

Orthogonal tensor decomposition

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Largely based on 2012 arXiv report “Tensor decompositions for learning latent variable models”, with Anandkumar, Ge, Kakade, and Telgarsky.

1. The basic decomposition problem

The basic decomposition problem

Notation: For a vector $\vec{x} = (x_1, x_2, \dots, x_n) \in \mathbb{R}^n$,

$$\vec{x} \otimes \vec{x} \otimes \vec{x}$$

denotes the 3-way array (call it a “tensor”) in $\mathbb{R}^{n \times n \times n}$ whose $(i, j, k)^{\text{th}}$ entry is $x_i x_j x_k$.

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Problem: Given $T \in \mathbb{R}^{n \times n \times n}$ with the promise that

$$T = \sum_{t=1}^n \lambda_t \vec{v}_t \otimes \vec{v}_t \otimes \vec{v}_t$$

for some orthonormal basis $\{\vec{v}_t\}$ of \mathbb{R}^n (w.r.t. standard inner product) and positive scalars $\{\lambda_t > 0\}$, approximately find $\{(\vec{v}_t, \lambda_t)\}$ (up to some desired precision).

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2. If so, is there an efficient algorithm for finding the decomposition?
3. What if T is perturbed by some small amount?

Perturbed problem: Same as the original problem, except instead of T , we are given $T + E$ for some “error tensor” E .

How “large” can E be if we want ε precision?

Analogous matrix problem

Matrix problem: Given $M \in \mathbb{R}^{n \times n}$ with the promise that

$$M = \sum_{t=1}^n \lambda_t \vec{v}_t \vec{v}_t^T$$

for some orthonormal basis $\{\vec{v}_t\}$ of \mathbb{R}^n (w.r.t. standard inner product) and positive scalars $\{\lambda_t > 0\}$, approximately find $\{(\vec{v}_t, \lambda_t)\}$ (up to some desired precision).

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Answer provided by **matrix perturbation theory**

(e.g., Davis-Kahan), which requires $\|E\|_2 < \min_{i \neq j} |\lambda_i - \lambda_j|$.

Back to the original problem

Problem: Given $T \in \mathbb{R}^{n \times n \times n}$ with the promise that

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Where the decompositions do exist, the Perturbed problem asks if they are “robust”.

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- ▶ Catch: We don't recover (\vec{v}_t, λ_t) exactly, so we actually can only replace T with

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- ▶ Therefore, must anyway deal with perturbations.

Rest of this talk

1. Identifiability of decomposition $\{(\vec{v}_t, \lambda_t)\}$ from \mathcal{T} .
2. A decomposition algorithm based on tensor power iteration.
3. Error analysis of decomposition algorithm.

2. Identifiability

Identifiability of the decomposition

Orthonormal basis $\{\vec{v}_t\}$ of \mathbb{R}^n , positive scalars $\{\lambda_t > 0\}$:

$$T = \sum_{t=1}^n \lambda_t \vec{v}_t \otimes \vec{v}_t \otimes \vec{v}_t$$

In what sense is $\{(\vec{v}_t, \lambda_t)\}$ uniquely determined?

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In what sense is $\{(\vec{v}_t, \lambda_t)\}$ uniquely determined?

Claim: $\{\vec{v}_t\}$ are isolated local maximizers of certain cubic form
 $f_T : \mathbb{S}^{n-1} \rightarrow \mathbb{R}^n$, and $f_T(\vec{v}_t) = \lambda_t$.

Aside: multilinear form

There is a natural trilinear form associated with T :

$$(\vec{x}, \vec{y}, \vec{z}) \mapsto \sum_{i,j,k} T_{i,j,k} x_i y_j z_k.$$

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$$(\vec{x}, \vec{y}) \mapsto \sum_{i,j} M_{i,j} x_i y_j = \vec{x}^T M \vec{y}.$$

Review: Rayleigh quotient

Recall Rayleigh quotient for matrix $M := \sum_{t=1}^n \lambda_t \vec{v}_t \vec{v}_t^\top$
(assuming $\vec{x} \in \mathbb{S}^{n-1}$):

$$R_M(\vec{x}) := \vec{x}^\top M \vec{x} = \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x})^2.$$

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Every \vec{v}_t such that $|\lambda_t| = \max!$ is a maximizer of R_M .

(These are also the only local maximizers.)

The natural cubic form

Consider the function $f_T: \mathbb{S}^{n-1} \rightarrow \mathbb{R}^n$ given by

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$$f_T(\vec{x}) = \sum_{t=1}^n \lambda_t \sum_{i,j,k} (\vec{v}_t \otimes \vec{v}_t \otimes \vec{v}_t)_{i,j,k} x_i x_j x_k$$

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Observation: $f_T(\vec{v}_t) = \lambda_t$.

Variational characterization

Claim: Isolated local maximizers of f_T on \mathbb{S}^{n-1} are $\{\vec{v}_t\}$.

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Objective function (with constraint):

$$\vec{x} \mapsto \inf_{\lambda \neq 0} \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x})^3 - 1.5\lambda(\|\vec{x}\|_2^2 - 1).$$

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First-order condition for local maxima:

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Second-order condition for isolated local maxima:

$$\vec{w}^\top \left(2 \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x}) \vec{v}_t \vec{v}_t^\top - \lambda I \right) \vec{w} < 0, \quad \vec{w} \perp \vec{x}.$$

Intuition behind variational characterization

May as well assume \vec{v}_t is t^{th} coordinate basis vector, so

$$\max_{\vec{x} \in \mathbb{R}^n} f_T(\vec{x}) = \sum_{t=1}^n \lambda_t x_t^3 \quad \text{s.t.} \quad \sum_{t=1}^n x_t^2 = 1.$$

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$$f_T(\vec{x}) = \lambda_1 x_1^3 + \lambda_2 x_2^3 < \lambda_1 x_1^2 + \lambda_2 x_2^2 \leq \max\{\lambda_1, \lambda_2\}.$$

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Better to have $|\text{supp}(\vec{x})| = 1$, *i.e.*, picking \vec{x} to be a coordinate basis vector. ■

Aside: canonical polyadic decomposition

Rank- K **canonical polyadic decomposition** (CPD) of T
(also called PARAFAC, CANDECAMP, or CP):

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N.B.: *Overcomplete* ($K > n$) CPD is also interesting *and a possibility* as long as $K(3n + 1) \ll n^3$.

3. Power iteration

The quadratic operator

Easy claim: Repeated application of a certain quadratic operator (based on T) recovers a single (λ_t, \vec{v}_t) up to any desired precision.

For any $A \in \mathbb{R}^{n \times n \times n}$ and $\vec{x} = (x_1, x_2, \dots, x_n) \in \mathbb{R}^n$, define the quadratic operator

$$\phi_A(\vec{x}) := \sum_{i,j,k} A_{i,j,k} x_j x_k \vec{e}_i \in \mathbb{R}^n$$

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$$\text{If } T = \sum_{t=1}^n \lambda_t \vec{v}_t \otimes \vec{v}_t \otimes \vec{v}_t, \text{ then } \phi_T(\vec{x}) = \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x})^2 \vec{v}_t.$$

An algorithm?

Recall: First-order condition for local maxima of

$$f_T(\vec{x}) = \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x})^3 \text{ for } \vec{x} \in \mathbb{S}^{n-1}:$$

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(Ignoring numerical issues, can just repeatedly apply ϕ_T and defer normalization until later.)

N.B.: [Gradient ascent also works](#) [Kolda & Mayo, '11].

Tensor power iteration

Start with some $\vec{x}^{(0)}$, and for $j = 1, 2, \dots$:

$$\vec{x}^{(j)} := \phi_T(\vec{x}^{(j-1)}) = \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x}^{(j-1)})^2 \vec{v}_t.$$

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Claim: For almost all initial $\vec{x}^{(0)}$, the sequence $(\vec{x}^{(j)} / \|\vec{x}^{(j)}\|)_{j=1}^\infty$ converges *quadratically fast* to some \vec{v}_t .

Review: matrix power iteration

Recall matrix power iteration for matrix $M := \sum_{t=1}^n \lambda_t \vec{v}_t \vec{v}_t^\top$:

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If $\lambda_1 > \lambda_2 \geq \dots$, then

$$\frac{(\vec{v}_1^\top \vec{x}^{(j)})^2}{\sum_{t=1}^n (\vec{v}_t^\top \vec{x}^{(j)})^2} \geq 1 - k \left(\frac{\lambda_2}{\lambda_1} \right)^{2j}.$$

i.e., converges *linearly* to \vec{v}_1 (assuming gap $\lambda_2/\lambda_1 < 1$).

Tensor power iteration convergence analysis

Let $c_t := \vec{v}_t^\top \vec{x}^{(0)}$ (initial component in \vec{v}_t direction); assume WLOG

$$\lambda_1 |c_1| > \lambda_2 |c_2| \geq \lambda_3 |c_3| \geq \dots$$

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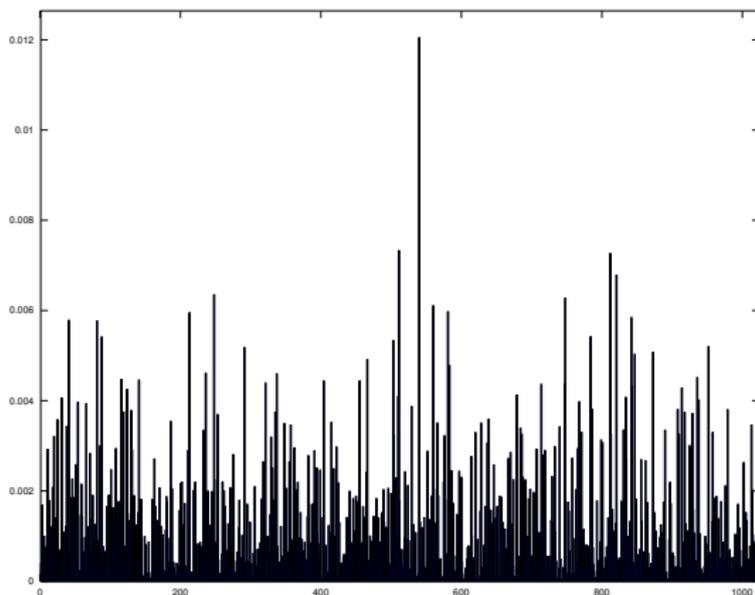
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Easy to show

$$\frac{(\vec{v}_1^\top \vec{x}^{(j)})^2}{\sum_{t=1}^n (\vec{v}_t^\top \vec{x}^{(j)})^2} \geq 1 - k \left(\frac{\lambda_1}{\max_{t \neq 1} \lambda_t} \right)^2 \left| \frac{\lambda_2 c_2}{\lambda_1 c_1} \right|^{2^{j+1}}.$$

Example

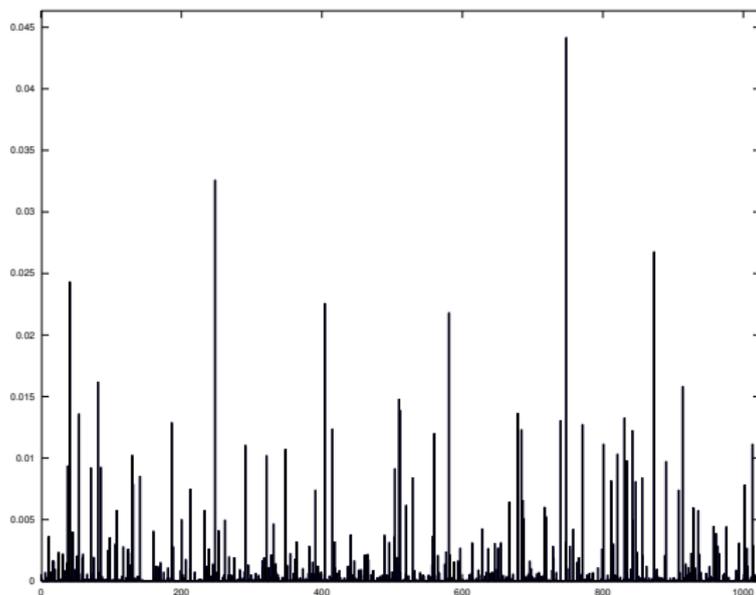
$$n = 1024, \lambda_t \sim_{\text{u.a.r.}} [0, 1].$$



Value of $(\vec{v}_t^\top \vec{x}^{(0)})^2$ for $t = 1, 2, \dots, 1024$

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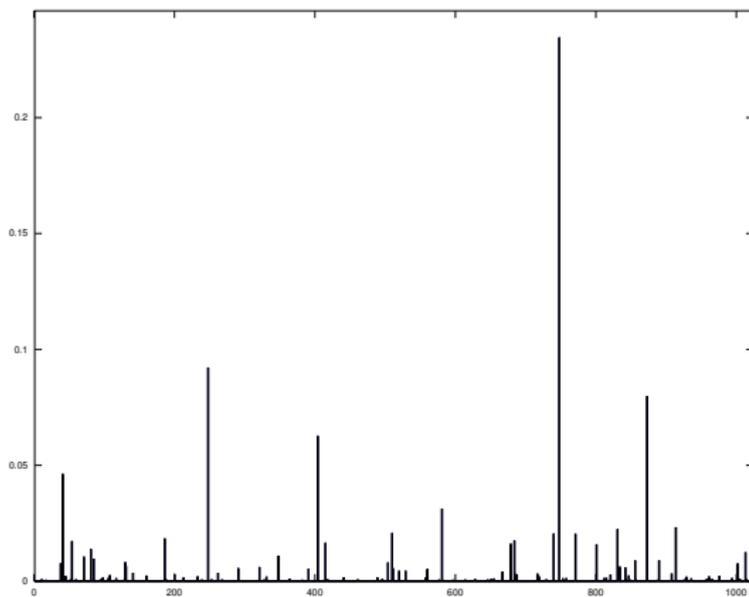
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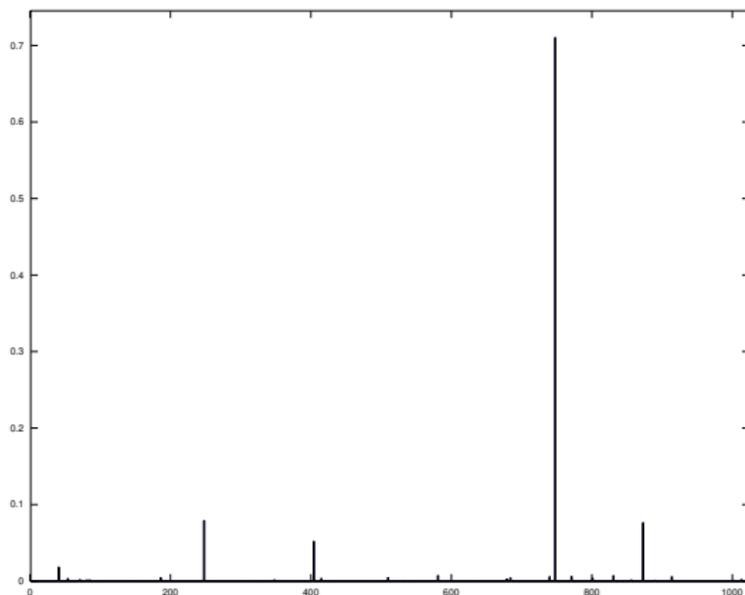
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Value of $(\vec{v}_t^\top \vec{x}^{(2)})^2$ for $t = 1, 2, \dots, 1024$

Example

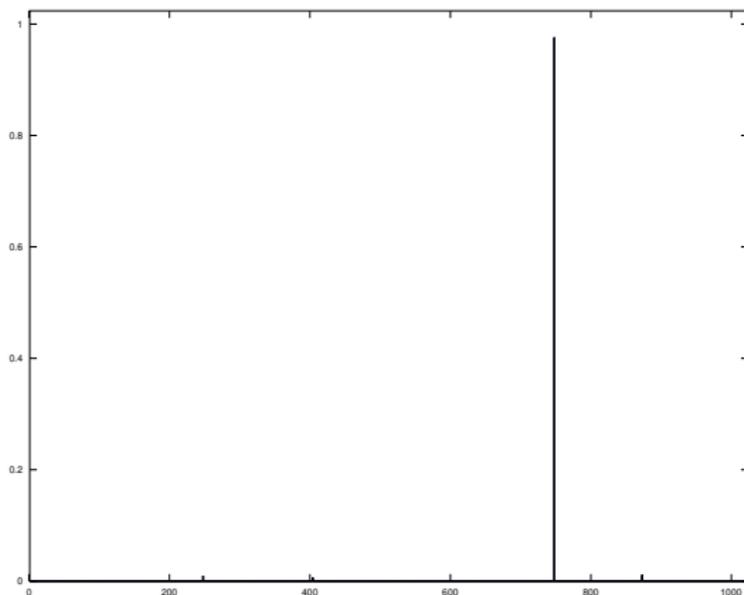
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Value of $(\vec{v}_t^\top \vec{x}^{(3)})^2$ for $t = 1, 2, \dots, 1024$

Example

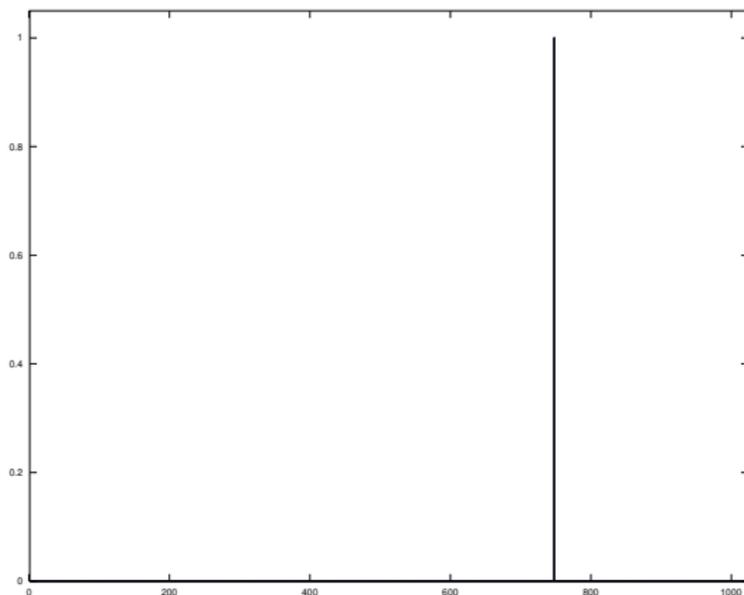
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Value of $(\vec{v}_t^\top \vec{x}^{(4)})^2$ for $t = 1, 2, \dots, 1024$

Example

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Value of $(\vec{v}_t^\top \vec{x}^{(5)})^2$ for $t = 1, 2, \dots, 1024$

Matrix vs. tensor power iteration

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1. Requires gap between largest and second-largest λ_t .
(Property of the matrix **only**.)

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3. **Linear** convergence. (Need $O(\log(1/\epsilon))$ iterations.)

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3. **Quadratic** convergence. (Need $O(\log \log(1/\epsilon))$ iterations.)

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Convergence of tensor power iteration requires **gap** between **largest** and **second-largest** $\lambda_t |\vec{v}_t^\top \vec{x}^{(0)}|$.

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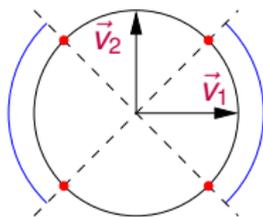
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Fortunately, bad initialization points are atypical.



Full decomposition algorithm

Input: $T \in \mathbb{R}^{n \times n \times n}$.

Initialize: $\tilde{T} := T$.

For $i = 1, 2, \dots, n$:

1. Pick $\vec{x}^{(0)} \in \mathbb{S}^{n-1}$ u.a.r.
2. Run tensor power iteration with \tilde{T} starting from $\vec{x}^{(0)}$ for N iterations.
3. Set $\hat{v}_i := \vec{x}^{(N)} / \|\vec{x}^{(N)}\|$ and $\hat{\lambda}_i := f_{\tilde{T}}(\hat{v}_i)$.
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Output: $\{(\hat{v}_i, \hat{\lambda}_i) : i \in [n]\}$.

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Actually: repeat Steps 1–3 several times, and take results of trial yielding largest $\hat{\lambda}_i$.

Aside: direct minimization

Can also consider directly minimizing

$$\left\| T - \sum_{t=1}^n \hat{\lambda}_t \hat{v}_t \otimes \hat{v}_t \otimes \hat{v}_t \right\|_F^2$$

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Decomposition algorithm via tensor power iteration can be viewed as **orthogonal greedy algorithm** for minimizing above objective [Zhang & Golub, '01].

Aside: implementation for bag-of-words models

Let $\vec{f}^{(i)}$ be empirical word frequency vector for document i :

$$\vec{f}_j^{(i)} = \frac{\text{\# times word } j \text{ appears in document } i}{\text{length of document } i}$$

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$$\widehat{\text{Pairs}} \approx \frac{1}{m} \sum_{i=1}^m \vec{f}^{(i)} \otimes \vec{f}^{(i)} \longrightarrow \sum_{t=1}^K \vec{\mu}_t \otimes \vec{\mu}_t.$$

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Use inner product system given by $\langle \vec{x}, \vec{y} \rangle := \vec{x}^\top \widehat{\text{Pairs}}^\dagger \vec{y}$.

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Final running time $\propto \# \text{ topics} \times (\text{model size} + \text{input size})$.

4. Error analysis

Effect of errors in tensor power iterations

Suppose we are given $\hat{T} := T + E$, with

$$T = \sum_{t=1}^n \lambda_t \vec{v}_t \otimes \vec{v}_t \otimes \vec{v}_t, \quad \varepsilon := \sup_{\vec{x} \in \mathbb{S}^{n-1}} \|\phi_E(\vec{x})\|.$$

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What can we say about the resulting \hat{v}_i and $\hat{\lambda}_i$?

Perturbation analysis

Theorem: If $\varepsilon \leq O(\frac{\min_t \lambda_t}{n})$, then with high probability, a modified variant of the full decomposition algorithm returns $\{(\hat{\mathbf{v}}_i, \hat{\lambda}_i) : i \in [n]\}$ with

$$\|\hat{\mathbf{v}}_i - \vec{\mathbf{v}}_i\| \leq O(\varepsilon/\lambda_i), \quad |\hat{\lambda}_i - \lambda_i| \leq O(\varepsilon), \quad i \in [n].$$

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Essentially third-order analogue of Wedin's theorem for SVD of matrices, but **specific to particular algorithm**.

Effect of errors in tensor power iterations

Quadratic operator $\phi_{\hat{T}}$ with \hat{T} :

$$\phi_{\hat{T}}(\vec{x}) = \sum_{t=1}^n \lambda_t (\vec{v}_t^\top \vec{x})^2 \vec{v}_t + \phi_E(\vec{x}).$$

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Claim: If $\varepsilon \leq O(\frac{\min_t \lambda_t}{n})$ and $N \geq \Omega(\log(n) + \log \log \frac{\max_t \lambda_t}{\varepsilon})$, then N steps of tensor power iteration on $T + E$ (with good initialization) gives

$$\|\hat{v}_j - \vec{v}_j\| \leq O(\varepsilon/\lambda_j), \quad |\hat{\lambda}_j - \lambda_j| \leq O(\varepsilon).$$

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Deflation danger: To find next \vec{v}_t , use

$$\begin{aligned} \hat{T} - \hat{v}_1 \otimes \hat{v}_1 \otimes \hat{v}_1 &= \sum_{t=2}^n \vec{v}_t \otimes \vec{v}_t \otimes \vec{v}_t \\ &\quad + E + \left(\vec{v}_1 \otimes \vec{v}_1 \otimes \vec{v}_1 - \hat{v}_1 \otimes \hat{v}_1 \otimes \hat{v}_1 \right). \end{aligned}$$

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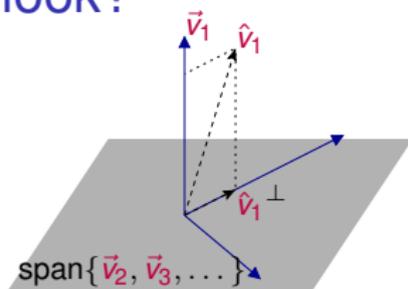
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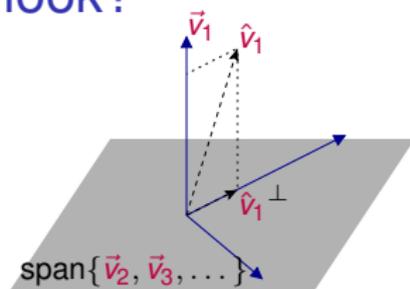
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How do the errors look?



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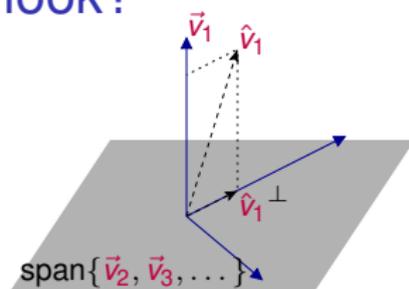
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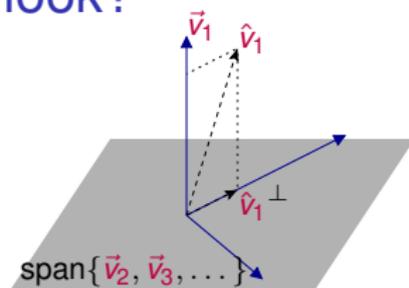


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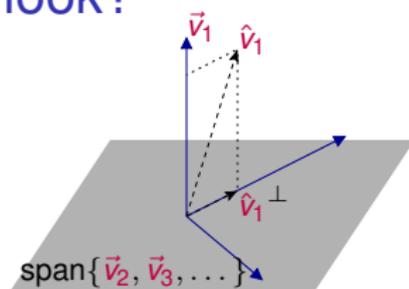


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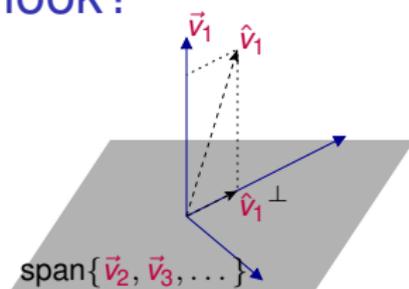


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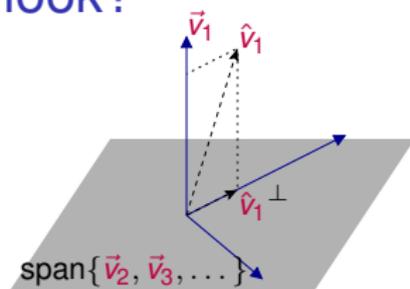


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- ▶ Effect of $E + E_1$ in directions orthogonal to \vec{v}_1 is just $(1 + o(1))\varepsilon$.

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Analogous statement for matrix power iteration is **not true**.

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- ▶ Non-orthogonal (*e.g.*, overcomplete) CP decomposition is active area of research.

Questions?

6. Tensor algebra

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Tensor algebra perspective

From tensor algebra: Since $\{\vec{v}_t : t \in [n]\}$ is a basis for \mathbb{R}^n ,

$\{\vec{v}_i \otimes \vec{v}_j \otimes \vec{v}_k : i, j, k \in [n]\}$ is a *basis* for $\mathbb{R}^n \otimes \mathbb{R}^n \otimes \mathbb{R}^n$

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Every tensor $T \in \mathbb{R}^n \otimes \mathbb{R}^n \otimes \mathbb{R}^n$ has a unique representation in this basis:

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N.B.: $\dim(\mathbb{R}^n \otimes \mathbb{R}^n \otimes \mathbb{R}^n) = n^3$.

Aside: general bases for $\mathbb{R}^n \otimes \mathbb{R}^n \otimes \mathbb{R}^n$

Pick any bases $(\{\vec{\alpha}_i\}, \{\vec{\beta}_i\}, \{\vec{\gamma}_i\})$ for \mathbb{R}^n
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Claim: A tensor T can be diagonal w.r.t. at most one basis.

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Rank- K **canonical polyadic decomposition** (CPD) of T
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N.B.: *Overcomplete* ($K > n$) CPD is also interesting *and a possibility* as long as $K(3n + 1) \ll n^3$.

7. Initialization

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Try $O(n^{1.3})$ initializers; chances are at least one is good.
(Very conservative estimate only; can be *much* better than this.)