

An Introduction to Applicative Functors

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What Is an Applicative Functor?

- An Applicative functor is a Monoid in the category of endofunctors, what's the problem?
- WAT?!

Functions in Haskell

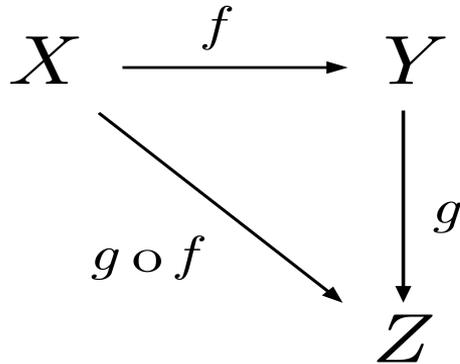
- Functions in Haskell are first-order citizens
- Functions in Haskell are curried by default
 - $f :: a \rightarrow b \rightarrow c$ is the curried form of $g :: (a, b) \rightarrow c$
 - $f = \text{curry } g, g = \text{uncurry } f$
- One type declaration, multiple interpretations
 - $f :: a \rightarrow b \rightarrow c$
 - $f :: a \rightarrow (b \rightarrow c)$
 - $f :: (a \rightarrow b) \rightarrow c$
 - Use parentheses when necessary:
 - $>>= :: \text{Monad } m \Rightarrow m \ a \rightarrow (a \rightarrow m \ b) \rightarrow m \ b$

Functors

- A **functor** is a type of mapping between categories, which is applied in category theory.
- What the heck is category theory?

Category Theory 101

- A category is, in essence, a simple collection. It has three components:
 - A collection of **objects**
 - A collection of **morphisms**
 - A notion of **composition** of these morphisms



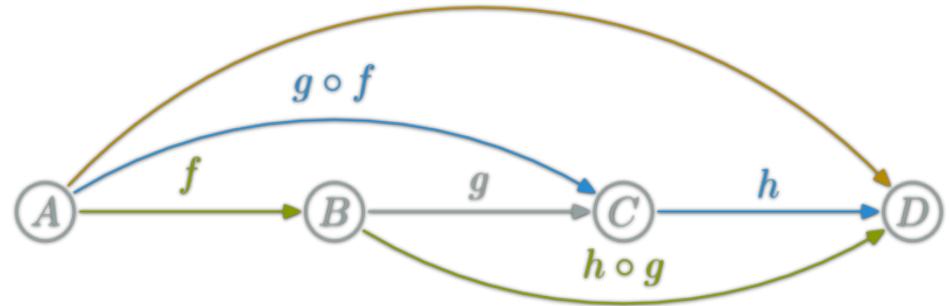
- Objects: X, Y, Z
- Morphisms: $f :: X \rightarrow Y, g :: Y \rightarrow Z$
- Composition: $g \circ f :: X \rightarrow Z$

Category Theory 101

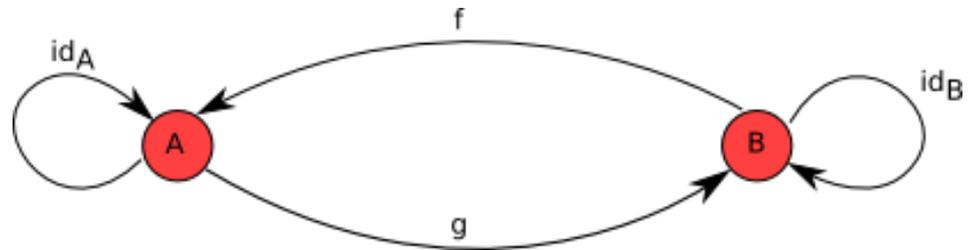
- Category laws:

$$f \circ (g \circ h) = (f \circ g) \circ h$$

$$(h \circ g) \circ f = h \circ (g \circ f)$$

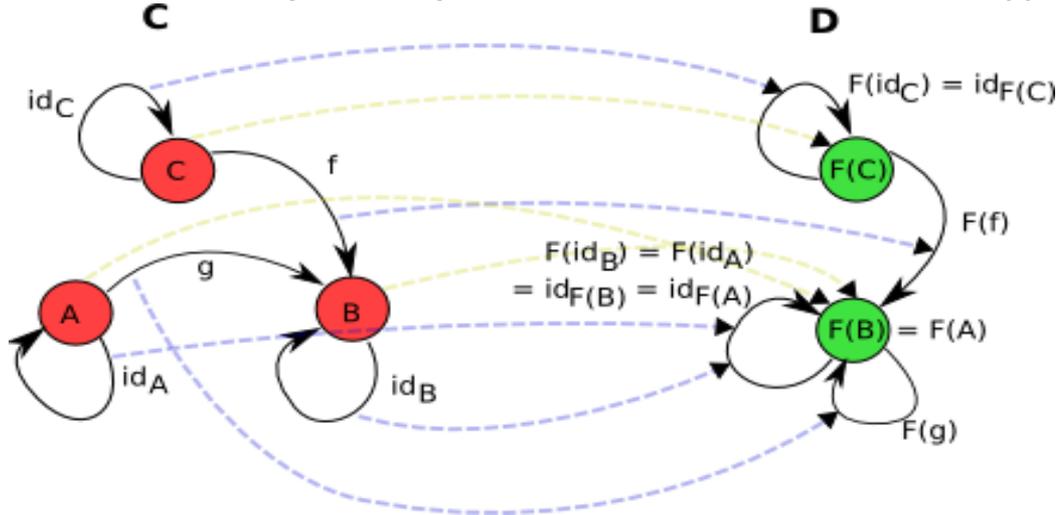


$$g \circ id_A = id_B \circ g = g$$



Functors Revisited

- Recall that a **functor** is a type of mapping between categories.
- Given categories **C** and **D**, a functor $F :: C \rightarrow D$
 - Maps any object A in **C** to $F(A)$ in **D**
 - Maps morphisms $f :: A \rightarrow B$ in **C** to $F(f) :: F(A) \rightarrow F(B)$ in **D**



Functors in Haskell

class Functor f where

fmap :: (a -> b) -> f a -> f b

- Recall that a functor maps morphisms $f :: A \rightarrow B$ in \mathbf{C} to $F(f) :: F(A) \rightarrow F(B)$ in \mathbf{D}
- morphisms \sim functions
- $\mathbf{C} \sim$ category of primitive data types like Integer, Char, etc.
- $\mathbf{D} \sim$ category of “functorized types” like Maybe Integer, Maybe Char, etc.
- fmap actually takes as parameter a function ($g :: a \rightarrow b$), and returns a function ($g' :: f a \rightarrow f b$)

Endofunctors

- A **functor** is a type of mapping between 2 categories.
- What if the 2 categories are the actually the same category? You got endofunctors
- Functors in Haskell are actually endofunctors

We have a category **Hask**, which treats ALL Haskell types as objects and Haskell functions as morphisms and uses (.) for composition

Applicative Functors

class (Functor f) => Applicative f where

pure :: a -> f a

<*> :: f (a -> b) -> f a -> f b

-- fmap

<\$> :: (a -> b) -> f a -> f b

Function-in-the-box

- Applicative functors are another mechanism for dealing with programming with effects(values wrapped in a context)
- Applicative functors are more powerful than functors because they are able to deal with functions in a context
- But how do functions get into a “box” in the first place?

fmap



Function-in-the-box

- How do functions get into a context?

- Just use `pure :: a -> f a`

- Use `fmap`:

```
fmap (+) [1] or (+) <$> [1]
```

```
>> [(+ 1)]
```

```
(+) <$> [1, 2] <*> [3, 4]
```

```
>> [4, 5, 5, 6]
```

A Use Case

```
data User = User { firstName :: Text,  
                  LastName :: Text,  
                  Email :: Text }
```

```
buildUser :: Profile -> Maybe User
```

```
buildUser p = User  
  <$> lookup "first_name" p  
  <*> lookup "last_name" p  
  <*> lookup "email" p
```

```
buildUser p = do  
  fn <- lookup "first_name" p  
  ln <- lookup "last_name" p  
  em <- lookup "email" p  
  return $ User fn ln em
```

Why Applicatives?

Q: We already got this Monad dude, who is, like, super awesome. Why do we need to hire you for this task?

A: I'm flexible on salary, and I get shit done faster

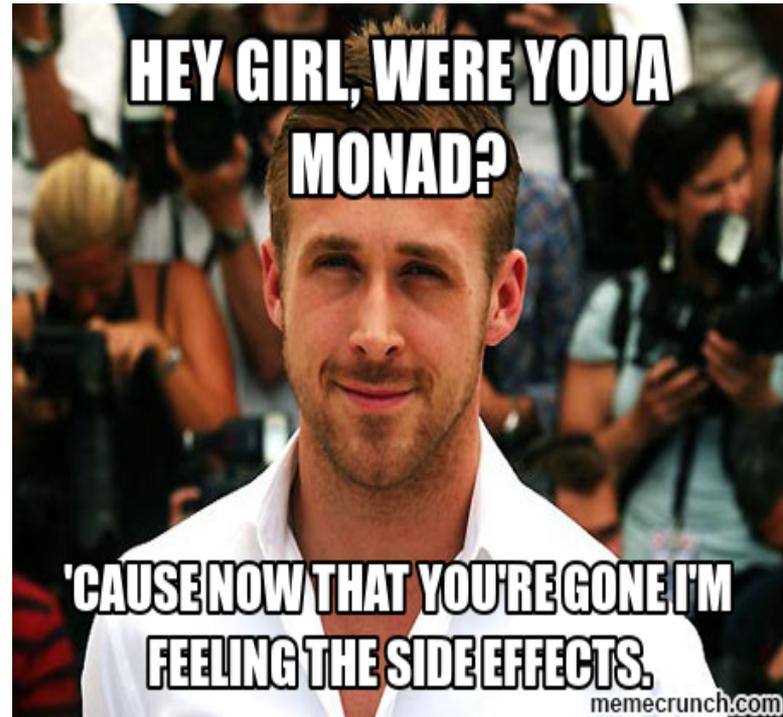
Q: Okay, what's your name again?

A: Applicative Functor

Q: Geez, that's a mouthful!

Applicatives vs. Monads

- Monads are about...
 - Effects
 - Composition
 - Sequence/Dependency
 - parsing context-sensitive grammar
 - branching on previous results
- Applicatives are about...
 - (less severe)Effects
 - Batching and aggregation
 - Concurrency/Independency
 - parsing context-free grammar
 - exploring all branches of computation



Disaster Averted (or Not)

- `miffy :: Monad m => m Bool -> m a -> m a -> m a`
`miffy mb mt me = do`
 `b <- mb`
 `if b then mt else me`
`>> miffy (Just True) (Just "Yay!") Nothing = Just "Yay!"`
- `iffy :: Applicative f => f Bool -> f a -> f a -> f a`
`iffy fb ft fe = cond <$> fb <*> ft <*> fe where`
 `cond b t e = if b then t else e`
`>> iffy (Just True) (Just "Yay!") Nothing = Nothing`

Should It Always Fail Early?

- Monads have this inherent property that they can branch on the results of previous computations, which implies they always fail early(short-circuited)
- What if you want to design a [signup page](#) for your website?
- What if you actually don't really care whether the computation should fail early or not?

Weaker But Sometimes Better

- Applicatives are weaker than Monads, which also means they are more common than Monads
- Applicative code is usually cleaner and shorter than its monadic counterpart, and lends itself to optimization
 - Facebook's Haxl provides a DSL that expose the monadic interfaces and converts them to applicatives when necessary
- Use the least powerful mechanism to get things done
- When there's no dependency issues or branching, just use applicatives

Like Father, Like Son

- All monads are applicatives, but not all applicatives are monads
 - ZipList
- Applicative is actually a superclass of monad
- Fun fact: Actually applicatives were discovered **later** than monads
- Due to historical reasons, applicative is NOT a superclass of monad in Haskell yet (but it soon will be)

Applicative => Monad Proposal (AMP)

- Applicative becomes a superclass of Monad
- Why?
 - lack of unity means there is a lot of duplication of API:
 - `liftA :: (Applicative f) => (a -> b) -> f a -> f b`
 - `liftM :: (Monad m) => (a -> b) -> m a -> m b`
 - `pure = return`, `<*> = ap`
 - `ap mf ma = do`
 - `f <- mf`
 - `a <- ma`
 - `return $ f a`
 - Enforce the use of the least restrictive functions

So an Applicative Functor Is...

- A Monoid in the category of endofunctors. That's it.
- Dammit! What the heck is a Monoid?
 - class Monoid m where
 mempty :: m
 mappend :: m -> m -> m
 - instance Monoid [a] where
 mempty = []
 la mappend lb = (++) <\$> la <*> lb

Resources

- <http://learnyouahaskell.com/functors-applicative-functors-and-monoids>
- [Applicative programming with effects](#)
- [Applicative Functors: Hidden in plain view](#)
- [Haskell/Category Theory](#)
- [Introduction to functional programming](#)
- [Beginning Haskell: A Project-Based Approach](#)
- [Haskell Ryan Gosling](#)