

Problem Set 8

Due: Thur, 04/02/09.

Reading: Chapter 3

- Prove that the class of Turing-recognizable languages is closed under the concatenation operation.
 - Prove that the class of Turing-decidable languages is closed under the complementation operation.
 - Does your proof of part (a) work to prove the same claim for the class of Turing-decidable languages? Does your proof of part (b) work to prove the same claim for the class of Turing-recognizable languages? Explain your answers.

- Let $NE = \{\langle M \rangle \mid M \text{ is a TM that accepts some string}\}$ (namely all $\langle M \rangle$ that recognize a non-empty language).

Prove that NE is Turing-recognizable (show an algorithm to recognize it, and argue the correctness of your algorithm).

- Prove that a language L is decidable if and only if there is some enumerator that enumerates the language in lexicographic order.
- Recall that a non-deterministic TM (NTM) is called a decider, if it always halts (all computations, on all inputs). In this problem you will show that having some halting computation on every input, does not necessarily mean the NTM is a decider.

Assume L is some language that is TM-recognizable but not TM-decidable. Construct a non-deterministic TM N recognizing L , such that N has the following properties: (1) For any input string w , there exists at least one computation of N on w that is halting; (2) N is not a decider.