

Last time:

- Introductions (you, me, topic)
- Admin. overview (web page)
- High level course overview (learning models)
- Some technical material: Key basic notions + terminology

Rocco OH:

Fri 1:30-

3:30

5th floor
CSB

- $X = \text{domain}$, $\mathcal{C} = \text{concept class} = \text{collection of concepts}$ (concept = subset of X)
- There's a known \mathcal{C} , + an unknown target concept $c \in \mathcal{C}$.
 - Learner has some source of info. about how c labels examples $x \in X$. Want to find/approximate c .

Today:

- some more examples of \mathcal{C} 's (DNFs, LTFs)
- FIRST LEARNING MODEL:

Online Mistake-Bound Learning (OLMB)

- Elimination alg. for monotone disjunctions (+ more)
- Decision lists, OLMB alg. for learning them

PS1: coming out
on Tues

Readings: see web
page

Questions?

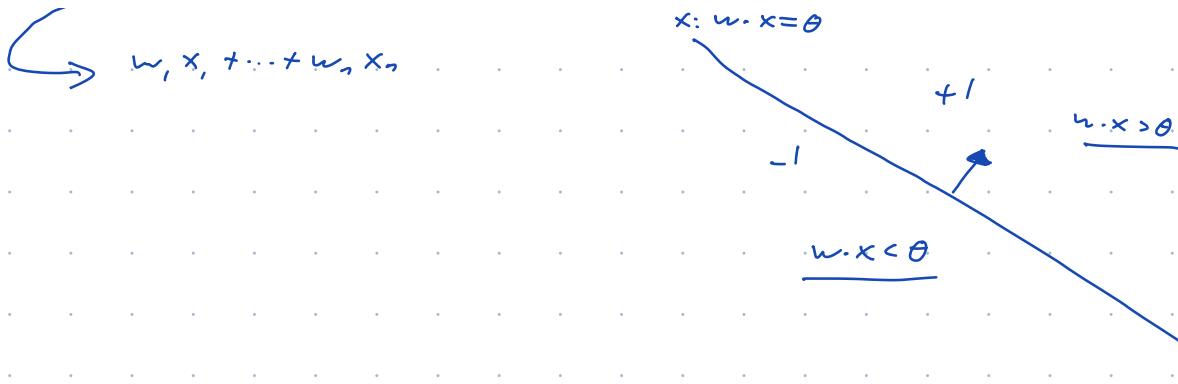
Last example:

LTFs
Linear Threshold Fns

- $X = \mathbb{R}^n$ $\mathcal{C} = \text{all } \underline{\text{halfspaces}}$
- ↳ or $\{\mathbb{R}^n\}$

A Bool fn $c: \mathbb{R}^n \rightarrow \{-1, 1\}$ is an LTF if
 \exists a weight vector $w \in \mathbb{R}^n$, threshold $\theta \in \mathbb{R}$ s.t.

$$c(x) = \text{sign}(w \cdot x - \theta), \quad \text{sign}(t) = \begin{cases} 1 & \text{if } t \geq 0 \\ -1 & \text{if } t < 0 \end{cases}$$



Ex: $n=4, c(x) = \text{sign}(-100x_1 + 2.7x_2 + 0.03x_3 + 1000x_4 - 7)$

First learning model: (OLMB)

Online Mistake-Bound Learning

"Learning session" consists of a seq. of trials.

Learner always has a hypothesis $h: X \rightarrow \{0, 1\}$

Know \mathcal{C} , don't know target concept $c \in \mathcal{C}$.

Each trial:

- Learner given unlabeled $x \in X$.
- Learner outputs $h(x) \in \{0, 1\}$.
- Learner given ^{then} true $c(x) \in \{0, 1\}$.

If $h(x) \neq c(x)$, learner is charged a mistake.



Learner can update h before next trial.

(this
is the
learning alg,
along w/ hinit.)

Performance on a seq. of trials: # mistakes made.

Def: A learning alg. A has mistake bd M for \mathcal{C} if, for any seq. of ex. from X , any target $c \in \mathcal{C}$, A makes $\leq M$ mistakes. (of any length!)

- No noise
- Worst case;
- no data distrib.

Obs #1: • If X finite ($X = \{x^1, x^2, \dots, x^N\}$)
can achieve mist. bd. $M \leq N$.

• If \mathcal{C} is finite ($\mathcal{C} = \{c^1, \dots, c^L\}$)
 $M \leq L - 1$.

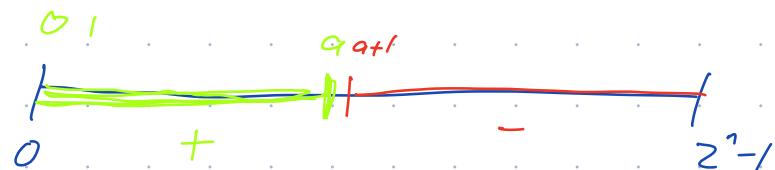
(try each concept in order)

Can often do better:

Ex: $X = \{0, 1, 2, \dots, 2^n - 1\}$

$\mathcal{C} = \text{class of initial intervals}$

$\mathcal{C} = \{\{0, 1, \dots, a\} : a \leq 2^n - 1\}$



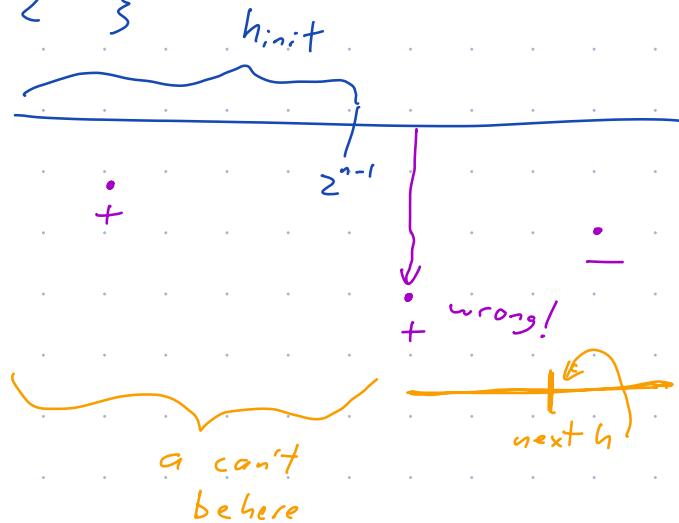
$$|X| = 2^n, |e| = 2^n$$

Better alg.: binary search.

Mist. bd.

$M \leq n$

$$h_{\text{init}} = \{0, 1, \dots, 2^{n-1}\}$$



Each mistake lets us halve the interval containing a , so after $\log_2(2^n)$ mistakes, know a exactly, & no more mistakes.

Ex: $X = \{0, 1\}$ (cont. interval):

\mathcal{C} = init. intervals: no finite
mst. bound.

$$c = [0, 0.237185620014\dots]$$

Some OCLMB Algorithms

① Learning Mon. Disjunctions

$$X = \{0, 1\}^n, \mathcal{C} = \text{all mon. disj.}$$

$$c(x) = x_5 \vee x_8 \vee x_9 \vee x_{10}$$

"Elim. alg."

- $h_{\text{init}} = x_1 \vee x_2 \vee \dots \vee x_n$ (all n vars)
- On false pos mistake ($h(x) = 1$ but $c(x) = 0$):
remove from h all x_i 's s.t. $x_i = 1$ in that example
 \equiv \equiv (elimination)
- On false neg mistake ($h(x) = 0, c(x) = 1$):
FAIL.

Ex: $n = 5, h_{\text{init}}(x) = x_1 \vee x_2 \vee x_3 \vee x_4 \vee x_5$

Get $x = 10010: h_{\text{init}}(x) = 1$

Told $c(x) = 0$: new $h(x) = x_2 \vee x_3 \vee x_5$.

Claim A: No var that's in c is ever removed from h .

Pf: x_i only rem. from h if $x_i = 1, c(x) = 0$;
but means x_i can't be in c . \blacksquare

Claim B: Alg only makes false pos on $c = \text{any}$
 mon. disj. (don't FAIL)

Pf By Claim A, h always includes all c 's vars.
So if $c(x) = 1$, also $h(x) = 1$. \blacksquare

Thm: Elim alg. has mistake bd $\leq n$ for
 $c = \text{mon. disj.}$

Pf: By B, don't fail.

By A, $(\text{vars in } h) \supseteq (\text{vars in } c)$ always.

Alg starts w/ n vars in h ;
each mistake removes ≥ 1 var from h ,
so $\leq n$ mistakes.

Notes: $n = \text{"size"};$

$\text{poly}(n)$ time / trial + }
" mistake bd } 

Noise would been fatal/...

- Can extend to learn arb. disj. (non-mon): run it over 2^n "meta-variables" $x_1, \dots, x_n, \bar{x}_1, \dots, \bar{x}_n$
- $n=2:$

$$h_{\text{init}}(x_1, x_2) = x_1 \vee x_2 \vee \bar{x}_1 \vee \bar{x}_2$$

must be $h_{\text{init}} = 1$,
 $c=0$

mist. bd. = $n+1$ (first mistake: cross off n of 2^n literals)

- Can extend to learn (non) conjunctions.
- negate each $c(x)$, + learn the disj.

$$c(x) = x_1 \wedge \bar{x}_2 \wedge x_3$$



$$\overline{c(x)} = \bar{x}_1 \vee x_2 \vee \bar{x}_3$$

② Learning Decision Lists

$$b_i = b_i \text{ if } 0/1$$

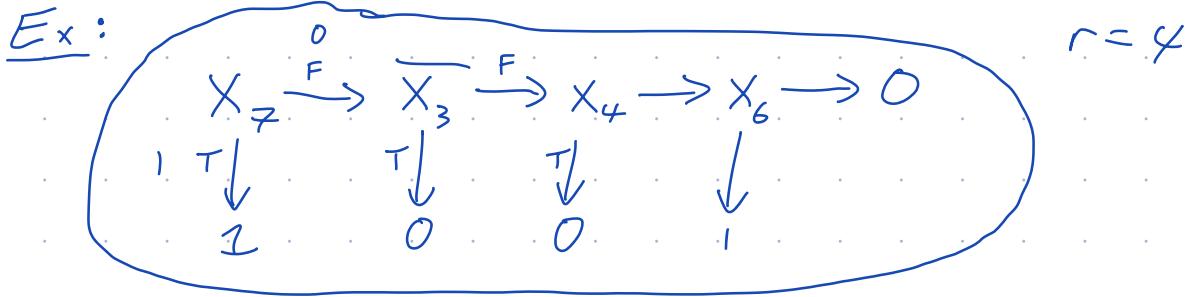
$$l_i = \text{literal}$$

A decision list is a fn $\{0, 1\}^r \rightarrow \{0, 1\}$ of the form

"if l_1 , then output b_1 ,
else if l_2)) b_2
⋮
else if l_r)) b_r ←
else output b_{r+1} "

$$r = \#$$

Ordered list
of rules

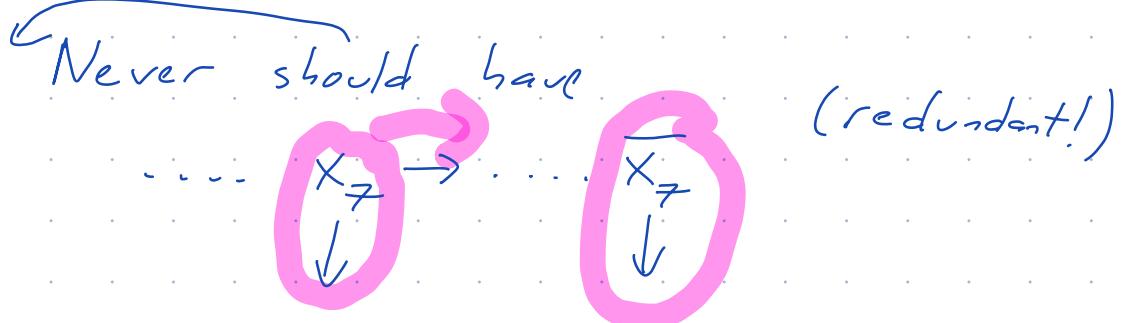


"if l then output b " : a rule.

$$\frac{4n}{\text{var. choice}} + 2 \quad \begin{matrix} \text{poss.} \\ \text{rules} \end{matrix}$$

\pm literal
 \pm bit

Length- r DL: $\approx (4n)^r$ poss.



Think of every DL over {0,1} as having length $\leq n$.

Every conj can be written as

a DL ... $x_1 \wedge \bar{x}_2 \wedge x_3$ is \equiv to $\bar{x}_1 \rightarrow x_2 \rightarrow \bar{x}_3 \rightarrow 1$

Next time: an efficient OLMB alg. for DLs...
