client sends the encrypted Image $[\mathbf{I}]$, and the filter values f to the server.

$$[I_{spt}(u,v)] = \frac{1}{m \times n} \otimes \sum_{u=1,v=1}^{m,n} f(u,v) \otimes [I(u,v)]$$

- Server computes the output image and sends it to the client.
- Client decrypts the result.

LPF



Input



PD Output

Edge Detection

- Client encrypts the source image \mathbf{I} .
- Servers computes the encrypted horizontal and vertical gradients of image

$$[[G_x(u, v)]] = \sum_{u=1, v=1}^{m, n} h_1(u, v) \otimes [[I(u, v)]]$$

$$[[G_y(u, v)]] = \sum_{u=1, v=1}^{m, n} h_2(u, v) \otimes [[I(u, v)]]$$

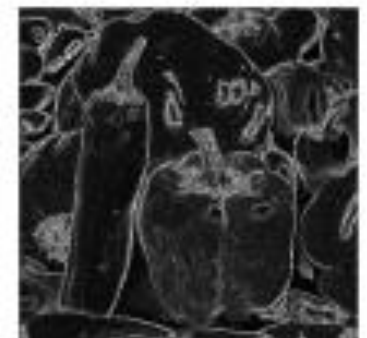
- Client decrypts the result to compute the gradient magnitude and direction

$$G = \sqrt{G_x^2 + G_y^2}$$

$$\Theta = \text{atan2}(G_y, G_x)$$



Input



PD Output

Morphological Operations

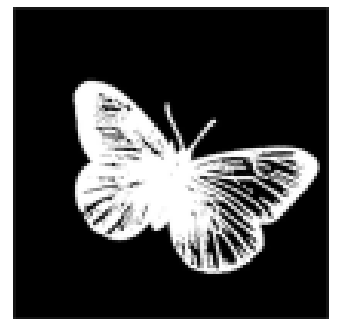
- Client encrypts the source image \mathbf{I} .

- Servers computes, and sends it to client.

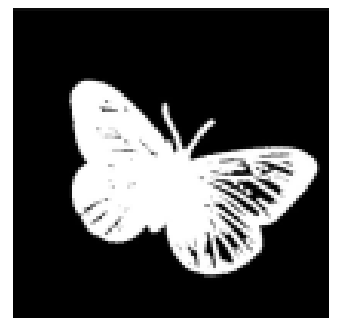
$$[[L(u, v)]] = \sum_{u=1, v=1}^{m, n} [[I(u, v)]]$$

- Client decrypts \mathbf{L} and applies a threshold \mathbf{T} to get the output image.

Dilation



Input



PD Output

Histogram Equalization

- Client computes and encrypts the image histogram $[\mathbf{H}]$.
- Server computes the brightness transformation $[\mathbf{T}(\mathbf{p})]$
$$[[H_c(0)]] = [[H(0)]]$$
$$[[H_c(p)]] = [[H_c(p-1)]] \oplus [[H(p)]], \text{ where } p = 1, 2, \dots, G-1$$
$$[[T(p)]] = (G-1)/(w \times \ell) \otimes [[H_c(p)]].$$
- Server sends $[\mathbf{T}(\mathbf{p})]$ to client.
- Client decrypts $\mathbf{T}(\mathbf{p})$ and applies it to get the output image.

Equalization



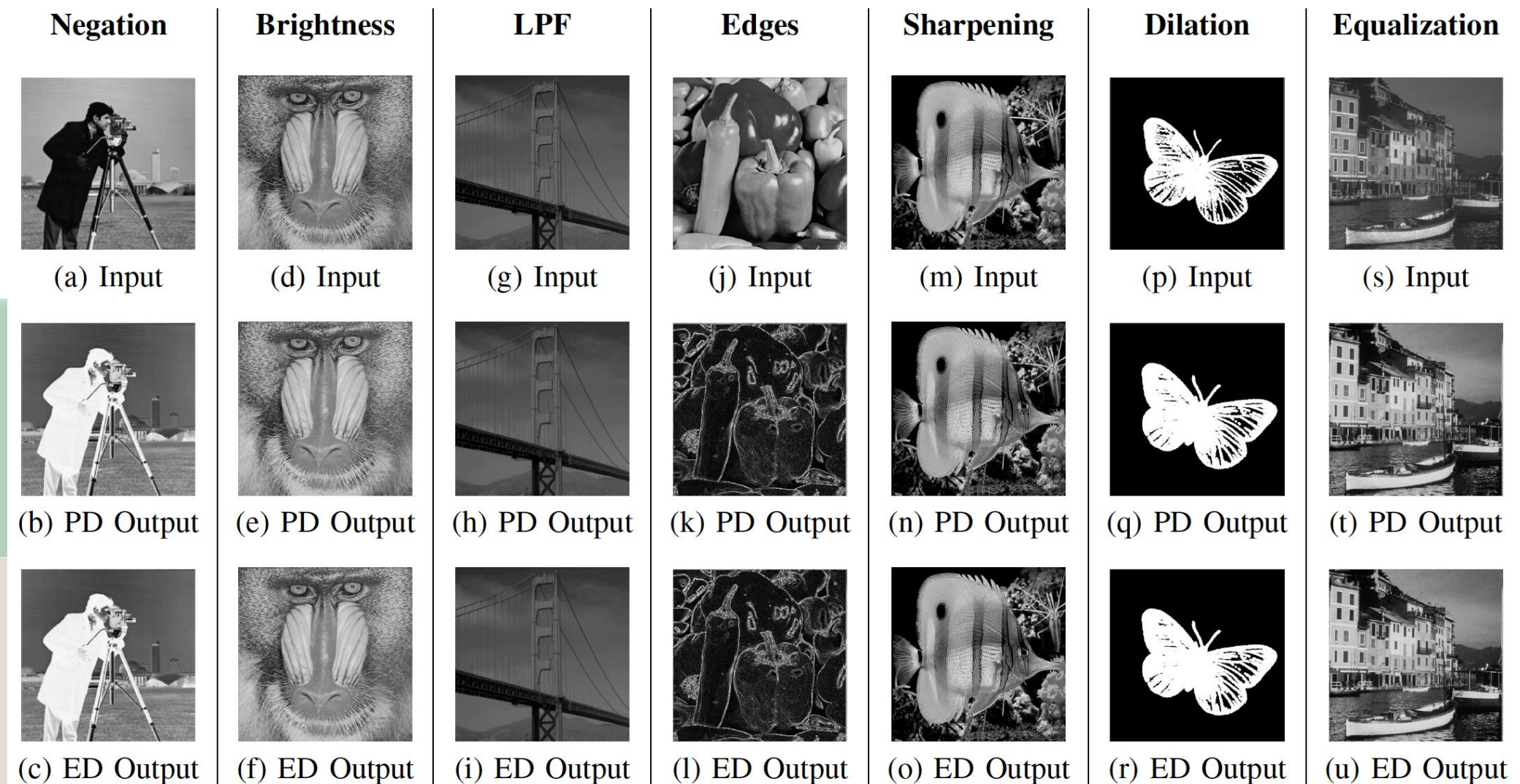
Input



PD Output

Evaluation

- **Cryptolmg** is implemented as an extension for OpenCV library.
- Evaluated using both desktop and mobile device clients.



Evaluation

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Encryption/ Decryption Time (Sec) of 256x256 Image on Mobile Device

	512 bit	1024 bit	2048 bit
Encryption	73	575	3701
Decryption	48	325	2268

Encryption/ Decryption Time (Sec) of 512x512 Image on Desktop Device

	512 bit	1024 bit	2048 bit
Encryption	156.905	1154.29	7670.49
Decryption	1.93	4.068	9.623

Evaluation

- **Cryptolmg** is implemented as an extension for OpenCV library.
- Evaluated using both desktop and mobile device clients.

Time for each operation (Sec.)

	PD	Pre-Proc		Server Time		Post-Proc	
		PC	Mob	1024 bit	2048 bit	PC	Mob
Brightness	0.00108	0	0	0.81	2.39	0	0
LPF	0.00763	0	0	180.50	609.199	0	0
Edge Detection	0.00642	0	0	147.56	482.195	0.0012	0.094
Erosion	0.00009	0	0	4.049	10.808	0.0006	0.0198
Histogram Equalization	0.00174	0.00182	0.177	0.0144	0.048	0.0007	0.029

Future Work

- Speedup the encryption / decryption time of an image using hardware accelerators, e.g. GPUs.

