



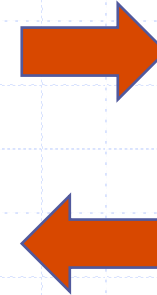
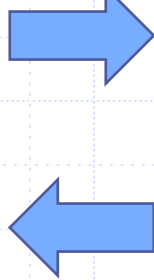
Network Security Protocols and Defensive Mechanisms

*Slides borrowed from John Mitchell

Network security

- ◆ What is the network for?
- ◆ What properties might attackers destroy?
 - Confidentiality : no information revealed to others
 - Integrity : communication remains intact
 - Availability : messages received in reasonable time

- Confidentiality
- Integrity
- Availability



Network Attacker

Intercepts and controls network communication

Plan for today

◆ Protecting network connections

- Wireless access– 802.11i/WPA2
- IPSEC

◆ Perimeter network defenses

- Firewall
 - ◆ Packet filter (stateless, stateful), Application layer proxies
- Intrusion detection
 - ◆ Anomaly and misuse detection



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Last lecture

◆ Basic network protocols

- IP, TCP, UDP, BGP, DNS

◆ Problems with them

■ TCP/IP

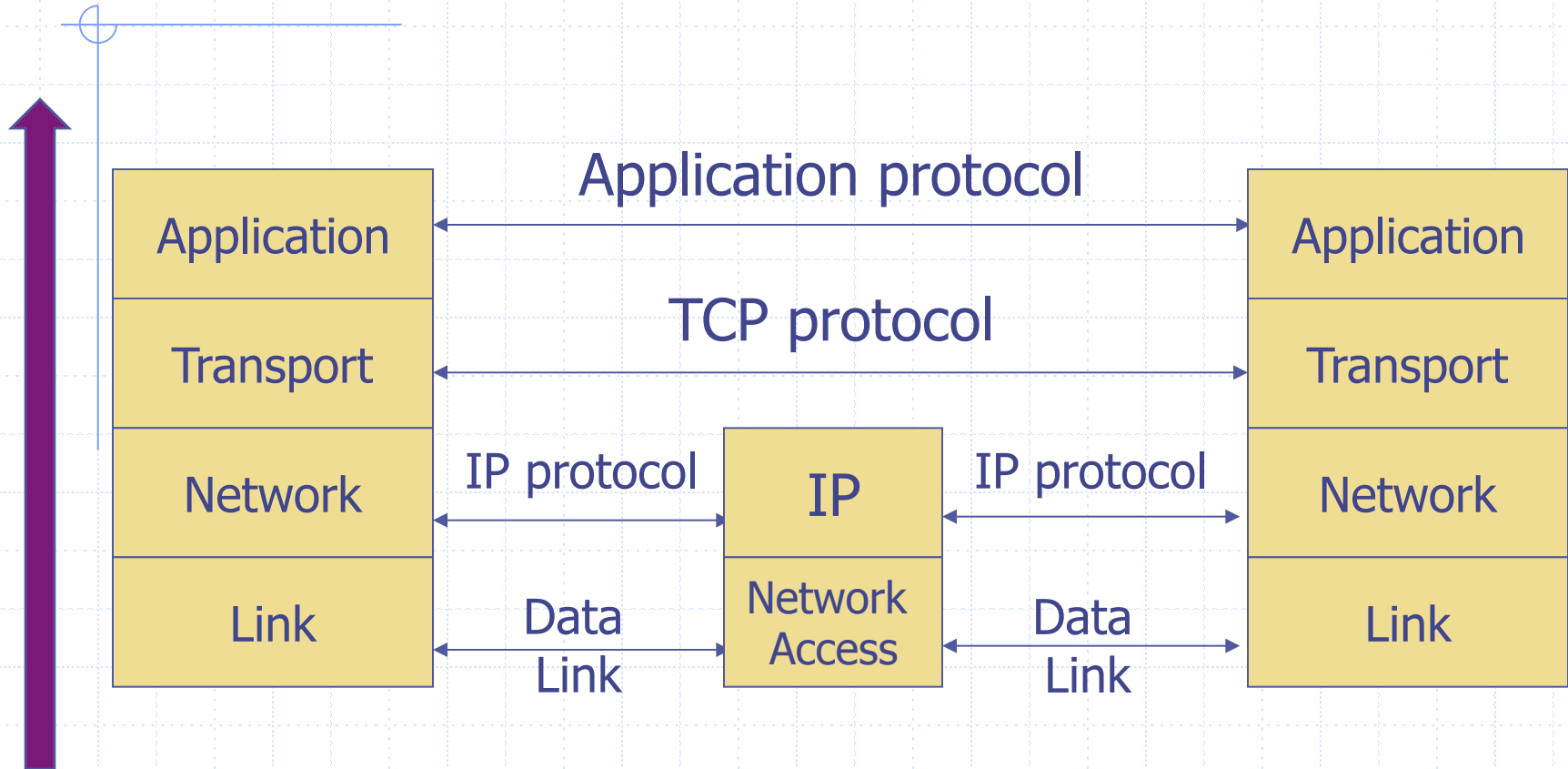
- ◆ No SRC authentication: can't tell where packet is from
- ◆ Packet sniffing
- ◆ Connection spoofing, sequence numbers

■ BGP: advertise bad routes or close good ones

■ DNS: cache poisoning, rebinding

- ◆ Web security mechanisms rely on DNS

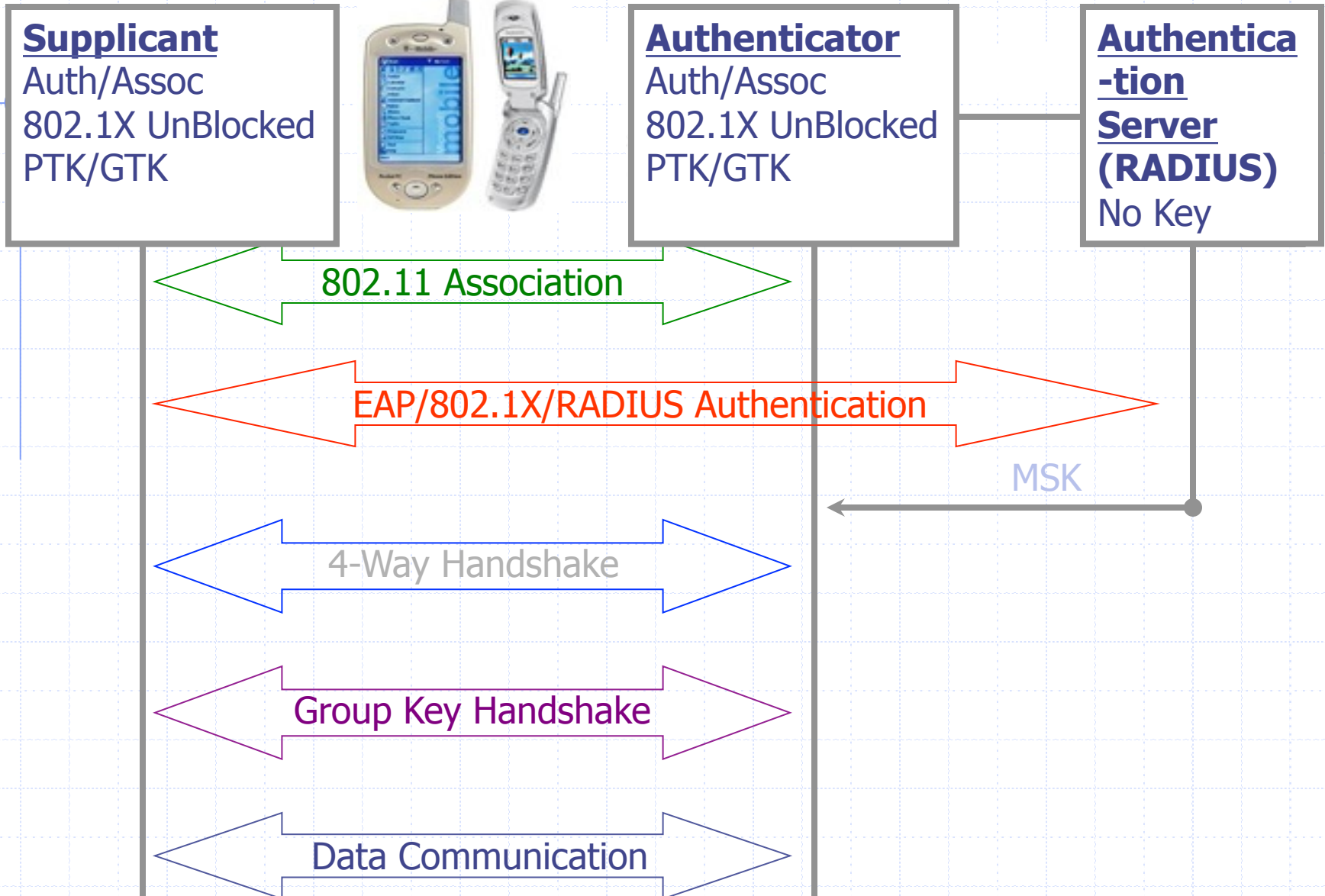
Network Protocol Stack



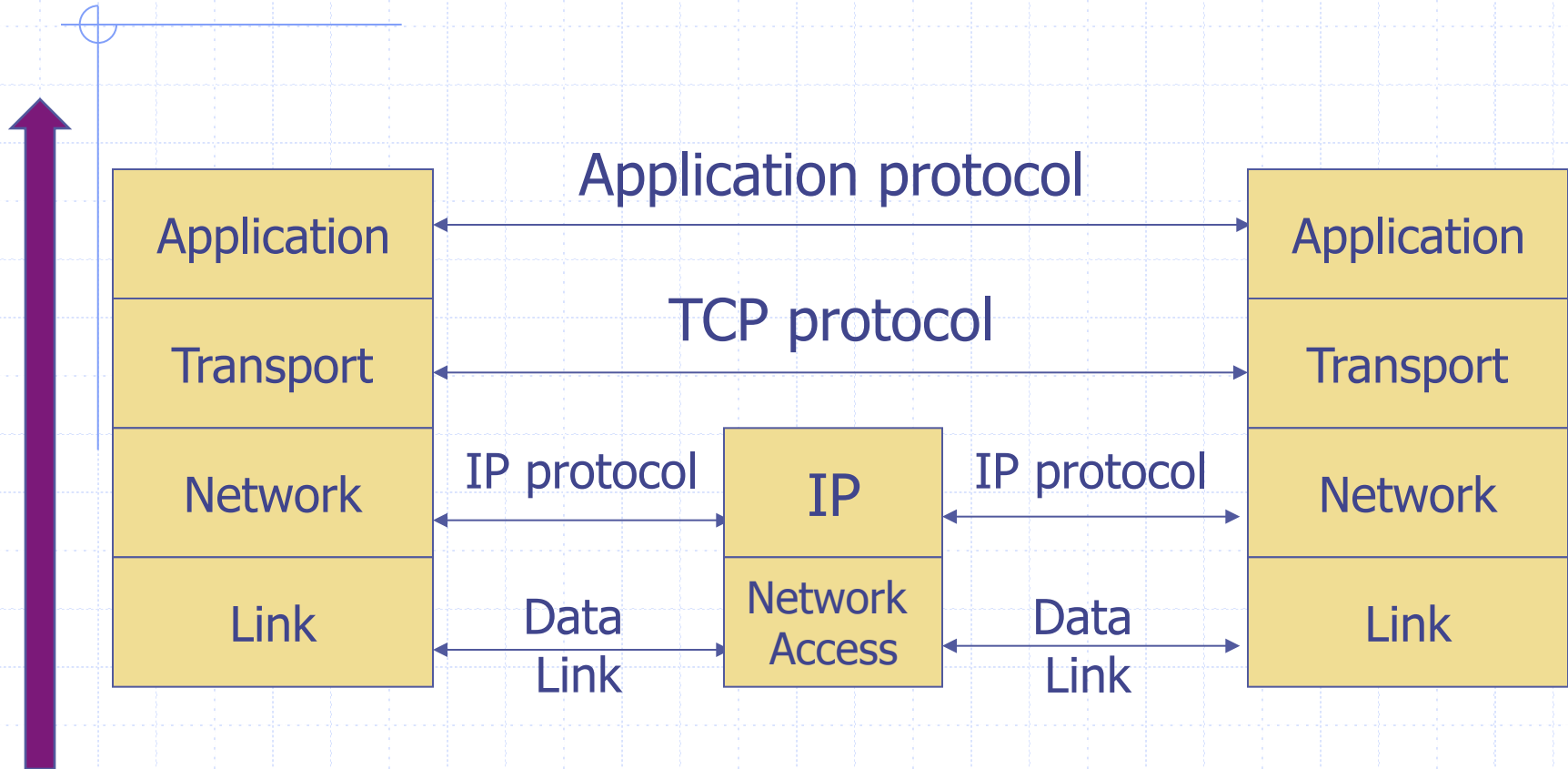


Protocol and link-layer connectivity

802.11i Protocol



Network Protocol Stack



TCP/IP CONNECTIVITY

How can we isolate our conversation from attackers on the Internet?

Basic Layer 2-3 Security Problems

- ◆ Network packets pass by untrusted hosts
 - Eavesdropping, packet sniffing
 - Especially easy when attacker controls a machine close to victim

- ◆ TCP state can be easy to guess
 - Enables spoofing and session hijacking

Virtual Private Network (VPN)

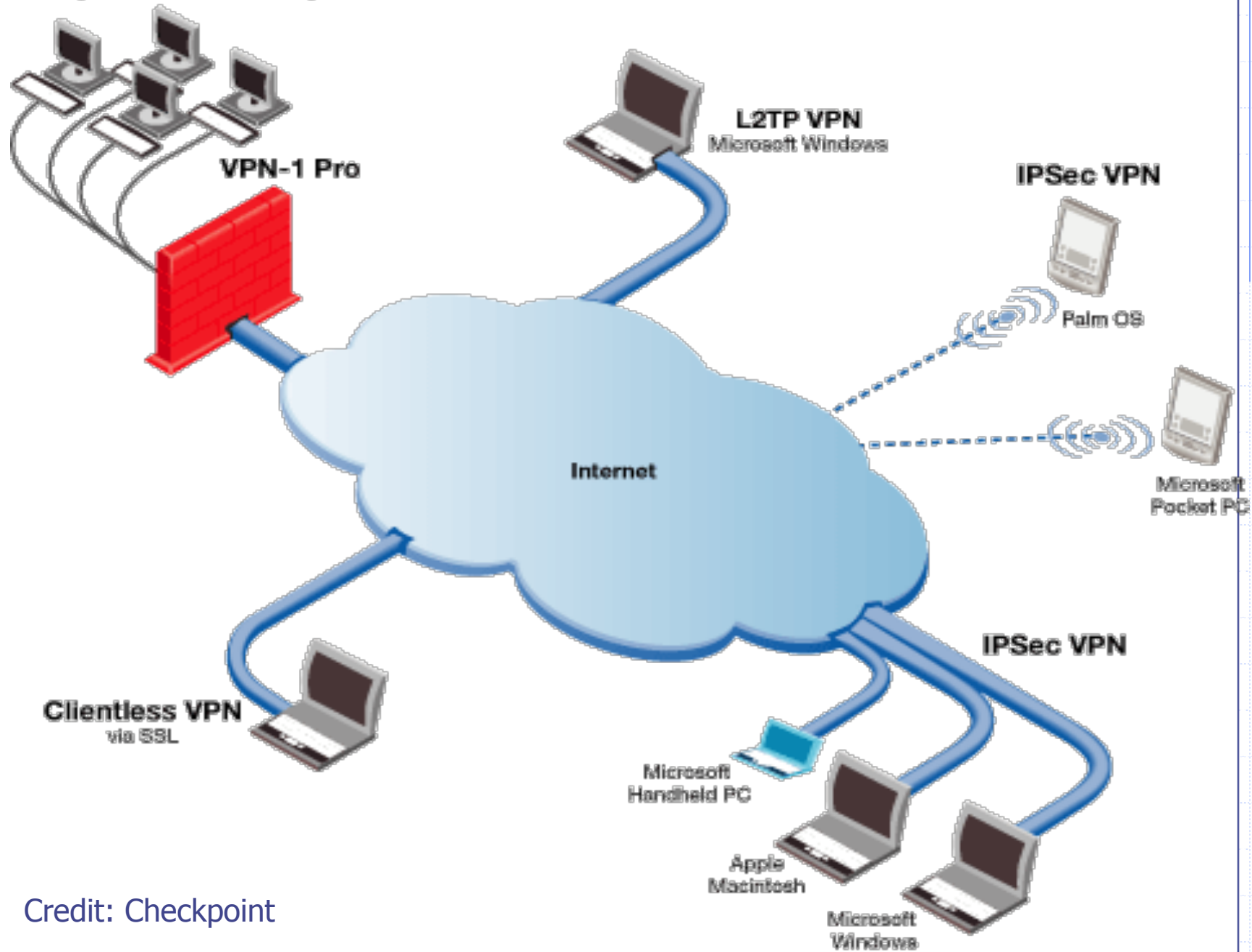
◆ Three different modes of use:

- Remote access client connections
- LAN-to-LAN internetworking
- Controlled access within an intranet

◆ Several different protocols

- PPTP – Point-to-point tunneling protocol
 - L2TP – Layer-2 tunneling protocol
 - IPsec (Layer-3: network layer)
- } Data layer

LAN (Trusted Network)

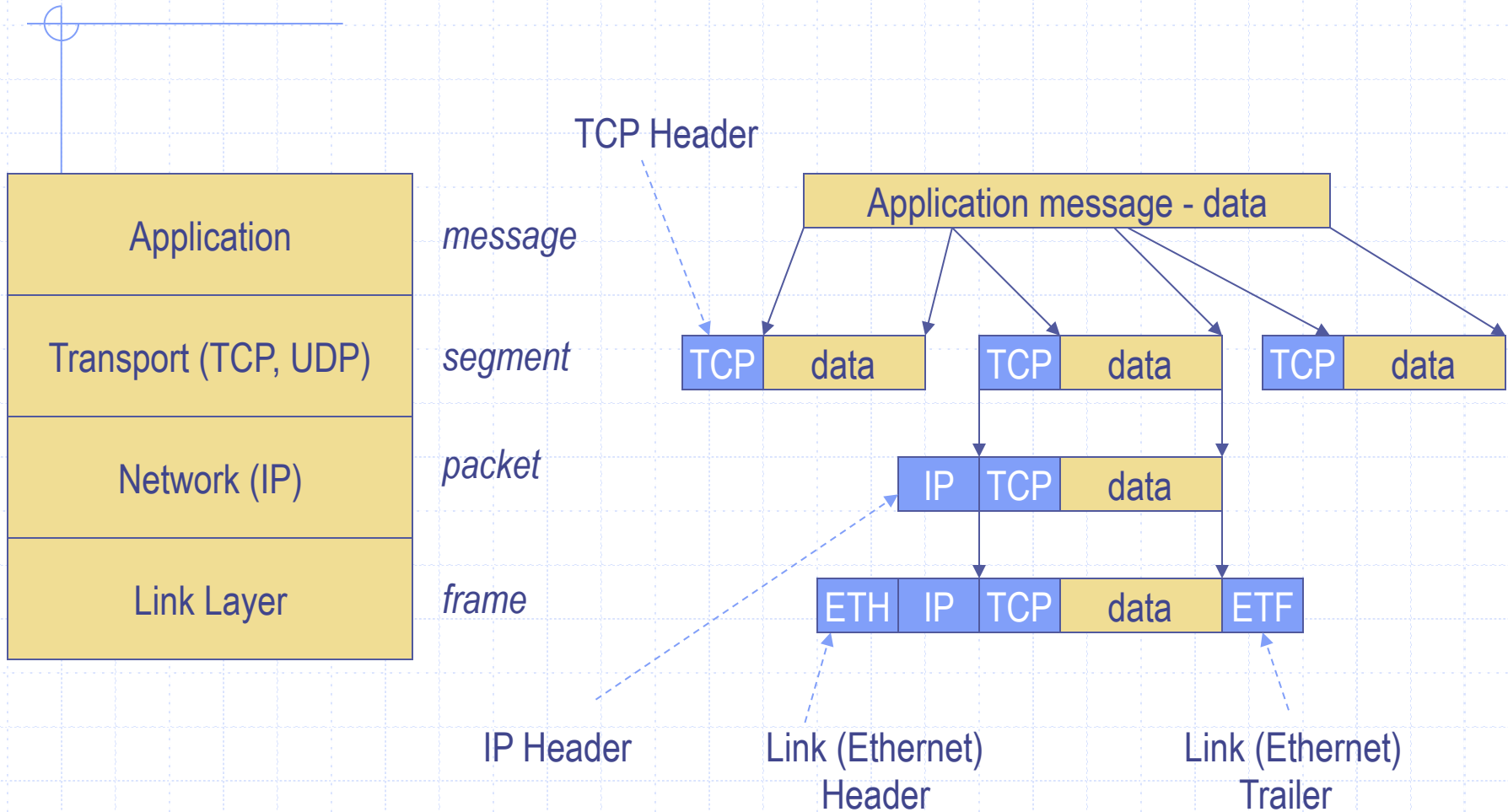


Credit: Checkpoint

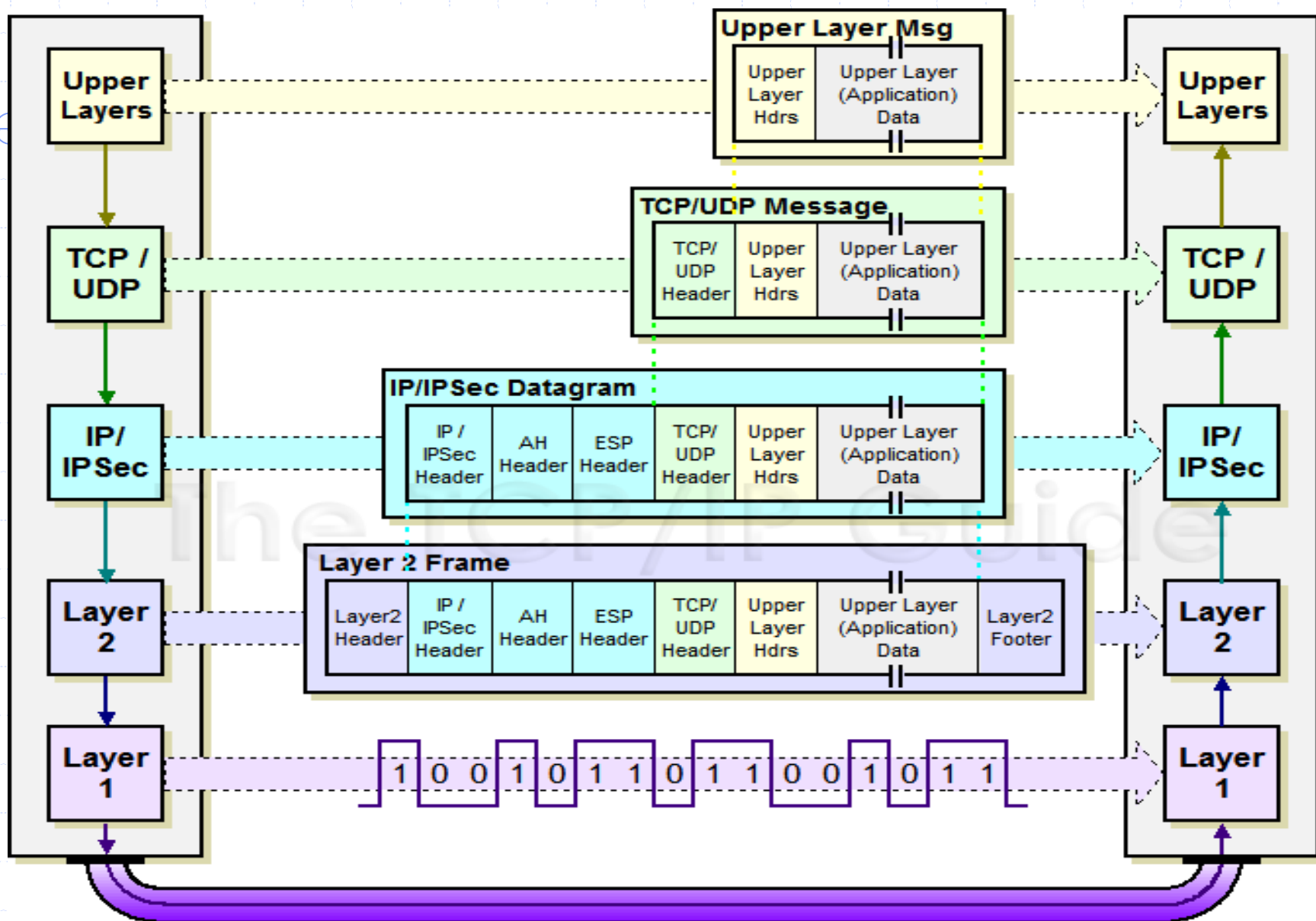
IPSEC

- ◆ Security extensions for IPv4 and IPv6
- ◆ IP Authentication Header (AH)
 - Authentication and integrity of payload and header
- ◆ IP Encapsulating Security Protocol (ESP)
 - Confidentiality of payload
- ◆ ESP with optional ICV (integrity check value)
 - Confidentiality, authentication and integrity of payload

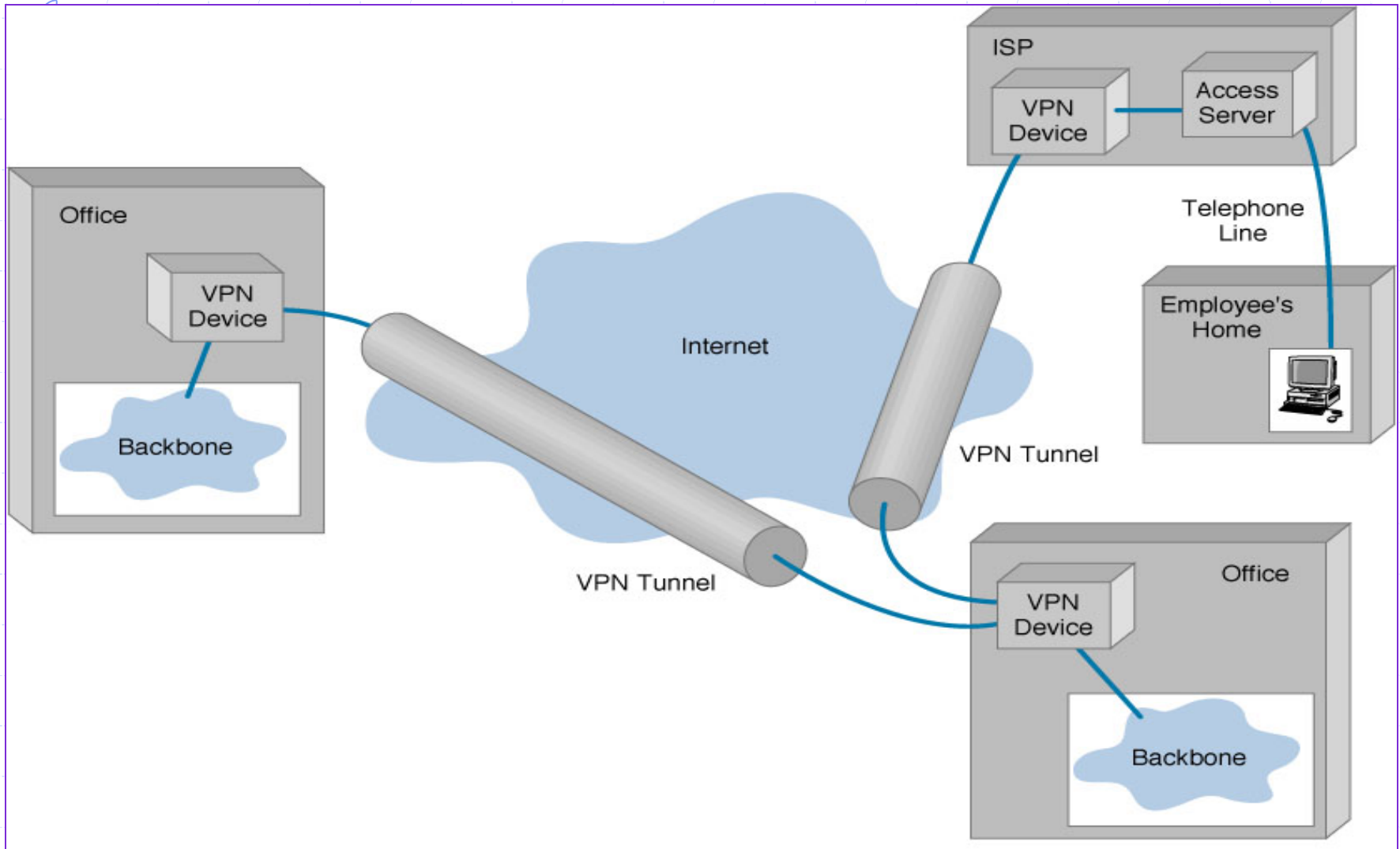
Recall packet formats and layers



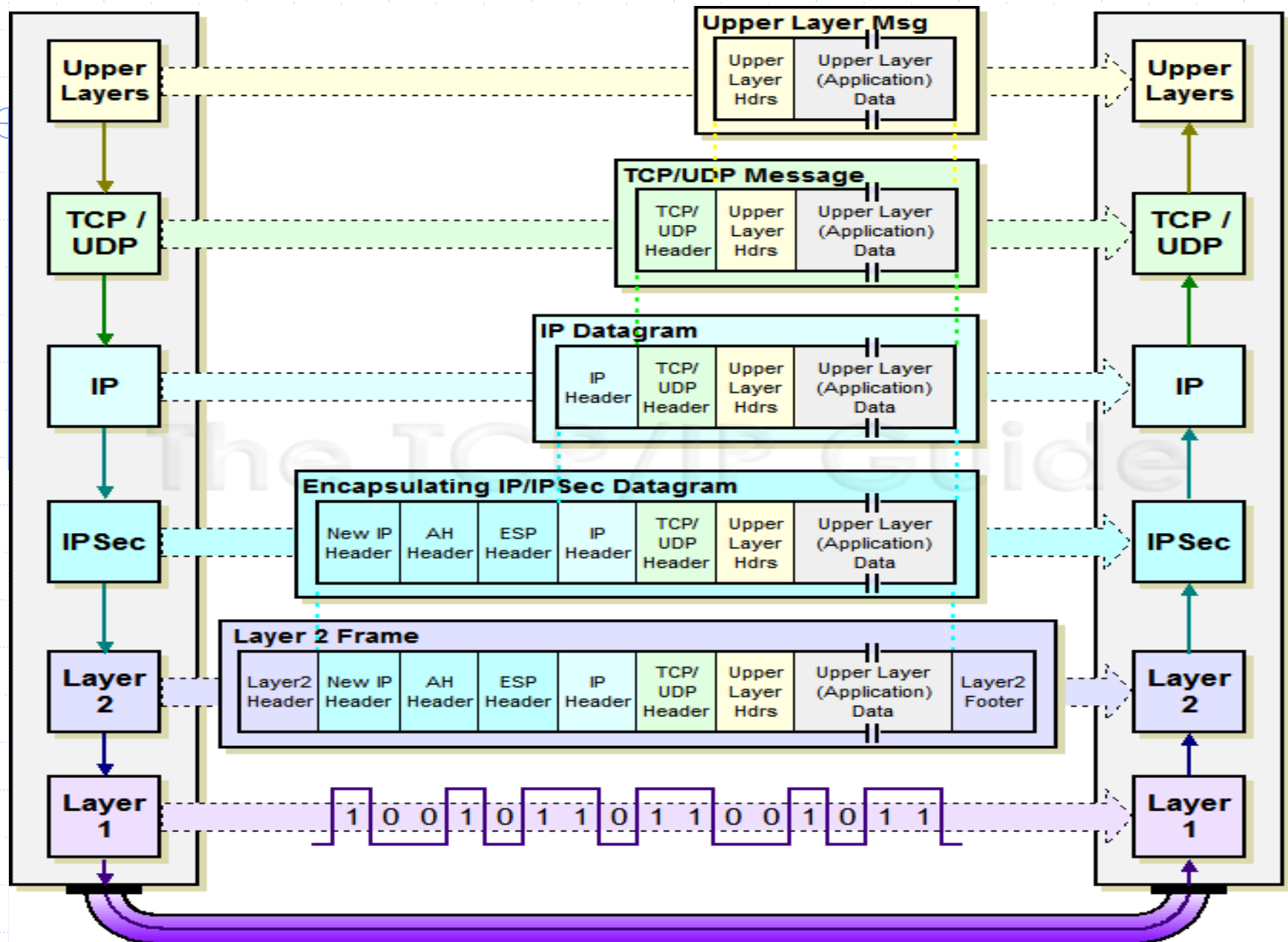
IPSec Transport Mode: IPSEC instead of IP header



IPSEC Tunnel Mode



IPSec Tunnel Mode: IPSEC header + IP header

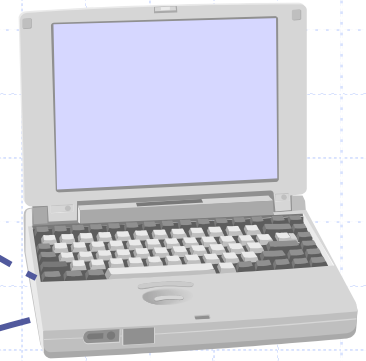


Mobile IPv6 Architecture

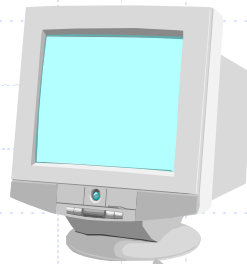
Mobile Node (MN)



Direct connection via binding update



Corresponding Node (CN)



Home Agent (HA)

- ◆ Authentication is a requirement
- ◆ Early proposals weak
- ◆ RFC 6618 – use IPsec

Summary

◆ Protecting network connections

- Wireless access– 802.11i/WPA2
 - ◆ Several subprotocols provide encrypted link between user device and wireless access point
 - ◆ Ideally – wireless attacker in range of access point has no better chance for attack than a remote attacker
- IPSEC
 - ◆ Give external Internet connections equivalent security to local area network connections
- Mobility
 - ◆ Preserve network connections when a device moves to different physical portions of the network
 - ◆ Ideally – no attacks other than against non-mobile user



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Second topic of today's lecture

◆ Perimeter defenses for local networks

- Firewall
 - ◆ Packet filter (stateless, stateful)
 - ◆ Application layer proxies
- Intrusion detection
 - ◆ Anomaly and misuse detection

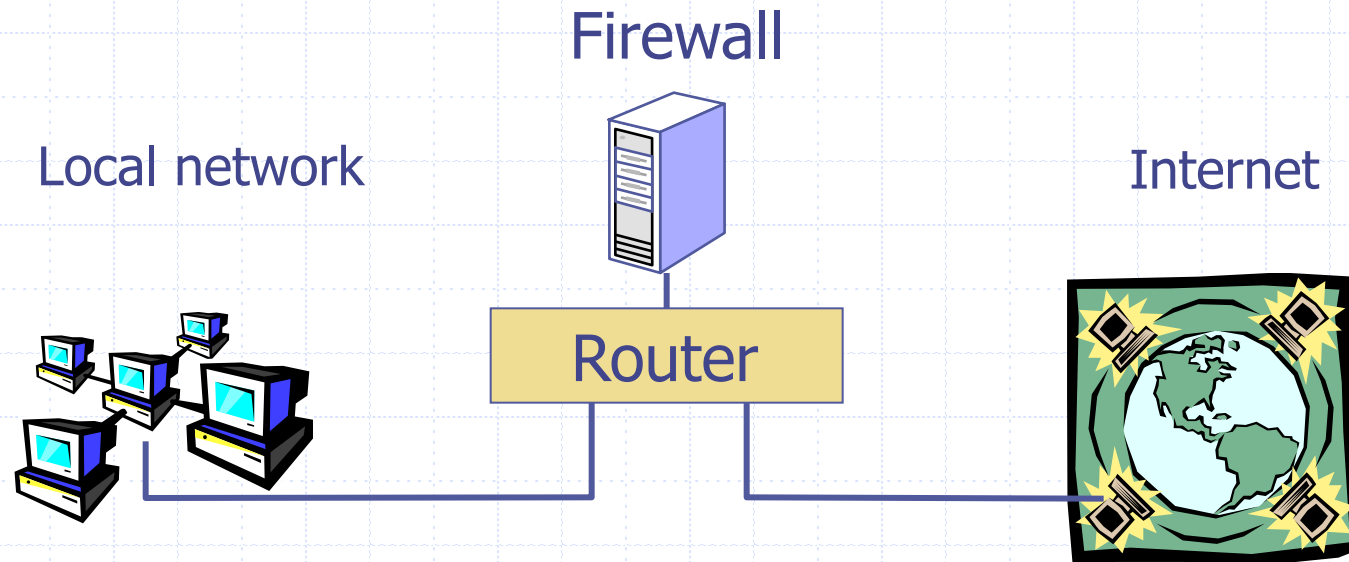


LOCAL AREA NETWORK

How can we protect our local area network from attackers on the external Internet?

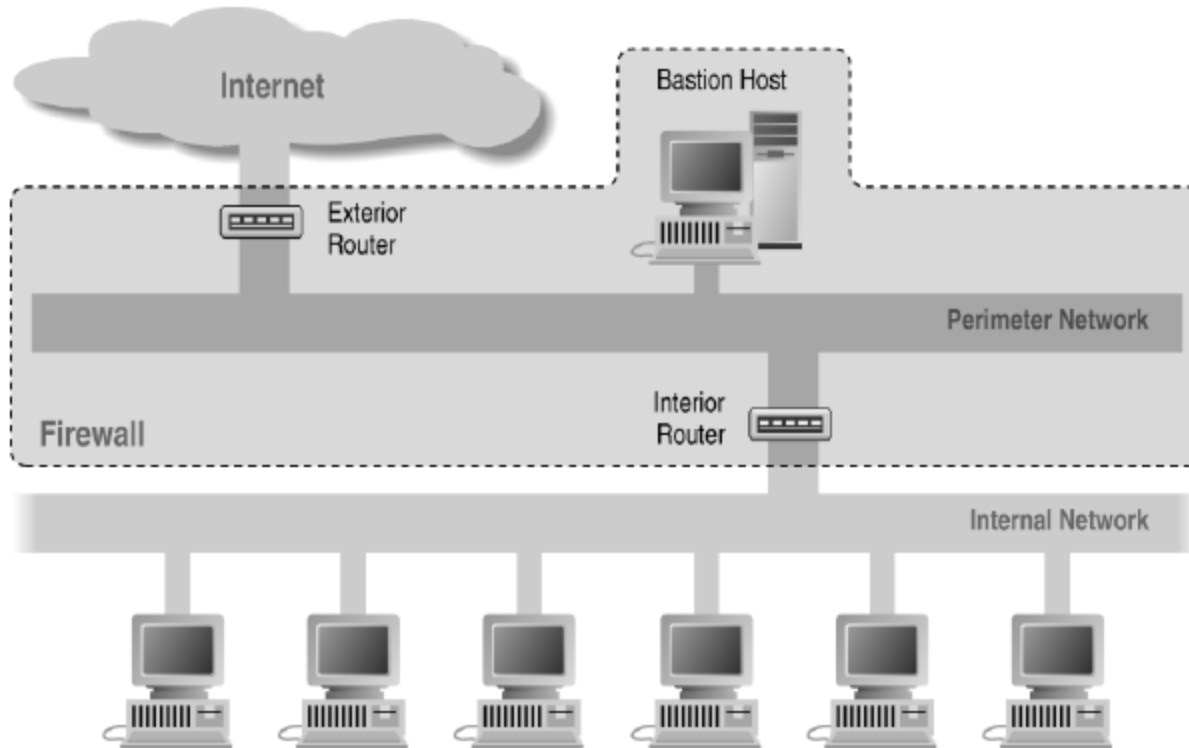
Basic Firewall Concept

- ◆ Separate local area net from internet

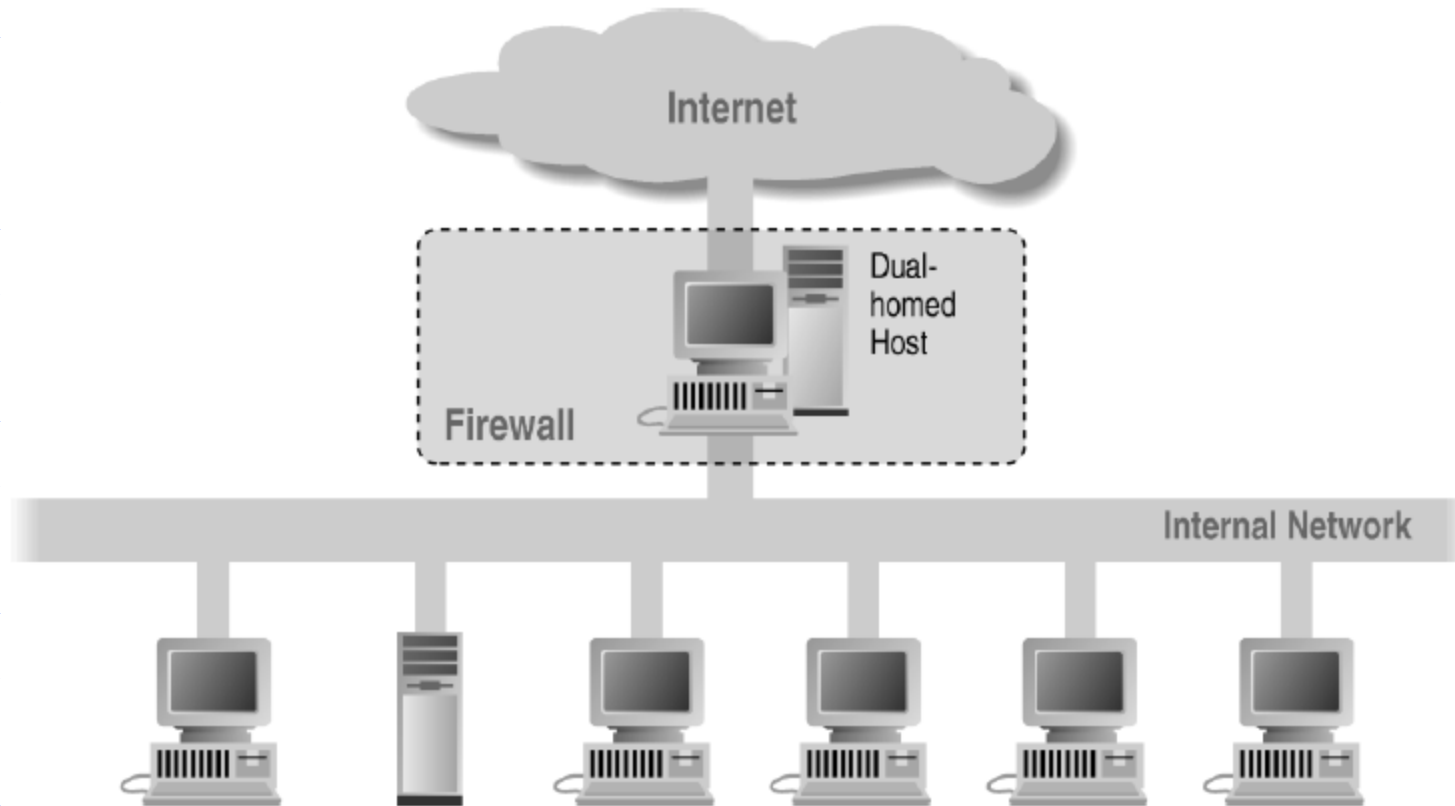


All packets between LAN and internet routed through firewall

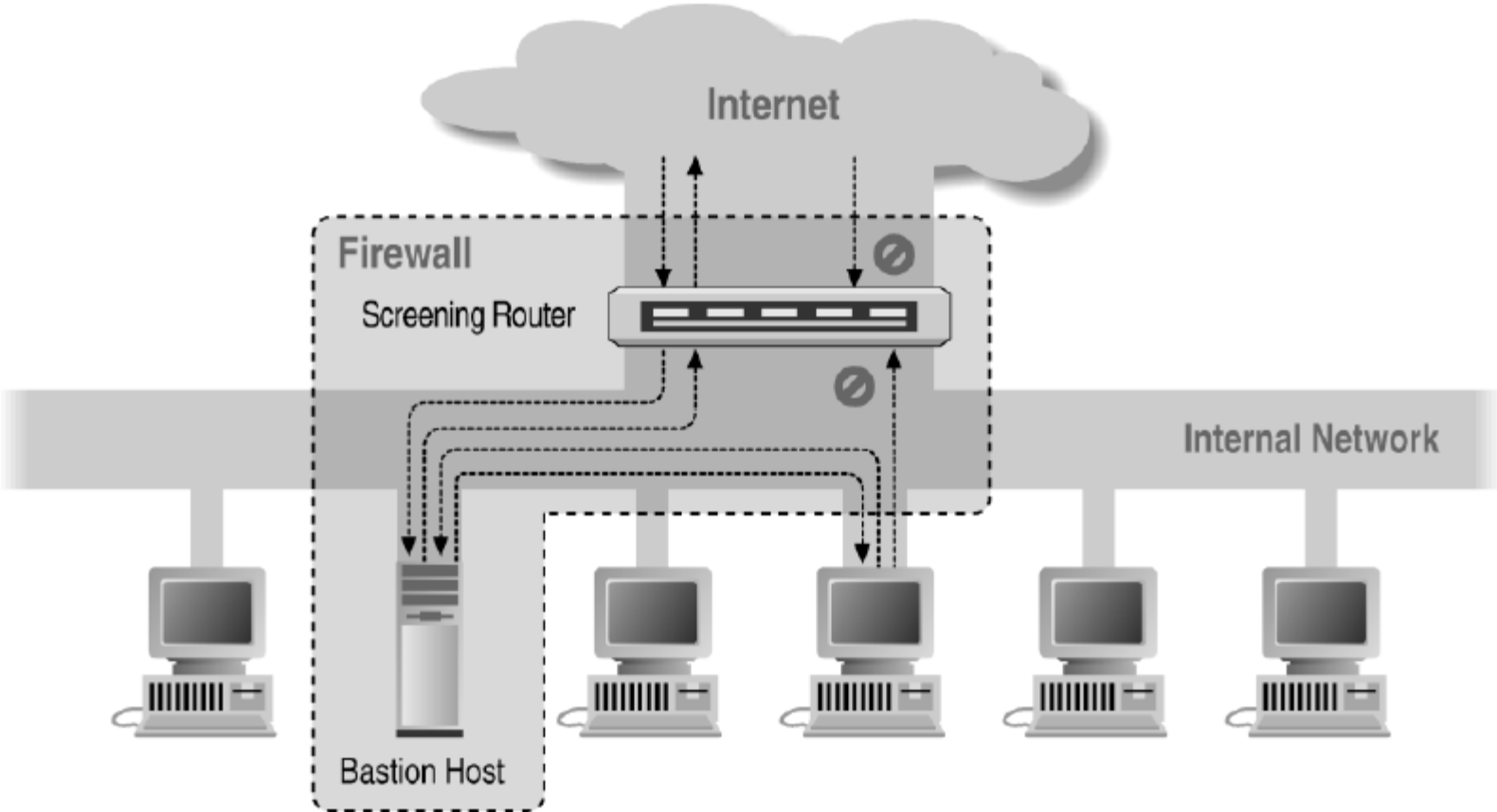
Screened Subnet Using Two Routers



Alternate 1: Dual-Homed Host



Alternate 2: Screened Host



Basic Packet Filtering

◆ Uses transport-layer information only

- IP Source Address, Destination Address
- Protocol (TCP, UDP, ICMP, etc)
- TCP or UDP source & destination ports
- TCP Flags (SYN, ACK, FIN, RST, PSH, etc)
- ICMP message type

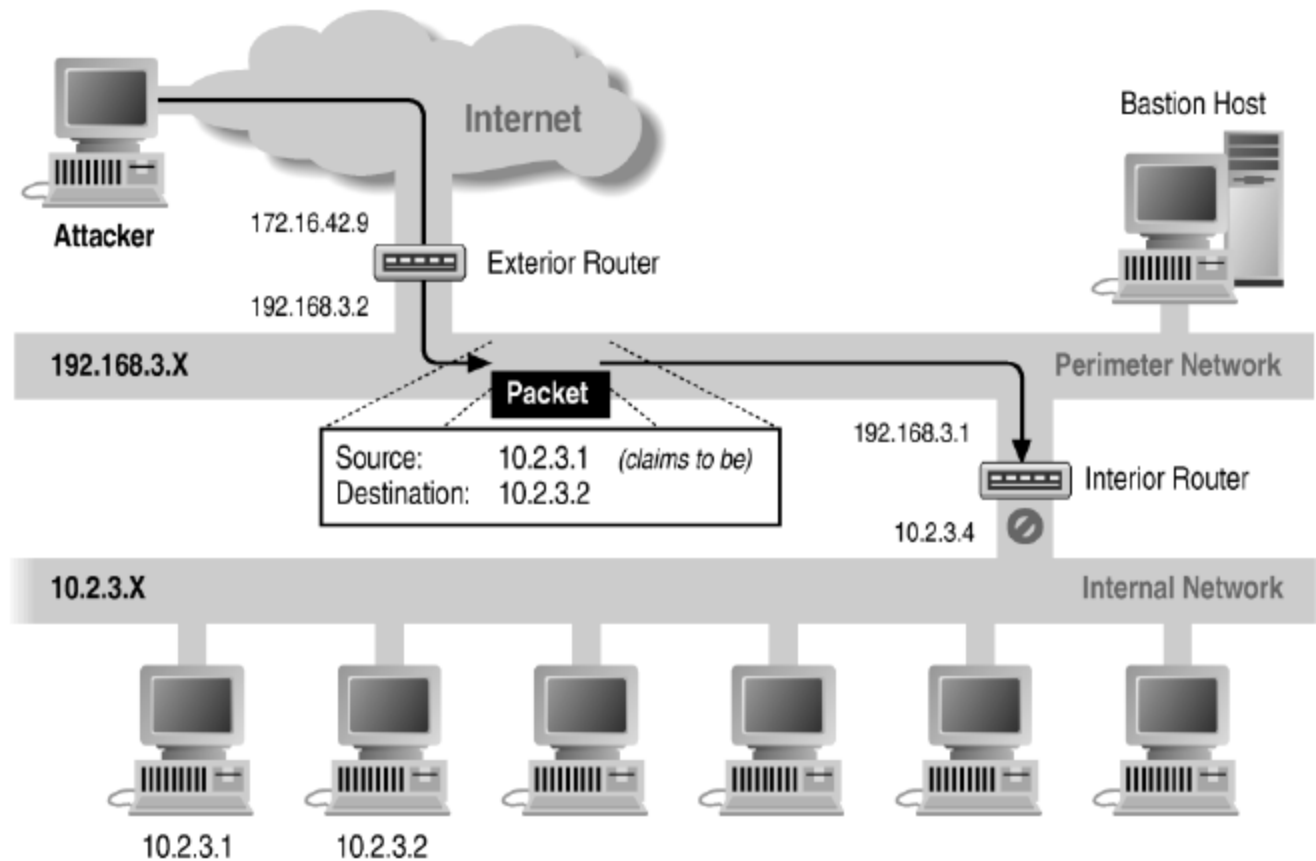
◆ Examples

- DNS uses port 53
 - ◆ Block incoming port 53 packets except known trusted servers

◆ Issues

- Stateful filtering
- Encapsulation: address translation, other complications
- Fragmentation

Source-Address Forgery



More about networking: port numbering

◆ TCP connection

- Server port uses number less than 1024
- Client port uses number between 1024 and 16383

◆ Permanent assignment

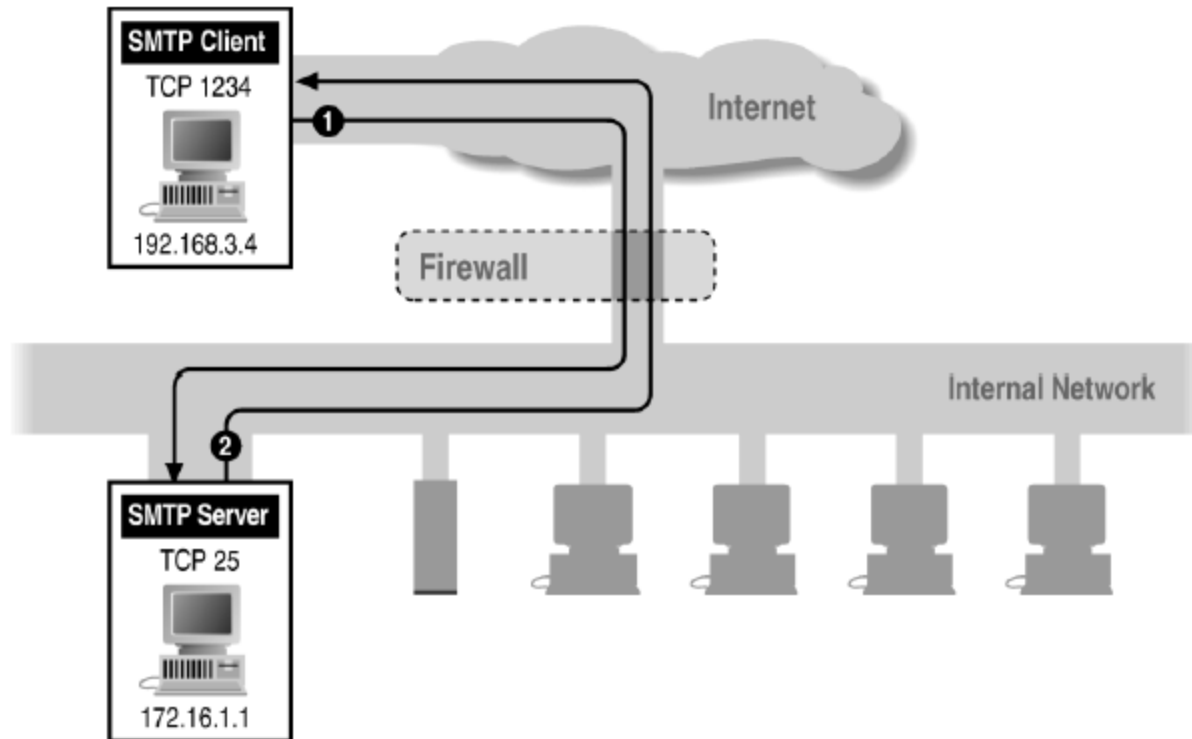
- Ports <1024 assigned permanently
 - ◆ 20,21 for FTP 23 for Telnet
 - ◆ 25 for server SMTP 80 for HTTP

◆ Variable use

- Ports >1024 must be available for client to make connection
- Limitation for stateless packet filtering
 - ◆ If client wants port 2048, firewall must allow incoming traffic
- Better: stateful filtering knows outgoing requests
 - ◆ Only allow incoming traffic on high port to a machine that has initiated an outgoing request on low port

Filtering Example: Inbound SMTP

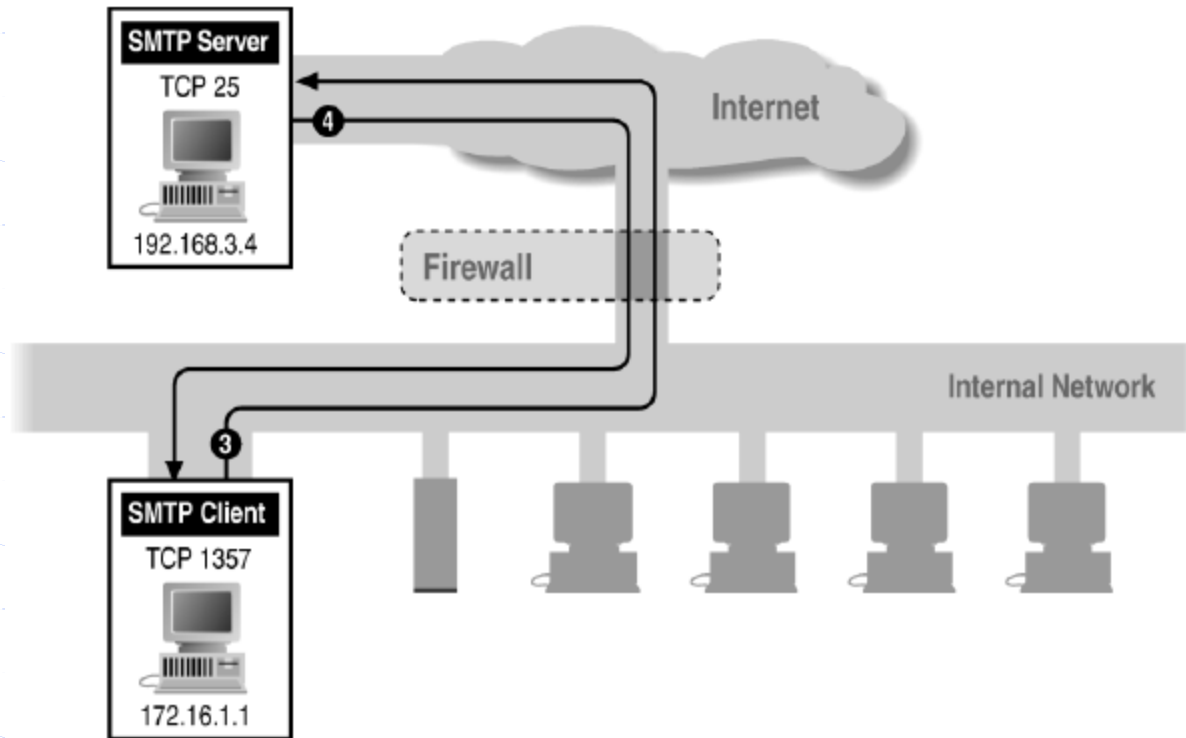
Assume we want to block internal server from external attack



Can block external request to internal server based on port number

Filtering Example: Outbound SMTP

Assume we want to allow internal access to external server

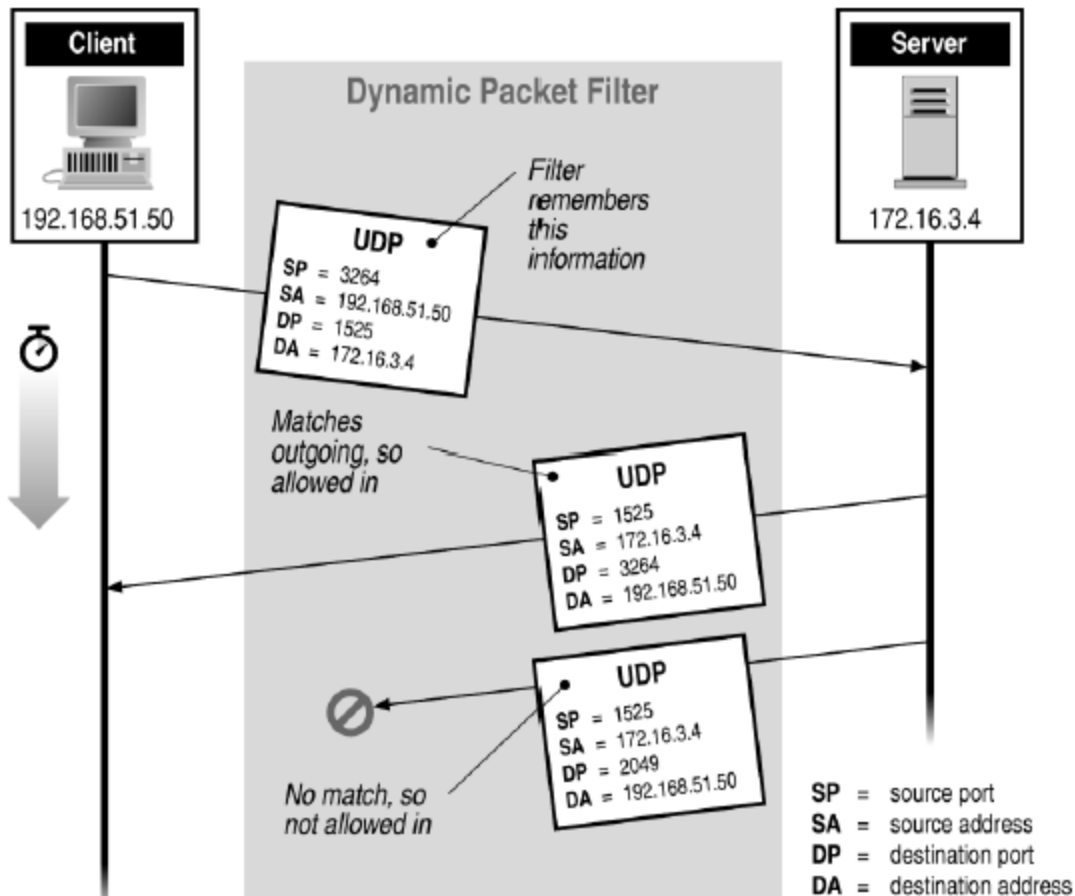


Known low port out, arbitrary high port in

If firewall blocks incoming port 1357 traffic then connection fails

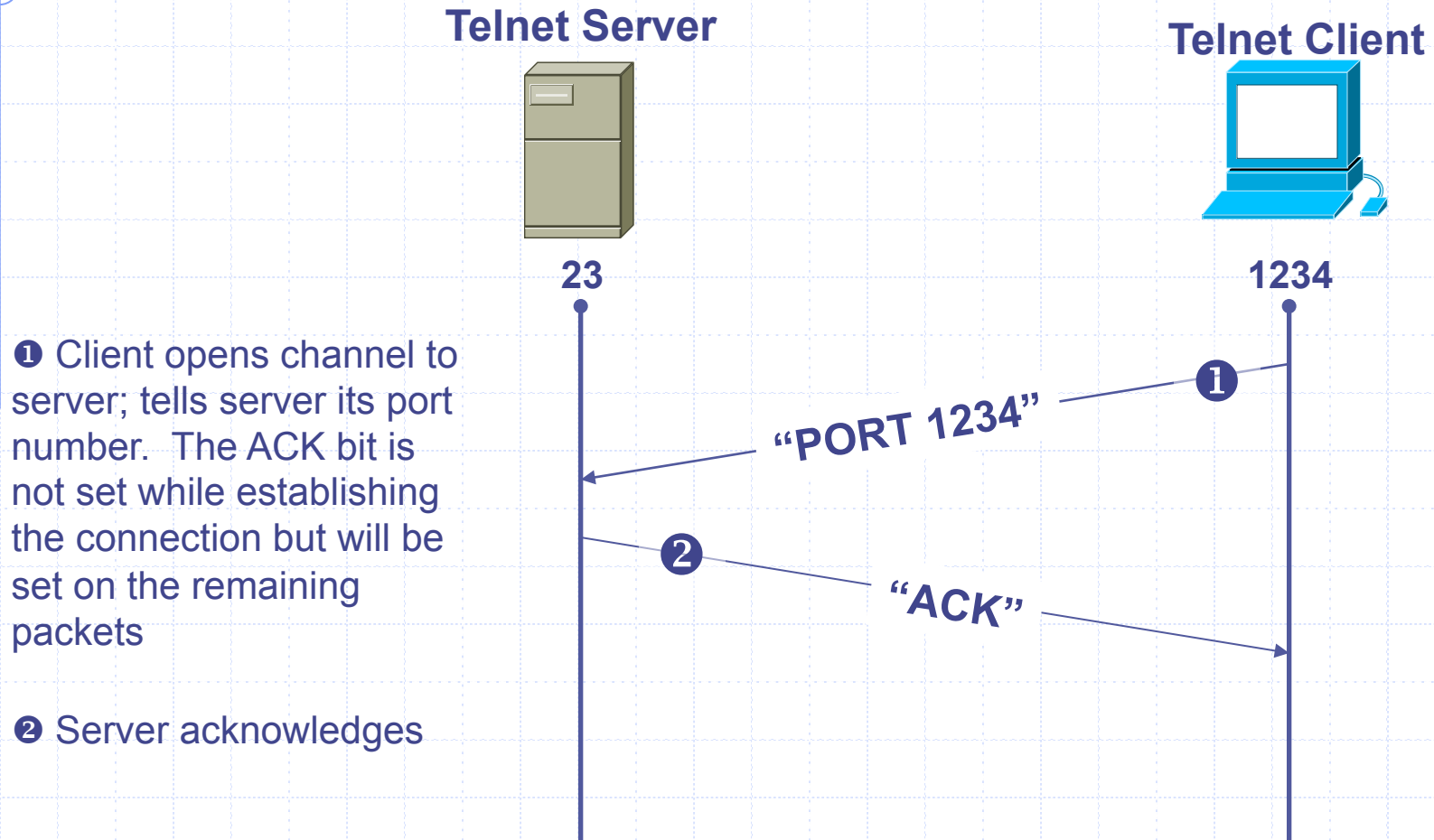
Stateful or Dynamic Packet Filtering

Assume we want to allow external UDP only if requested



Telnet

How can stateful filtering identify legitimate session?

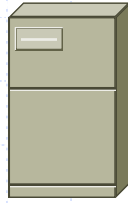


Stateful filtering can use this pattern to identify legitimate sessions

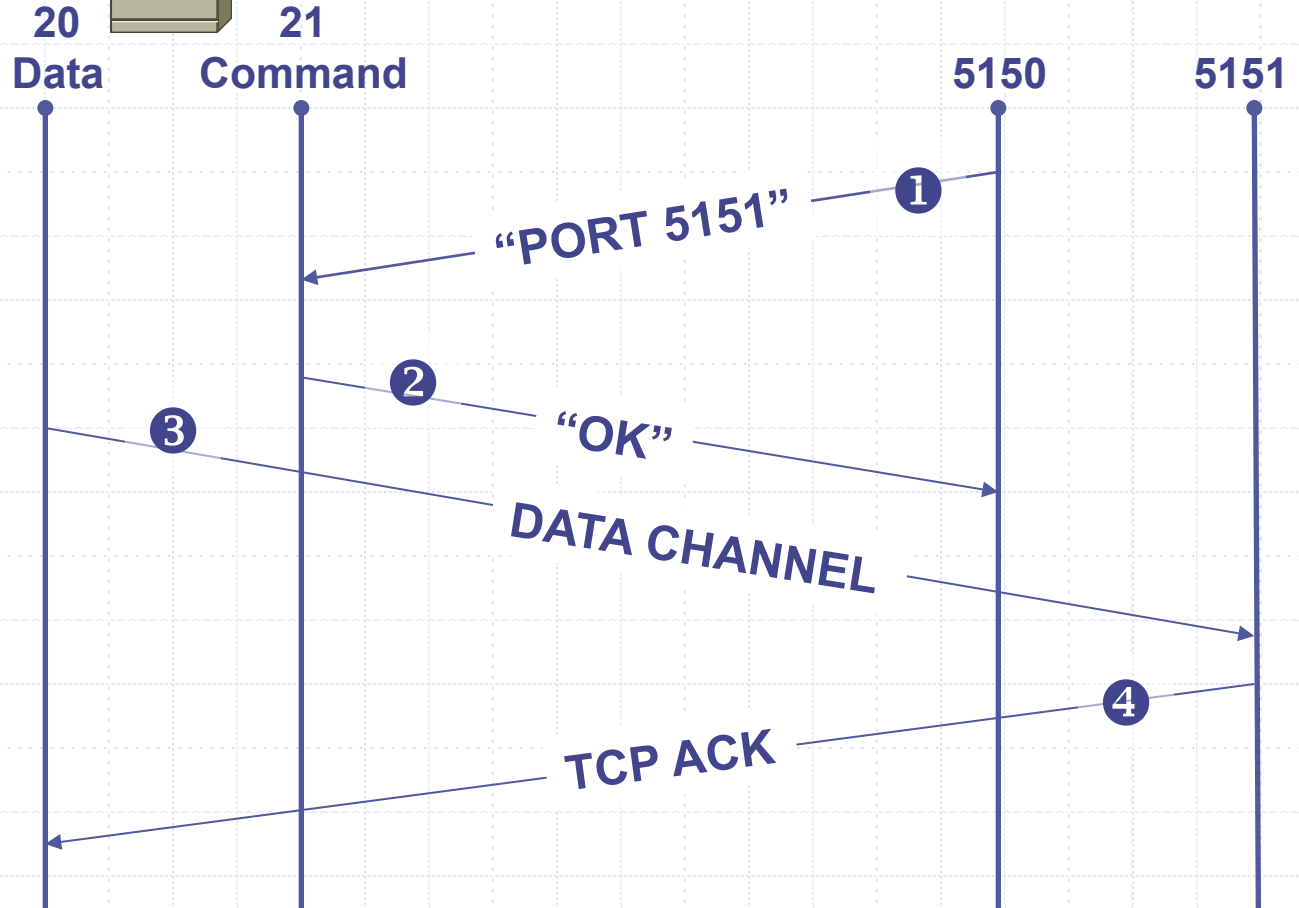
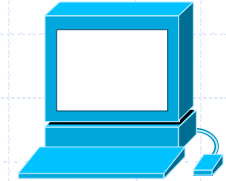
FTP

How can stateful filtering identify legitimate session?

FTP Server



FTP Client



① Client opens command channel to server; tells server second port number

② Server acknowledges

③ Server opens data channel to client's second port

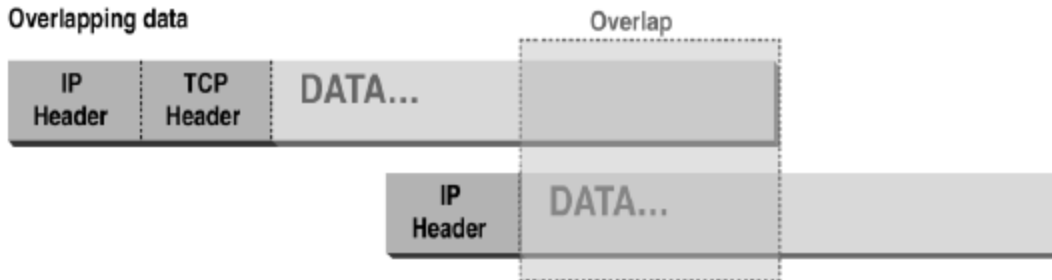
④ Client acknowledges

Abnormal Fragmentation

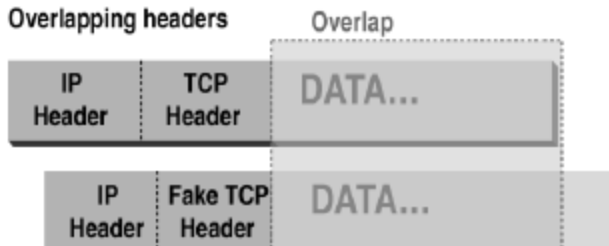
Normal



Overlapping data



Overlapping headers

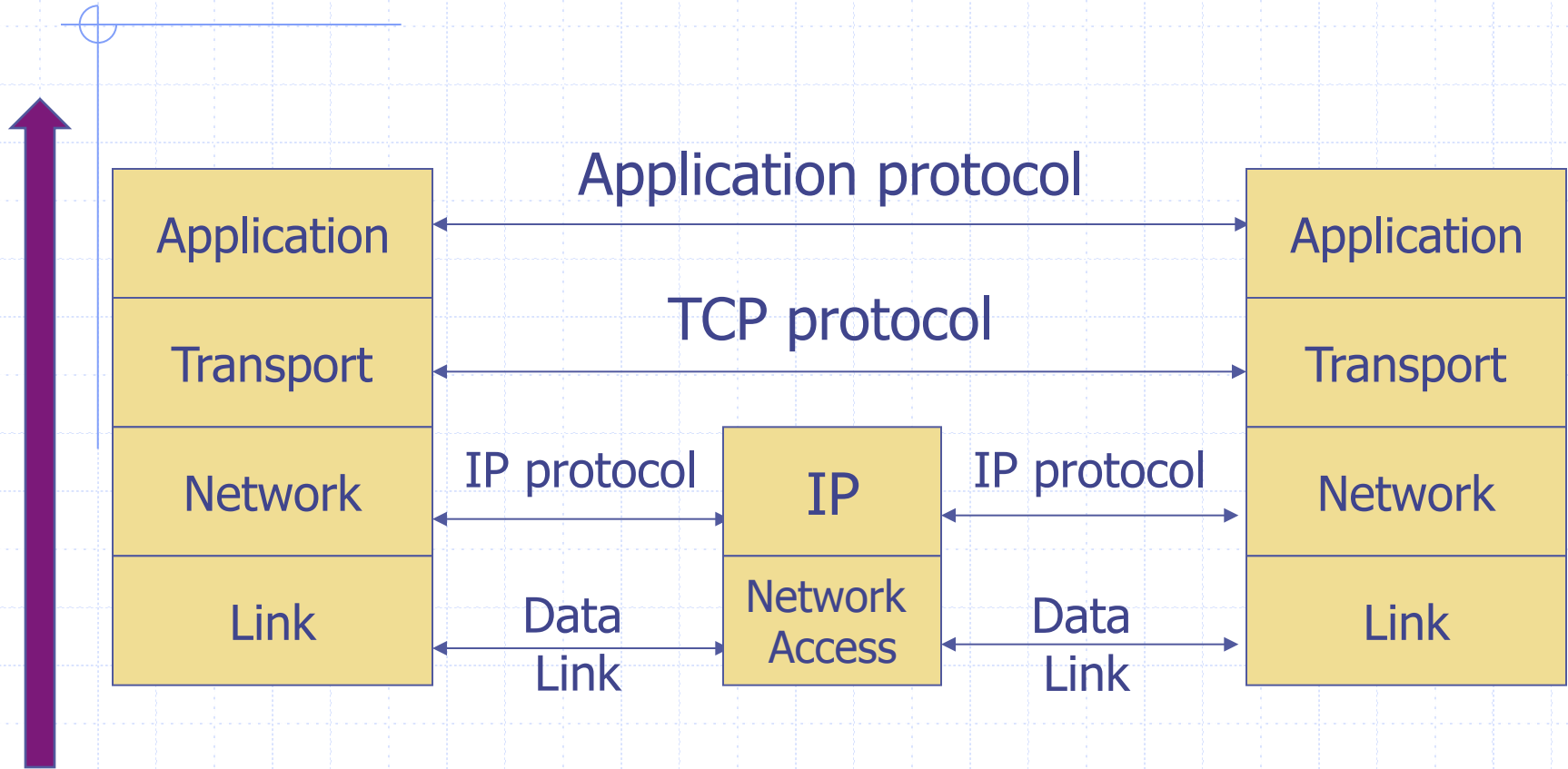


Low offset allows second packet to overwrite TCP header at receiving host

Packet Fragmentation Attack

- ◆ Firewall configuration
 - TCP port 23 is blocked but SMTP port 25 is allowed
- ◆ First packet
 - Fragmentation Offset = 0.
 - DF bit = 0 : "May Fragment"
 - MF bit = 1 : "More Fragments"
 - Destination Port = 25. TCP port 25 is allowed, so firewall allows packet
- ◆ Second packet
 - Fragmentation Offset = 1: second packet overwrites all but first 8 bits of the first packet
 - DF bit = 0 : "May Fragment"
 - MF bit = 0 : "Last Fragment."
 - Destination Port = 23. Normally be blocked, but sneaks by!
- ◆ What happens
 - Firewall ignores second packet "TCP header" because it is fragment of first
 - At host, packet reassembled and received at port 23

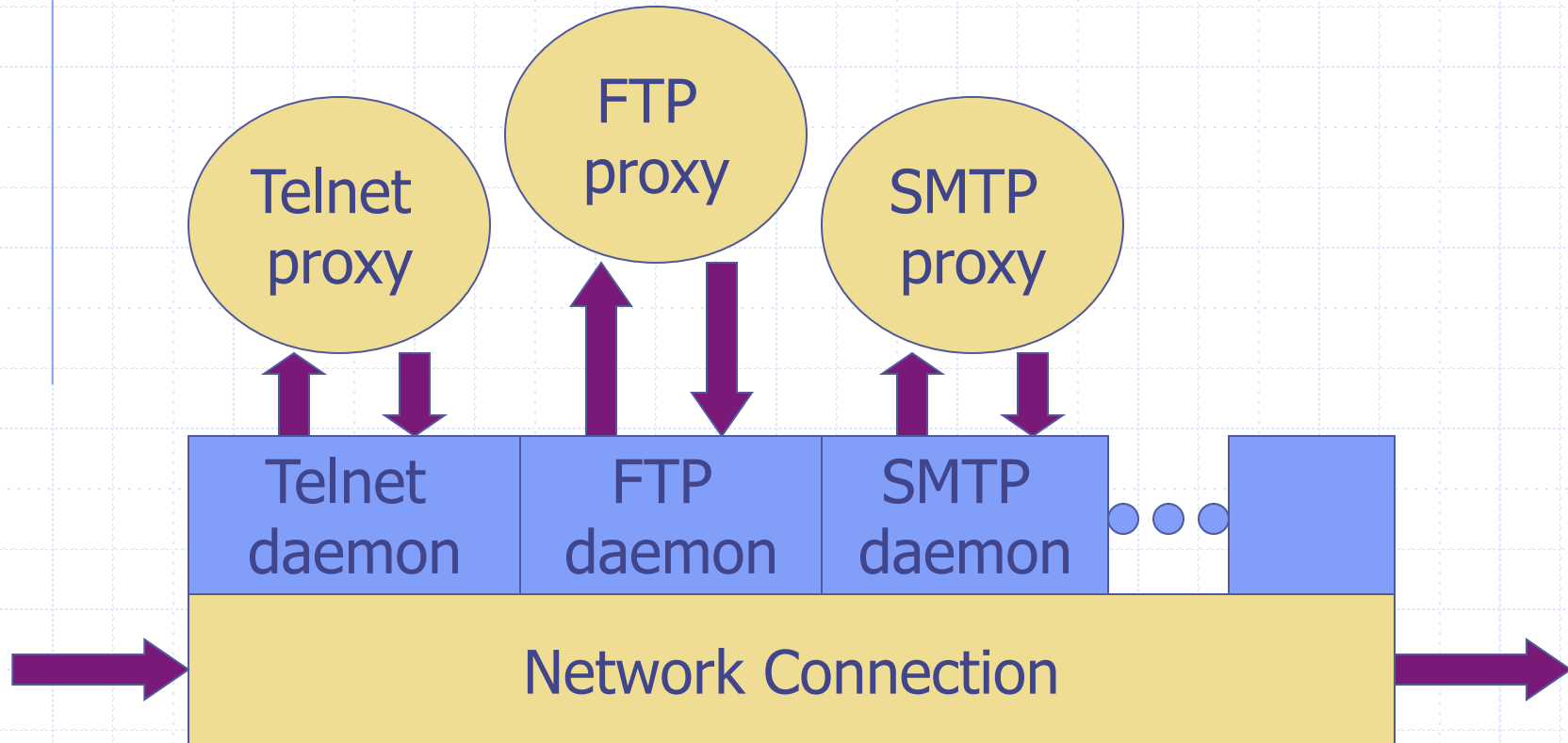
TCP Protocol Stack



Proxying Firewall

- ◆ Application-level proxies
 - Tailored to http, ftp, smtp, etc.
 - Some protocols easier to proxy than others
- ◆ Policy embedded in proxy programs
 - Proxies filter incoming, outgoing packets
 - Reconstruct application-layer messages
 - Can filter specific application-layer commands, etc.
 - ◆ Example: only allow specific ftp commands
 - ◆ Other examples: ?
- ◆ Several network locations – see next slides

Firewall with application proxies



Daemon spawns proxy when communication detected ...

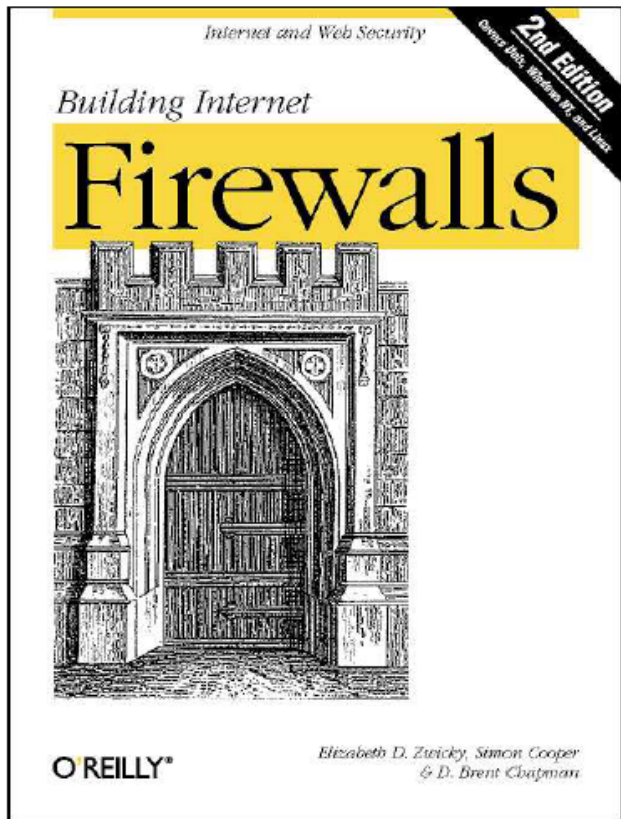
Application-level proxies

- ◆ Enforce policy for specific protocols
 - E.g., Virus scanning for SMTP
 - ◆ Need to understand MIME, encoding, Zip archives
 - Flexible approach, but may introduce network delays
- ◆ “Batch” protocols are natural to proxy
 - SMTP (E-Mail) NNTP (Net news)
 - DNS (Domain Name System) NTP (Network Time Protocol)
- ◆ Must protect host running protocol stack
 - Disable all non-required services; keep it simple
 - Install/modify services you want
 - Run security audit to establish baseline
 - Be prepared for the system to be compromised

Web traffic scanning

- ◆ Intercept and proxy web traffic
 - Can be host-based
 - Usually at enterprise gateway
- ◆ Block known bad sites
- ◆ Block pages with known attacks
- ◆ Scan attachments
 - Virus, worm, malware, ...

Firewall references



Elizabeth D. Zwicky
Simon Cooper
D. Brent Chapman

Firewalls and Internet Security Second Edition

Repelling the Wily Hacker

William R. Cheswick
Steven M. Bellovin
Aviel D. Rubin



William R Cheswick
Steven M Bellovin
Aviel D Rubin

ADDISON-WESLEY PROFESSIONAL COMPUTING SERIES

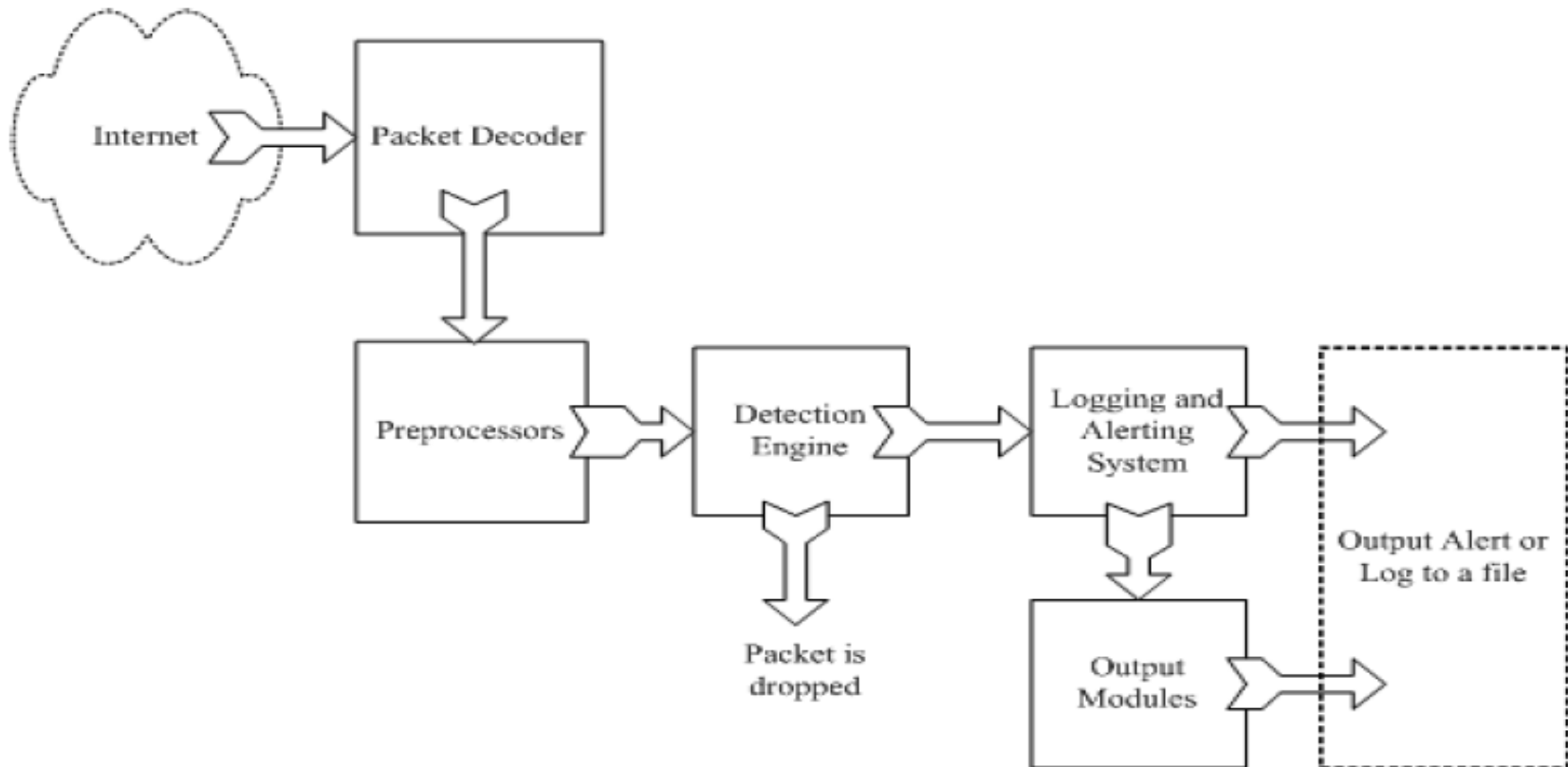
Intrusion detection

- ◆ Many intrusion detection systems
 - Network-based, host-based, or combination
- ◆ Two basic models
 - Misuse detection model
 - ◆ Maintain data on known attacks
 - ◆ Look for activity with corresponding signatures
 - Anomaly detection model
 - ◆ Try to figure out what is “normal”
 - ◆ Report anomalous behavior
- ◆ Fundamental problem: too many false alarms



<http://www.snort.org/>

Example: Snort

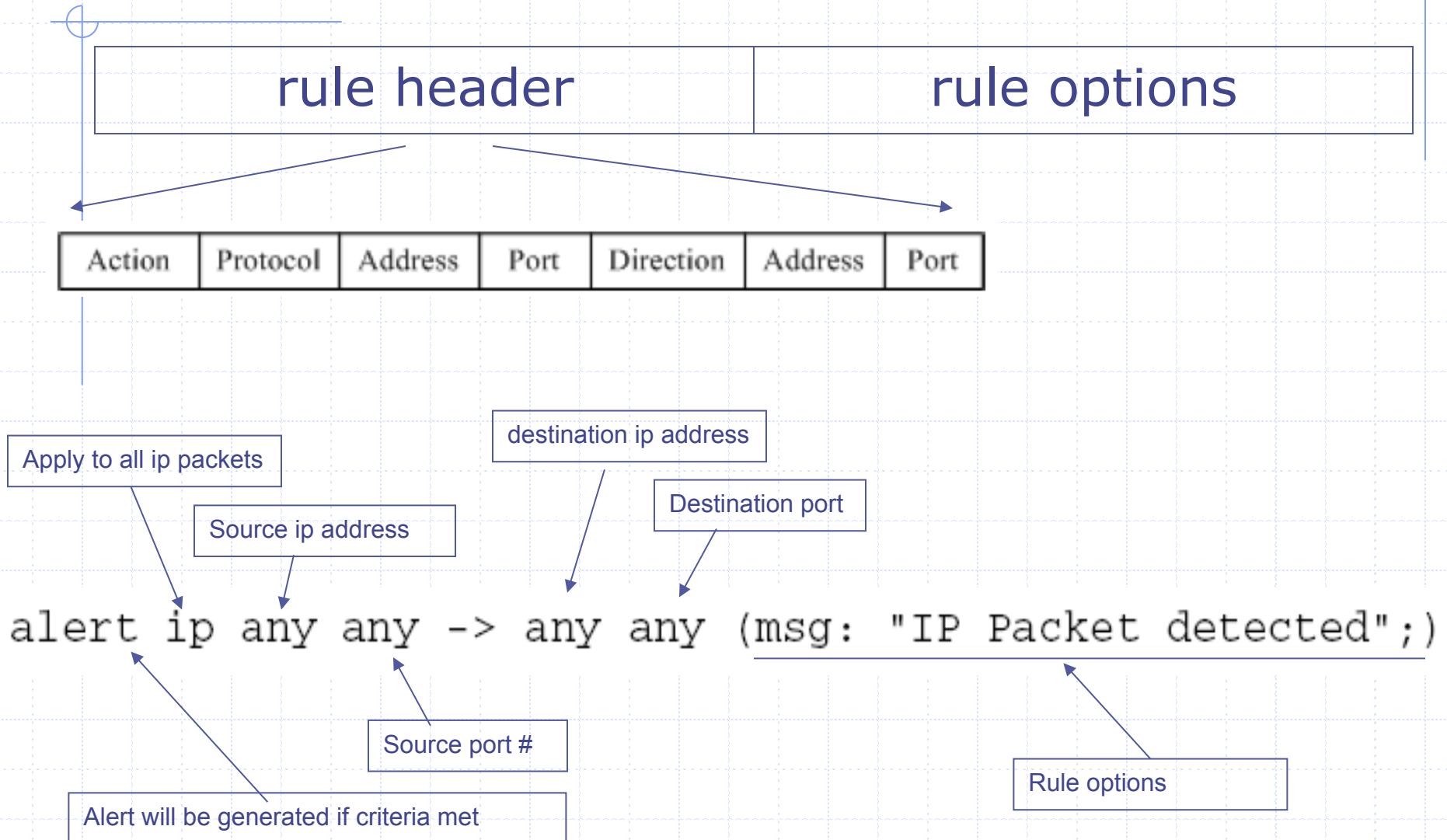


From: Rafeeq Ur Rehman, *Intrusion Detection Systems with Snort: Advanced IDS Techniques with Snort, Apache, MySQL, PHP, and ACID.*

Snort components

- ◆ Packet Decoder
 - input from Ethernet, SLIP, PPP...
- ◆ Preprocessor:
 - detect anomalies in packet headers
 - packet defragmentation
 - decode HTTP URI
 - reassemble TCP streams
- ◆ Detection Engine: applies rules to packets
- ◆ Logging and Alerting System
- ◆ Output Modules: alerts, log, other output

Snort detection rules



Additional examples

```
alert tcp any any -> 192.168.1.0/24 111  
(content:"|00 01 86 a5|"; msg: "mountd access");)
```

```
alert tcp !192.168.1.0/24 any -> 192.168.1.0/24 111  
(content: "|00 01 86 a5|"; msg: "external mountd access");)
```

! = negation operator in address

content - match content in packet

192.168.1.0/24 - addr from 192.168.1.1 to 192.168.1.255

<https://www.snort.org/documents/snort-users-manual>

Snort challenges

- ◆ Misuse detection – avoid known intrusions
 - Database size continues to grow
 - ◆ Snort version 2.3.2 had 2,600 rules
 - Snort spends 80% of time doing string match
- ◆ Anomaly detection – identify new attacks
 - Probability of detection is low

Difficulties in anomaly detection

◆ Lack of training data

- Lots of “normal” network, system call data
- Little data containing realistic attacks, anomalies

◆ Data drift

- Statistical methods detect changes in behavior
- Attacker can attack gradually and incrementally

◆ Main characteristics not well understood

- By many measures, attack may be within bounds of “normal” range of activities

◆ False identifications are very costly

- Sys Admin spend many hours examining evidence

Summary

- ◆ Protecting network connections
 - Wireless security – 802.11i/WPA2
 - IPSEC
- ◆ Perimeter network perimeter defenses
 - Firewall
 - ◆ Packet filter (stateless, stateful),
 - ◆ Application layer proxies
 - Intrusion detection
 - ◆ Anomaly and misuse detection
- ◆ Network infrastructure security
 - BGP vulnerability and S-BGP
 - DNSSEC, DNS rebinding