

COMS 3101

Programming Languages: MATLAB

Lecture 6

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Lecture Outline

- Review: HW4
- Practical math / optimization (continued)
- Advanced functionality:
 - `convhull()`, `inpolygon()`, `polyarea()`
 - `triplot()`, `fill()`
 - `mldivide (\)`, `linprog()`
 - `sortrows()`
 - Convexity

HW4

- Normalize a vector 'V' to the [0,1] range:
 - $V = (V - \min(V)) / (\max(V) - \min(V))$
- Check a set of fields for a struct:
 - `ck = isfield(specs, {'f1', 'f2', 'f3'})`
- Check for a non-empty string variable 'fdir':
 - `str_ck = isstr(fdir) & ~isempty(fdir);`
 - `str_ck = ischar(fdir) & ~isempty(fdir);`
- Check for a positive integer variable 'N':
 - `int_ck = isnumeric(N) & (round(N) == N) & (N > 0);`

textscan()

- Read formatted data from a text file
- `textread()` is another (deprecated) function with similar functionality
- How does it work?
 - `C = textscan (FID, 'format')`
 - Recall that to open a text file, we need to create an identifier:
`FID = fopen ('myfile.text', 'r');`
 - `'format'` is a string specifying how each line should be read
 - `'C'` is a cell array. The number of specifiers (`'format'`) determines the number of cells

More Plotting – 3D

- Very similar to 2D plotting
- Keyword: `plot3()`
- Example:
 1. `X = 0:0.1:1; Y = X; Z = X.^2+Y.^2;`
 2. `plot3(X,Y,Z);`
 3. `xlabel('locX'); ylabel('locY'); zlabel('power');`
 4. `xlim([0 1]); zlim([0 40]);`

More Plotting – Surface

- 'plot' functions produce line plots. Suppose we wish to display a 3D surface
- Keywords: `meshgrid()`, `mesh()`, `surf()`
- The command `meshgrid(X,Y)` creates a grid (domain) with all combinations of $\{X,Y\}$ elements. `mesh(X,Y,Z)` / `surf()` are then used to plot the surface $Z = f(X,Y)$ over the grid

3D Surface – Example

- Example:

1. `X = 1:4; Y = 1:3;`

2. `[Xg, Yg] = meshgrid(X,Y);` % produce a grid

3. `Z = Xg.^2 + Yg.^2;` % Z = f(X,Y)

4. `mesh(X,Y,Z);` % plot

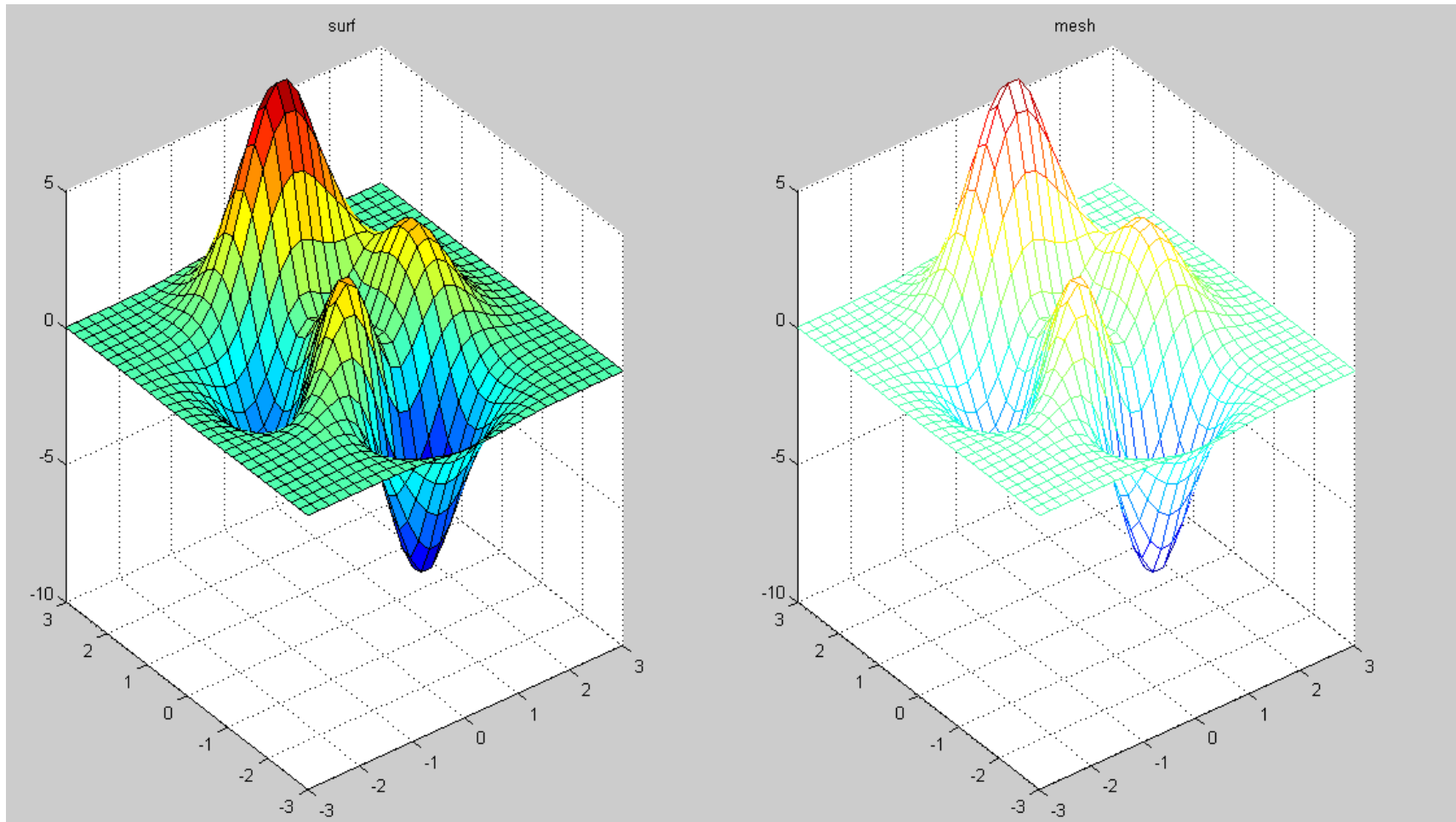
5. `surf(X,Y,Z);` % plot

- ‘mesh’ plots wireframe. ‘surf’ plots with shading.

- Can pass a 4th parameter specifying color e.g. “`surf(X,Y,Z,C)`”.
Otherwise, color is proportional to mesh height.

3D Surface – Example

- surf vs. mesh



Exercise (In Class)

- Sets & handles
- Write a function with 3 input parameters {A, B, fh}, which applies the function stored in the handle 'fh' to the vectors/matrices stored in {A,B}. You can assume that the function in 'fh' can only be one of the “set” methods (union, intersect, setdiff, ismember)

Sets

- Relevant functions: `unique()`, `intersect()`, `union()`, `setdiff()`, `ismember()`
- General form:
 - `<function>(A, B)` % Applies function to vectors {A,B}
 - `<function>(A, B, 'rows')` % Applies function to matrices {A,B}
 - `[C, IA, IB] = <function>(...)` % Returns combination & indices
- Rules:
 - If {A,B} are matrices, must supply the 'rows' parameter
 - {A,B} must have the same number of columns
- `ismember(A, B)`: returns a 0/1 array with the size of A. Where element is set to 1 if element is in B

Sets – Examples

- Define: $V1 = [1,2,3,4]$; $V2 = [0,1,2]$;
- Define: $A1 = [1,0;1,1;1,2]$; $A2 = [1,1;1,1;1,2]$;
- Examples:
 1. `ans = intersect (V1, V2)`
 2. `ans = union (V1, V2)`
 3. `ans = union (A1, A2)`
 4. `ans = setdiff(V1, V2)`
 5. `ans = ismember (A1, 1)`
 6. `ans = ismember (V1, [-1,2,4])`

Function Handles @

- Function handle: a variable that stores an identifier for a function
 - `h1 = @min;` % 'h1' can now be used instead of 'min'
 - `val = h1(rand(3))` % Same as `val = min(rand(3))`
- Can use handles in cell arrays or structs, but not in regular arrays:
 - `C = {@min, @max, @mean};`
 - `S.a = @min; S.b = @max; S.c = @mean;`
 - ~~`A = [@min, @max];`~~ **WRONG!**

Function Handles @

- Purpose of function handles:
 - Suppose the user should be able to choose which sub-routine to use. Thus, a mechanism to pass a parameter which specifies the sub-routine is required
 - Can't pass a function, so pass a handle instead
 - Also possible to define functions on the fly (anonymous functions):
 1. `sqr = @(x) x.^2;`
 2. `a = sqr(5);`

Exercise (In Class)

- Variable input arguments
- Write a function with *at least* 3 input parameters {A, B, fh1, fh2, fh3...}, which applies the functions stored in the handles 'fh1' (fh2, fh3...if present) to the vectors/matrices stored in {A,B}. You can assume that the function in 'fh#' can only be one of the “set” methods (union, intersect, setdiff, ismember)

Variable Input Arguments

- Function which can have any number of input args
- Rules:
 - Keyword: `varargin`
 - 'varargin' is a cell array containing the optional arguments ONLY
 - Must be declared as the last input argument
 - Must be lowercase
 - Example: "function my_fun(x,y,z,varargin)"
 - Similarly, use `varargout` to create a function with any number of output args. The same rules apply (last, lowercase)

The MATLAB Heist

- Suppose you are trying to steal a valuable object from a room. The room is protected by a laser grid which functions as a cloak of invisibility. Each laser dot is connected to every other dot, and each triple of dots ensures that the triangle defined by them is invisible
- You have a map which specifies the location of each laser dot in the room. Your goal is to turn off parts of the grid until the object becomes visible. To turn off a part, you need to “unravel” each dot in the triangle

The MATLAB Heist

- However, unraveling some dots requires more time than others. This is reflected by a value (in seconds) attached to each dot.
- Moreover, you need to unravel a given dot each time you wish to turn off a triangle for which it is a vertex
- Not to be misled, you should check whether the area you are searching is indeed visible
- You can stop when the object is found. That is, you may not need to turn off every triangle

Heist – Problem Formulation

- More specifically:
 - The room is defined by the 2D square $[0.0,1.0],[0.0,1.0]$
 - The grid is defined by a set of N laser dots ($N \geq 3$)
 - Each dot 'n', requires $C(n)$ seconds to unravel
 - The location of each dot 'n' is specified by two coordinates $\{x_n, y_n\}$
Both coordinates are in the range $[0.0,1.0]$
 - To visualize the grid, we can connect every pair of dots with a straight line
- Given the complete grid, how do you choose which parts of the grid to turn off? How much time does it take to find the object?

Heist – Problem Formulation

- Assumptions:
 - M is an N -by-2 ($N \geq 3$) matrix specifying the grid
 - Columns $\{1,2\}$ of M are the $\{x,y\}$ coordinates of the dots respectively
 - C (time) is an N -by-1 (column) vector of positive real values
 - M consists of real values in the range $[0.0,1.0]$
- All the relevant info is given. That is, $\{M,N,C\}$ must be known
- Note: no guarantee that the parameters are set correctly
- Note: there is no a priori information about the location of the object in the room

1st Step – Approach

- Approach:

1. Collect and verify all parameters {M,N,C}
2. Generate some plots: grid (2D), grid vs. time (3D), progress (2D triangles)
3. Decide on a search strategy:
 - I. Determine whether the starting point is in the convex hull for the grid.
 - II. Choose which triangle to turn off (e.g. minimize effort or maximize triangle area)
 - III. Check whether the area you are searching is visible
 - IV. Update your strategy if necessary and keep track of progress
4. Evaluate your strategy in hindsight

Convex Set

- *Definition:* a set of points C is convex, if the line segment between any two points in C lies in C
- Formally, if for any $\{x,y\}$ in C and any $0 \leq w \leq 1$, $wx + (1-w)y$ is in C , then C is convex
- Implicitly assuming that C is a subset of the n -dimensional real space
- We can generalize the definition to hold for more than two points

Convex Hull

- *Definition:* we call a point 'p' of the form $p = w_1x_1 + w_2x_2 + \dots + w_kx_k$, where $w_1 + w_2 + \dots + w_k = 1$ and $w_i \geq 0$ ($\forall i$), a *convex combination* of the points $\{x_1, \dots, x_k\}$.
- *Convex combination:* weighted average of the points
- *Definition:* the set of all convex combinations of points in C is called the *convex hull* of C
- *Convex hull:* The smallest convex set that contains C
- These notions can be applied to infinite sums and non-Euclidean spaces as well

2nd Step – Concrete Tasks

- Problems & relevant functionality:
 1. Generate plots: grid (`plot`, `triplot`), grid vs. time (`plot3`)
 2. Compute convex hull: use grid data (`convhull`)
 3. Choose a triangle to turn off: maximize area (`polyarea`) or minimize effort by recording every triangle and value and then sorting (`sortrows`)
 4. Check if area is visible: verify whether a given location is inside a triangle (`inpolygon`)
 5. Keep track of progress: mark visible triangles (`fill`)

2nd Step – Remark

- Built-in functionality simplifies work considerably
- The alternative is to define and solve linear equations (`mldivide`, `linprog`)

Plotting Grid & Triangles

- Recall that M is an N -by-2 ($N \geq 3$) matrix specifying the coordinates of the grid. Every pair of dots is connected
- To plot every triangle, we first need to create a list of triples. Each triple specifies three (non-collinear) dots on the grid
- Built-in function to plot triangles: `triplot()`
- Built-in function to compute triples without loops: `combntns()`

combntns()

- **Syntax:** `combos = combntns(set, subset)`
- **Action:** returns a matrix whose rows are the various combinations of elements from vector 'set'. Each combo (row) is of length 'subset'
- **Example:**
 - `combos = combntns (1:3,2) % "3 choose 2"`
 - `combos:`

```
[ 1 2  
  1 3  
  2 3 ]
```

triplot()

- **Syntax:** `triplot(TRI, x, y)`
 - `triplot(TRI, x, y, color)`
 - `triplot(TRI, x, y, 'param', 'value')`
- **Action:** displays the triangles defined in the N-by-3 matrix TRI. A row of TRI contains indices into the vectors x,y that define a single triangle. The default line color is blue
- **Example:**
 - `x = rand(5,1); y = rand(5,1);`
 - `combos = combntns(1:5,3);` % "5 choose 3"
 - `triplot(combos, x, y);`

3rd Step (Code) – Plots

- Code to generate plots:

1. `X = M(:,1); Y = M(:,2);` `% grid dots`
2. `figure(1);` `% Create figure`
3. `plot(X,Y, 'r.', 'MarkerSize', 20);` `% plot dots`
4. `figure(2);`
5. `plot3(X,Y, M(:,3), 'r.');` `% plot dots vs. time`
6. `figure(1); hold on;` `% Reset to figure 1`
7. `LT = combntns(1:length(X), 3);` `% List all triangles`
8. `triplot(LT, X, Y, 'black');` `% Plot triangles`

Computing Convex Hull

- Given a set of points in N-D, computing the hull is not a trivial task. Several algorithms are available
- The optimal algorithm may depend on the properties of the points
- Built-in functions: `convhull()`, `convhulln()`
- The hull is specified by a set of 'outer boundary' points

convhull()

- **Syntax:** `[CH,V] = convhull(X,Y)`
- **Action:** returns the 2D convex hull CH of the points (X,Y), where X and Y are column vectors, and the corresponding area/volume V bounded by K. 'CH' is a vector of point indices arranged in a counter-clockwise cycle around the hull
- **Example:**
 - `x = [0,1,0.5,1,0]'; y = [0,0,0.5,1,1]';`
 - `[CH, A] = convhull(x,y); % CH = [1,2,4,5,1]', A = 1.0 (area of square)`
 - `plot(x(CH),y(CH),'r-', x,y,'b+')`

Polygon Interior Check

- *Polygon*: a plane shape consisting of straight lines that are joined together to form a circuit
- The convex hull for a set of 2D points is a polygon. A triangle is the simplest polygon
- To check whether a point is inside the hull/triangle we need to define the polygonal region and check if the point is in the interior
- Built-in function to check: `inpolygon()`
- Another approach: solve a set of linear equations using `linprog()`

inpolygon()

- **Syntax:** `IN = inpolygon(X,Y, Px,Py)`
 - `[IN ON] = inpolygon(X,Y, Px,Py)`
- **Action:** returns a 0/1 matrix 'IN' the same size as X & Y. `IN(k) = 1` if `{X(k), Y(k)}` is inside the polygon or on the boundary. The polygon vertices are specified by the vectors `{Px, Py}`
- **Example:**
 - `Px = [0,1,1,0]'`; `Py = [0,0,1,1]'`; % Polygon vertices
 - `X = [0.1,0.2,0.5,1]`; `Y = [0.1,0.2,1.1,0.8]`; % Points to test
 - `IN = inpolygon (X,Y,Px,Py)`; % `IN = [1,1,0,1]`

3rd Step (Code) – Interior Check

- Code to compute/plot convex hull and check interior:
 1. `X = M(:,1); Y = M(:,2);` % grid dots
 2. `figure(1); hold on; prop = 'MarkerSize';` % create figure
 3. `plot(X,Y, 'r.', prop, 20);` % plot dots
 4. `[CH, A] = convhull (X,Y);` % compute hull
 5. `Px = X(CH); Py = Y(CH);` % polygon vertices
 6. `plot(Px,Py, 'g-', X,Y, 'r.', prop, 20);` % plot hull
 7. `IN_H = inpolygon (lx,ly,Px,Py);` % {lx,ly} starting point
 8. `LT = combntns(1:length(X), 3);` % LT is the list of TRG
 9. `Tx = X(LT(k,:)); Ty = Y(LT(k,:));` % triangle vertices
 10. `IN_T = inpolygon (lx,ly,Tx,Ty);` % check if inside triangle

Hull Interior Check – SLE

- Notice that every point inside the convex hull is a convex combination (weighted avg.) of the points for which the hull was computed
- In other words, if $pt = \{ptx, pty\}$ is inside the hull, then for some set of weights $\{w_1 \dots w_k\}$:
 1. $ptx = w_1x_1 + w_2x_2 + \dots + w_kx_k$
 2. $pty = w_1y_1 + w_2y_2 + \dots + w_ky_k$
 3. $w_1 + w_2 + \dots + w_k = 1$
 4. $\forall i, w_i \geq 0$
- Need to solve a system of linear equations (SLE) with lower bounds on the variables

Hull Interior Check – Notation

- Standard notation:

1. $ptx = w_1x_1 + w_2x_2 + \dots + w_kx_k$
2. $pty = w_1y_1 + w_2y_2 + \dots + w_ky_k$
3. $w_1 + w_2 + \dots + w_k = 1$
4. $\forall i, w_i \geq 0$

- Vector notation:

1. $X^TW = ptx$
2. $Y^TW = pty$
3. $\mathbf{1}^TW = 1$
4. $W \geq 0$

- Solver form:

1. $[X; Y; \mathbf{1}]^TW = [ptx; pty; 1]$
2. $0 \leq W$

Solving SLE

- Given a system of linear equations $Ax = b$, does it have a solution?
- Depends on the matrix A :
 1. If A is a square matrix: $x = A^{-1}b$
 2. Else we have an under/over-determined system
 3. Could have multiple/no solutions for either case
- To compute the inverse of A efficiently, can decompose the matrix. Multiple ways of doing this (LU,QR)
- When A is not square, can minimize $\text{norm}(A^*X - B)$ i.e., the length of the vector $AX - B$. This is the least squares solution

mldivide (\)

- **Syntax:** $x = A \backslash b$

- $x = \text{mldivide}(A,b)$
- $x = \text{inv}(A)*b$

- **Action:** If A is a square matrix, $A \backslash b$ is roughly the same as $\text{inv}(A)*B$. Otherwise, $x = A \backslash b$ is the least squares solution. A warning message is displayed if A is badly scaled or nearly singular

- **Example:**

- $A = [X; Y; 1]^T$; $b = [ptx; pty; 1]$; % Hull interior equations
- $x = A \backslash b$; % Solution to SLE

Choosing a Strategy

- We need a strategy (set of rules) to select which triangle to turn off next
- *Greedy algorithm* makes the decision that gives the maximum benefit in the immediate next step (locally optimal). This decision might not be the best considering more (all) steps (globally optimal)
- Greedy strategy 1: choose the largest triangle
 - Built-in function to compute area: `polyarea()`
- Greedy strategy 2: choose the least-effort triangle
 - Sum the time to unravel the vertices for each triangle

polyarea()

- **Syntax:** $A = \text{polyarea}(X,Y)$
- **Action:** returns the area of the polygon specified by the vertices in the vectors X and Y . If X and Y are matrices of the same size, then the area is computed for each column (polygon) of $\{X,Y\}$
- **Example:**
 - $\text{TSx} = X(\text{LT}); \text{TSy} = Y(\text{LT});$ % Vertices for every triangle
 - $\text{T_area} = \text{polyarea}(\text{TSx}', \text{TSy}');$ % Area for every triangle

3rd Step (Code) – Strategy

- Code to implement strategy:

1. `X = M(:,1); Y = M(:,2); T = M(:,3);` % grid dots & time
2. `LT = combntns(1:length(X), 3);` % LT is the list of TRG
3. `TSx = X(LT); TSy = Y(LT);` % vertices \forall TRG
4. `T_area = polyarea(TSx', TSy');` % area \forall TRG
5. `Tot_effort = sum(T(LT),2);` % total effort \forall TRG
6. `choice = [LT, T_area', Tot_effort];` % Combine into one mtx
7. `top_choice = sortrows (choice,-4);` % sort rows by area
8. `top_choice = sortrows (choice,[-4,5]);` % sort by area then by time

fill()

- **Syntax:** `A = fill(X,Y,C)`
- **Action:** fills the polygon whose vertices are specified in `{X,Y}` with the constant color specified in `C` (`C` can be a single character string chosen from the list `{r,g,b,c,m,y,w,k}` or an RGB row vector triple, `[r g b]`)
- **Example:**
 - `fill(X(LT(1,:)),Y(LT(1,:)), 'm');` % fill one triangle
 - `fill(X(LT)',Y(LT)','m');` % fill every triangle

3rd Step (Code) – Tracking

- Code to track progress:

1. `X = M(:,1); Y = M(:,2); T = M(:,3);` % grid dots & time
2. `LT = combntns(1:length(X), 3);` % LT is the list of TRG
3. `figure(1); hold on;` % create figure
4. `fill(X(LT)',Y(LT)', 'm');` % fill every triangle
5. `triplot(LT, X, Y, 'black');` % Plot triangles
6. `[CH, A] = convhull (X,Y);` % compute hull
7. `plot(X(CH),Y(CH), 'g-', X,Y, 'r.', prop, 20);` % plot hull

Conclusion

- *High-level language*: built-in functionality at your fingertips. Easy to bypass even the most trivial mathematical problems e.g. how do you compute triangle area? I don't know, but I don't care, can just use `polyarea()`
- Convenient platform for *prototyping*. May wish to design an elaborate game. Can add routines one after another while verifying outcome using plots. All this does not require considerable time investment