



# Breaking Bidder Collusion in Large-Scale Spectrum Auctions

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University of California, Santa Barbara

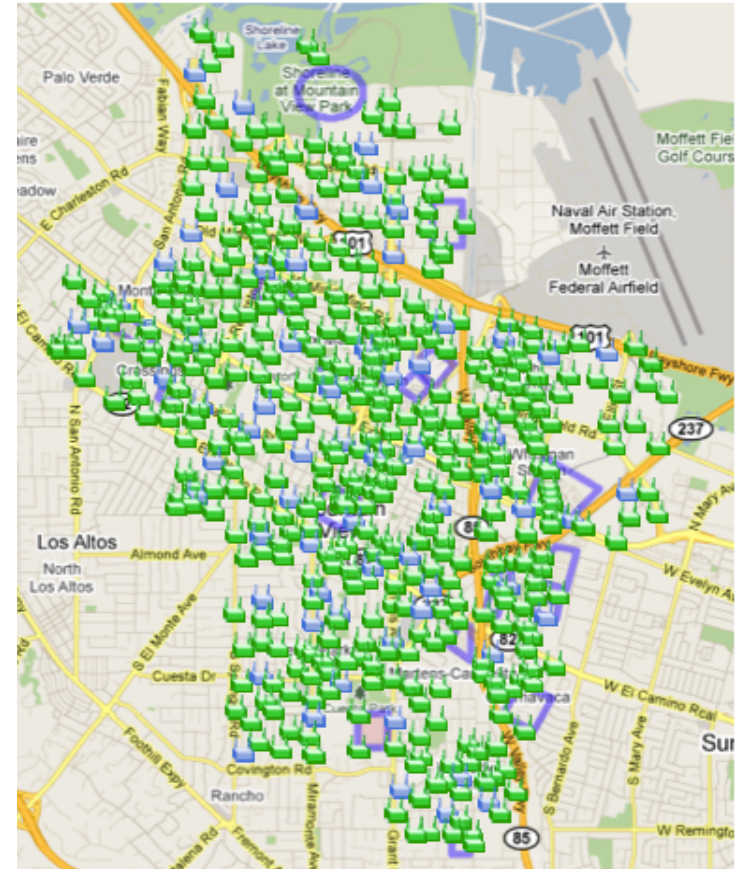


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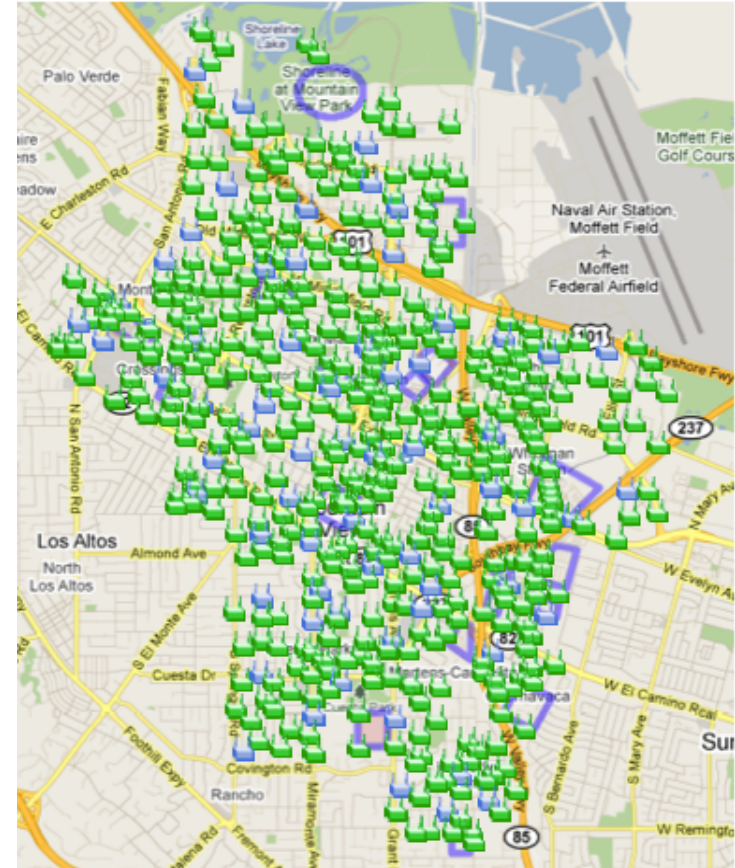
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  - Dynamical demands, local wireless service



*GoogleWiFi Network*

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- Obtaining spectrum is difficult
  - Unlicensed band → too crowded! ☹️
  - Licensed band → long-term usage, pricy! ☹️

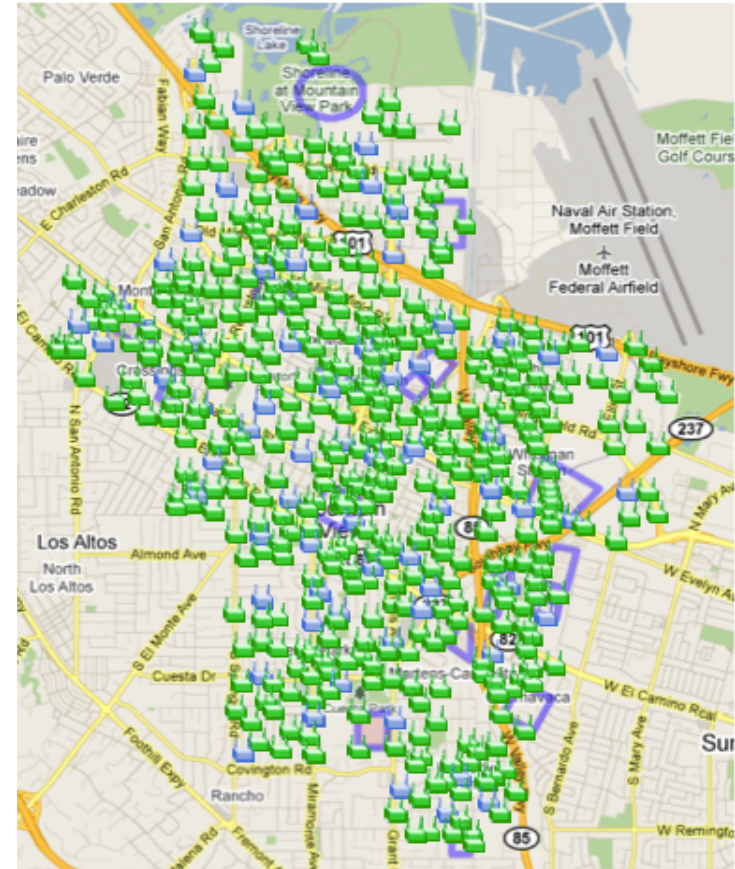


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- Obtaining spectrum is difficult
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  - Licensed band → long-term usage, pricy! ☹
- Dynamic spectrum distribution with spatial reuse




*GoogleWiFi Network*



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  - Bidders bid for spectrum to match their needs

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- Key requirements:

- Maximize spectrum distribution efficiency
  - Enabling spectrum reuse
- Resist bidder cheating

 Networking

 Economics



# A Closer Look at Bidder Cheating



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- Individual cheating
  - Change bid to gain unfair advantage
  - Solution: truthful spectrum auction designs
    - VERITAS [zhou08], TRUST[zhou09], [jia09]...



# A Closer Look at Bidder Cheating

- Individual cheating
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  - Solution: truthful spectrum auction designs
    - VERITAS [zhou08], TRUST[zhou09], [jia09]...
- Collusion
  - Cheat in groups, improving the group's utility
  - Popular in large-scale networks
    - Example: P2P networks
  - Few studies in dynamic spectrum auctions



# Our Contributions

- **Understand the impact of bidder collusion in dynamic spectrum auctions**
- **Propose a collusion-resistant design for large scale spectrum auctions**



# Outline

- Is bidder collusion a serious threat to spectrum auction?
- How to address bidder collusion?
- Evaluation
- Conclusion and future works



# (Truthful) Spectrum Auctions 101



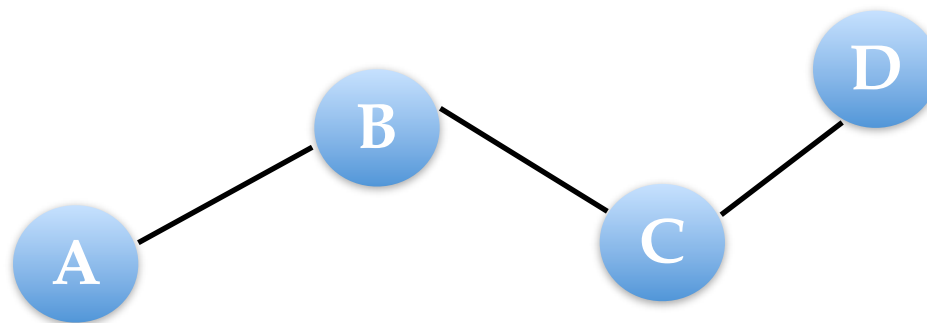
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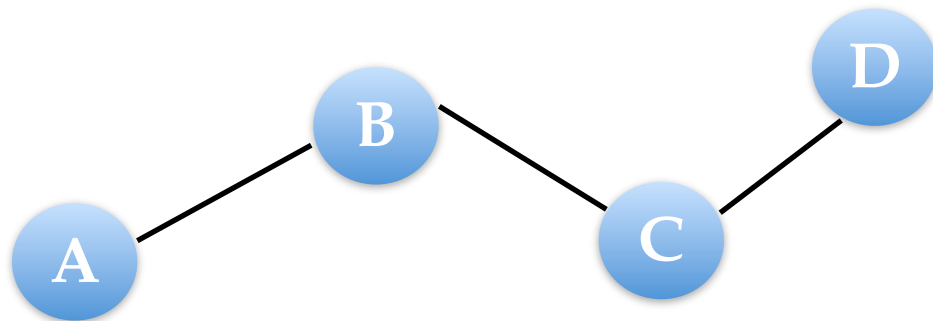
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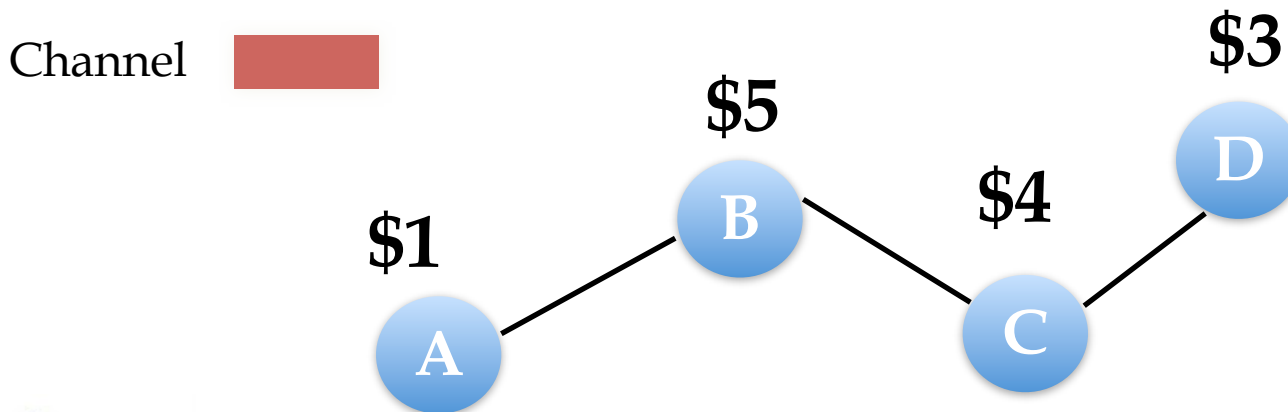
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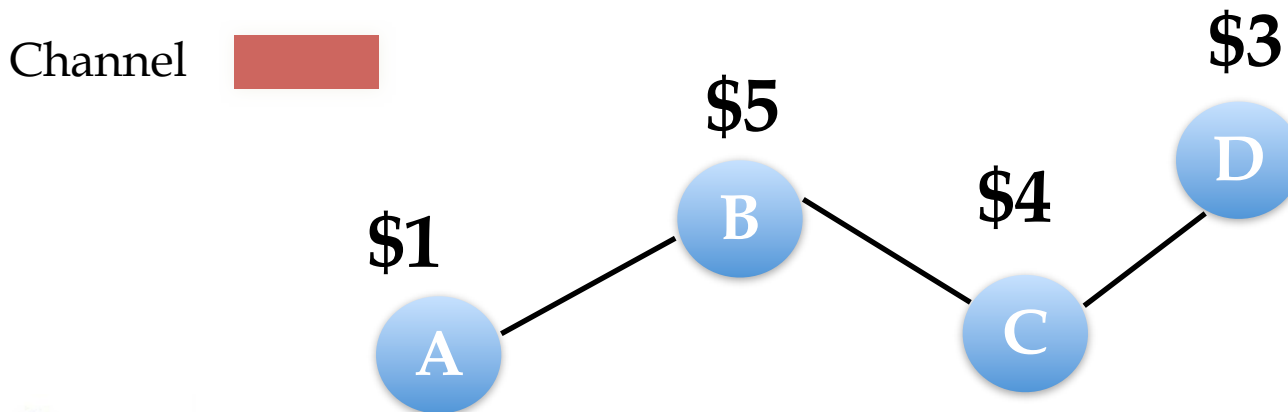
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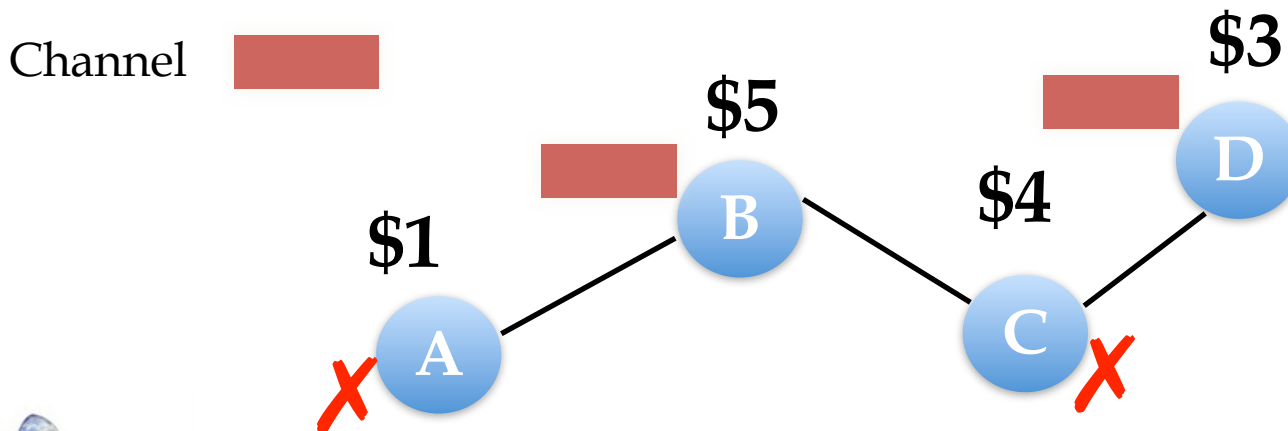
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  - Bid-dependent greedy allocation



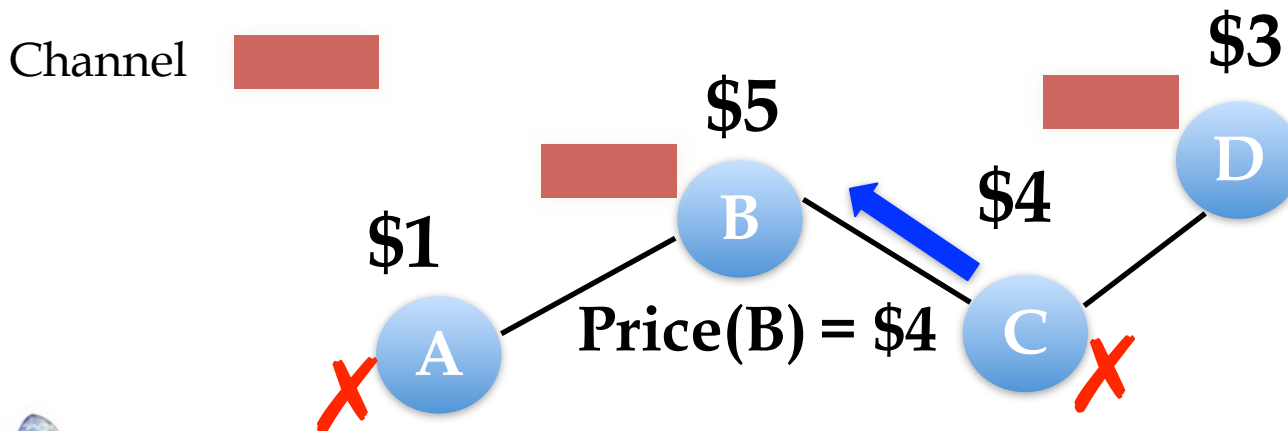
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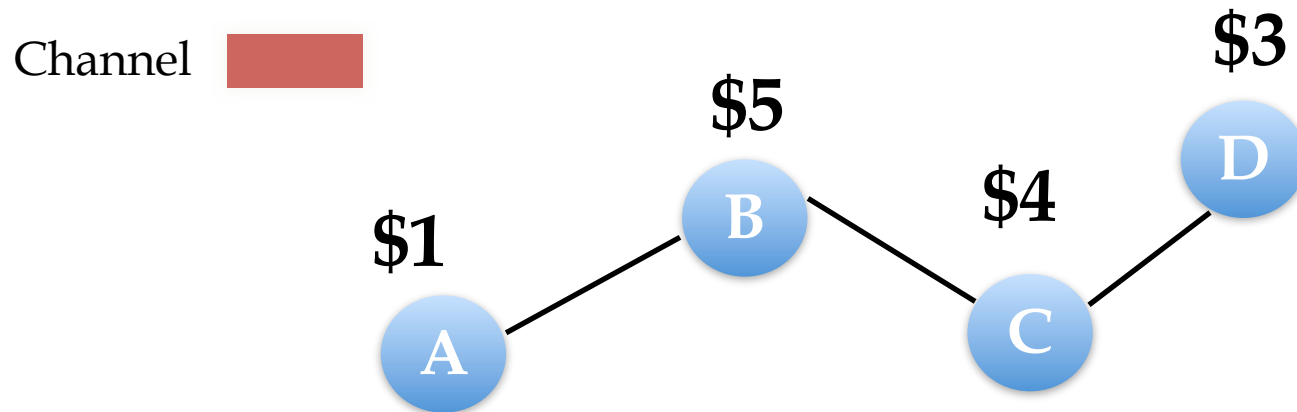


# (Truthful) Spectrum Auctions 101

- Must enable spatial reuse
- **VERITAS**: A representative truthful spectrum auction
- Allocation
  - Bid-dependent greedy allocation
- Pricing
  - **Critical neighbor**: for bidder  $i$ , if  $i$  bids lower than its critical neighbor, then  $i$  cannot win the auction; otherwise it wins.

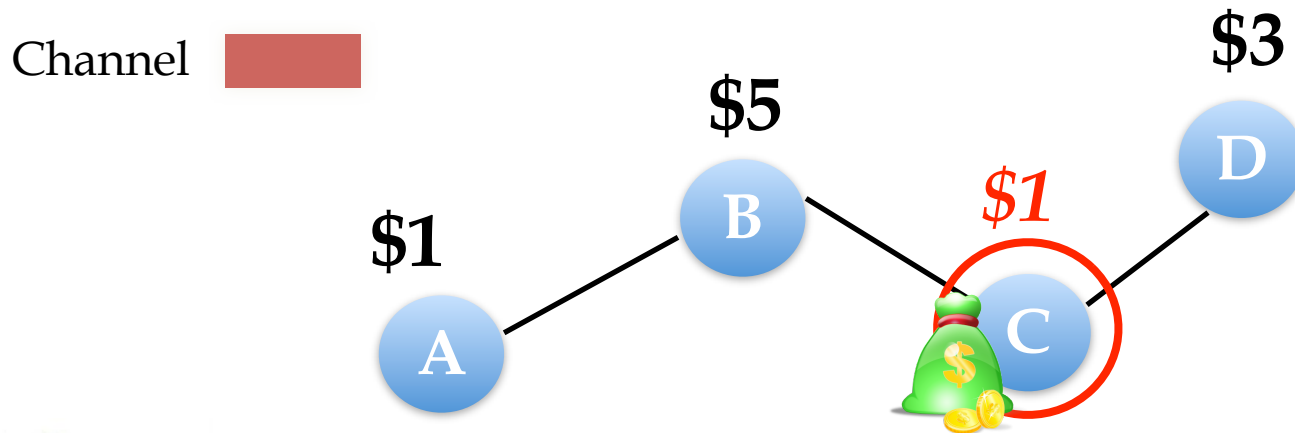


# An Illustrative Collusion Example



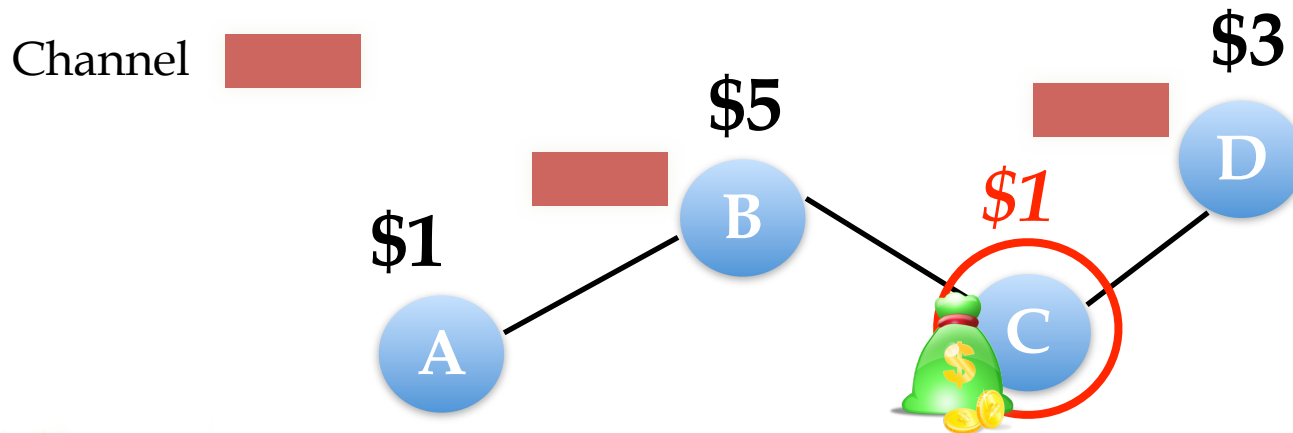
# An Illustrative Collusion Example

- Winner-Critical Neighbor (WCN) Collusion
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  - B pays C to bid lower
  - B wins and pays ONLY \$1
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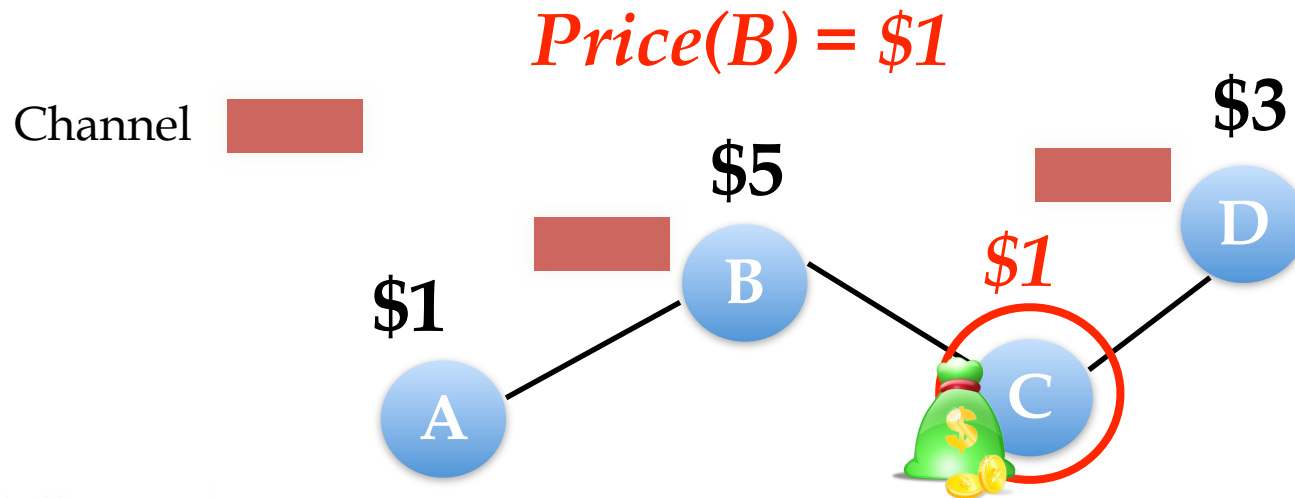
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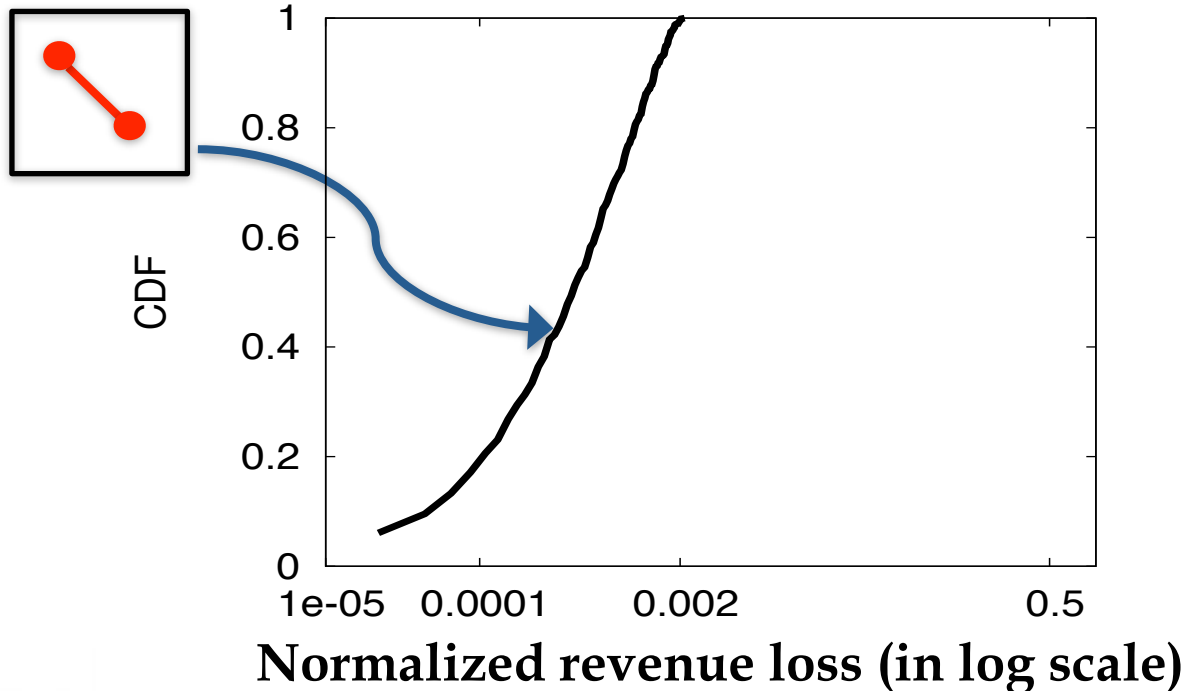
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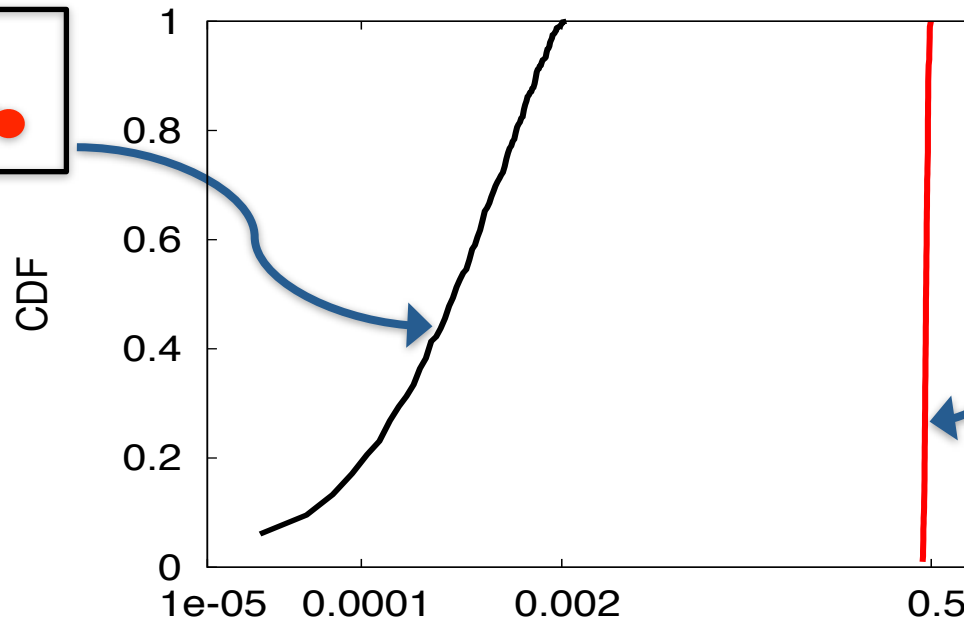
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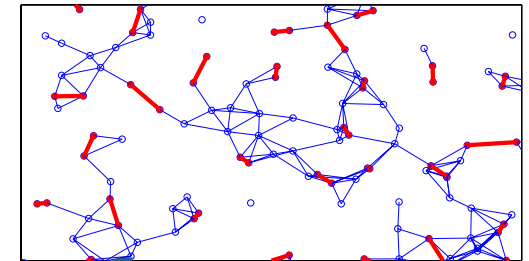
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## Single Collusion group



## Multiple Collusion groups



*Up to 50%  
revenue loss!*

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- Is bidder collusion a serious threat to spectrum auction? – Yes, small-size bidder collusion is a huge threat

## How to address bidder collusion?



- Evaluation
- Conclusion and future works

# Our Methodology



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- **Prevention** rather than detection
  - ‘Needle in a hay’: hard to detect small size collusion group
  - **Prevention**  $\equiv$  nullify collusion gain  $\rightarrow$  no gain, no collusion





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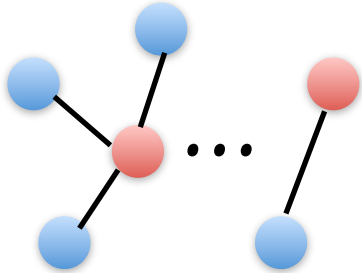
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- **Soft prevention while enabling spectrum reuse**
  - Existing designs assume “all conflict” or “none conflicts”
  - Need new design



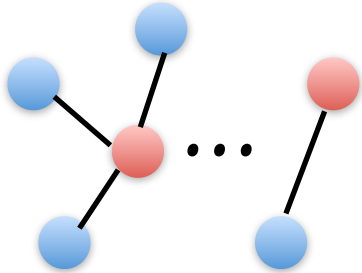
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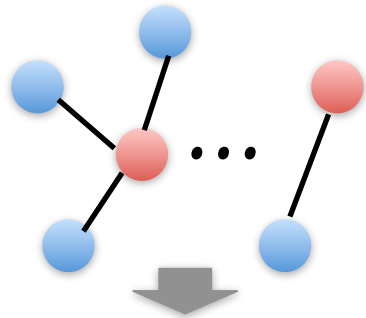


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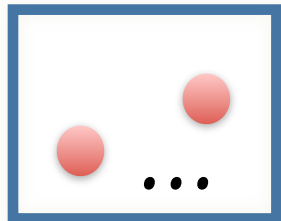
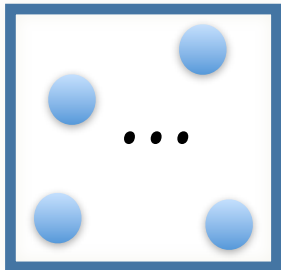
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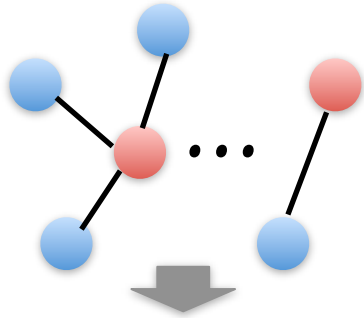
- Form bidder segments
- Bidders in each segment do not conflict



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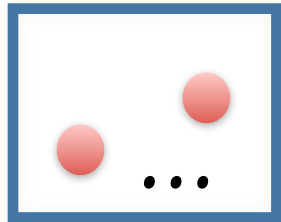
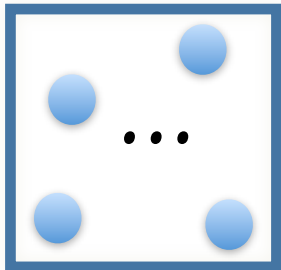
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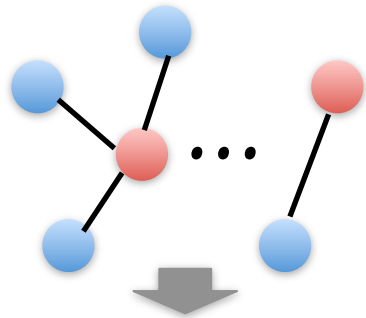
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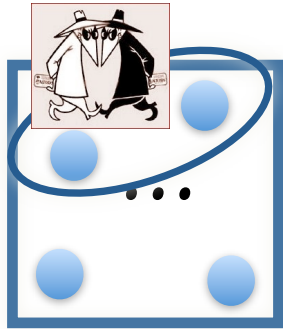


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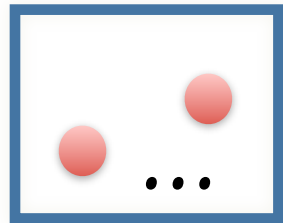


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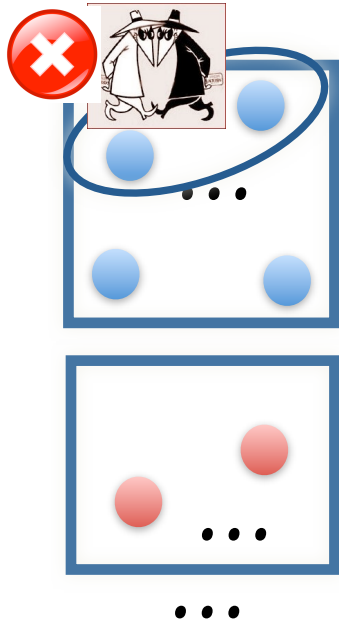
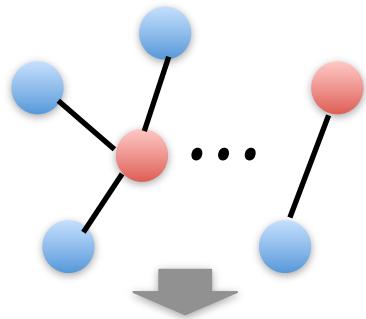


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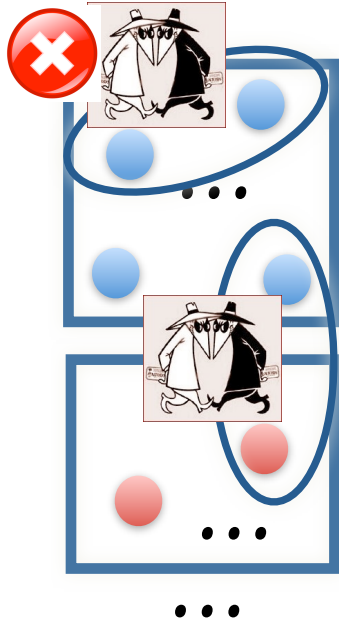
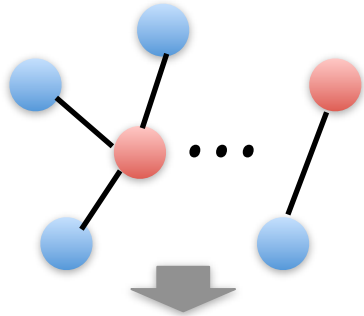
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- Tackle collusion within a segment
  - Use collusion-resistant design (tCP) to choose potential winners in each segment



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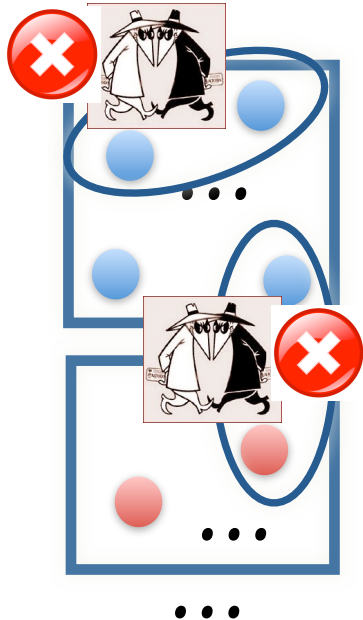
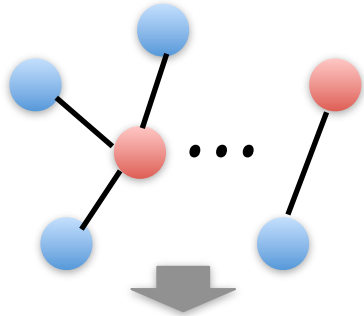
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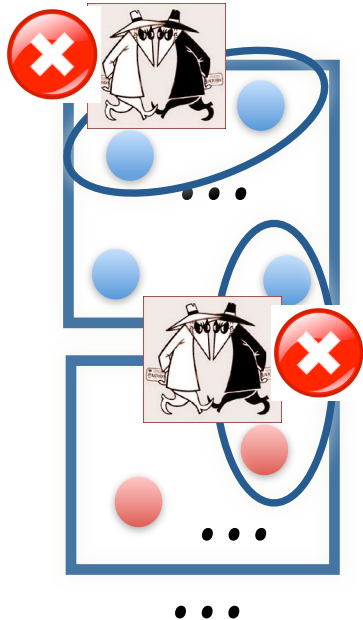
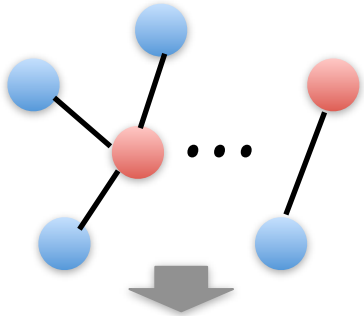
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- Tackle collusion across segments
  - Add **randomness** to winning segment selection

# Athena Spectrum Auctions



## Enabling spectrum reuse

- Form bidder segments
- Bidders in each segment do not conflict

Divide

## Diminishing collusion gain

- Tackle collusion within a segment
  - Use collusion-tolerant design (tTCP) to choose potential winners in each segment
- Tackle collusion across segments
  - Add randomizing segment selection

Conquer

Combine

# Detailed Design

Divide

Conquer

Combine



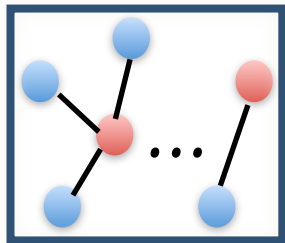
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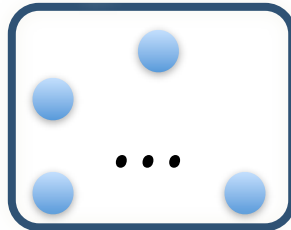
Conquer

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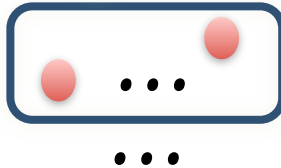
- Divide bidders into segments
- Bidders in each segment do not conflict
- **Partition is bid-independent**



Segment 1



Segment 2

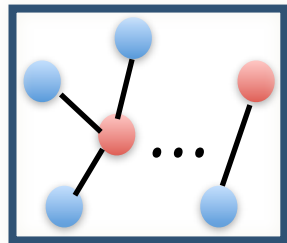


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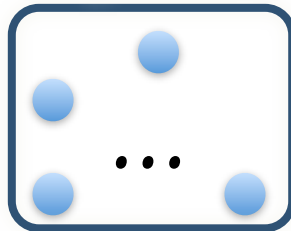
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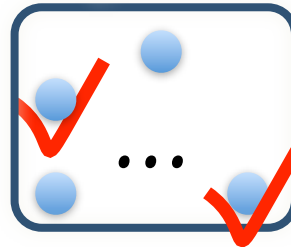


Segment 2



...

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Segment 2



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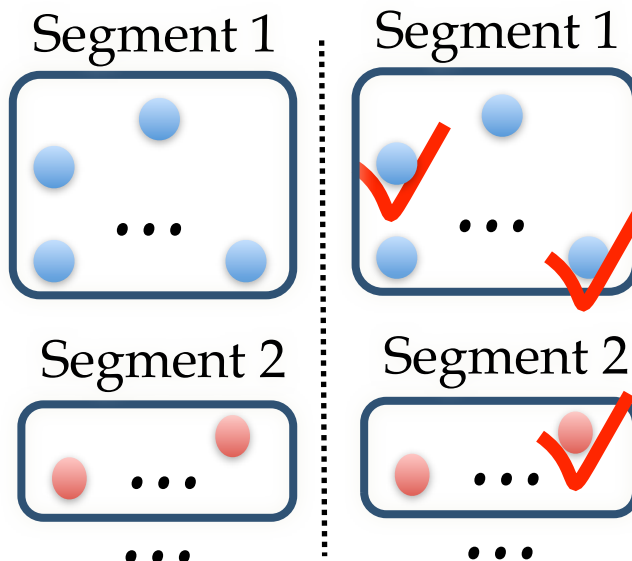
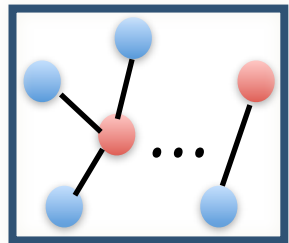
# Detailed Design

Divide

Conquer

Combine

- For each segment  $i$ , select potential winners using a uniform price  $p_i$
- **tCP method<sup>1</sup>**: make  $p_i$  **insensitive** to bid changes within segment  $i$   
→ no gain for intra-segment collusion



*1: [goldberg03]*





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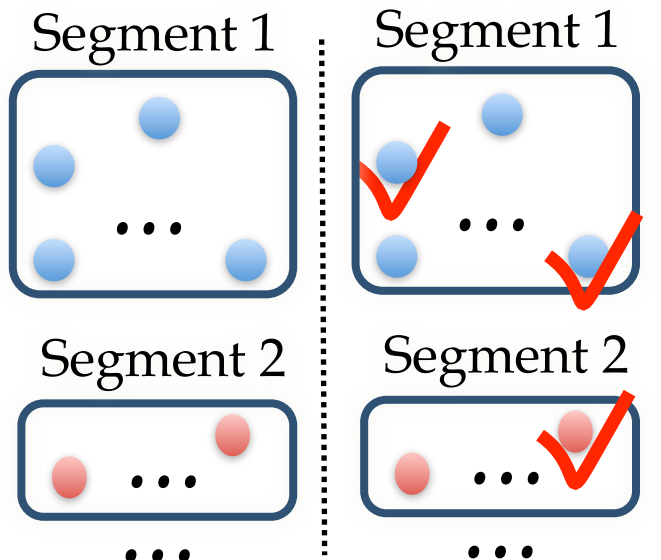
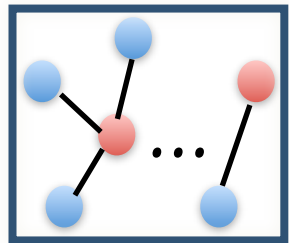
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Conquer

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- Estimate each segment's revenue
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- **Add randomness in revenue estimation to diminish the impact of inter-segment collusion**



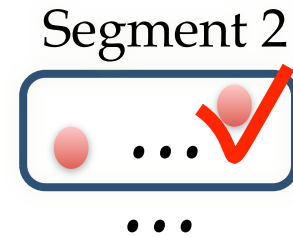
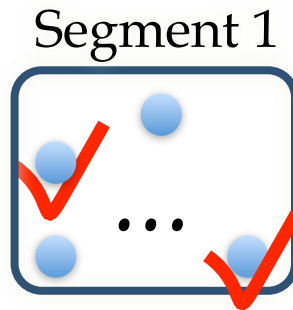
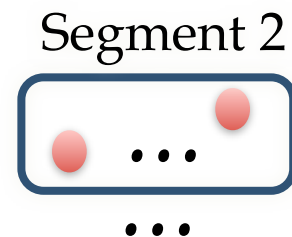
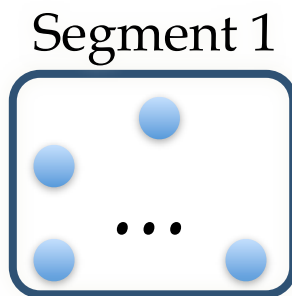
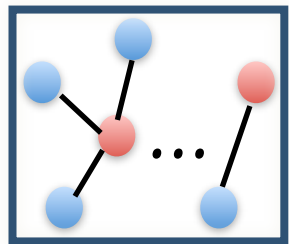
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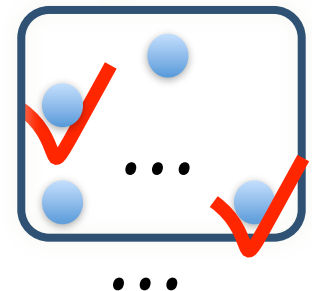
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Estimated Revenue \$

Estimated Revenue \$

Final winners & prices



# Summary

**Divide**



**Spatial reuse by bid-independent partition**

**Conquer**



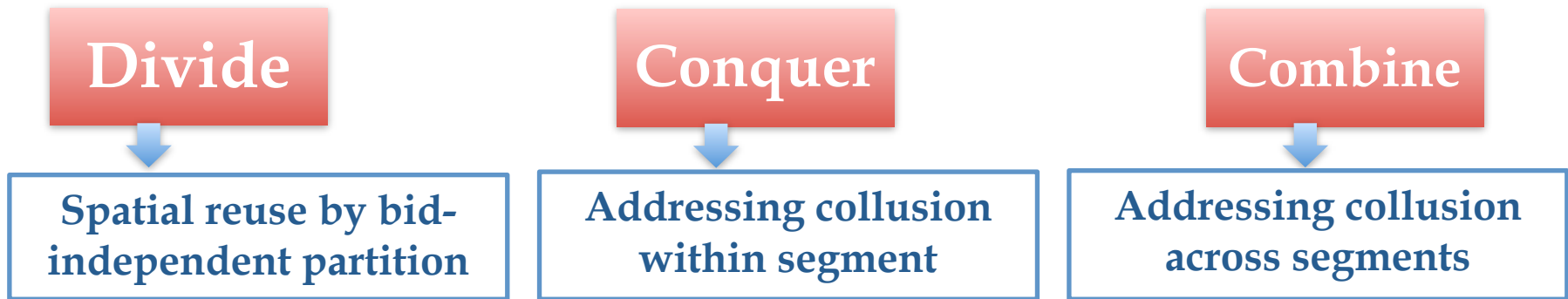
**Addressing collusion within segment**

**Combine**





**Addressing collusion across segments**

# Summary



- Athena's collusion resistance
  - *(t, p)-truthfulness*: with probability  $\geq p$ , no collusion group of  $\leq t$  bidders can improve group utility by collusion
  - Athena achieves *(t, p)-truthfulness*,  $p$  depends on  $t$  and the #winners in the smallest segment

# Fine-Tuning Athena

- *Segment sizes* affect the choice of the pricing scheme in 'Conquer' in order to maximize revenue given  $(t, p)$ 
  - Uniform segment sizes 
  - Non-uniform segment sizes 
    - Carefully select segments running tCP and their configurations
- Athena's revenue bound
  - When all segments run tCP, the distance of Athena's revenue to the optimal is a function of  $t$ ,  $p$ , and *segment sizes*



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- Is bidder collusion a serious threat to spectrum auction? – Yes, small-size bidder collusion is a huge threat
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# Evaluating Athena



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- **Case study 1:** Effectiveness on resisting collusion
  - Can Athena diminish collusion group gain?



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  - Experiment with **typical** bid patterns;
- **Case study 1:** Effectiveness on resisting collusion
  - Can Athena diminish collusion group gain?
- **Case study 2:** The cost of collusion resistance
  - How much revenue Athena needs to sacrifice for collusion-resistance?
  - Compare to VERITAS (truthful auctions)

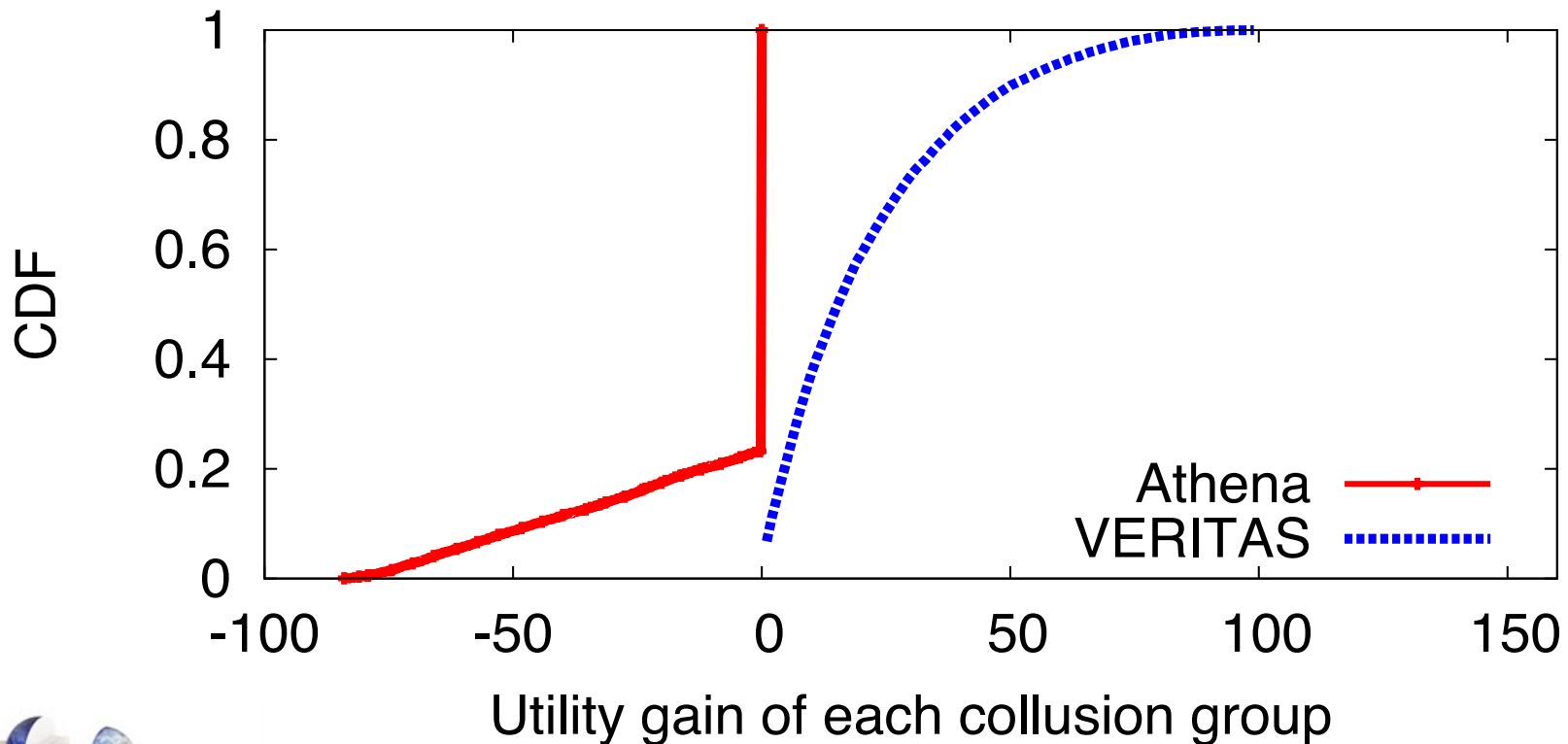


# Athena's Collusion Resistance



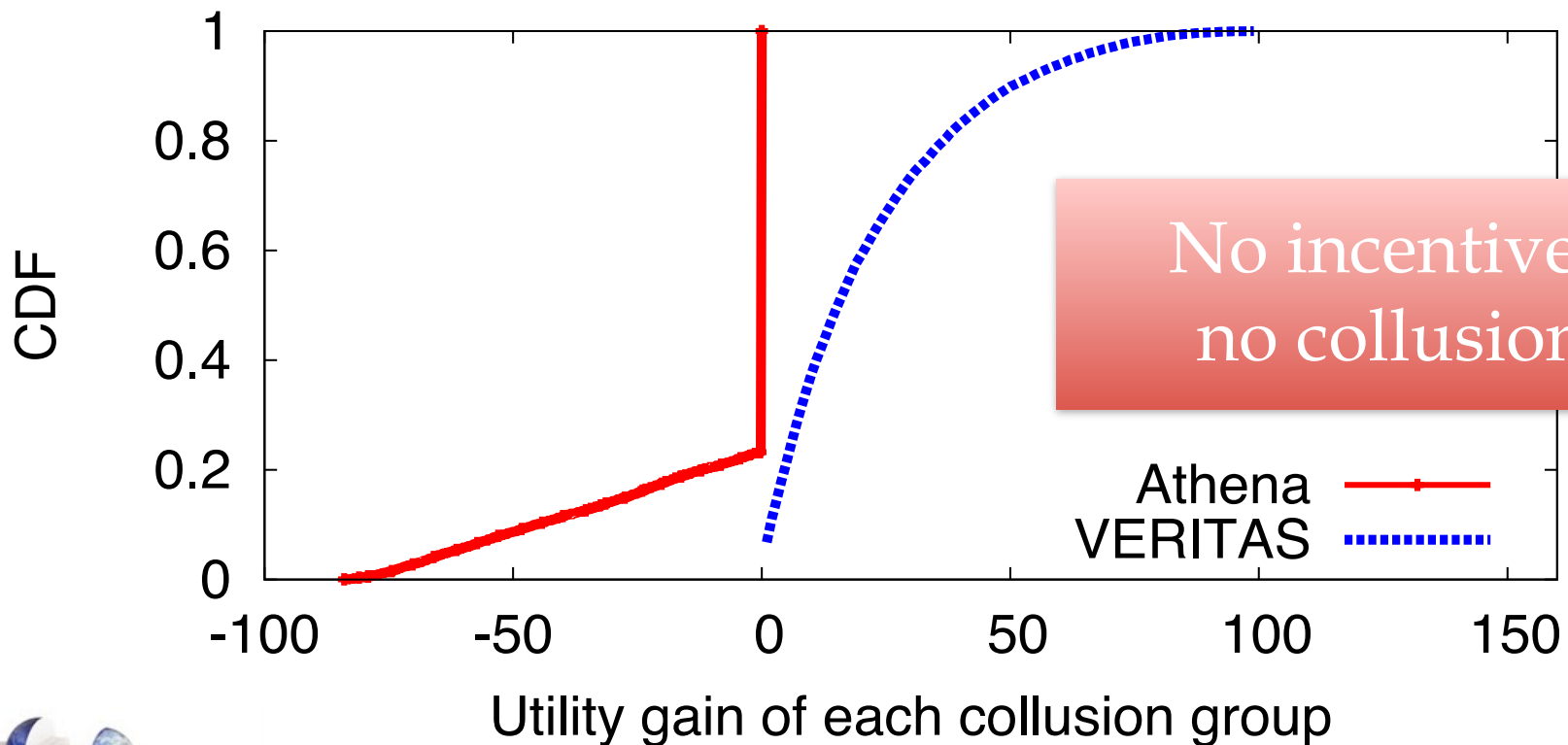
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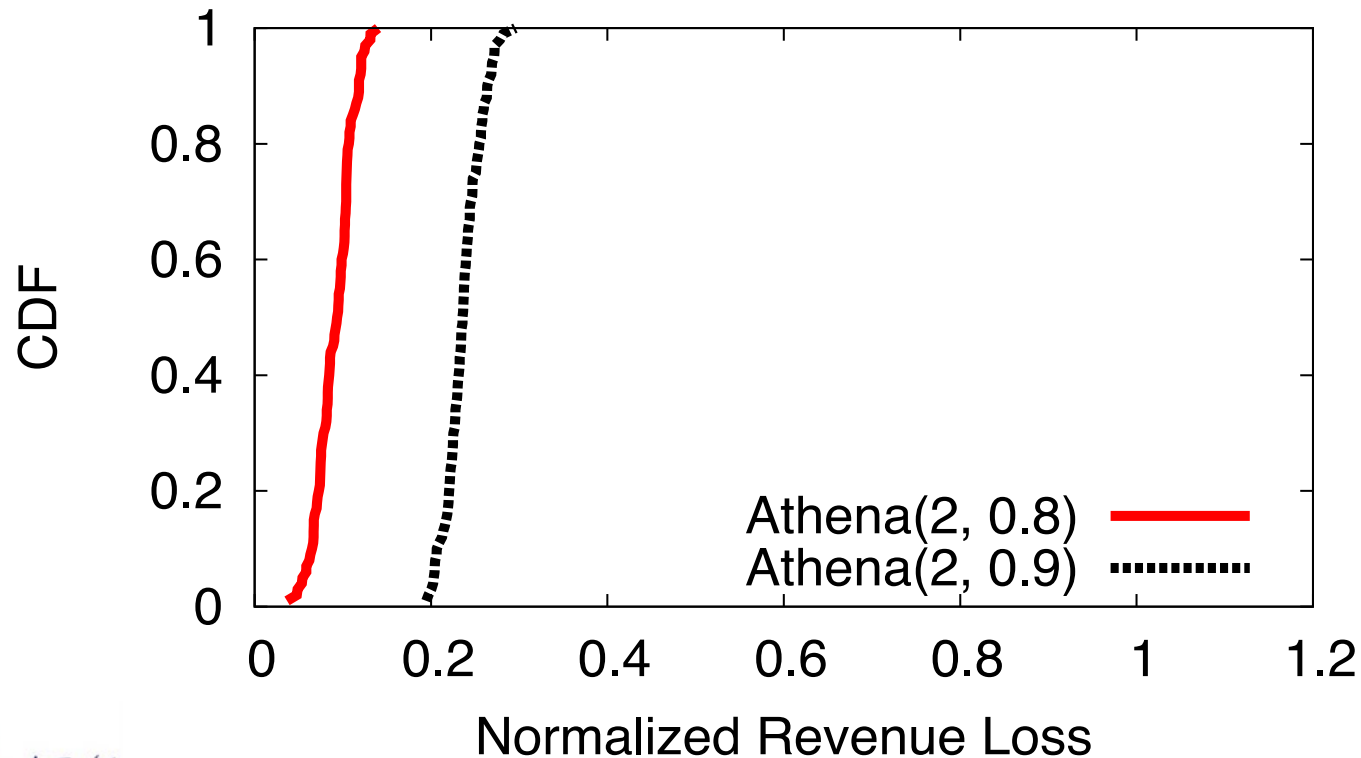
# The Cost of Collusion-Resistance

$$\text{Normalized revenue loss} = 1 - \frac{\text{Revenue}}{\text{VERITAS revenue}}$$



# The Cost of Collusion-Resistance

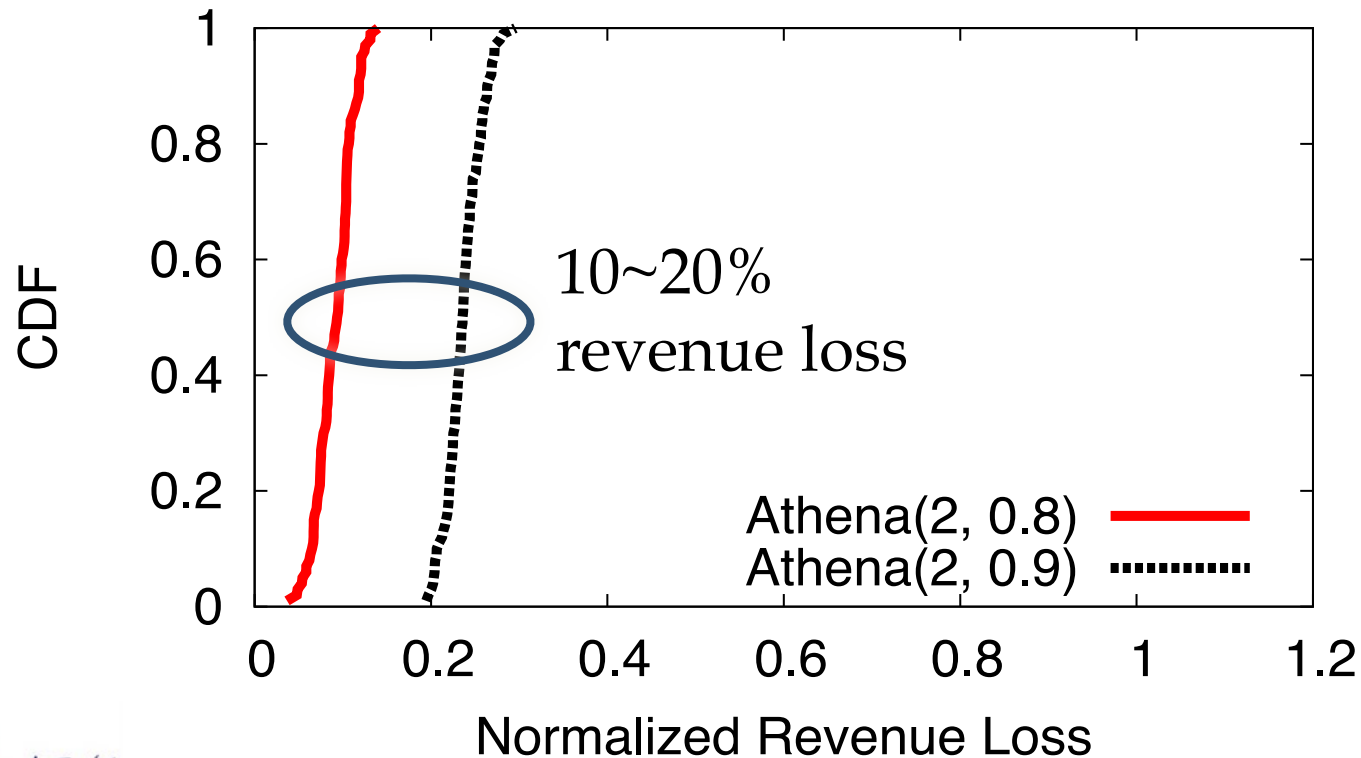
$$\text{Normalized revenue loss} = 1 - \frac{\text{Revenue}}{\text{VERITAS revenue}}$$





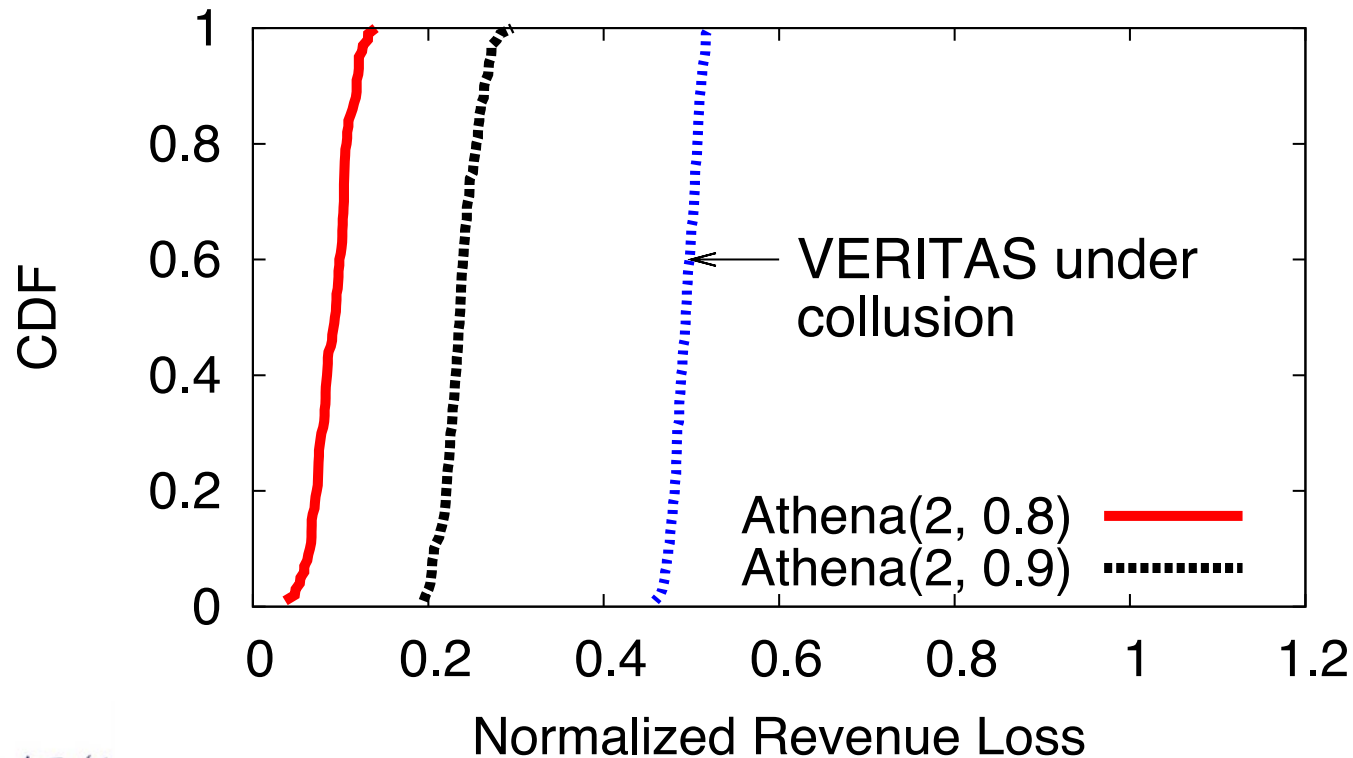
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  - Customizable collusion-resistance
- Future work
  - Extend to multi-channel request
  - Explore the optimal segment formation



# Q & A

- Thanks!

For more information, please visit:

<http://link.cs.ucsb.edu/project/mercury.html>

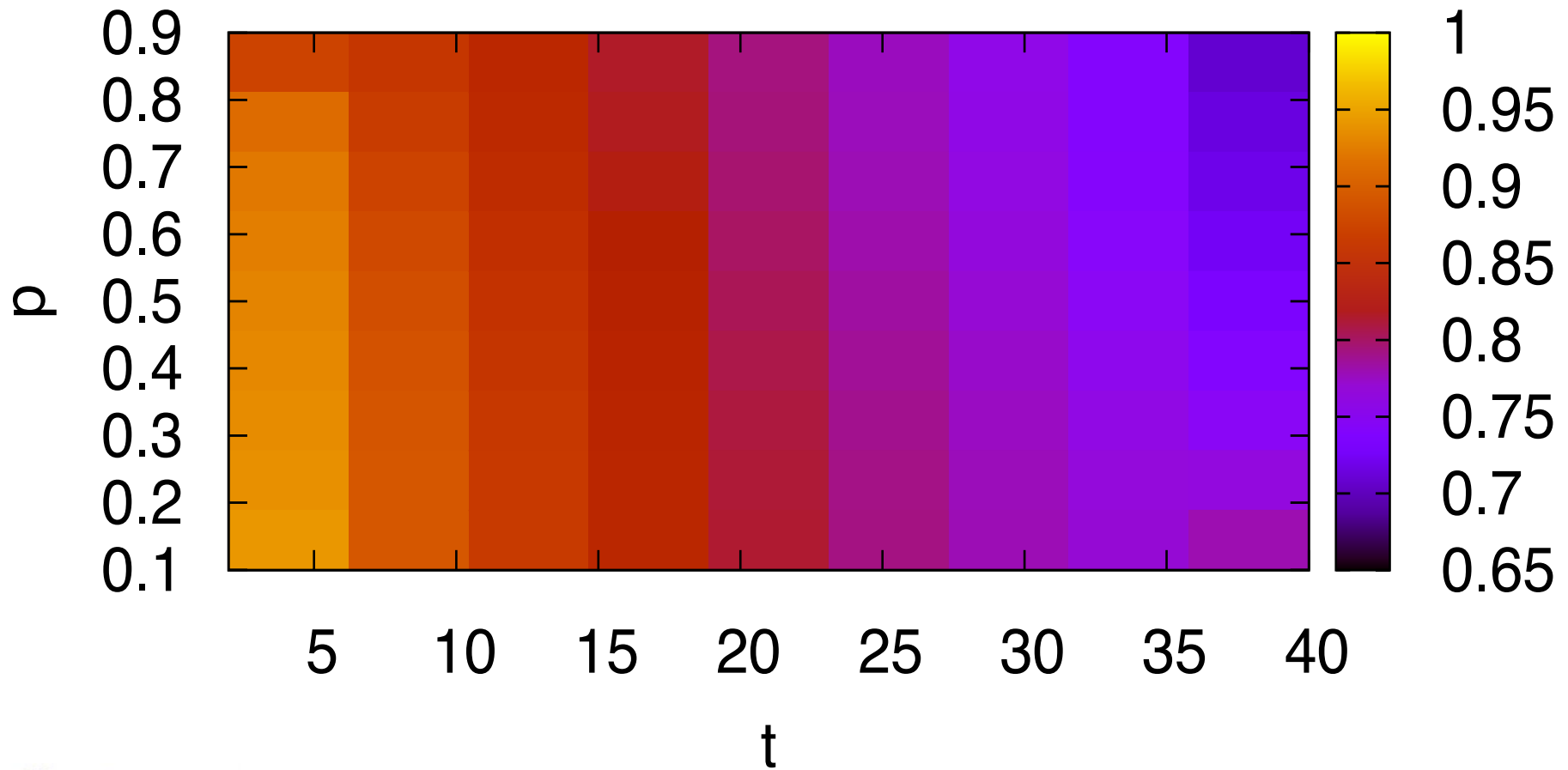


# BACK-UP SLIDES





# For all $(t, p)$



# Comparing to Posted Price

- Assuming no bidders collude due to the awareness of the design's collusion resistance

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