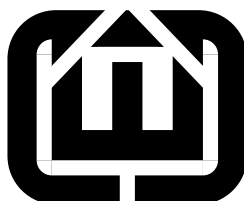


CEC GRC-to-PDG Converter



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Abstract

This converts the control-flow graph portion of the GRC graph into a program dependence graph using the algorithm described by Cytron et al. in their 1991 TOPLAS article.

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1 Utilities

1.1 contains

Return true if the set contains the object.

2a	$\langle \text{utilities 2a} \rangle \equiv$	(30)	2b
	template <class T> bool contains(set<T> &s, T o) {		
	return s.find(o) != s.end();		
	}		
2b	$\langle \text{utilities 2a} \rangle + \equiv$	(30)	<2a
	template <class T, class U> bool contains(map<T, U> &m, T o) {		
	return m.find(o) != m.end();		
	}		

2 The STDPS class

```

3  <stdps class 3>≡ (30)
    class STDPS {
        EnterGRC *entergrc;
        set<GRCNode *> visited;
        map<GRCNode *, set<GRCNode *> > enter_nodes; //enter nodes under, for node with multi-par only

    public:

        STDPS(EnterGRC *entergrc): entergrc(entergrc) {}
        ~STDPS() {}

        Status execute() {
            visited.clear();
            variable_dfs(entergrc);
            return Status();
        }

    private:

        set<GRCNode *> variable_dfs(GRCNode *n)
        {
            set<GRCNode *> RET;

            if (!n)
                return RET;
            if (visited.count(n) > 0){
                assert(enter_nodes.count(n) > 0);
                return enter_nodes[n];
            }

            visited.insert(n);

            for (vector<GRCNode *>::iterator i = n->successors.begin();
                i != n->successors.end(); i++){
                set<GRCNode *> ch_set = variable_dfs(*i);
                if (ch_set.size() > 0){ // if find some enters under
                    RET.insert(ch_set.begin(), ch_set.end());
                }
                if (dynamic_cast<Enter *>(n)){ // if it's an enter, decide whether to add dps
                    for (set<GRCNode *>::iterator j = ch_set.begin();
                        j != ch_set.end(); j++){
                        if (same_sstp(n,*j)){
                            **j << n;
                        }
                    }
                    RET.insert(n);
                }
            }
            if ((dynamic_cast<STSuspend *>(n))

```

```

        ||(dynamic_cast<Switch *>(n))) { //if it's suspend or switch, decide whether to add dp
        for (set<GRCNode *>::iterator j = ch_set.begin();
            j != ch_set.end(); j++){
//        if (same_stp(n,*j)){
            if(st_ancestor(n,*j)){
                **j << n;
            }
        }
    }
}

if (n->predecessors.size()>1){
    enter_nodes[n].insert(RET.begin(), RET.end());
}

return RET;
}

//test if two nodes have the same st pointer, n1-suspend, n2-enter
bool same_stp(GRCNode *n1, GRCNode *n2)
{
    STSuspend *s;
    Enter *e;

    s = dynamic_cast<STSuspend *>(n1);
    e = dynamic_cast<Enter *>(n2);

    if (s->st == e->st)
        return true;

    return false;
}

bool same_sstp(GRCNode *n1, GRCNode *n2)
{
    Enter *e1, *e2;

    e1 = dynamic_cast<Enter *>(n1); assert(e1);
    e2 = dynamic_cast<Enter *>(n2); assert(e2);

    //if they point to the same stnode, not need to add constrain btw them
    if(e1->st == e2->st)
        return false;

    if (e1->st->parent == e2->st->parent) {
        if (dynamic_cast<STexcl *>(e1->st->parent))
            return true;
    }
    return false;
}

```

```

bool st_ancestor(GRCNode *p, GRCNode *c)
{
    GRCSTNode *pp = dynamic_cast<GRCSTNode *>(p); assert(p);
    GRCSTNode *cc = dynamic_cast<GRCSTNode *>(c); assert(cc);

    STNode *stp = pp ->st;
    STNode *stc = cc ->st;

    while(stc != NULL){
        if(stp==stc) return true;
        stc = stc->parent;
    }
    return false;
}

};

```

3 Signal Dependency Calculator Class

This class removes & re-computes dependencies between signal emissions and tests.

```

5  <dependency class 5>≡ (30)
    class Dependencies : public Visitor {
protected:
    set<GRCNode *> visited;
    GRCNode *current;
    map<GRCNode *, bool> par_label;

    struct SignalNodes {
        set<GRCNode *> writers;
        set<GRCNode *> readers;
    };

    map<SignalSymbol *, SignalNodes> dependencies;

    <dependency methods 8b>
    void mark_par(GRCNode* n);
    bool have_comm_pp_gen(GRCNode* n, GRCNode* m);
    bool have_comm_pp(GRCNode* n, GRCNode* m);

public:
    Dependencies() {}
    virtual ~Dependencies() {}
    void compute(GRCNode *);
};

```

6a \langle dependency method definitions 6a $\rangle \equiv$ (30) 6b \triangleright

```

void Dependencies::compute(GRCNode *root)
{
    assert(root);

    Dependencies depper;

    depper.dfs(root);

    for ( map<SignalSymbol *, SignalNodes>::const_iterator i =
          depper.dependencies.begin() ; i != depper.dependencies.end() ;
          i++ ) {
        const SignalNodes &sn = (*i).second;
        if (!sn.writers.empty() && !sn.readers.empty()) {
            for ( set<GRCNode*>::const_iterator j = sn.writers.begin() ;
                  j != sn.writers.end() ; j++ ){
                visited.clear();
                par_label.clear();
                mark_par(*j);
                for ( set<GRCNode*>::const_iterator k = sn.readers.begin() ;
                      k != sn.readers.end() ; k++ ){
                    visited.clear();
                    if (have_comm_pp_gen((*k),(*j)))
                        **k << *j;
                }
            }
        }
    }
}

```

6b \langle dependency method definitions 6a $\rangle + \equiv$ (30) <6a 7a \rangle

```

void Dependencies::mark_par(GRCNode* n)
{
    int sz, i;

    if (visited.count(n) > 0)
        return;

    sz = n->predecessors.size();
    for (i = 0; i < sz; i++){
        if ( par_label[n->predecessors[i]] == false){
            par_label[n->predecessors[i]] = true;
            mark_par(n->predecessors[i]);
        }
    }

    //also mark n itself as its parent
    par_label[n] = true;
    visited.insert(n);
}

```

7a \langle dependency method definitions 6a $\rangle + \equiv$ (30) \langle 6b 7b \rangle

```

//test if two nodes n & m have parallel first-comm-parent
//where m's parents have been labeled
bool Dependencies::have_comm_pp_gen(GRCNode* n, GRCNode* m)
{
    if (par_label[n])
        return true;

    return have_comm_pp(n,m);
}

```

7b \langle dependency method definitions 6a $\rangle + \equiv$ (30) \langle 7a 8a \rangle

```

bool Dependencies::have_comm_pp(GRCNode* n, GRCNode* m)
{

    assert(n);
    if (visited.count(n)>0)
        return false;

    visited.insert(n);

    if (par_label[n]){//found a first_comm_parent
        if ((dynamic_cast<Fork *>(n)) //is it a parallel node?
            || (n == m))//or, is it the emitter corasp?
            return true;
    }
    else {
        for (vector<GRCNode *>::iterator it =n->predecessors.begin();
            it != n->predecessors.end(); it++){
            if (have_comm_pp(*it, m))
                return true;
        }
    }

    return false;
}

```

3.1 DFS

This is the core dispatch procedure for the walker. It verifies it has not already visited the given node, visits it, then calls itself recursively on its successors.

```

8a  <dependency method definitions 6a>+≡ (30) <7b 8c>
    void Dependencies::dfs(GRCNode *n)
    {
        if (!n || visited.find(n) != visited.end() ) return;

        visited.insert(n);

        current = n;
        n->welcome(*this);

        for (vector<GRCNode*>::const_iterator i = n->successors.begin() ;
             i < n->successors.end() ; i++ ) dfs(*i);
    }

8b  <dependency methods 8b>≡ (5) 8d>
    void dfs(GRCNode *);

```

3.2 Action

An action may be an emit or exit statement, which emit signals.

```

8c  <dependency method definitions 6a>+≡ (30) <8a 9a>
    Status Dependencies::visit(Action &act)
    {
        Emit *emt = dynamic_cast<Emit *>(act.body);
        if (emt) {
            dependencies[emt->signal].writers.insert(current);
            current->dataSuccessors.clear();
        } else {
            // Traps are treated the same as signals
            Exit *ex = dynamic_cast<Exit *>(act.body);
            if (ex) {
                dependencies[ex->trap].writers.insert(current);
                current->dataSuccessors.clear();
            }
        }
        return Status();
    }

8d  <dependency methods 8b>+≡ (5) <8b 9b>
    Status visit(Action &);

```


3.3 DefineSignal

The DefineSignal node is like an emit.

- 9a $\langle \text{dependency method definitions 6a} \rangle + \equiv$ (30) $\langle 8c$
- ```

 Status Dependencies::visit(DefineSignal &ds)
 {
 assert(ds.signal);
 dependencies[ds.signal].writers.insert(current);
 current->dataSuccessors.clear();
 return Status();
 }

```
- 9b  $\langle \text{dependency methods 8b} \rangle + \equiv$  (5)  $\langle 8d \ 9c \rangle$
- ```

    Status visit(DefineSignal &);

```

3.4 Test

This descends down its predicate, possibly adding signal testers

- 9c $\langle \text{dependency methods 8b} \rangle + \equiv$ (5) $\langle 9b \ 10a \rangle$
- ```

 Status visit(Test &t) {
 t.predicate->welcome(*this); return Status();
 }

```

### 3.5 Expressions

10a  $\langle \text{dependency methods 8b} \rangle + \equiv$  (5)  $\langle 9c \ 10b \rangle$

```

Status visit(LoadSignalExpression &e) {
 dependencies[e.signal].readers.insert(current);
 current->dataPredecessors.clear();
 return Status();
}

Status visit(LoadSignalValueExpression &e) {
 dependencies[e.signal].readers.insert(current);
 current->dataPredecessors.clear();
 return Status();
}

Status visit(BinaryOp &e) {
 e.source1->welcome(*this);
 e.source2->welcome(*this);
 return Status();
}

Status visit(UnaryOp &e) {
 e.source->welcome(*this);
 return Status();
}

Status visit(CheckCounter &e) {
 e.predicate->welcome(*this);
 return Status();
}

Status visit(Delay &d) {
 d.predicate->welcome(*this);
 return Status();
}

```

#### 3.5.1 Vacuous Expression Nodes

10b  $\langle \text{dependency methods 8b} \rangle + \equiv$  (5)  $\langle 10a \ 11 \rangle$

```

Status visit(Literal &) { return Status(); }
Status visit(LoadVariableExpression &) { return Status(); }
Status visit(FunctionCall &) { return Status(); }

```

### 3.6 Trivial visitors

These nodes have no dependency implications and hence do nothing when visited.

```

11 <dependency methods 8b>+≡ (5) <10b
 Status visit(EnterGRC &) { return Status(); }
 Status visit(ExitGRC &) { return Status(); }
 Status visit(Nop &) { return Status(); }
 Status visit(Switch &) { return Status(); }
 Status visit(STSuspend &) { return Status(); }
 Status visit(Fork &) { return Status(); }
 Status visit(Terminate &) { return Status(); }
 Status visit(Enter &) { return Status(); }
 Status visit(Sync &s) { return Status(); }

```

## 4 The GRCPDG class

12     $\langle \text{grcpdg class 12} \rangle \equiv$  (30)  
       class GRC2PDG {  
           CFGmap &dotrefmap;  
           map<GRCNode \*, int> nodenum; // RDFS numbering (index) of each node  
           vector<GRCNode\*> vert; // nodes in RDFS order  
           vector<int> parent; // index of the RDFS spanning tree parent of  
                               // each node  
           vector<int> ancestor;  
           vector<int> semi; // Semi-dominator of each node  
           vector<int> idom; // The immediate dominator of each node  
           vector<set<int> > ichild; // The nodes immediately dominated by each node  
           vector<set<int> > df; // Dominance frontier for each node  
           vector<set<int> > cd; // Nodes control dependent on each node  
           map<int, vector<int> > succmap;  
           map<int, vector<int> > predmap;  
           map<int, bool> reachability;  
           set<int> visited;  
           int N; // Total number of nodes  
           int nullnum;  
           EnterGRC \*enternode;  
           ExitGRC \*exitnode;  
       public:  
            $\langle \text{method declarations 13} \rangle$   
       };

## 5 The Constructor

This uses the algorithm described in Cytron et al. [1] to calculate control dependence relationship and transform the GRC concurrent control-flow graph into a program dependence graph.

```

13 <method declarations 13>≡ (12) 14>
 GRC2PDG(GRCNode *top, CFGmap &dotrefmap) : dotrefmap(dotrefmap)
 {
 assert(top);
 enternode = dynamic_cast<EnterGRC *>(top);
 assert(enternode);
 exitnode = dynamic_cast<ExitGRC *>(enternode->successors[0]);
 assert(exitnode);

 N = 0; // Used to number the nodes during reverse DFS
 reverse_dfs(NULL, exitnode);

 build_dominance_tree();

 df.resize(N);
 compute_dominance_frontier(nodenum[exitnode]);
 //print_df();

 cd.resize(N);
 compute_control_dependence();
 //print_CD();

 //cerr<<"start building pdg\n";
 build_pdg();
 //print_PDG();

 visited.clear();
 opt_dfs(enternode);
 //cerr<<"finished\n";
 }

```

## 6 Depth-first search on the reverse graph

Depth-first search on the reverse graph. Number all the nodes.

```

14 <method declarations 13>+≡ (12) <13 15>
 void reverse_dfs(GRCNode *p, GRCNode *n)
 {
 if (!n || contains(nodenum,n)) return;

 nodenum[n] = N;
 vert.push_back(n);
 parent.push_back(p ? nodenum[p] : -1);
 N++;

 if (n != enternode)
 for (vector<GRCNode*>::iterator i = n->predecessors.begin() ;
 i != n->predecessors.end() ; i++)
 reverse_dfs(n, *i);
 }

```

## 7 Build Dominance Tree

Build the dominance tree for the reverse graph.

```

15 <method declarations 13>+≡ (12) <14 16>
 void build_dominance_tree()
 {
 ancestor.resize(N,-1);
 semi.resize(N,-1);
 idom.resize(N,-1);
 vector<int> samedom;
 samedom.resize(N,-1);

 vector<set<int> > bucket;
 bucket.resize(N);

 ichild.resize(N);

 for (int n = N-1 ; n > 0 ; n--) {

 assert(dotrefmap.count(vert[n])>0); // FIXME: ??

 int p = parent[n];
 int s = p;

 for(vector<GRNode*>::iterator iv = vert[n]->successors.begin() ;
 iv != vert[n]->successors.end() ; iv++) {
 if (*iv) {
 int v = nodenum[*iv];
 int s1 = (v <= n) ? v : semi[ancestor_lowest_semi(v)];
 if (s1 < s) s = s1;
 }
 }

 semi[n] = s;
 if (!contains(bucket[s], n)) bucket[s].insert(n);
 ancestor[n] = p;

 for(set<int>::iterator iv = bucket[p].begin() ;
 iv != bucket[p].end() ; iv++) {
 int v = *iv;
 int y = ancestor_lowest_semi(v);
 if (semi[y] == semi[v]) idom[v] = p;
 else samedom[v] = y;
 }

 bucket[p].clear();
 }

 for (int n = 1 ; n < N ; n++)

```

```

 if (samedom[n] != -1)
 idom[n] = idom[samedom[n]];

 for (int n = 1 ; n < N ; n++)
 if (idom[n] != -1)
 ichild[idom[n]].insert(n);
}

```

## 7.1 ancestor lowest semi

16     $\langle \text{method declarations } 13 \rangle + \equiv$     (12)  $\langle 15 \ 17a \rangle$

```

 int ancestor_lowest_semi(int v)
 {
 int u = v;
 while (ancestor[v] != -1) {
 if (semi[v] < semi[u]) u = v;
 v = ancestor[v];
 }

 return u;
 }

```



## 8 Compute Dominance Frontier

This is Fig. 10 from Cytron et al. [1]. It builds the `df` sets.

```
17a <method declarations 13>+≡ (12) <16 17b>
 void compute_dominance_frontier(int n)
 {
 for(set<int>::iterator iz = ichild[n].begin(); iz != ichild[n].end() ; iz++)
 compute_dominance_frontier(*iz);

 int enternodeidx = nodenum[enternode];

 if (n != enternodeidx) {
 for (vector<GRNode*>::iterator i = vert[n]->predecessors.begin() ;
 i != vert[n]->predecessors.end(); i++) {
 assert(contains(nodenum, *i));
 int y = nodenum[*i];
 if (idom[y] != n && !contains(df[n], y)) {
 assert(contains(dotrefmap, *i));
 df[n].insert(y);
 }
 }
 }

 for(set<int>::iterator iz = ichild[n].begin() ;
 iz != ichild[n].end() ; iz++) {
 int z = *iz;
 for(set<int>::iterator iy = df[z].begin() ; iy != df[z].end() ; iy++) {
 int y = *iy;
 if(idom[y] != n && !contains(df[n], y)) df[n].insert(y);
 }
 }
 }
}
```

## 9 Compute control dependence

This is Fig. 11 from Cytron et al. [1]. It builds the `cd` sets.

```
17b <method declarations 13>+≡ (12) <17a 18>
 void compute_control_dependence()
 {
 for(int y = 0 ; y < N ; y++)
 for(set<int>::iterator ix=df[y].begin() ; ix!=df[y].end() ; ix++) {
 int x = *ix;
 if (!contains(cd[x], y)) cd[x].insert(y);
 }

 //a trick - force EnterGRC's child[1] to be CD of EnterGRC
 cd[nodenum[enternode]].insert(nodenum[enternode->successors[1]]);
 }
}
```

## 10 Build PDG

```

18 <method declarations 13>+≡ (12) <17b 21a>
 void build_pdg()
 {
 copy_conn();
 remove_conn();

 int counter = N;

 //for each node i
 for (int i = 0; i < N; i++) {
 //cerr<<"for node "<<dotrefmap[vert[i]]<<"\n";
 GRNode *n = vert[i];

 assert(dotrefmap.count(vert[i])>0);

 if (n == exitnode) {

 // n is ExitGRC; ignore it

 } else if ((dynamic_cast<Fork *>(n))
 ||
 (n == enternode && (cd[i].size() < 2))) {

 // A parallel node or EnterGRC with a single child:
 // Make each CD member a child, disregard its original child number
 // If n is EnterGRC with 1 child, take it as a parallel node
 // **** something may happen, if one can exit in two branches

 for(set<int>::iterator iy = cd[i].begin() ; iy != cd[i].end() ; iy++) {
 GRNode *y = vert[*iy];
 if (y != exitnode && ((*iy) != i)) {
 n->successors.push_back(y);
 y->predecessors.push_back(n);
 }
 }

 } else if (n == enternode) {

 // EnterGRC with more than 1 child

 Fork *reg = new Fork();
 for (set<int>::iterator iy = cd[i].begin() ; iy != cd[i].end() ; iy++) {
 GRNode *y = vert[*iy];
 if (y != exitnode && ((*iy) != i)) {
 reg->successors.push_back(y);
 y->predecessors.push_back(reg);
 }
 }
 }
 }
 }

```

```

//new region node
nodenum[reg] = counter++;
vert.push_back(reg);
n->successors.push_back(reg);
reg->predecessors.push_back(n);

} else {

// else, for each successor ic of i, make a region node reg

//cerr<<" build regions for ic succ:\n";
for(vector<int>::iterator ic = succmap[i].begin();
 ic != succmap[i].end(); ic++) {

// NULL node
if (*ic == -1){
 n->successors.push_back(NULL);
 //cerr<<" null succ\n";
 continue;
}
if (dynamic_cast<ExitGRC *>(vert[*ic])){
 //cerr<<" exit grc succ\n";
 continue;
}

Fork *reg = new Fork();

//cerr<<" real succ IC "<<dotrefmap[vert[*ic]]<<"\n";

//for each node iy in CD set of node i,
// check if iy is reachable from brunch ic
for(set<int>::iterator iy=cd[i].begin(); iy!=cd[i].end(); iy++){
 if ((dynamic_cast<ExitGRC *>(vert[*iy])) || ((*iy) == i))
 continue;

//cerr<<" IY "<<dotrefmap[vert[*iy]]<<"\n";

reachability.clear();
//cerr<<"testing reachablility...";
if (reachable((*ic), (*iy))) {
 // if yes, add it as a child of the brunch region node reg
 reg->successors.push_back(vert[*iy]);
 vert[*iy]->predecessors.push_back(reg);
}
//cerr<<" finshed\n";
}

//place the region node reg as n's child
// if reg only has one child, add this child directly

```

```

switch (reg->successors.size()){
case 0:
 //if n is sync|switch|test, instead of reg, place a null node there
 if ((dynamic_cast<Switch *>(n)) || (dynamic_cast<Sync *>(n))
 || (dynamic_cast<Test *>(n)))
 n->successors.push_back(NULL);
 break;
case 1:
 n->successors.push_back(reg->successors[0]);
 reg->successors[0]->predecessors.pop_back();
 reg->successors[0]->predecessors.push_back(n);
 reg->successors.clear();
 break;
default:
 //cerr<<"add new reg node: "<<sz<<"\n";
 nodenum[reg] = counter++;
 vert.push_back(reg);
 n->successors.push_back(reg);
 reg->predecessors.push_back(n);
 break;
}
}
//cerr<<"N"<<dotrefmap[vert[i]]<<" is finished\n";
}
}
}

```

## 11 copy conn

21a  $\langle \text{method declarations } 13 \rangle + \equiv$  (12)  $\langle 18 \ 21b \rangle$

```

void copy_conn()
{
 nullnum = 0;

 for (int i = 0; i < N; i++){
 for (vector<GRNode *>::iterator ic = vert[i]->successors.begin();
 ic != vert[i]->successors.end(); ic++){
 if (*ic)
 succmap[i].push_back(nodenum[*ic]);
 else{
 succmap[i].push_back(-1);
 nullnum++;
 }
 }
 for (vector<GRNode *>::iterator ip = vert[i]->predecessors.begin();
 ip != vert[i]->predecessors.end(); ip++){
 predmap[i].push_back(nodenum[*ip]);
 }
 }
}

```

## 12 remove conn

21b  $\langle \text{method declarations } 13 \rangle + \equiv$  (12)  $\langle 21a \ 22 \rangle$

```

void remove_conn()
{
 for (int i = 0; i < N; i++){
 vert[i]->successors.clear();
 if (vert[i] != enternode)
 vert[i]->predecessors.clear();
 }
}

```

## 13 reachable

```

22 <method declarations 13>+≡ (12) <21b 24>
 bool reachable(int from, int to)
 {
 //cerr<<" dfs "<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<"\n";

 if (reachability.count(from) > 0)
 return reachability[from];

 if (from == 0){
 //cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" NO2\n";
 reachability[from] = false;
 return false;
 }

 if (from == -1){
 //cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" YES2\n";
 reachability[from] = true;
 return true;
 }

 if (to == from){
 //cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" YES1\n";
 reachability[from] = true;
 return true;
 }

 assert(vert[from]);

 //for fork node, reachable from any one of the children is reable
 if (dynamic_cast<Fork *>(vert[from])){
 for (vector<int>::iterator ic = succmap[from].begin();
 ic != succmap[from].end(); ic++){
 if (reachable((*ic), to)){
 //cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" YES3\n";
 reachability[from] = true;
 return true;
 }
 }
 //cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" NO3\n";
 reachability[from] = false;
 return false;
 }

 //for else node, reachable means can be reached from all of the children
 for (vector<int>::iterator ic = succmap[from].begin();
 ic != succmap[from].end(); ic++){
 if (!reachable((*ic), to)){
 //cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" NO4\n";

```

```
 reachability[from] = false;
 return false;
 }
}
//cerr<<dotrefmap[vert[from]]<<"->"<<dotrefmap[vert[to]]<<" YES4\n";
reachability[from] = true;
return true;
}
```

## 14 opt dfs

```

24 <method declarations 13>+≡ (12) <22 25>
 void opt_dfs(GRCNode *n)
 {
 vector<GRCNode *>::iterator i,j,k;
 GRCNode *ch;
 bool pis_fork = false;

 if (!n)
 return;

 if (visited.count(nodenum[n]) > 0){
 return;
 }

 if (dynamic_cast<Fork *>(n)){
 pis_fork = true;
 }

 i = n->successors.begin();
 while (i != n->successors.end()){

 if ((*i) == NULL){
 i++;
 continue;
 }

 opt_dfs(*i);

 // remove node with only null children
 if (all_child_null(*i)){
 //for sync, test, switch node, replace i with null child
 if ((dynamic_cast<Sync *>(n))
 || (dynamic_cast<Switch *>(n)) || (dynamic_cast<Test *>(n))){
 rm_predecessor(*i, n);
 (*i) = NULL;
 i++;
 }
 //else remove child i ::FIXME -- is it always true?
 } else{
 if ((*i)->predecessors.size() == 1) //FIXME
 rm_datadps(*i);
 rm_predecessor(*i, n);
 i = n->successors.erase(i);
 //i--;
 }
 continue;
 }
 }

```



```

// & continuous fork nodes with the same control dependence
if ((dynamic_cast<Fork *>(*i)) && (pis_fork)){
 //cerr<<dotrefmap[n]<<"'s child "<<dotrefmap[*i]<<" is also a fork\n";
 if((*i)->predecessors.size() == 1){
 //cerr<<" SAME control flow, merge them\n";
 ch = *i;
 rm_predecessor(*i, n);
 k = n->successors.erase(i);

 for (j = ch->successors.begin(); j != ch->successors.end(); j++){
 k = n->successors.insert(k, *j);
 replace_par(*j, ch, n);
 k++;
 }
 if (ch->predecessors.size() == 0)
 ch->successors.clear();
 i = k-1;
 }
 //else, it has other control dependences, do nothing currently
}
i++;
}

visited.insert(nodenum[n]);

}

```

## 15 rm predecessor

25     $\langle$ method declarations 13 $\rangle + \equiv$  (12)  $\langle$ 24 26a $\rangle$

```

void rm_predecessor(GRCNode *n, GRNode *par)
{
 vector<GRCNode *>::iterator i;

 assert(n);
 assert(par);
 for (i = n->predecessors.begin(); i != n->predecessors.end(); i++)
 if (*i == par){
 n->predecessors.erase(i);
 return;
 }
}

```

## 16 rm datadps

26a     $\langle \text{method declarations } 13 \rangle + \equiv$  (12)  $\langle 25 \ 26b \rangle$

```

void rm_datadps(GRCNode *n)
{
 vector<GRCNode *>::iterator i,j;

 for (i = n->dataPredecessors.begin(); i != n->dataPredecessors.end(); i++){
 for (j = (*i)->dataSuccessors.begin();
 j != (*i)->dataSuccessors.end(); j++){
 if ((*j) == n){
 (*i)->dataSuccessors.erase(j);
 break;
 }
 }
 }

 for (i = n->dataSuccessors.begin(); i != n->dataSuccessors.end(); i++){
 for (j = (*i)->dataPredecessors.begin();
 j != (*i)->dataPredecessors.end(); j++){
 if ((*j) == n){
 (*i)->dataPredecessors.erase(j);
 break;
 }
 }
 }

 n->dataPredecessors.clear();
 n->dataSuccessors.clear();
}

```

## 17 all child null

26b     $\langle \text{method declarations } 13 \rangle + \equiv$  (12)  $\langle 26a \ 27a \rangle$

```

bool all_child_null(GRCNode *n)
{
 int sz;

 assert(n);
 sz = n->successors.size();

 if (sz == 0)
 return false;

 for(vector<GRCNode *>::iterator i = n->successors.begin();
 i != n->successors.end(); i++){
 if (*i)
 return false;
 }

 return true;
}

```

## 18 replace par

27a  $\langle \text{method declarations 13} \rangle + \equiv$  (12)  $\triangleleft 26b$

```

void replace_par(GRCNode *n, GRCNode *org_par, GRCNode *new_par)
{
 vector<GRCNode *>::iterator i;

 for (i = n->predecessors.begin(); i != n->predecessors.end(); i++)
 if ((*i) == org_par){
 // *i = new_par;
 i = n->predecessors.erase(i);
 n->predecessors.insert(i, new_par);
 return;
 }
}

```

## 19 Printing methods

27b  $\langle \text{printing method declarations 27b} \rangle \equiv$  27c  $\triangleright$

```

void print_df()
{
 int i;

 cerr<<"DF\n";
 for(i=0; i<N ;i++){
 cerr<<"Node"<<dotrefmap[vert[i]]<<": ";
 for(set<int>::iterator iy=df[i].begin(); iy!=df[i].end(); iy++){
 cerr<<dotrefmap[vert[*iy]]<<" ";
 }
 cerr<<"\n";
 }
}

```

27c  $\langle \text{printing method declarations 27b} \rangle + \equiv$   $\triangleleft 27b \ 28a \triangleright$

```

void print_CD()
{
 int i;

 cerr<<"CD\n";
 for(i=0; i<N ;i++){
 cerr<<"Node"<<dotrefmap[vert[i]]<<": ";
 for(set<int>::iterator iy=cd[i].begin(); iy!=cd[i].end(); iy++){
 cerr<<dotrefmap[vert[*iy]]<<" ";
 }
 cerr<<"\n";
 }
}

```

28a     *<printing method declarations 27b>+≡* *<27c 28b>*

```

void print_conn()
{
 cerr<<"Connectivity:\n";
 for(int i=0; i<N ;i++){
 cerr<<"Node"<<dotrefmap[vert[i]]<<": ";
 for(vector<int>::iterator iy=succmap[i].begin(); iy!=succmap[i].end(); iy++)
 if (*iy > -1)
 cerr<<dotrefmap[vert[*iy]]<<" ";
 else
 cerr<<"NULL ";
 cerr<<"\n";
 }
}

```

28b     *<printing method declarations 27b>+≡* *<28a*

```

void print_PDG()
{
 cerr<<"PDG:\n";

 for (int i = 0; i < (int)(vert.size()); i++){

 if (!(vert[i]))
 continue;

 cerr<<"Node"<<dotrefmap[vert[i]]<<": ";
 for(vector<GRNode *>::iterator iy=vert[i]->successors.begin();
 iy!=vert[i]->successors.end(); iy++)
 cerr<<dotrefmap[*iy]<<" ";
 cerr<<"\n";
 }
}

```

## 20 Main function

29     $\langle \text{main function 29} \rangle \equiv$  (30)

```

int main(int argc, char* argv[])
{
 IR::XMListream f(std::cin);
 IR::Node *n;
 f >> n;

 Modules *mods = dynamic_cast<AST::Modules*>(n);
 if (!mods) {
 std::cerr<<"Root node is not a module object\n";
 exit(-2);
 }

 for(vector<AST::Module*>::iterator i = mods->modules.begin();
 i != mods->modules.end(); i++){
 assert(*i);

 GRCgraph *gf = dynamic_cast<GRCgraph*>((*i)->body);
 assert(gf);
 GRCNode *top = gf->control_flow_graph;

 CFGmap dotrefmap;
 STmap strefmap;

 EnterGRC *engrc = dynamic_cast<EnterGRC*>(top);
 assert(engrc);

 // compute the data dependencies between Enter & STsuspend nodes
 // remove & recompute dps between variables
 Dependencies vardps;
 vardps.compute(engrc);

 STDPS compdps(engrc);
 compdps.execute();

 // Convert the GRC graph into a PDG
 gf->enumerate(dotrefmap, strefmap);
 GRC2PDG converter(top, dotrefmap);
 }

 IR::XMLostream o(std::cout);
 o << n;

 return 0;
}

```

```
30 <cec-grcpdg.cpp 30>≡
 #include "IR.hpp"
 #include "AST.hpp"

 #include <iostream>
 #include <fstream>
 #include <set>
 #include <map>
 #include <vector>

 using namespace AST;
 using namespace std;

 typedef map<GRCNode *, int> CFGmap;
 typedef map<STNode *, int> STmap;

 <utilities 2a>

 <stdps class 3>

 <dependency class 5>

 <dependency method definitions 6a>

 <grcpdg class 12>

 <main function 29>
```

## References

- [1] Ron Cytron, Jeanne Ferrante, Barry K. Rosen, Mark N. Wegman, and F. Kenneth Zadeck. Efficiently computing static single assignment form and the control dependence graph. *ACM Transactions on Programming Languages and Systems*, 13(4):451–490, October 1991.