

# Fundamentals of Computer Systems

## Thinking Digitally

Stephen A. Edwards

Columbia University

Summer 2020

## The Subject of this Class

0

## The Subjects of this Class

0

1

But let your communication be, Yea, yea; Nay, nay: for whatsoever is more than these cometh of evil.

— Matthew 5:37



*SUPERCODER 2000*

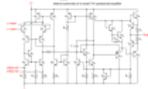
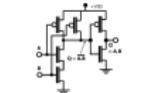
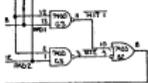
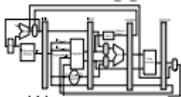
Air cooled coding keyboard for professional use.



# Engineering Works Because of Abstraction



```
;; voice 1 wave select
ld a, (#CH1_W_NUM)
and a
ld a, (#CH1_W_SEL)
jr nz, #00b4
ld a, (#CH1_E_TABLE0)
```



Application Software

Operating Systems

Architecture

Micro-Architecture

Logic

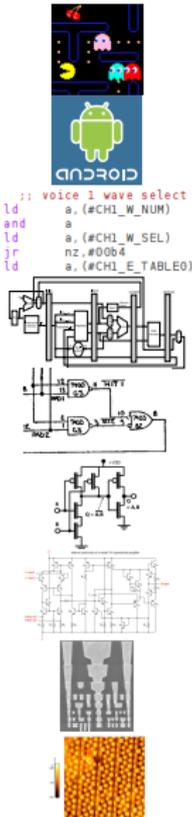
Digital Circuits

Analog Circuits

Devices

Physics

# Engineering Works Because of Abstraction



Application Software COMS 3157, 4156, et al.

Operating Systems COMS W4118

Architecture Second Half of 3827

Micro-Architecture Second Half of 3827

Logic First Half of 3827

Digital Circuits First Half of 3827

Analog Circuits ELEN 3331

Devices ELEN 3106

Physics ELEN 3106 et al.

# Boring Stuff

<http://www.cs.columbia.edu/~sedwards/classes/2020/3827-summer/>

Prof. Stephen A. Edwards  
sedwards@cs.columbia.edu

Lectures 1:00 – 4:00 PM, Mondays and Wednesdays  
May 27–July 1

Weight	What	When
40%	Homeworks	See Webpage
60%	Final exam	<del>July 1st</del> July 1st

Submit homework online via Courseworks

# Software You Need



INKSCAPE

The Inkscape SVG File Editor [inkscape.org](https://inkscape.org)

Do homework by downloading an SVG file from the class website, edit it in Inkscape, and upload it to Courseworks

The Digital Circuit Simulator [github.com/hneemann/Digital](https://github.com/hneemann/Digital)

Circuit design problems: download (class website) .zip file with .dig files, edit with Digital, upload to Courseworks

SPIM: A MIPS32 Simulator [spimsimulator.sourceforge.net](https://spimsimulator.sourceforge.net)

MIPS assembly coding:, download .zip file with .s files, edit in favorite text editor, test and debug in SPIM, upload to Courseworks

# Rules and Regulations

Each assignment turned in must be unique; work must ultimately be your own.

*Don't cheat: Columbia Students Aren't Cheaters*

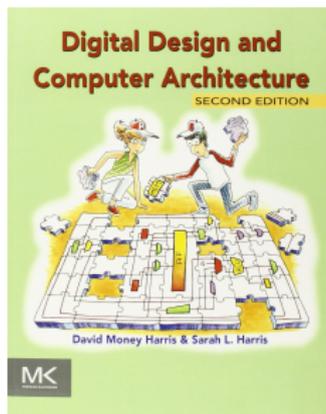
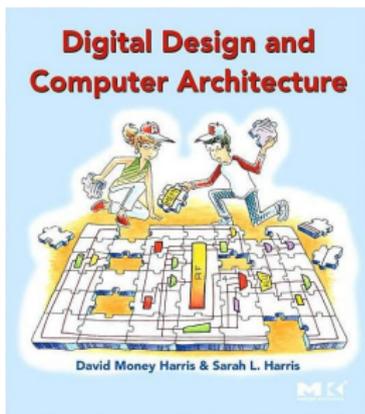
Test will be closed-book; you may use a single sheet of your own notes

# Optional Texts: Alternative 1

No required text. One option:

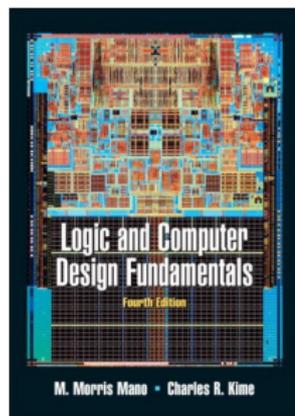
- ▶ David Harris and Sarah Harris. *Digital Design and Computer Architecture*. Either 1st or 2nd ed.

Almost precisely right for the scope of this class: digital logic and computer architecture.

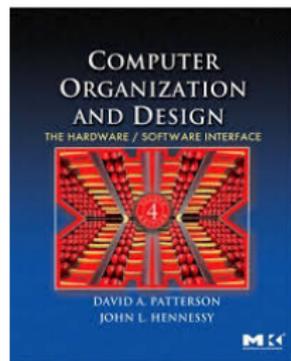


## Optional Texts: Alternative 2

- ▶ M. Morris Mano and Charles Kime. *Logic and Computer Design Fundamentals*. 4th ed.



- ▶ David A. Patterson and John L. Hennessy. *Computer Organization and Design, The Hardware/Software Interface*. 4th ed.



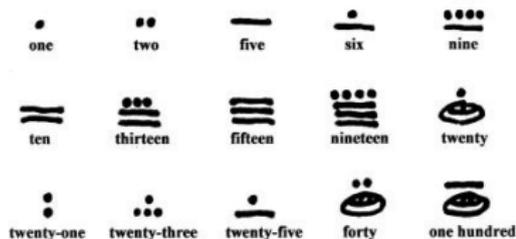
GILDAN  
ULTRA  
COTTON™

There are only 10 types  
of people in the world:  
Those who understand binary  
and those who don't.

# Which Numbering System Should We Use?



Roman: I II III IV V VI VII VIII IX X



Mayan: base 20, Shell = 0

1	𐎀	11	𐎁𐎂	21	𐎁𐎂𐎁	31	𐎁𐎂𐎁𐎂	41	𐎁𐎂𐎁𐎂𐎁	51	𐎁𐎂𐎁𐎂𐎁𐎂
2	𐎀𐎀	12	𐎁𐎂𐎀	22	𐎁𐎂𐎁𐎀	32	𐎁𐎂𐎁𐎂𐎀	42	𐎁𐎂𐎁𐎂𐎁𐎀	52	𐎁𐎂𐎁𐎂𐎁𐎂𐎀
3	𐎀𐎀𐎀	13	𐎁𐎂𐎀𐎀	23	𐎁𐎂𐎁𐎀𐎀	33	𐎁𐎂𐎁𐎂𐎀𐎀	43	𐎁𐎂𐎁𐎂𐎁𐎀𐎀	53	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀
4	𐎀𐎀𐎀𐎀	14	𐎁𐎂𐎀𐎀𐎀	24	𐎁𐎂𐎁𐎀𐎀𐎀	34	𐎁𐎂𐎁𐎂𐎀𐎀𐎀	44	𐎁𐎂𐎁𐎂𐎁𐎀𐎀𐎀	54	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀𐎀
5	𐎀𐎀𐎀𐎀𐎀	15	𐎁𐎂𐎀𐎀𐎀𐎀	25	𐎁𐎂𐎁𐎀𐎀𐎀𐎀	35	𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀	45	𐎁𐎂𐎁𐎂𐎁𐎀𐎀𐎀𐎀	55	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀
6	𐎀𐎀𐎀𐎀𐎀𐎀	16	𐎁𐎂𐎀𐎀𐎀𐎀𐎀	26	𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀	36	𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀	46	𐎁𐎂𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀	56	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀
7	𐎀𐎀𐎀𐎀𐎀𐎀𐎀	17	𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀	27	𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀𐎀	37	𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀	47	𐎁𐎂𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀𐎀	57	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀
8	𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	18	𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀𐎀	28	𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀	38	𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀𐎀	48	𐎁𐎂𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀	58	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀𐎀
9	𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	19	𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	29	𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	39	𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	49	𐎁𐎂𐎁𐎂𐎁𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀	59	𐎁𐎂𐎁𐎂𐎁𐎂𐎀𐎀𐎀𐎀𐎀𐎀𐎀𐎀
10	𐎀	20	𐎁	30	𐎁𐎂	40	𐎁𐎂𐎀	50	𐎁𐎂𐎀𐎀		

Babylonian: base 60

# The Decimal Positional Numbering System

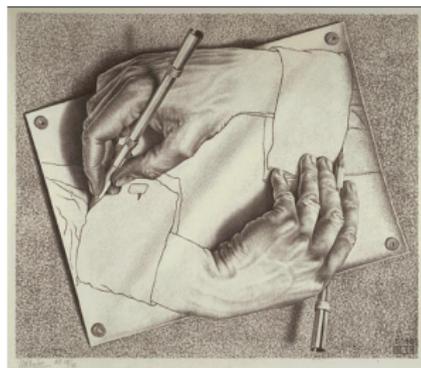


Ten figures: 0 1 2 3 4 5 6 7 8 9

$$730_{10} = 7 \times 10^2 + 3 \times 10^1 + 0 \times 10^0$$

$$990_{10} = 9 \times 10^2 + 9 \times 10^1 + 0 \times 10^0$$

Why base ten?



## Hexadecimal, Decimal, Octal, and Binary

Hex	Dec	Oct	Bin
0	0	0	0
1	1	1	1
2	2	2	10
3	3	3	11
4	4	4	100
5	5	5	101
6	6	6	110
7	7	7	111
8	8	10	1000
9	9	11	1001
A	10	12	1010
B	11	13	1011
C	12	14	1100
D	13	15	1101
E	14	16	1110
F	15	17	1111

# Binary and Octal: Electronics Likes Powers of Two



DEC PDP-8/I, c. 1968

Oct	Bin
0	0
1	1
2	10
3	11
4	100
5	101
6	110
7	111

$$\begin{aligned} \text{PC} &= 010110111101_2 \\ &= 0 \times 2^{11} + 1 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + 1 \times 2^7 + 0 \times 2^6 + \\ &\quad 1 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= 2675_8 \\ &= 2 \times 8^3 + 6 \times 8^2 + 7 \times 8^1 + 5 \times 8^0 \\ &= 1469_{10} \end{aligned}$$

# Hexadecimal Numbers

Base 16: 0 1 2 3 4 5 6 7 8 9 A B C D E F

Instead of groups of 3 bits (octal), Hex uses groups of 4.

$$\begin{aligned}\text{CAFEEF00D}_{16} &= 12 \times 16^7 + 10 \times 16^6 + 15 \times 16^5 + 14 \times 16^4 + \\ &\quad 15 \times 16^3 + 0 \times 16^2 + 0 \times 16^1 + 13 \times 16^0 \\ &= 3,405,705,229_{10}\end{aligned}$$

C	A	F	E	F	0	0	D	Hex			
1100	1010	1111	1110	1110	0000	0000	1101	Binary			
3	1	2	7	7	5	7	0	0	1	5	Octal

# Computers Rarely Manipulate True Numbers

Infinite memory still very expensive

Finite-precision numbers typical

32-bit processor: naturally manipulates 32-bit numbers

64-bit processor: naturally manipulates 64-bit numbers

How many different numbers can you

represent with 5

binary	$2^5 = 32$	digits?
octal	$8^5 = 32768$	
decimal	$10^5 = 100,000$	
hexadecimal	$16^5 = 1,048,576$	

Handwritten diagram showing 5 bits grouped together with a brace, equated to  $2^5 = 32_{10}$ .

# Jargon



Drill

Bit Binary digit: 0 or 1



Godzilla

Byte Eight bits



Prison Tattoo

Word Natural number of bits for the processor, e.g., 16, 32, 64



LSB

LSB

101001

Least Significant Bit ("rightmost")



MSB

MSG

Most Significant Bit ("leftmost")

# Decimal Addition Algorithm

$$\begin{array}{r} 434 \\ +628 \\ \hline \end{array} \quad 2$$

$$4 + 8 = 12$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

# Decimal Addition Algorithm

$$\begin{array}{r} 1 \\ 434 \\ + 628 \\ \hline 2 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

# Decimal Addition Algorithm

$$\begin{array}{r} 1 \\ 434 \\ + 628 \\ \hline 62 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

# Decimal Addition Algorithm

$$\begin{array}{r} 1 \ 1 \\ 434 \\ + 628 \\ \hline 062 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

# Decimal Addition Algorithm

$$\begin{array}{r} 1\ 1 \\ 434 \\ +628 \\ \hline 1062 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

# Binary Addition Algorithm

$$\begin{array}{r} 10011 \\ +11001 \\ \hline \end{array}$$

$$1 + 1 = 10$$

+	0	1
0	00	01
1	01	10
10	10	11

# Binary Addition Algorithm

$$\begin{array}{r} \phantom{0}1 \\ 10011 \\ +11001 \\ \hline \phantom{0}0 \end{array}$$

$$1 + 1 = 10$$

$$1 + 1 + 0 = 10$$

+	0	1
0	00	01
1	01	10
10	10	11

# Binary Addition Algorithm

$$\begin{array}{r} \phantom{100}11 \\ 10011 \\ +11001 \\ \hline \phantom{100}00 \end{array}$$

$$1 + 1 = 10$$

$$1 + 1 + 0 = 10$$

$$1 + 0 + 0 = 01$$

+	0	1
0	00	01
1	01	10
10	10	11

# Binary Addition Algorithm

$$\begin{array}{r} 011 \\ 10011 \\ +11001 \\ \hline 100 \end{array}$$

$$1 + 1 = 10$$

$$1 + 1 + 0 = 10$$

$$1 + 0 + 0 = 01$$

$$0 + 0 + 1 = 01$$

+	0	1
0	00	01
1	01	10
10	10	11

# Binary Addition Algorithm

$$\begin{array}{r} 0011 \\ 10011 \\ +11001 \\ \hline 1100 \end{array}$$

$$1 + 1 = 10$$

$$1 + 1 + 0 = 10$$

$$1 + 0 + 0 = 01$$

$$0 + 0 + 1 = 01$$

$$0 + 1 + 1 = 10$$

+	0	1
0	00	01
1	01	10
10	10	11

# Binary Addition Algorithm

$$\begin{array}{r} 10011 \\ 10011 \\ +11001 \\ \hline 101100 \end{array}$$

$$1 + 1 = 10$$

$$1 + 1 + 0 = 10$$

$$1 + 0 + 0 = 01$$

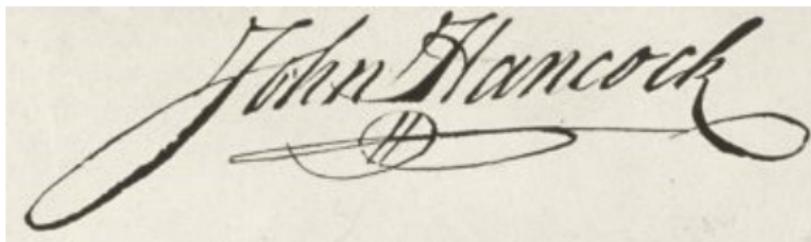
$$0 + 0 + 1 = 01$$

$$0 + 1 + 1 = 10$$

+	0	1
0	00	01
1	01	10
10	10	11

# Signed Numbers: Dealing with Negativity

How should we represent negative numbers?

A handwritten signature in cursive script, reading "John Hancock". The signature is written in dark ink on a light-colored, aged paper background. The letters are fluid and connected, with a prominent flourish under the "H" and a long, sweeping underline that extends to the right.

## Binary Signed Magnitude Numbers

The familiar notation: negative numbers have a leading –

Binary signed-magnitude encoding: leading 1 indicates negative; remaining bits treated as binary.

$$0000_2 = 0$$

$$0010_2 = 2$$

$$1010_2 = -2$$

$$1111_2 = -7$$

$$1000_2 = -0?$$

Can be made to work, but addition is annoying:

If the signs match, add the magnitudes and use the same sign.

If the signs differ, subtract the smaller number from the larger; return the sign of the larger.

# One's Complement Numbers

Like Signed Magnitude, a leading 1 indicates a negative One's Complement number. However, number magnitude is *complement* of remaining bits interpreted as binary.

To negate a number, complement (flip) each bit.

$$0000_2 = 0$$

$$0010_2 = 2$$

$$1101_2 = -2$$

$$1000_2 = -7$$

$$1111_2 = -0?$$

Addition is nicer: just add the one's complement numbers as if they were normal binary.

Really annoying having a  $-0$ : two numbers are equal if their bits are the same or if one is 0 and the other is  $-0$ .



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ZEROS  
ARE CREATED  
EQUAL**

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## Two's Complement Numbers



Really neat trick: just make only the most significant bit represent a *negative* number instead of positive; treat the rest as binary.

$$1101_2 = -8 + 4 + 1 = -3$$

$$1111_2 = -8 + 4 + 2 + 1 = -1$$

$$0111_2 = 4 + 2 + 1 = 7$$

$$1000_2 = -8$$

Easy addition: just add in binary and discard any carry.

Negation: complement each bit (as in one's complement) then add 1.

Subtraction done with negation and addition.

Very good property: no  $-0$

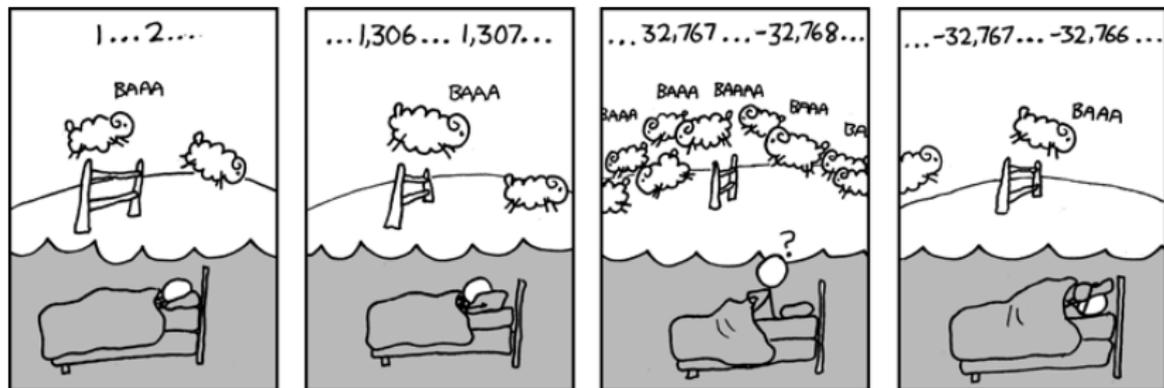
Two's complement numbers are equal if and only if all their bits are the same.

# Number Representations Compared

Code	Binary	Signed Mag.	One's Comp.	Two's Comp.
0000	0	0	0	0
0001	1	1	1	1
⋮				
0111	7	7	7	7
1000	8	-0	-7	-8
1001	9	-1	-6	-7
⋮				
1110	14	-6	-1	-2
1111	15	-7	-0	-1

Smallest number

Largest number



<https://xkcd.com/571/>

How many bits in his brain?

# Fixed-point Numbers



How to represent fractional numbers? In decimal, we continue with negative powers of 10:

$$31.4159 = 3 \times 10^1 + 1 \times 10^0 + \\ 4 \times 10^{-1} + 1 \times 10^{-2} + 5 \times 10^{-3} + 9 \times 10^{-4}$$

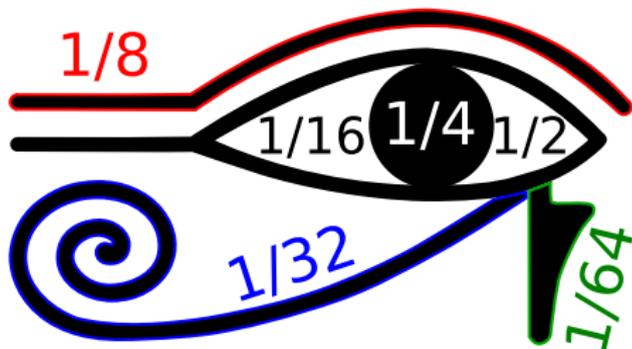
Also works in binary:

$$1011.0110_2 = 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 + \\ 0 \times 2^{-1} + 1 \times 2^{-2} + 1 \times 2^{-3} + 0 \times 2^{-4} \\ = 8 + 2 + 1 + 0.25 + 0.125 \\ = 11.375$$

Addition and subtraction algorithms the same.

F F  
u a  
c  
Interesting

The ancient Egyptians used binary fractions:



The Eye of Horus

# Binary-Coded Decimal



thinkgeek.com

Humans prefer reading decimal numbers; computers prefer binary.

BCD is a compromise: every four bits represents a decimal digit.

Dec	BCD
0	0000 0000
1	0000 0001
2	0000 0010
⋮	⋮
8	0000 1000
9	0000 1001
10	0001 0000
11	0001 0001
⋮	⋮
18	0001 1000
19	0001 1001
20	0010 0000
⋮	⋮



# BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 158 \\ +242 \\ \hline \end{array}$$

$$\begin{array}{r} 000101011000 \\ +001001000010 \\ \hline \phantom{000}1010 \text{ First group} \\ \phantom{000}+0110 \text{ Correction} \\ \hline \end{array}$$

$$\begin{array}{r} \hline \\ \hline \end{array}$$

# BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 1 \\ 158 \\ +242 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 1 \\ 000101011000 \\ +001001000010 \\ \hline 1010 \quad \text{First group} \\ +0110 \quad \text{Correction} \\ \hline 10100000 \quad \text{Second group} \\ \hline \\ \hline \end{array}$$

# BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 1 \\ 158 \\ +242 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 1 \\ 000101011000 \\ +001001000010 \\ \hline 1010 \quad \text{First group} \\ +0110 \quad \text{Correction} \\ \hline 10100000 \quad \text{Second group} \\ +0110 \quad \text{Correction} \\ \hline \hline \end{array}$$

# BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 11 \\ 158 \\ +242 \\ \hline 00 \end{array}$$

$$\begin{array}{r} \phantom{000}1\phantom{00}1\phantom{00}1\phantom{00}000 \\ +001001000010 \\ \hline \phantom{000}1010 \text{ First group} \\ \phantom{000}+ 0110 \text{ Correction} \\ \hline \phantom{000}10100000 \text{ Second group} \\ \phantom{000}+ 0110 \text{ Correction} \\ \hline 01000000 \text{ Third group} \\ \hline \end{array}$$

# BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

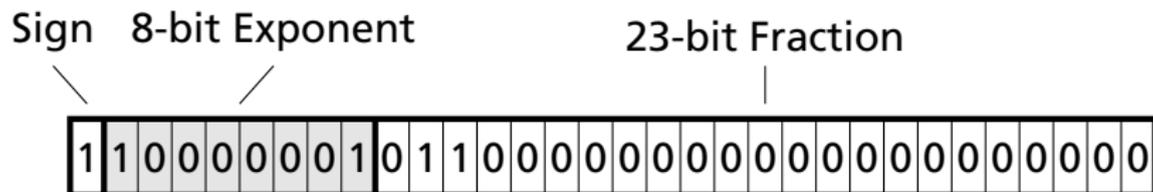
$$\begin{array}{r} 11 \\ 158 \\ +242 \\ \hline 400 \end{array}$$

$$\begin{array}{r} \phantom{000}1\phantom{000}1\phantom{000} \\ 000101011000 \\ +001001000010 \\ \hline \phantom{000}1010 \text{ First group} \\ \phantom{000}+0110 \text{ Correction} \\ \hline \phantom{000}10100000 \text{ Second group} \\ \phantom{000}+0110 \text{ Correction} \\ \hline 01000000 \text{ Third group} \\ \phantom{000} \text{ (No correction)} \\ \hline 010000000000 \text{ Result} \end{array}$$

# Floating-Point Numbers: "Scientific Notation"

Greater dynamic range at the expense of precision  
Excellent for real-world measurements

IEEE 754 Single-Precision (32-bit)



implicit "excess 127"

$$= - 1.0110000_2 \times 2^{10000001_2 - 127}$$
$$= -1.375 \times 2^2$$
$$= -5.5$$

## ASCII For Representing Characters and Strings

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	<b>0</b>	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>	<b>6</b>	<b>7</b>
<b>0</b>	NUL	DLE	SP	0	@	P	'	p
<b>1</b>	SOH	DC1	!	1	A	Q	a	q
<b>2</b>	STX	DC2	"	2	B	R	b	r
<b>3</b>	ETX	DC3	#	3	C	S	c	s
<b>4</b>	EOT	DC4	\$	4	D	T	d	t
<b>5</b>	ENQ	NAK	%	5	E	U	e	u
<b>6</b>	ACK	SYN	&	6	F	V	f	v
<b>7</b>	BEL	ETB	'	7	G	W	g	w
<b>8</b>	BS	CAN	(	8	H	X	h	x
<b>9</b>	HT	EM	)	9	I	Y	i	y
<b>A</b>	LF	SUB	*	:	J	Z	j	z
<b>B</b>	VT	ESC	+	;	K	[	k	{
<b>C</b>	FF	FS	,	<	L	\	l	
<b>D</b>	CR	GS	-	=	M	]	m	}
<b>E</b>	SO	RS	.	>	N	^	n	~
<b>F</b>	SI	US	/	?	O	_	o	DEL

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