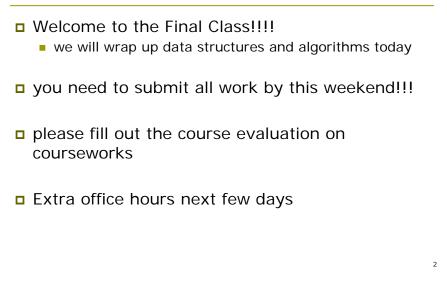
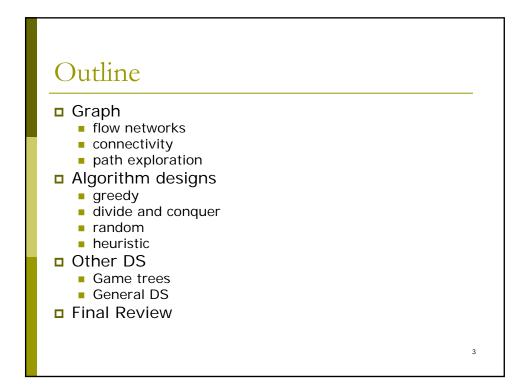
## 3137 Data Structures and Algorithms in C++

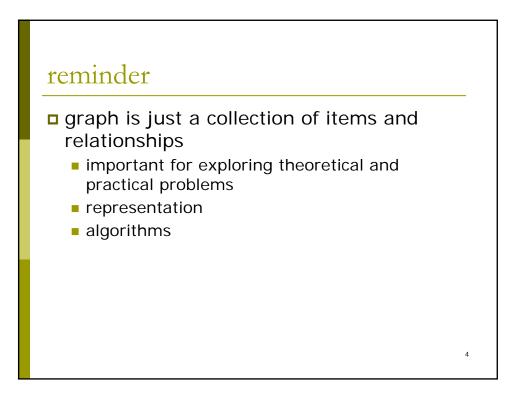
Lecture 10 Aug 9 2006 Shlomo Hershkop

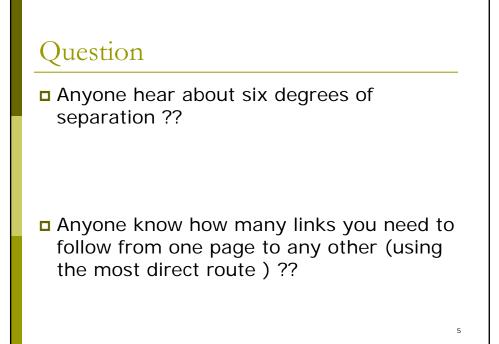
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### Announcements

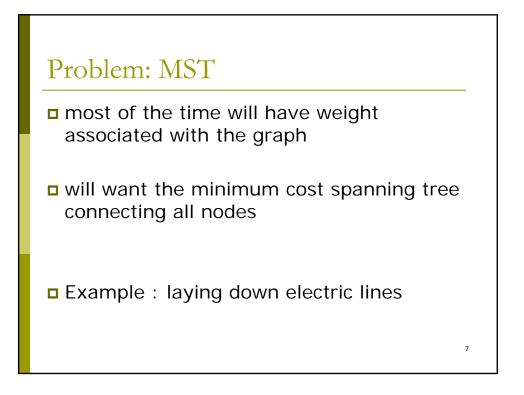


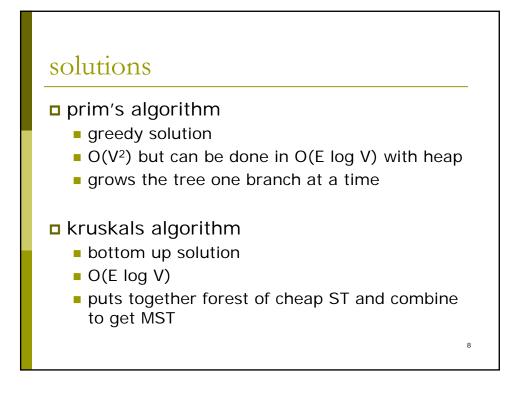


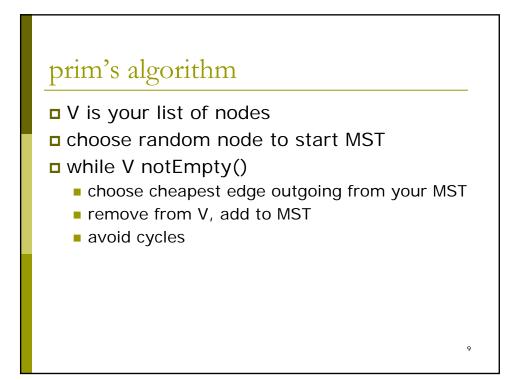


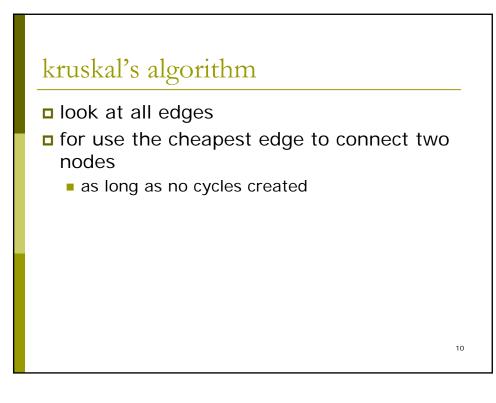


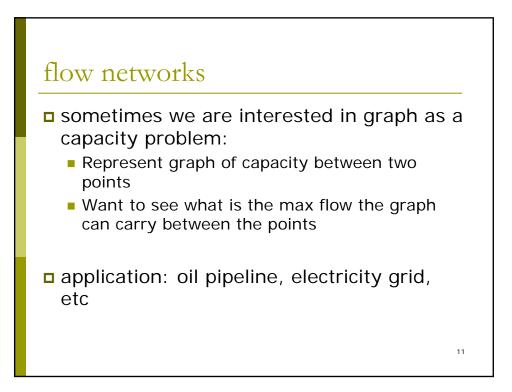
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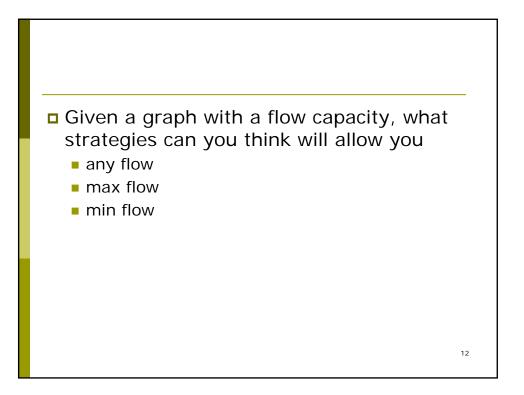














hand out from last class

• will be using extra graph for book keeping

# undoing • an augment the algorithm by adding back edges • handout



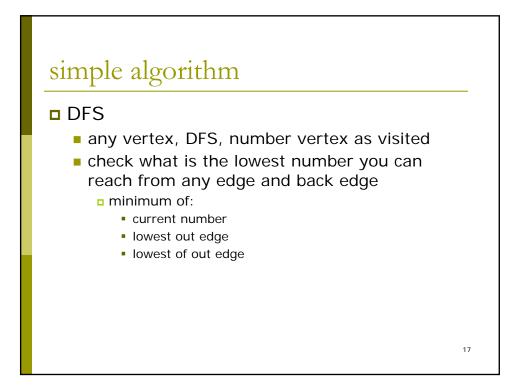
- say you have a bunch of tasks which have to be done over time
- some can be done in parallel
- some have to wait for a specific task to be done before being able to be started
- think building construction, can't paint the inside without putting up the walls...etc

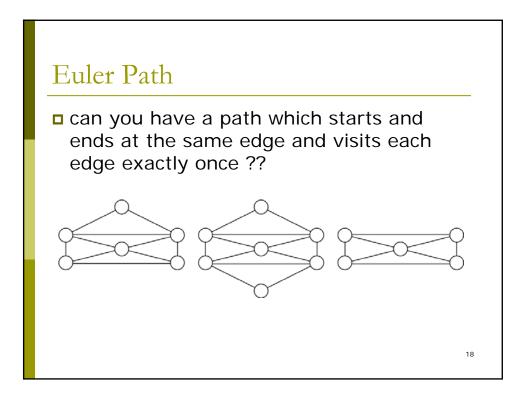


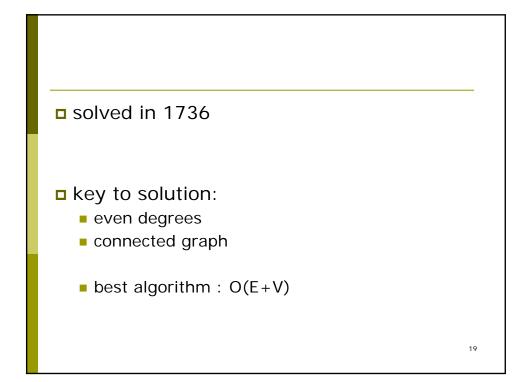
biconnected:

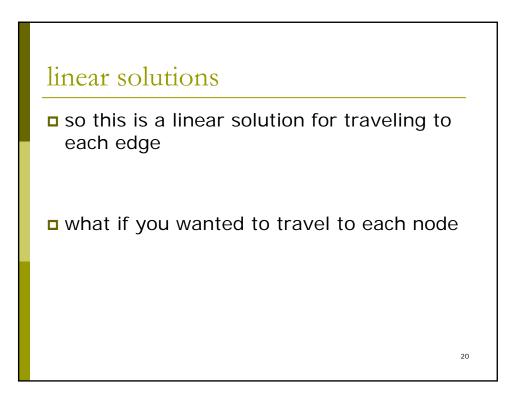
- graph is biconnected is there are no vertices which can disconnect the graph
- i.e no articulation points
  - points of failurehow to find them all??

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### TSP

□ Traveling salesman problem:

- shortest route through n cities that visits each city once and returns to starting one
- can be stated as finding the shortest Hamiltonian circuit of the graph
- how many combinations of tours available ??

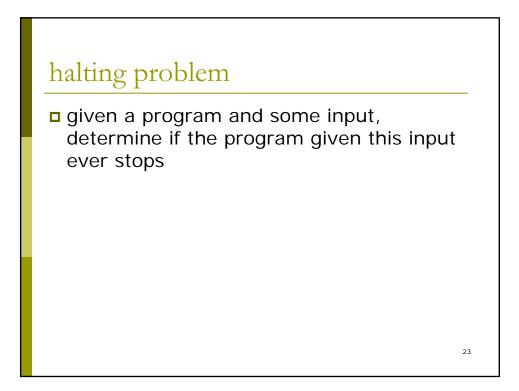
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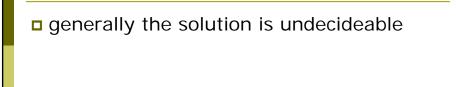
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### computing limitations

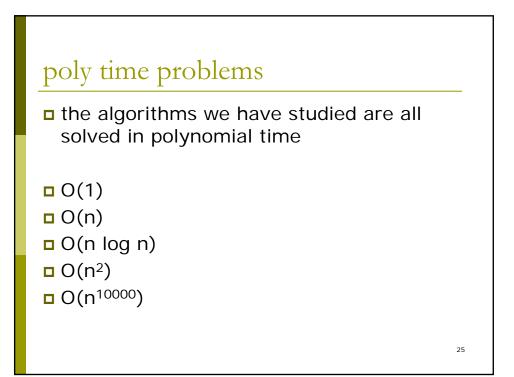
at the beginning of the 1900's there was an open problem to show what would be the theoretical limitation of computing

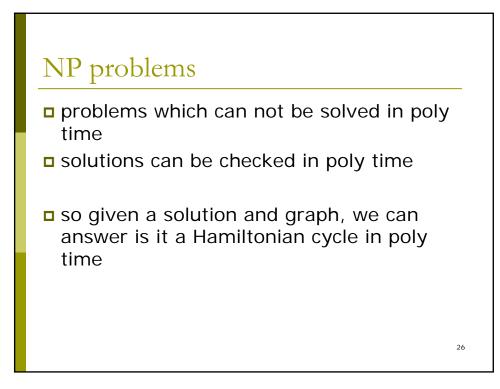
can a computer solve every problem

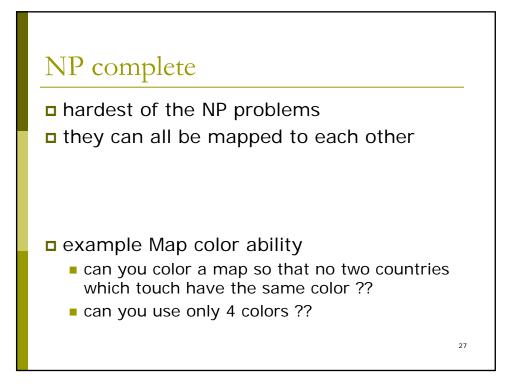


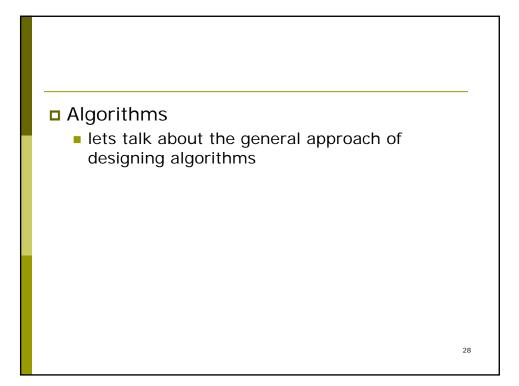


but specific instances of this problem can be solved





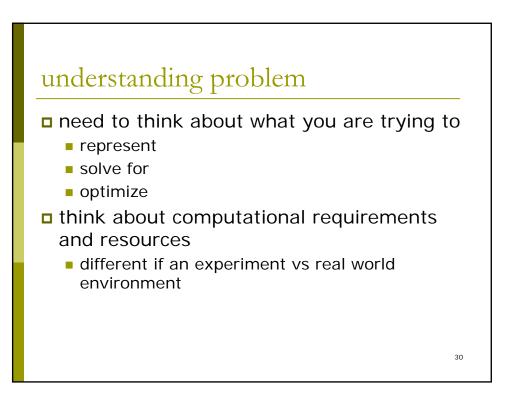






- 1. Understand the problem
- 2. Decide on exact vs approximate solution
- 3. Design your algorithm
- 4. Prove its correct
  - 1. might have to go back to 2 or 3
- 5. Analyze the algorithm1. might have to go back to 2 or 3
- 6. Code the algorithm

DON'T START WITH 6!!





Why wouldn't you want an exact solution ??

- can't be done
- might be possible, very very slow

might be part of an exact solution

### Design issues

- Data structures can make or break your algorithm
  - sometimes can use the standard, more often than not cut-paste various models
- look for similar problems, how were they solved
- feel free to approach the problem in an unusual way
- documentation is essential
- try not to reinvent the wheel (Everywhere)



check basic cases

we've done very easy proofs, as mentioned some algorithms avg cases took years to find a proof

what does it mean for approximate algorithm...how do you prove it works ?

Analyze it

- time resources
- space resources
- can it be simplified ?
- can you generalize it a little ?
- can it handle a broad set of inputs ?

not always all steps necessary!

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### coding

■ have lots of experience from this class ☺

most of the time proof approximated by testing...remember to code this!!

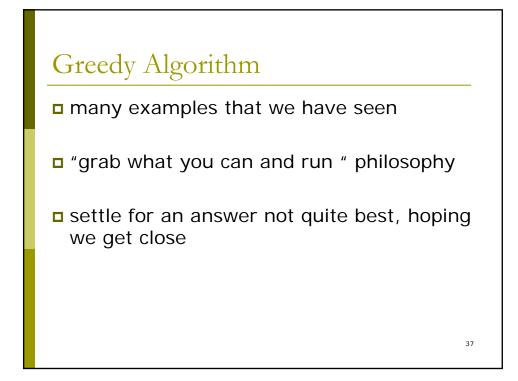
- packages available!
- choose language carefully
  - consider available language resources!
- don't loose efficiency of the algorithm with poor implementation

### Classes of algorithms

Lets talk about different types of algorithms

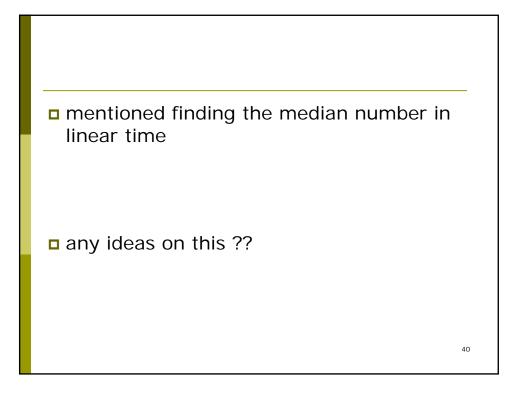
- greedy
- divide and conquer
- dynamic algorithms
- backtracking
- decrease and conquer
- randomized algorithms

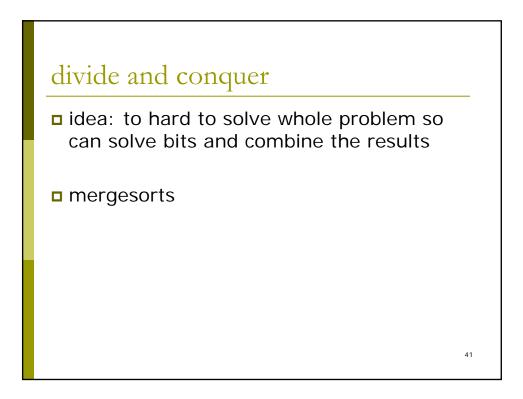
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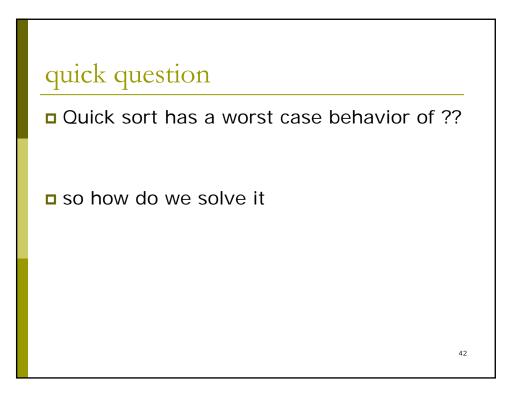


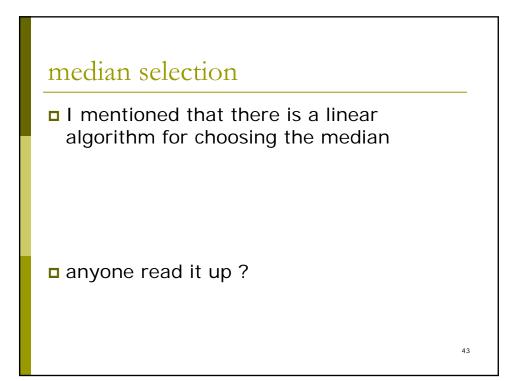
## problem a given a set of points in 2 dimensional space a can you find closest pair of points to any point ? b mmm sounds familiar ??





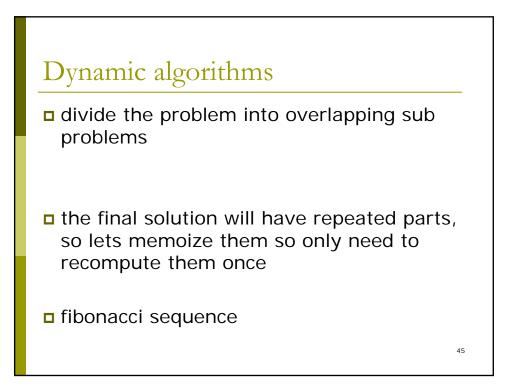


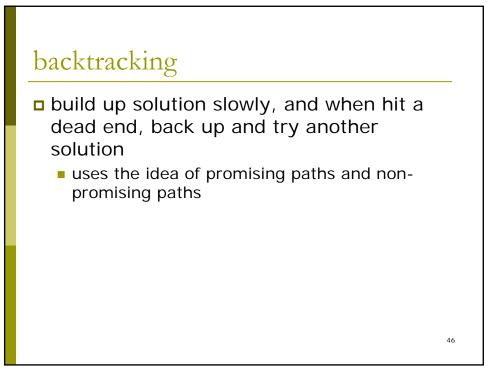




### basic idea:

- divide your elements into groups of small size say 5
- **•** find the median in each group...running time ???
- so we have N/2 medians
- find the medians of these medians
- how close are we to the median ??
- what does it do to quicksort ?







how to place n queens on a n by n board

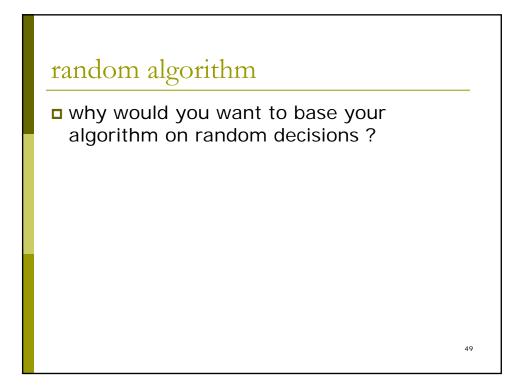
lets do 4 queens

### decrease and conquer

idea: problem of size n can be solved as problem size n-1 + the i'th item

think of some of the sorting routines

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□ for some hard problems

- can not solve directly
- we can get the correct answer with high probability
- run efficiently in expectation!

### Example

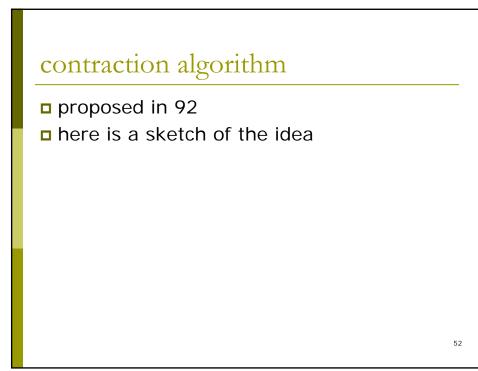
- computer networks....anyone use an ethernet network
- anyone know how they work
- how the tcp/ip protocol works ?

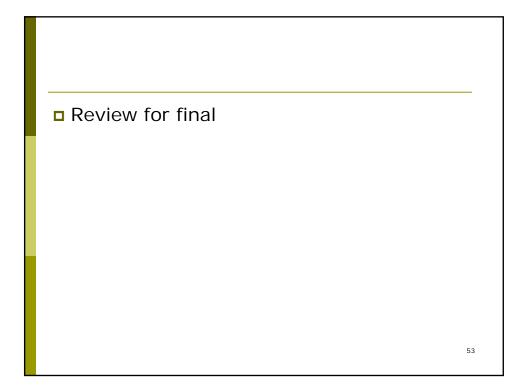


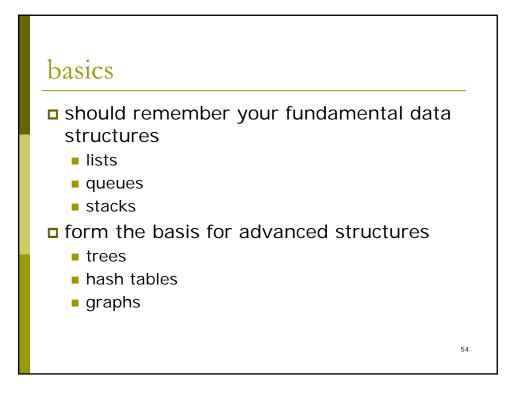
- given a Graph (undirected) we want to find the minimum number of edge deletions so that we can divide the graph into two distinct non-empty sets A and B
- real life example: say you have a complex network, what would be the minimum number of damages to it to split the network

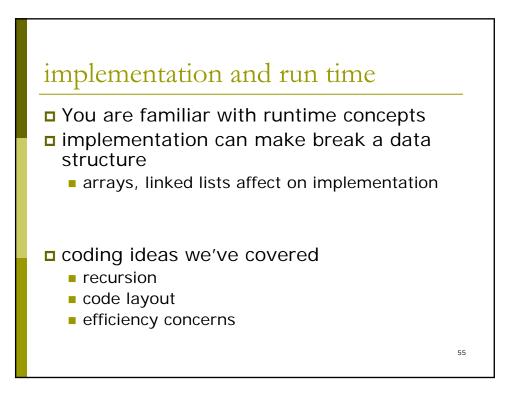
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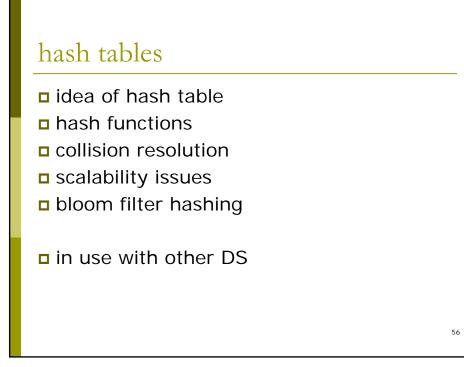
any ideas on this ?











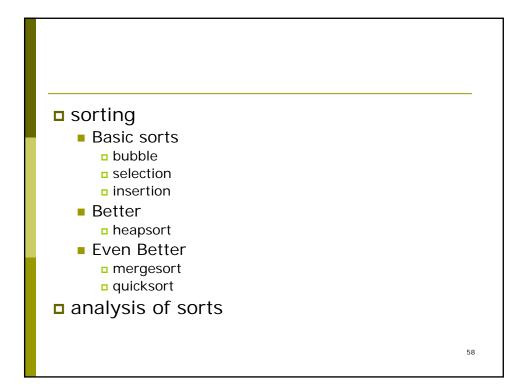
### Trees

implementation details

BST

- Traversals
- AVL
- □ B+ trees

how are trees related to graphs ??



### sets

idea of sets
union / find
path compression
union by rank

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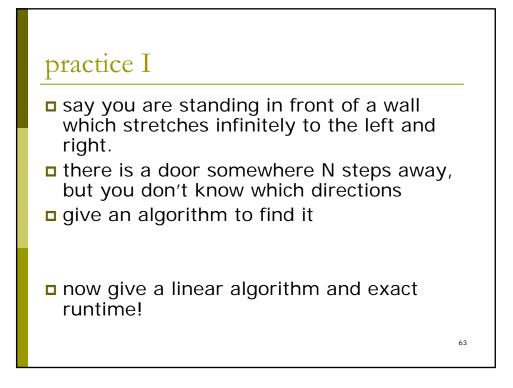
## Algorithms

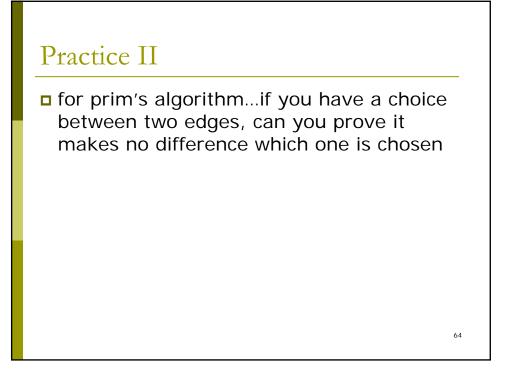
analyzing general run times
comparing classes of algorithms
getting a feel if it's a good algorithm

Lessons Learned

Think then code!

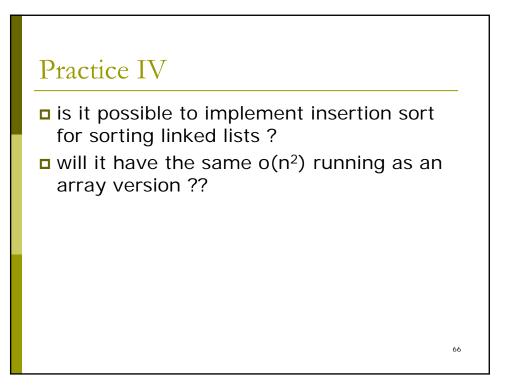
Try not to reinvent wheel or do things in n<sup>3</sup> or n! land <sup>(C)</sup>

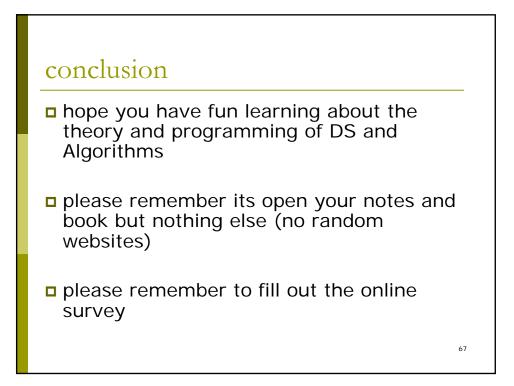


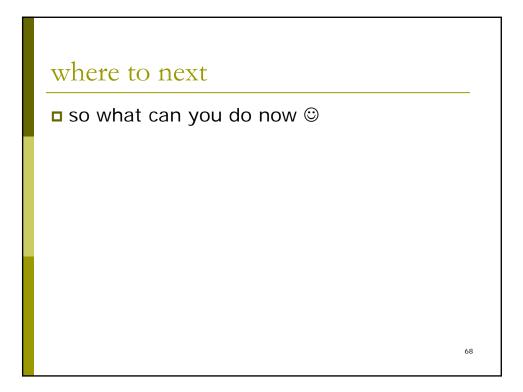


### Practice III

- You are given an array that contains N numbers.
- We want to determine if there are two numbers whose sum equals a given number K.
- for instance if the input is 8,4,1,6 and K=10, the answer is yes (use 4 and 6).
- A number may be used twice.
- Do the following:
  - Give an O(N<sup>2</sup>) algorithm to solve this problem
  - Give an O(N log N) algorithm to solve this problem.







Thank you

□ good luck on the final....