

Last time:

- introductions (me, you, topic)
- administrative overview (web page)
- **high level overview of course** → learning models
- some technical material: key basic notions & terminology

Learner knows  $X, \mathcal{C}$ , goal is to learn unknown target concept  $c \in \mathcal{C}$

$X =$  instance space  $x \in X$   
 $c: X \rightarrow \{0,1\}$  concept  $c \in \mathcal{C}$   
 $\mathcal{C} =$  concept class = set of concepts

Today:

- a few more examples of  $\mathcal{C}$ 's (DNFs, LTFs)

• FIRST LEARNING MODEL:

Online Mistake-Bound Learning

specific algs. for specific  $\mathcal{C}$

- Elimination alg. for monotone disjunctions, extensions
- Decision lists, OCMB alg. for learning them

Admin:

- PS 1 out on Wed

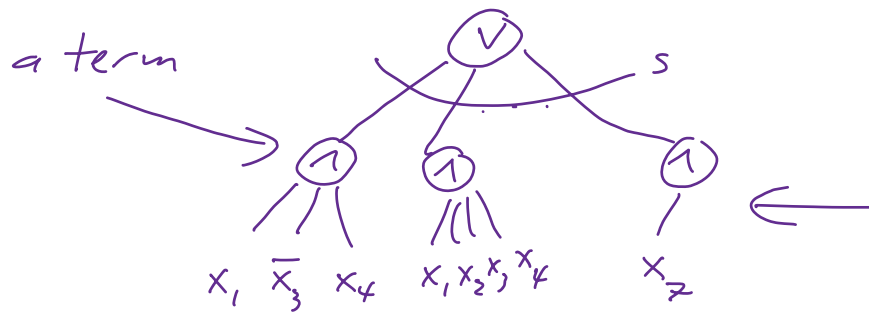
Readings: web page

Questions?

Some more  $\mathcal{C}$ 's of interest:

$X = \{0,1\}^n$ ,  $\mathcal{C} =$  all s-term DNF formulas  
s-clause Conjunctive Disjunctive Normal Form CNF

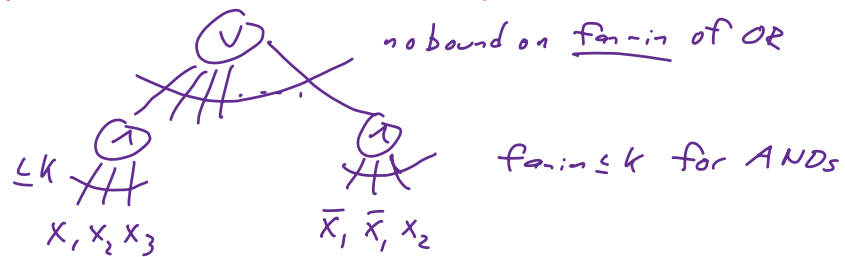
s-term DNF:  $\swarrow$  OR disj. of AND conj of literals.  
s-way AND ORs



K-CNFs

$$X = \{0,1\}^n, \mathcal{C} = \text{K-DNFs}$$

K-DNF: OR of K-literal ANDs.  
 K-CNF AND ORs



halfspace

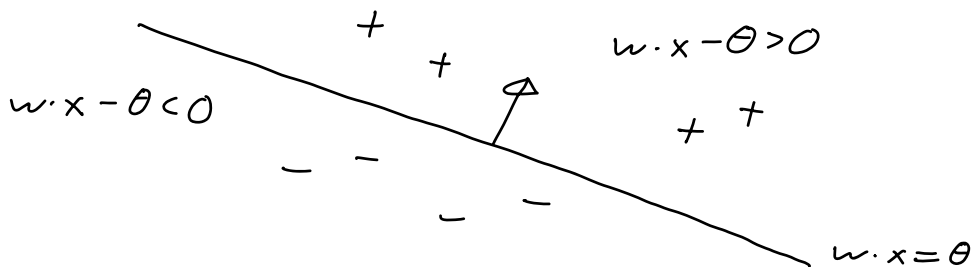
(LTF)

$$X = \mathbb{R}^n, \mathcal{C} = \text{linear threshold functions}$$

A function  $c: \mathbb{R}^n \rightarrow \{-1,1\}$  is an LTF if  
 if there are  $w_1, \dots, w_n, \theta \in \mathbb{R}$  s.t.

$$c(x) = \text{sign}(w \cdot x - \theta)$$

$$\text{sign}(t) = \begin{cases} 1 & t \geq 0 \\ -1 & t < 0 \end{cases}$$



$\mathbb{1}[w \cdot x \geq \theta] \in \{0,1\}$   
 $\mathcal{C}$  indicator function

ex:

$$c(x) = \text{sign}(100x_1 - 7x_2 + 8x_3 - 16)$$

Can also consider LTFs over  $X = \{0, 1\}^3$ .

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First learning model:

## ONLINE MISTAKE-BOUND

### LEARNING MODEL

A learning session consist of a seq. of trials.  
Throughout, learner maintains some hypothesis  
 $h: X \rightarrow \{0, 1\}$ .

Learning an unknown  $c \in \mathcal{C}$  goes like this.  
Each trial:

- Learner is given an unlabeled  $x \in X$ .
- Learner outputs  $h(x) \in \{0, 1\}$
- ⊛ • Learner is given  $c(x) \in \{0, 1\}$ .  
If  $c(x) \neq h(x)$ , learner is charged a mistake.
- ⊛ • Learner can update  $h$  before next trial.  
↳ this update rule, & initial hyp, constitute the learning alg.

Perf. measure: worst-case tot. # mistakes.

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Def: <sup>or</sup> Learning alg  $A$  has mistake bound  $M$  for  $\mathcal{C}$  if for any sequence of examples from  $X$ , any  $c \in \mathcal{C}$ ,  $A$  makes  $\leq M$  mistakes when learning  $c$ .  
(could be arbitrarily long!)

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Observe:

- If  $X$  finite, for any  $\mathcal{C}$  over  $X$ , can achieve m.b.  $\leq |X|$ . (memorization)

- If  $\mathcal{C}$  finite, can achieve m.b.  $|\mathcal{C}|-1$  (try each concept till it causes a mistake).

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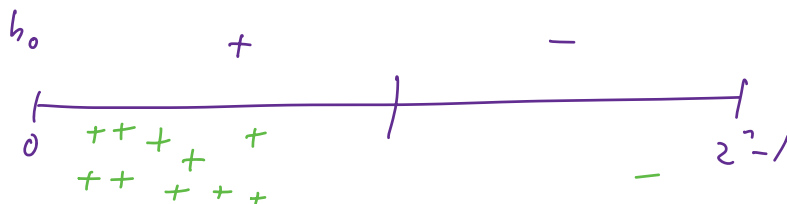
Let's do better:

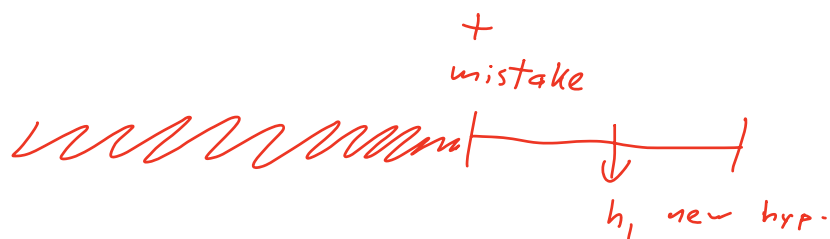
EX  $X = \{0, 1, \dots, 2^n - 1\}$

$\mathcal{C} =$  initial intervals  $c = \{0, 1, 2, \dots, a\}$

Alg. w/ m.b.  $n$ : binary search

$h_0$  init. hyp:  $\{0, \dots, 2^{n-1}\}$  (midpoint of "uncertainty region")





$h$  have uncertainty at each mistake  $\Rightarrow \leq n$  mistakes.

$X = [0, 1]$  continuous,  $\mathcal{C} =$  initial intervals  
 $c = [0, 0.371428556001\dots]$

(no finite mistake bound...)

### Learning monotone disjunctions

$X = \{0, 1\}^n$ ,  $\mathcal{C} =$  mon. disj.  $x_1 \vee x_5 \vee x_6 \vee x_7$

Algorithm ("elimination alg."):

- Initial hyp is  $h(x) = x_1 \vee \dots \vee x_n$
- Get  $z \in \{0, 1\}^n$  as example, predict  $h(z)$ , given  $c(z)$
- If  $h(z) = 1$ ,  $c(z) = 0$  (false positive): ~~⊖~~  
 remove  $x_i$  from  $h$  for all  $i$  s.t.  $z_i = 1$ .
- If  $h(z) = 0$ ,  $c(z) = 1$ , stop + FAIL (won't happen!)
- If  $h(z) = c(z)$ , keep  $h$  the same.

Ex  $n = 5$   $h_0 = x_1 \vee x_2 \vee x_3 \vee x_4 \vee x_5$

Get  $z = 01001$

$h(z) = 1$ . Suppose  $c(z) = 0$ .

Remove  $x_2, x_5$ : new  $h(x)$  is  $x_1 \vee x_3 \vee x_4$ .

MB  $n$ .

Claim 1: if  $x_i$  is in  $c$ ,  $x_i$  never removed from  $h$ .

Pf: only way  $x_i$  removed is if ex  $z$  has  $c(z) = 0$ ,  
 $z_i = 1$ ; can't happen if  $x_i$  is in  $c$ .  $\blacksquare$

Claim 2: alg only makes false pos. mistakes (won't FAIL).

Pf: by prev claim,  $h$  always includes all vars in  $c$ .  
So if  $c(z) = 1$  b/c of  $z_i = 1$ ,  $z_i$  also causes  $h(z) = 1$ .  
So no false neg.  $\blacksquare$

Thm: Elim alg. has mistake bound  $n$  for  $\mathcal{C} = \text{mon disj}$ .

Pf: Doesn't FAIL.

Vars in  $h \supseteq$  Vars in  $c$ .

Every mistake removes  $\geq 1$  var from  $h$ .

$h$  initially  $n$  variables, never has  $< 0$  vars;

So # mistakes can't be greater than  $n$ .

Note: •  $n = \text{"size" of input } x \in \{0,1\}^n$ .  
can describe any  $c \in \mathcal{C}$  with  $n$  bits.

- Alg is efficient  $\rightarrow n$  mistakes  
 $O(n)$  time per trial
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- Noise? OLMB Model: no noise.

This alg would be very susceptible to noise: elimination is irrevocable.

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- Extension to alg for  $\mathcal{C} = \text{all disj}$ :

initial  $h = \cancel{x_1} \vee \bar{x}_1 \vee x_2 \vee \cancel{\bar{x}_2} \vee \dots \vee x_n \vee \bar{x}_n$

First mistake  $z = 10100$ :  $c(z) = 0$ ,  
 $h(z) = 1$ :

updated  $h = \bar{x}_1 \vee x_2 \vee \bar{x}_3 \vee x_4 \vee x_5$

After 1 mistake: learning mon. disj. "in disguise"

$n+1$  for .

- Easy ext. to  $\mathcal{C} = \text{mon. conj.}$ , general conj.  
(think about it...)
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# Decision Lists

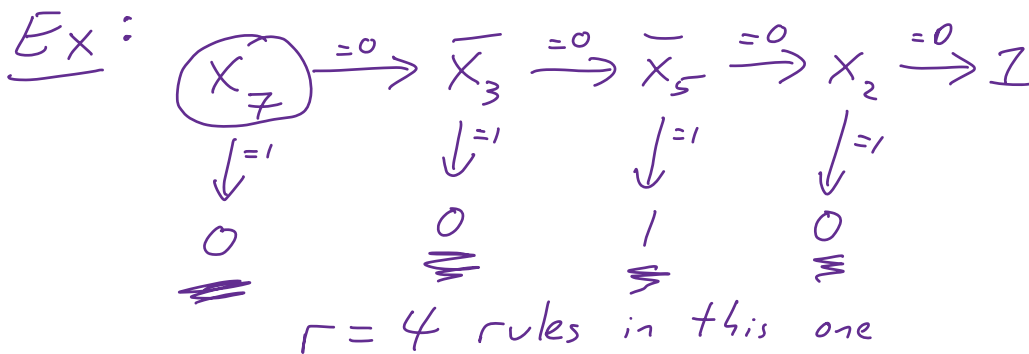
A 1-decision list (1-DL) is a concept

$$c: \{0,1\}^n \rightarrow \{0,1\} \quad \text{which is an ordered}$$

list of if-then-else rules:

literal *one literal per rule*

if  $l_1$  true then output  $b_1$   
 else "  $l_2$  " " "  $b_2$   $\geq$  poss  
 :  
 : "  $l_r$  " " "  $b_r$   
 else output  $b_{r+1}$   $\geq$  poss.

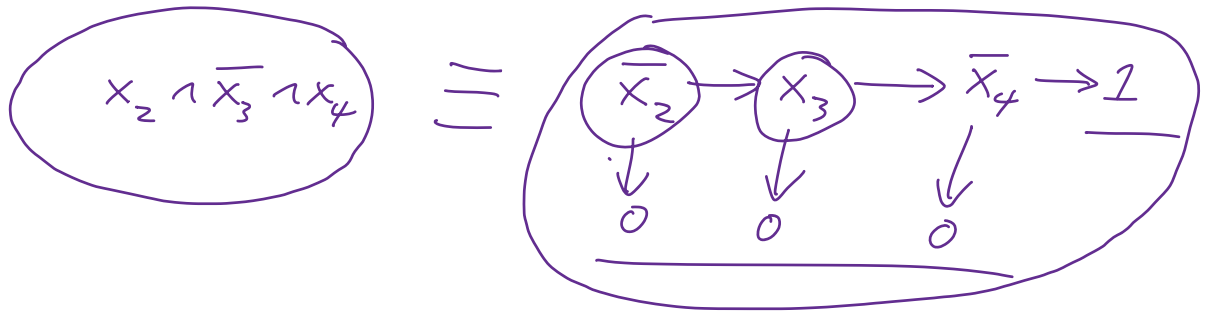


• Would never need to repeat lits or vars:





- Any conj is expressible as a DL:



Also any disj.

$$C_{\text{conj}} \subseteq C_{\text{DL}}$$

Note: # of 1-DL of length  $r$

$$\text{is } \approx 2^{r+1} \cdot (2n)^r \approx (4n)^r.$$

Total # of possible rules, including the two "default rules", is  $4n+2$

Next time:  $O(nr)$ -mistake bound  
 alg. for class of 1-DL's of  
 length  $r$ . Based on reordering  
 rules in careful way.