COMS W4115 Fall 2012
Course Project
EasKey Final Report

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Chapter 1

An Introduction to EasKey

EasKey aims to facilitate programming of keyboard and mouse actions. It consists of most fundamental operations and functions of regular programming languages, such as branches and loops. Besides, all the potential keyboard and mouse actions are realized by EasKey internal functions and integrated tightly with the compiler, which makes the process of keyboard and mouse programming super easy for users.

1.1 Motivation

Technology makes people's life easier. Computer is supposed to help people complete iterative and boring calculations. However, with the high-speed development of software industry, software installed in a computer becomes more and more complicated and often requires thousands of mouse clicks and keyboard inputs, most of which are repetitive and boring.

Take a few examples, in the course registration period, you may wish to register a class that is already full. So the most common way may be that you constantly press the “Class Search” button to refresh the website in order to see that if someone has just dropped the course so you are lucky enough to fill in the position. Or perhaps you need to update a database system with thousands of operations that are similar to each other, so you have to input all the instructions one by one. Or you may want to click your mouse as fast as possible to win a game which needs immediate reaction.

In a circumstance like any of these, it is definitely a waste of time to do finish the task in normal ways presented above. So, is there something we can do to make it easier and faster and more efficient? If what we do is simply a combination of keyboard inputs and mouse clicks, why don't we let the computer itself to finish the task? These results in our idea to design a language to facilitate people doing iterative keyboard and mouse input. And the solution is our EasKey Language!

1.2 Description

Our EasKey language is developed using OCaml language, and it is text-based and has a lot of similarity with C language. We support all the basic types and functionality defined in microC, which is the mainly reference language that we turned to, and we have built lots of built-in functions to facilitate the basic operation of mouse and keyboard. In EasKey, we define the same types of statements as in C but with different grammar,
which help to better organize the structure of program. The basic idea and rules in EasKey is to discard the use of parenthesis “{ }” as the notation of block and replace it with key word “end”, besides, make the best convenience for users to manipulate the windows API operations with just sparse code. With EasKey language, user can get the position of cursor and keep it clicking any button thousands of time. Also, just like many other primary language, EasKey supports user self-defined function and the basic loop or recursive function.

1.3 Features

Functional
The mainly advantage that EasKey has over other language is the integration of several operations of Mouse and keyboard in windows environment. It makes programming easy for the user who doesn’t have much programming experience.

Abundant Data Types
EasKey is, to some extent, derived from microC, but far more powerful than microC. We have several data types, including basic Integer, Float, String, Char, Boolean and the extended data type, point, color and key.

Statically-scoped
EasKey is statically-scoped which means that the scope of a defined variable in EasyKey will be extended when a function is called later.

Applicative-ordered
Just like C, EasKey uses applicative-order during the evaluation.

Robust
EasKey has a very detailed-checking analyzer, which helps to minimize the potential bug in the program.

Convenient and Interesting
As we have stated, EasKey is designed to facilitate people to get rid of boring routine computer operation, which we think is very interesting and commercial promising.
Chapter 2
Language Tutorial

EasKey is a really simple language that can easily get started with. User really don’t need to have much knowledge about programming and can easily use the inherited function to build the program they want. Now, let’s start the tutorial with a simple example.

2.1 Introductory Sample

This is a very simple function that print a string “EasKey” to the terminal:
function void out(string s)
  __print(s);
  return;
end function

function int main()
  string s=”EasKey”;
  out(s);
  return 0;
end function

2.2 Functionalities

2.2.1 Variable declaration

We have 8 kinds of variables, and there declaration should follow the form as bellow:

    int a=1;
    string b=”EasKey”;
    char c=”E”;
    float d=3.1415;
    boolean e=true;
    point f={50,50};
    key g=$F$;
    color h=#123,456,789#;
2.2.2 Loop

There are two kinds of loop can be defined in our language:
For loop:
   for (i=0;i<10;i=i+1)
      forstatement;
   end for

While loop:
   while (i !=0)
      whilestatement;
   end while

2.2.3 recursion

Our language support recursion, which means a function can call itself in its function body

function int fib (int n)
   if(n == 1)
      return 1;
   elseif(n == 2)
      return 1;
   else
      return (fib(n-1)+fib(n-2));
   end if
end function

2.2.4 condition

There are two types of condition control
if :
   if (a==b)
      a=a+1;
   else b=b+1;
   end if

switch:
   int i;
   switch (i)
      case i: a = 1;
case i: a = 2;
case i: a = 3;
default: a = 4;
end switch

2.2.5 function
We have defined a lot of built in functions which can help user better achieve their goal:

void __leftClickPoint(point p);  -------perform mouse left click at point p
void __rightClickPoint(point p);  -------perform mouse right click at point p
void __movePoint(point p);        -------move cursor to point p
void __leftDown();               -------perform mouse left click at current point
void __rightDown();              -------perform mouse right click at current point
void __keyStroke (key k);        -------press key k on the keyboard
point __getPoint();              -------get the point value of current cursor location
color __getColor(point p);       -------get the color value of certain point p
void __print(string s);          -------print out a string s in the terminal
string __stringOfInt(int i);     -------convert a variable of type int to type string
void __delay(int t);             -------delay t time

Besides, user can build their own function with the form like this:

function foo(int i, int j)

        if (i==j)
            __print("equal");
        else __print("nonequal");
end if

end function
Chapter 3
Language Reference Manual

1. Lexical conventions
Four kinds of tokens are included in EasKey: identifiers, keywords, constants and expression operators. Blanks, tabs, newlines and comments can only serve to separate tokens, otherwise they will be ignored. Tokens are distinguished by adding at least one of the characters above to separate them.

1.1. Comments
Any character between a /* and the first */ after that are considered as a comment. Nested /* and */ pairs are illegible.

1.2. Identifiers
An identifier is a sequence of letters and digits; the first character must be alphabetic. The underscore ‘_’ counts as alphabetic. Upper and lower case letters are considered different.

1.3. Keywords
The following identifiers are reserved for use as keywords, and may not be used otherwise:

- `void`
- `int`
- `float`
- `char`
- `string`
- `boolean`
- `point`
- `color`
- `key`
- `if`
- `elseif`
- `else`
- `while`
- `for`
- `continue`
- `break`
- `switch`
- `case`
- `default`
- `function`
- `return`
- `end`
- `true`
- `false`

1.4. Constants
There are several kinds of constants as follows.
1.4.1. Integer constants
An integer constant is a sequence of digits.

1.4.2. Character constants
A character constant is one or two characters enclosed in single quotes ‘‘’. Within a
character constant a single quote must be preceded by a backslash ‘‘\’’. Certain
non-graphic characters, and ‘‘\’’ itself, may be escaped according to the following table:

- back space \b
- newline \n
- carriage return \r
- horizontal tab \t
- null byte \0
- ddd \ddd
- \\ 

The escape ‘‘\ddd’’ consists of the backslash followed by exact 3 digits which are taken
to specify the value of the desired character. The number ddd must fall in the region 000 to 255.

1.4.3. Floating constants
A floating constant consists of an integer part, a decimal point, a fraction part, an e, and
an optionally signed integer exponent. The integer and fraction parts both consist of a
sequence of digits. Either the integer part or the fraction part (not both) may be missing;
either the decimal point or the e and the exponent (not both) may be missing.

1.4.4. Boolean constants
A boolean constant only has two possible values: true or false.

1.4.5. Point constants
A point constant is two integer constants separated by a comma ‘,’ and enclosed in
parentheses ‘‘{’’, ‘’}’’, indicating the coordinates of a pixel in the computer screen.

1.4.6. Key constants
A key constant is one or more characters enclosed in ‘‘$’’, indicating the most of the keys
of an ordinary keyboard. The key constants are as follows:

<table>
<thead>
<tr>
<th>$a$</th>
<th>$b$</th>
<th>$c$</th>
<th>$d$</th>
<th>$e$</th>
<th>$f$</th>
<th>$g$</th>
<th>$h$</th>
<th>$i$</th>
<th>$j$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$k$</td>
<td>$l$</td>
<td>$m$</td>
<td>$n$</td>
<td>$o$</td>
<td>$p$</td>
<td>$q$</td>
<td>$r$</td>
<td>$s$</td>
<td>$t$</td>
</tr>
<tr>
<td>$u$</td>
<td>$v$</td>
<td>$w$</td>
<td>$x$</td>
<td>$y$</td>
<td>$z$</td>
<td>$A$</td>
<td>$B$</td>
<td>$C$</td>
<td>$D$</td>
</tr>
<tr>
<td>$E$</td>
<td>$F$</td>
<td>$G$</td>
<td>$H$</td>
<td>$I$</td>
<td>$J$</td>
<td>$K$</td>
<td>$L$</td>
<td>$M$</td>
<td>$N$</td>
</tr>
</tbody>
</table>
1.4.7. Color constants
A color constant consists of three integer constants enclosed in “#”, indicating the color of pixel in the computer screen. The three integers must fall in the region 000 to 255, representing the three parameters in a RGB color model.

1.4.8. String constants
A string is a sequence of characters surrounded by double quotes “"”. String itself is a type. In a string, the character “"” must be preceded by a “\"”; in addition, the same escapes as described for character constants may be used.

2. Conversions
EasKey supports the following type conversions based on built-in functions:

- `string __stringOfInt(int i);`
- `string __stringOfFloat(float f);`
- `string __stringOfChar(char c);`
- `string __stringOfBoolean(boolean b);`
- `string __stringOfPoint(point p);`
- `string __stringOfKey(key k);`
- `string __stringOfColor(color c);`

3. Expressions
3.1. Primary expressions
Primary expressions include the following.

3.1.1. Identifier
An identifier is a primary expression, provided it has been suitably declared as discussed below. Its type is specified by its declaration.

3.1.2. Constant
Any constant is a primary expression.

3.1.3. ( expression )
A parenthesized expression is a primary expression whose type and value are identical to those of the unadorned expression.

3.1.4. primary-expression ( expression-list opt )
A function call is a primary expression followed by parentheses containing a possibly empty, comma separated list of expressions which constitute the actual arguments to the function. The result of the function call is of return type of that function.

3.2. Unary operators
Expressions with unary operators group right-to-left.

3.2.1. –expression
The result is the negative of the expression, and has the same type. The type of the expression must be int or float.

3.2.2. !expression
The result of the logical negation operator ! is true if the value of the expression is false, false if the value of the expression is true. Note that this operator is applicable to int or char type.

3.2.3. expression @x
The result is the x-coordinate of the expression, and has the integer type. The type of the expression must be point.

3.2.4. expression @y
The result is the y-coordinate of the expression, and has the integer type. The type of the expression must be point.

3.2.5. expression @r
The result is the red parameter of the expression, and has the integer type. The type of the expression must be color.

3.2.6. expression @g
The result is the green parameter of the expression, and has the integer type. The type of the expression must be color.

3.2.7. expression @b
The result is the blue parameter of the expression, and has the integer type. The type of the expression must be color.

3.2.8. expression @k
The result is the string of the expression, and has the string type. The type of the expression must be key.

3.3 Additive operators
3.3.1 expression + expression
The result is the sum of the expressions. Both operands must have the same type, which should be int, char, float. If the expressions are of type point or color, then each field is added up respectively. If the expressions are of type string, then the string are concatenated.

3.3.2 expression – expression
The result is the difference of the operands. Same type consideration is applied here, except that type string cannot be used here.

3.3 Multiplicative operators
3.3.1 expression * expression
The binary * operator indicates multiplication. Both operands must be int or float, and the result is of the same type.

3.3.2 expression / expression
The binary / operator indicates division. The same type considerations as for multiplication apply.

3.3.3 expression % expression
The binary % operator obtains the remainder from the division of the first expression by the second. Both operands must be int, and the result is int.

4. Relational operators
The relational operators group left or right. Relational operators cannot be used sequentially. For example, “a<b<c” is not permitted by EasKey.
expression < expression;
expression > expression;
expression <= expression;
expression >= expression
The operators < (less than), > (greater than), <= (less than or equal to) and >= (greater than or equal to) return boolean type. Relational operators can apply to int, char, float, string types.

5. Equality operators
EasKey do not have pointer, the equality operators only operates the contents.
6.1 expression == expression; expression != expression
The == (equal to) and the != (not equal to) operators are exactly analogous to the relational operators except for their lower precedence. Equality operators can be used to compare all the types defined in EasKey.

6.2 expression && expression
The && operator returns true if both its operands are true, false otherwise. The precedence of && is higher than ||, but lower than relational operators and equality operators.

6.3 expression || expression
The || operator returns true if either of its operands is true, false otherwise.

7. Assignment operators
There is one assignment operator, which is ‘=’. The value of the right group is assigned to left. The assignment operator requires an identifier as their left operand, and the type of an assignment expression is that of its left operand. The value is the value stored in the left operand after the assignment has taken place.
7.1 identifier = expression
The value of the expression replaces that of the identifier. The operands must have the same type. The value of the expression is simply stored into the object referred to by the identifier.

8. Declarations
Declarations are used within function definitions to specify the interpretation of each identifier. Declarations have the form

    declaration:
    decl-specifiers declarators;

The declaratory is the identifier being declared. Unlike C or Java or other languages, EasKey does not have storage class specifier; it only consists of type-specifier and it only consists of at most one type-specifier.
8.1 The type-specifiers are:
type-specifier:
  int
  char
  float
  string
  boolean
  key
  point
  color

Examples:

```c
int a          // (Uninitialized Integer identifier will be set to 0 as default value)
char b         // (Uninitialized char identifier will be set to '\000' as default value)
string c       // (Uninitialized string identifier will be set to "" as default value)
float d        // (Uninitialized float identifier will be set to 0.0 as default value)
boolean e      // (Uninitialized boolean identifier will be set to false as default value)
color f        // (Uninitialized color identifier will be set to white, which represents a
                // RGB form #0, 0, 0#, as default value)
key g          // (Uninitialized key identifier will be set to $null$, as default value)
point h        // (Uninitialized point identifier will be set to tuple {0, 0}, as default value)
i int i=10      // (This will initialize the value of h to 0)
j key j=$A$     // (This will initialize the value of j to $A$)
```

8.2 Declarators
The specifier in the declaration indicates the type of the objects to which the declarator refer.
Declarators have the syntax:

```
declarator:
  identifier
  declarator ( )
```

8.3 Meaning of declarators
Each declarator is taken to be an assertion that when a construction of the same form as
the declarator appears in an expression, it yields an object of the indicated type. Each
declarator contains exactly one identifier; it is this identifier that is declared.
If an unadorned identifier appears as a declarator, then it has the type indicated by the
specifier heading the declaration.
For example,

```c
int a;
```
means that the identifier a, which is to be declared, has the type of int.
If a declarator has the form

\[ D( ) \]

then the contained identifier has the type “function returning \( \ldots \)”, where \( \ldots \) is the type
which the identifier would have had if the declarator had been simply \( D \).

9. Constant expressions
In one case it requires expressions which evaluate to a constant: after case. In this case,
the expression can involve only constants of any type except float, possibly connected by
the binary operators

\[ + \quad - \quad * \quad / \quad \% \]

or by the unary operators:

\[ ! \quad @x \quad @y \quad @r \quad @g \quad @b \quad @k \]

10. Built-in functions
The built-in functions are classified into three categories: (1) Mouse operations; (2) Key
operations; (3) Other functions.
10.1 Mouse operations
10.1.1 void __leftClickPoint(point p);
Click the left mouse button on a designated point.
10.1.2 void __rightClickPoint(point p);
Click the right mouse button on a designated point.
10.1.3 void __movePoint(point p);
Move the mouse from current point to a designated point.
10.1.4 double_click (point);
Double click the mouse on a designated point.
10.1.5 void __leftDown();
Press down the left mouse button.
10.1.6 void __rightDown();
Press down the right mouse button.
10.1.7 void __leftUp();
Let go the left mouse button.
10.1.8 void __rightUp();
Let go the right mouse button.
10.2 Keyboard operations
10.2.1 void __keyDown(key k);
Press down a designated key on the keyboard.
10.2.2 void __keyUp(key k);
Let go the key.
10.2.3 void __keyStroke (key k);
Stroke a designated key on the keyboard.
10.2.4 void __comboStroke(key k1, key k2, key k3, key k4, key k5, key k6, key k7);
This function supports at most 4 different keys as arguments at a time by the means of overload. It presses the first three keys by calling function \texttt{key\_down(key)}, then strokes the fourth key by calling function \texttt{key\_stroke(key)} to implement the whole process.

10.3 Other functions
10.3.1 \texttt{point \_getPoint();}
10.3.2 \texttt{color \_getcolor(point p);}  
Get the color of a designated point.
10.3.3 \texttt{void \_print(string s);}  
Print out designated content.

11. Statements
Except as indicated, statements will be executed in sequence within certain Blocks, and the execution of Statements in different Block will also follow the sequence of block appearance.

11.1. Expression Statement
Most statements are expression statements, which have the form
\begin{verbatim}
expression ;
\end{verbatim}
Usually expression statements are assignments or function calls.
Expression statement is a construct consisting of variables, operators, certain real value and even functions.
Examples:
\begin{verbatim}
int a=32  
int a=b+32+c;  
point p={2,4}+g;
\end{verbatim}

11.2. Conditional Statement
There are 4 forms of conditional statement:
\begin{verbatim}
if (expression) then (statement) end if  
if (expression) then (statement) else (statement) end if  
if (expression) then (statement) elseif(expression) then (statement) end if  
if (expression) then (statement) elseif(expression) then (statement) else (statement) end if
\end{verbatim}
Each keyword if must end with up with a keyword end if. The elseif and the else after the keyword if are optional. With these forms, the body of a conditional statement are covered in a pair of if and end if, with no exception. Besides, each keyword end if will be paired with the last encountered if.
If the outcome of first expression is Boolean true then the first statement will be executed, else the following statement with Boolean true expression outcome will be executed.
Example:
1. \texttt{if (a==b) then a=0 end if;}
2. if (a==b) then a=0 else a=1 end if;
3. if (a==b) then a=0 elseif (a<b) then a=1 end if;
4. if (a==b) then a=0 elseif (a<b) then a=0 else a=1 end if;
5. if (a==b) then
   if (a>c) then a=0
   end if
end if;

11.3. While Statement
The while statement has the form:

\[
\text{while ( expression )} \\
\text{statement} \\
\text{end while}
\]

The substatement will be executed repeatedly so long as the value of the expression remains true.
Example:

\[
\text{while (a>b)} \\
\quad c=a; \quad a=b; \quad b=c; \\
\text{end while}
\]

The statement means a single statement or a sequence of statements, which are placed in the block between while and end while, will be executed consecutively if the condition in expression is true.

11.4. For Statement
The for statement has the form

\[
\text{for ( expression1; expression2; expression3) } \\
\text{statement} \\
\text{end for}
\]

It is equivalent to

\[
\text{expression1;} \\
\text{while ( expression2) } \\
\text{statement} \\
\text{expression3;} \\
\text{end while}
\]

The first expression specifies initialization for the loop. The second expression specifies a conditional judgment, which is tested before each iteration. The loop exits when the second expression is false; the third expression specifies an incrementation or decrementation which is performed after each iteration.
Example:

\[
\text{for(int i=0; i<100; i = i+1)}
\]
\[ a = a + i; \]

\textit{end for}

11.5. Switch Statement

The switch statement causes control to be transferred to one of several statements depending on the value of an expression. It has the form:

\[
\text{switch (expression)} \\
\text{case constant\_statement: statement} \\
\ldots \\
\text{(default: statement)}_{\text{opt}} \\
\text{end switch}
\]

The expression outcome can be any type other than float. The statement is typically a list of statements and each statement within the statement may follow a case prefixes as follows:

\[
\text{case constant-expression:}
\]

where the constant expression can be any type other than float. No two of the case constants in a switch may have the same value.

There may also be at most one statement prefix of the form

\[
\text{default:}
\]

When the switch statement is executed, its expression is evaluated and compared with each case constant in an undefined order. If one of the case constants is equal to the value of the expression, control is passed to the statement following the matched case prefix. If no case constant matches the expression, and if there is a default prefix, control passes to the prefixed statement. In the absence of a default prefix none of the statements in the switch is executed.

11.6. Break Statement

\[
\text{break;}
\]

causes termination of the smallest enclosing while, for, or switch statement; in other words, it will jump out of the inner loop.

11.7. Return Statement

A function returns to its caller according to the return statement, which has the following form:

\[
\text{return expression;}
\]

Example:

\[
\text{return 0;} \cr \text{return a+b+c;} \cr \text{return a==b;}
\]

12. Scope
In EasKey, local variables can shadow the presence of global variables of the same name. In each block statement, such as a while loop, etc. a new local scope will be built. In that scope, all the variables outside this scope are visible except a local variable is declared with the same name. In the scope of a function declaration, the names of formals and locals mustn’t be the same, since they belong to the same scope.

13. Examples
13.1. Basic function implementation
In this function, a dividend calculation function is implemented. Two doubles are passed into the function as parameters. Then the function calculates the dividend of the two doubles and returns the dividend which is a double value. If the second parameter is zero, the function will return an error message and return -1.0.

```%
calculate i divided by j
function float dividend(float i, float j)
if j==0 then __print("error");return -1.0;
else
    return i/j;
end if;
end function
```

13.2. Basic array operation
In this example, an integer array is passed to the function to calculate the sum of the elements inside the array. This function takes two arguments, the first one is the array and the second one is the size of the array. At the end of the function, the sum of the array is returned.

```function sum(int s[], int num)
int i=0;
int sum=0;
for(;i<num;i=i+1)
    sum=sum+s[i];
end for
end function
```

13.3. Drag window to a different position
In this function, two point type parameters are passed to the function to implement the action that move a dialog window from point A to point B. This function left presses the mouse and keeps the mouse down to drag the window to another position.

```function void dragwin(point origin, point desti)
    __movePoint(origin);
    __leftDown();
    __movePoint(desti);
```
13.4. Enter a string into a textbox
In this function, a point textb and a string msg are passed to the function to implement the function that enters the string msg into the textbox pointed by the textb point. The function left click the region inside the textbox to focus on the textbox and then enter the string msg into the box by traversing each character in the string.

```plaintext
function void enter(point textb, string msg)
    %point textb is the point within the area of the textbox
    __leftClickPoint(point);
    while msg.i!=eof do  //msg.i is not implemented)
        __keyStroke(msg.i);  //__keyStroke is not implemented for string yet
        __delay(100);
    end while
end function
```

13.5. Color dependent mouse click
In this function, a color type variable and a point pointing to the button are passed to the function to implement the action that if the color at or is the same as the red variable then click the button, else do not press the button.

```plaintext
function void press(color red, point or, point but)
    if(_getColor(or)==red)
        then
            __leftClickPoint(but);
        end if
end function
```
Chapter 4
Project Plan

4.1 Project Management process

Our team is formed with 5 members, and each member has their own specialty. So, we just try to make the best use of our intelligence through uneven task distribution. After we confirmed the project title, each member start to gather some useful reference and try to get used to OCaml language as soon as possible. Also, we have several conversions with TA Jiabin Hu, who has given us a lot of advices.

We held our group meeting once per week, discussing the knowledge that we learned in the class and trying to convert this knowledge into our language. During Oct and Dec, we have several long meetings and finally confirmed the LRM, which specified what we should do in our Language.

Whenever we came up with a new problem either about OCaml or our own language, we will reserve a room in butler the next day, and tried to figure it out as soon as possible.

We started to work on the test when we finally merged the source code and successfully produced our first bytecode file. Every team member tries to come up with his or her own test to find the bug in our compiler.

Finally, we really don't have time to further discuss in the final week, and we just work on the demo and presentation after 12/15/2012.

4.2 Programming Style guide

✓ For every code file, the author name must be notated, as well as the modified time and the version, where may potentially has bug.
✓ The OCaml code should try to follow the style of microC and the good project done by former team, and the C code should be neat and clean, making it easy to read.
✓ Almost no IDE is used, we just try to ensure that we work on the same path and on the same environment, and that’s why the plain text editor was chosen.
✓ Commit the code on Dropbox whenever it is modified and without bug, but have to make sure that the version and data are notated.
✓ Each team member has their own file directory in Dropbox, and a public directory is used to store the bug-free code.
✓ Two main directory were built in the public directory, one is EasKey_Compiler, and the other one is EasKey_Interpreter, each stored the files as their name imply.
4.3 Project Timeline

<table>
<thead>
<tr>
<th>Started Data</th>
<th>Event</th>
<th>Finished Data</th>
</tr>
</thead>
<tbody>
<tr>
<td>09/10/2012</td>
<td>Team Formation, some rough ideas about project</td>
<td>09/10/2012</td>
</tr>
<tr>
<td>09/12/2012</td>
<td>OCaml language learning, project title discussion</td>
<td>09/24/2012</td>
</tr>
<tr>
<td>09/23/2012</td>
<td>Project title confirmed and talked with TA</td>
<td>10/02/2012</td>
</tr>
<tr>
<td>09/26/2012</td>
<td>Learning the structure of compiler, and design our own compiler, task distributed</td>
<td>10/27/2012</td>
</tr>
<tr>
<td>10/03/2012</td>
<td>Holding several meetings, finally confirmed the grammar and submitted LRM</td>
<td>10/29/2012</td>
</tr>
<tr>
<td>11/05/2012</td>
<td>Interpreter is complete</td>
<td>11/16/2012</td>
</tr>
<tr>
<td>11/05/2012</td>
<td>Scanner is complete</td>
<td>11/26/2012</td>
</tr>
<tr>
<td>11/05/2012</td>
<td>Ast and parser are complete</td>
<td>12/01/2012</td>
</tr>
<tr>
<td>12/06/2012</td>
<td>Sast is complete</td>
<td>12/11/2012</td>
</tr>
<tr>
<td>12/06/2012</td>
<td>Analyzer and compile are complete</td>
<td>12/16/2012</td>
</tr>
<tr>
<td>12/16/2012</td>
<td>Test cases applied and demo program is complete</td>
<td>12/18/2012</td>
</tr>
<tr>
<td>12/16/2012</td>
<td>Final report is complete, and prepare for presentation</td>
<td>12/18/2012</td>
</tr>
<tr>
<td>12/19/2012</td>
<td>Presentation</td>
<td>12/19/2012</td>
</tr>
</tbody>
</table>

4.4 Roles and Responsibilities

4.4.1 Keqiu Hu
Leader of this team, and he is responsible for distributing the work and worked on every part of the project. Besides, he finished the back-end of the project independently.

4.4.2 Jinqi Huang
Mainly worked on the front-end of the project, including scanner and ast. Besides, participated in the writing of test cases.

4.4.3 Xiaoyu Huang
Worked on the record of our team Log, participated in the implementation of analyzer and sast. Also, she wrote most of the test cases.

4.4.4 Zongheng Wang
He has contributed a lot to this project, participating in the whole design of front-end, including scanner, ast, parser, sast and analyzer.
4.4.5 Lizhong Zhang
Most of his work is focused on the debug and designing byte code, as well as the implementation of the compile. Every time when our program collapse, he will check the bytecode line by line, and tell us where goes wrong.

4.5 Software Development Environment

4.5.1 Operating System
Actually, we had to use both Windows and Linux operating system, because we need to use the winAPI, which can only be accessed via windows platform. However, most of our codes are written in Linux environment except the interpreter part, which is written in windows with visual studio.

4.5.2 Language Used
We finish the front-end, including scanner, parser, ast, sast, analyzer and compile, with OCaml language, and the back-end, including interpreter and mainly API function, is built with C language.

4.5.3 Software Used

4.5.3.1 Dropbox
We used Dropbox as working directory, and it turns out to be a bad idea. Because it is really hard to control the version.

4.5.3.2 Gedit
Most of our group members use gedit in either Linux or Windows

4.5.3.3 Visual Studio
We finished the interpreter part with visual studio in windows.

4.6 Project Log

<table>
<thead>
<tr>
<th>Data</th>
<th>Events</th>
</tr>
</thead>
<tbody>
<tr>
<td>09/10/2012</td>
<td>Team formed, 5 members get to know with each other</td>
</tr>
<tr>
<td>09/12/2012</td>
<td>Discussed about the project title and exchange opinion</td>
</tr>
<tr>
<td>09/20/2012</td>
<td>Project title confirmed, begin to evaluate the difficulty of title we choose</td>
</tr>
<tr>
<td>09/23/2012</td>
<td>Begin writing the Project Proposal, and ask TA Jiabin Hu for advice</td>
</tr>
<tr>
<td>09/26/2012</td>
<td>Submit Project Proposal, installed OCaml environment for each member</td>
</tr>
<tr>
<td>Date</td>
<td>Event Description</td>
</tr>
<tr>
<td>-----------</td>
<td>-----------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>10/02/2012</td>
<td>Feedback from TA and discussed about the feasibility of our project</td>
</tr>
<tr>
<td>10/03/2012</td>
<td>1st meeting after project submitted, discussed about OCaml</td>
</tr>
<tr>
<td>10/15/2012</td>
<td>2nd meeting, discussed about the parser and scanner</td>
</tr>
<tr>
<td>10/20/2012</td>
<td>3rd meeting, discussed the basic structure of our compiler</td>
</tr>
<tr>
<td>10/24/2012</td>
<td>4th meeting, started to work on the LRM, confirmed the grammar and data types</td>
</tr>
<tr>
<td>10/29/2012</td>
<td>5th meeting, discussed and confirmed the LRM, then submitted</td>
</tr>
<tr>
<td>11/05/2012</td>
<td>6th meeting, started to work on project, and set up the coding style</td>
</tr>
<tr>
<td>11/16/2012</td>
<td>7th meeting, a team member has finished the interpreter and C API parts</td>
</tr>
<tr>
<td>11/26/2012</td>
<td>8th meeting, discussed about AST, the scanner is finished</td>
</tr>
<tr>
<td>12/01/2012</td>
<td>9th meeting, discussed about compile, Sast and analyzer, all the other parts are finished</td>
</tr>
<tr>
<td>12/06/2012</td>
<td>10th meeting, still haven’t begin the work of compile, will work on it after the final</td>
</tr>
<tr>
<td>12/11/2012</td>
<td>11th meeting, group coding and finished Sast</td>
</tr>
<tr>
<td>12/16/2012</td>
<td>12th meeting, finished analyzer but still some issue with compile</td>
</tr>
<tr>
<td>12/17/2012</td>
<td>13th meeting, finished compile and started to debug and test cases</td>
</tr>
<tr>
<td>12/18/2012</td>
<td>14th meeting, finished test cases and final report, prepared for the representation</td>
</tr>
<tr>
<td>12/19/2012</td>
<td>Representation</td>
</tr>
</tbody>
</table>
Chapter 5
Architecture Design

5.1 Block Diagram

5.2 Interface between Components
As we can see from the diagram above, the source file of EasKey language which has suffix (.ek), will be first processed by the scanner and produce a sequence of tokens. Parser will recognize these tokens and fill them into the AST tree, ensure the correctness of syntax. Then, analyzer will translate the AST to the SAST, and check any potential inconsistence that may result in the failure of compiler. If the semantic is correct, the compile will translate the source file into bytecode file with suffix (.e), which uses almost the same definition as microC. Finally, the interpreter will run the bytecode and call winAPI to perform Mouse or Keyboard operations.

5.3 Roles and Responsibility

<table>
<thead>
<tr>
<th>Components</th>
<th>Implementer(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Scanner.mll</td>
<td>Jinqi Huang, Zongheng Wang</td>
</tr>
<tr>
<td>Parser.mly</td>
<td>Zongheng Wang, Lizhong Zhang</td>
</tr>
<tr>
<td>Ast.mli</td>
<td>Zongheng Wang</td>
</tr>
<tr>
<td>Analyzer.ml</td>
<td>Zongheng Wang, Xiaoyu Huang</td>
</tr>
<tr>
<td>Sast.mli</td>
<td>Xiaoyu Huang</td>
</tr>
<tr>
<td>Bytecode.ml</td>
<td>Jinqi Huang</td>
</tr>
<tr>
<td>Compile.ml</td>
<td>Keqiu Hu, Lizhong Zhang</td>
</tr>
<tr>
<td>Interpreter</td>
<td>Keqiu Hu</td>
</tr>
<tr>
<td>Testcases</td>
<td>Xiaoyu Huang, Jinqi Huang</td>
</tr>
</tbody>
</table>

We can roughly divide our compiler into 6 components, and distribute the tasks evenly to each team member. Keqiu Hu is good at C language, so he is responsible for the building of interpreter and all the winAPI functions. Scanner, bytecode and some test cases are written by Jinqi Huang. Parser, AST, Analyzer, SAST are finished by Zongheng Wang and Xiaoyu Huang. Compile is implemented by Lizhong Zhang.
Chapter 6

Test Plan

6.1 Representative Source Code

The input of EasyKey is the source code with suffix (.ek). The output of EasyKey compiler is bytecode file, which will be further translated by the interpreter written in C, and then a serial of mouse and keyboard operation will be executed.

As we have discussed about, we are not able to present the output of program here. Besides, because these operations have to call the winAPI function, which can only be invoked in windows environment, we may need to transfer the output bytecode file from Linux to windows.

Hence, we will show the bytecode we generate here. Since we use the same bytecode format as microC, we can verify whether our output is right or wrong through the comparison of bytecode between our output and microC’s output.

6.2 Test Suits

Part of the test cases we chose are derived from microC test suite, and the rest are designed to test our special build in function. These tests is a mixture of both simple and complex program, which consists of arithmetic, logic, control structures, different data types, and our built-in functions.

6.2.1 Arithmetic

<table>
<thead>
<tr>
<th>Test case</th>
<th>Operation</th>
<th>Output</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>int a=(1+2);</td>
<td>Integer addition</td>
<td>3</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(__stringOfInt(a));</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>int c=1-2;</td>
<td>Integer subtraction</td>
<td>-1</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(__stringOfInt(c));</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>int b=1*2;</td>
<td>Integer multiplication</td>
<td>2</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(__stringOfInt(b));</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>int d=4/2;</td>
<td>Integer division</td>
<td>2</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(__stringOfInt(c));</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>float f=5%2;</td>
<td>Integer Modulus</td>
<td>2.5</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(__stringOfFloat(f));</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>float g=1+3*3-7;</td>
<td>Order of operations</td>
<td>3</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(__stringOfFloat(g));</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### 6.2.2 Logic

<table>
<thead>
<tr>
<th>Test case</th>
<th>Operation</th>
<th>Output</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>if(true</td>
<td></td>
<td>false)</td>
<td>Test</td>
</tr>
<tr>
<td>__print(&quot;true&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>end if</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>if(true&amp;&amp;false)</td>
<td>Test &amp;&amp;</td>
<td>false</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(&quot;true&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>else</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>__print(&quot;false&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>end if</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>if(true</td>
<td></td>
<td>(false</td>
<td></td>
</tr>
<tr>
<td>__print(&quot;true&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>end if</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>if((true&amp;&amp;false)</td>
<td></td>
<td>true)</td>
<td>Test</td>
</tr>
<tr>
<td>__print(&quot;true&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>else</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>__print(&quot;false&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>end if</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

### 6.2.3 if

<table>
<thead>
<tr>
<th>Test case</th>
<th>Motivation</th>
<th>Output</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>int i=20; if (i==20)</td>
<td>Test if and if/elseif/else</td>
<td>if</td>
<td>Correct</td>
</tr>
<tr>
<td>__print(&quot;if&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>elseif (i&lt;20)</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>__print(&quot;elseif&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>else</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>__print(&quot;else&quot;);</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>end if</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
### 6.2.3 for

<table>
<thead>
<tr>
<th>Test case</th>
<th>Motivation</th>
<th>Output</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>int i; for (i=0;i&lt;10;i=i+1) __print(__stringOfInt (i)); end for</td>
<td>Test for</td>
<td>0 1 2 3 4 5 6 7 8 9</td>
<td>Correct</td>
</tr>
</tbody>
</table>

### 6.2.4 Derived Types

<table>
<thead>
<tr>
<th>Test case</th>
<th>Operation</th>
<th>Output</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>point p={200,400}; int a=p@x; __print(__stringOfInt(a));</td>
<td>Test type point @x operator</td>
<td>200</td>
<td>Correct</td>
</tr>
<tr>
<td>color c=#1,2,4#; int a=c@r; __print(__stringOfInt(a));</td>
<td>Test type color @r operator</td>
<td>1</td>
<td>Correct</td>
</tr>
<tr>
<td>key k1=$R$; string a=k1@k; __print(a);</td>
<td>Test type key @k operator</td>
<td>R</td>
<td>Correct</td>
</tr>
</tbody>
</table>
## 6.2.5 EasyKey Built in function

<table>
<thead>
<tr>
<th>Test case</th>
<th>Operation</th>
<th>Output</th>
<th>Result</th>
</tr>
</thead>
<tbody>
<tr>
<td>color c=#0,54,99#; point p1={600,320}; color c1; while(true) c1=__getcolor(p1); if(c1!=c) __movePoint(p1); __leftClickPoint(p1); c1=c; end if end while</td>
<td>This is used to test built in function __movePoint(point p), __leftClickPoint(point p)</td>
<td>Mouse will move to the given point and click the point</td>
<td>Correct</td>
</tr>
</tbody>
</table>

### 6.3 Testing Method

We use black box testing; write box testing and unit testing.
Chapter 7

Lessons Learned

7.1.1 Keqiu Hu

OCaml is such a language that you can never discover its elegance that until you deeply get into it. Around 1500 lines C++ are needed to implement our interpreter (because we have to call Windows API) while only 50 lines are needed for a similar interpreter in OCaml to work. For a team to finish a project peacefully and gracefully, I think at least one lady is needed to be a part of a team so that the team environment will be much softer and make in-team discussion much more rational.

For lessons, a team must start the project earlier and divide work more wisely; so that every team member can work at the same time in parallel, instead of dividing work into a sequence of jobs that one can only work after another member finish his job.

This project does impress me with the beauty of OCaml and gain a much much deeper understanding of Programming Languages and Translators.

7.1.2 Jinqi Huang

Compiler course is always torture, but you can surely learn a lot of stuff. Scanner, parser, ast, analyzer and compile will teach you a lot not only about OCaml, but also about how a language is built and what we should think about when we program.

Also, advice from me is the same as most of former students. Start early! Or you will have to stay up late every night in the final week.

Trust me, OCaml is not worse than java, just make sure you know how to write a decent, neat and powerful OCaml program.

7.1.3 Xiaoyu Huang

The most important thing I learned is to start early. You never know what you will meet next time. Even if there are only a few hours before the deadline, there may still have a lot of problems to fix. Therefore, start early, you may fix the problems early.

And designing period is a very hard and difficult time. However, a good design will help to improve the whole project quality. Sometimes, it will save time. Since, a bad design may lead to reconfiguration the whole project. And I learned a lot about the how the scanner, parser, analyzer, compiler and interpreter work.

Last, but not the least, I learned a lot from my teammates. They are warm-hearted and passion. Zonghen Wang is a careful consideration person. Keqiu Hu is a good
Lizhong Zhang always fixes the problems. And Jinqi Huang helps to keep all our project done successfully.

7.1.4 Zongheng Wang

In this project, it is better to have a focus when implementing your language, rather than try to include every function, every type and every good characteristic that you have seen in other languages. The reason behind it is obvious. First of all, the time limitation is always a great constraint. Although everyone knows it is very necessary to start this project early, but the truth is that when you have time, you know very little about this project and this course; when you know what to do, you will find there is not much time left. The second reason is that the more you try to implement a lot of general functions, perhaps the less the innovation of your programming language is. In addition, a big project will lead to a large number of debugging processes for sure. At the end, you probably mess up something and leave the professor with a bad impression, while in fact, your group might have done a lot of hard work.

7.1.5 Lizhong Zhang

The most important thing is to start doing something. Ideas keep coming out when having a group discussion. However, it is not possible, or not possible in a limited amount of time, to implement every feature generated from discussion. Some ideas which take infinite time to discuss may not be practical at all. So try to keep everything in a narrow scope at first and start to implement a small demo program. During the process, you'll find that greatest ideas come from practice.

Learn to compromise with teammates. It is even more important than the project to learn to work with others. Everyone has his own working and learning style. Most of the time there is no best ideas; it is just trade-off. So when you find yourselves in a argument, insist on yourself may sometimes not be very efficient.

The last but not the least, start early. It is not only because the whole project is huge, but also it takes a lot of time to debug when complete the first version of the project. You may be surprised to find there are so many bugs in you code. But don't worry, it will get better.
7.2 Advice for future teams

✓ Start early! Or you will stay up late every night in the final week.

✓ Try to come up with a good proposal, and don’t bother talk to professor and TA about the feasibility of your project.

✓ LRM should be written carefully, it will make your work easier.

✓ Don’t use Dropbox as version control system, it makes things worse.

✓ Holding meeting every week and telling your teammate if you run into trouble

✓ Compile and analyzer should be designed carefully and start early, it will take most of your time.

✓ Don’t worship microC as the god, try to write your own style.
Reference

Appendix

./Makefile: easkey: scanner.cmo parser.cmo analyzer.cmo compile.cmo bytecode.cmo easkey.ml
./Makefile: ocamlc -o easkey scanner.cmo parser.cmo analyzer.cmo bytecode.cmo compile.cmo easkey.ml
./Makefile:  cp easkey ..
./Makefile: scanner_test: scanner.cmo parser.cmo scanner_test.ml
./Makefile:  ocamlc -o scanner_test scanner.cmo parser.cmo scanner_test.ml
./Makefile: analyzer_test: scanner.cmo parser.cmo analyzer.cmo analyzer_test.ml
./Makefile:  ocamlc -o analyzer_test scanner.cmo parser.cmo analyzer.cmo analyzer_test.ml
./Makefile: scanner.ml: scanner.mll
./Makefile:  ocamlex scanner.mll
./Makefile: parser.ml parser.mli: parser.mly ast.cmi
./Makefile:  ocamlyacc parser.mly
./Makefile: ast.cmi: type.cmi
./Makefile: sast.cmi: type.cmi
./Makefile: bytecode.cmo: type.cmi
./Makefile: scanner.cmo: parser.cmi
./Makefile: parser.cmo: parser.cmi ast.cmi
./Makefile: compile.cmo: sast.cmi bytecode.cmo
./Makefile: analyzer.cmo: ast.cmi sast.cmi
./Makefile: generator.cmo: sast.cmi bytecode.cmi
./Makefile: %.cmo : %.ml
./Makefile:  ocamle -c $<
./Makefile: %.cmi : %.mli
./Makefile:  ocamle -c $<
./Makefile: PHONY: clean
./Makefile: clean:
./Makefile:  rm -f *.cmi *.cmo scanner.ml parser.ml parser.ml easkey scanner_test analyzer_test
./Makefile:  -rm ../easkey
./Makefile: Generated by ocamldep
./Makefile:#compiler.cmo: scanner.cmo parser.cmi
./Makefile:#compiler.cm: scanner.cm parser.cm
./Makefile:#parser.cmo: parser.cmi
./Makefile:
./Makefile:#scanner.cm: parser.cm
./Makefile:#scanner_test.cmo: scanner.cmo parser.cmi
./Makefile:#scanner_test.cm: scanner.cm parser.cm
./Makefile:#parser.cmi:
./analyzer.ml:(* 11:37 Dec 16, 2012 hxy, wzh *)
./analyzer.ml:
./analyzer.ml:open Type
./analyzer.ml:open Sast
./analyzer.ml:
./analyzer.ml:type symbol_table = {
./analyzer.ml:    parent : symbol_table option;
./analyzer.ml:    variables : (t * string) list; (* variable type * variable name *)
./analyzer.ml:    in_loop : bool; (* true if in a loop *)
./analyzer.ml:    in_switch : bool (* true if in a switch *)
./analyzer.ml:}
./analyzer.ml:
./analyzer.ml:type function_table = {
./analyzer.ml:    functions : (t * string * t list) list; (* reture type * function name * formals *)
./analyzer.ml:}
./analyzer.ml:
./analyzer.ml:let string_of_type = function
./analyzer.ml:  Void_T -> "void"
./analyzer.ml:| Int_T -> "int"
./analyzer.ml:| Char_T -> "char"
./analyzer.ml:| Float_T -> "float"
./analyzer.ml:| String_T -> "string"
./analyzer.ml:| Boolean_T -> "boolean"
./analyzer.ml:| Point_T -> "point"
./analyzer.ml:| Color_T -> "color"
./analyzer.ml:| Key_T -> "key"
./analyzer.ml:
./analyzer.ml:let string_of_unop = function
./analyzer.ml:  Neg -> "-"  
./analyzer.ml:| Not -> "!"
./analyzer.ml:| Atx -> "@x"
./analyzer.ml:| Aty -> "@y"
let string_of_binop = function
  | Add -> "\+"
  | Sub -> "\-
  | Mul -> "\*"
  | Div -> "\/
  | Mod -> "\%
  | Gt -> "\>
  | Lt -> "\<"
  | Geq -> "\>="
  | Leq -> "\<="
  | Eq -> "\=="
  | Neq -> "\!="
  | And -> "\&&"
  | Or -> "\||"

let oprand_check_unop t unop = (* unop : Type.unop -> unit = <fun>, check operand type of unary operator *)
  match unop with
    | Neg -> (match t with Int_T | Float_T -> true | _ -> false)
    | Not -> (t = Boolean_T)
    | Atx | Aty -> (t = Point_T)
    | Atr | Atg | Atb -> (t = Color_T)
    | Atk -> (t = Key_T) in
    if not correct then
      raise (Failure("The oprand of " ^ string_of_unop unop ^ " can not be " ^ string_of_type t ^ ""))

let oprand_check_binop t leftside binop = (* binop : Type.binop -> unit = <fun>,
  check operand type of binary operator *)
  match binop with
    (* '+' for string is to concatenate two strings *)
    | Add -> (match t with Int_T | Char_T | Float_T | Point_T | Color_T | String_T
      -> true | _ -> false)
    | Sub -> (match t with Int_T | Char_T | Float_T | Point_T | Color_T -> true | _ ->
      false)
```plaintext
// analyzer.ml: Mul | Div -> (match t with Int_T | Float_T -> true | _ -> false)
// analyzer.ml: Mod -> (t = Int_T)
// analyzer.ml: Gt | Lt | Geq | Leq -> (match t with Int_T | Char_T | Float_T | String_T -> true | _ -> false)
// analyzer.ml: Eq | Neq -> (match t with Int_T | Char_T | Float_T | String_T -> true | _ -> false)
// analyzer.ml: And | Or -> (t = Boolean_T) in
// analyzer.ml: if not correct then
// analyzer.ml: if leftside = true then raise (Failure("The left operand of " ^ string_of_binop binop ^ " can not be " ^ string_of_type t ^ "")
// analyzer.ml: else raise (Failure("The right operand of " ^ string_of_binop binop ^ " can not be " ^ string_of_type t ^ "))

let rec find_variable scope name =
  try
  List.find (fun (_, s) -> s = name) scope.variables
  with Not_found ->
  match scope.parent with
  Some(parent) -> find_variable parent name
  | _ -> raise Not_found

let rec expr scope funcs = function
  (* AST.expr -> SAST.expr *)
  Ast.Id(id)
  let (var_type, _) = find_variable scope id in
  (Sast.Id(id), var_type, false)
  with Not_found ->
  raise (Failure("" ^ id ^ "' was not declared in this scope"))
  Ast.Literal(lit)
  let lit_type =
  match lit with
  Integer(i) -> Int_T
  Char(c) -> Char_T
  Float(f) -> Float_T
  String(str) -> String_T
  Boolean(b) -> Boolean_T
  Point(x,y) -> Point_T
  Color(r,g,b) -> Color_T
  Key(key) -> Key_T
```
in (Sast.Literal(lit), lit_type, true)

let new_e1 = expr scope funcs (Ast.Id(vname)) and new_e2 = expr scope funcs e2 in
let (_, t1, _) = new_e1 and (_, t2, _) = new_e2 in
if t1 = t2 then
  (Sast.Assign(vname, new_e2), t1, false)
else
  raise(Failure("type mismatch of " ^ string_of_type t1 ^ " and " ^ string_of_type t2 ^ " in the assignment of " ^ vname ^ ")

let new_e1 = expr scope funcs e1 and new_e2 = expr scope funcs e2 in
let (_, t1, c1) = new_e1 and (_, t2, c2) = new_e2 in
let _ = oprand_check_binop t1 true binop and _ = oprand_check_binop t2 false binop in (* check left & right oprands *)
if t1 = t2 then
  (Sast.Binop(new_e1, binop, new_e2), (fun t ->
    match binop with
    Gt | Lt | Geq | Leq | Eq | Neq -> Boolean_T
    | _ -> t) t1, c1 && c2) (* constant expr if both oprands are constant expr *)
else
  raise (Failure("type mismatch of " ^ string_of_type t1 ^ " and " ^ string_of_type t2 ^ " for the binary operator " ^ string_of_binop binop))

let new_e = expr scope funcs e in
let (_, t, c) = new_e in
let _ = oprand_check_unop t unop in
(Sast.Unop(unop, new_e), (fun t ->
  match unop with
  Neg | Not -> t
  | Atx | Aty | Atr | Atg | Atb -> Int_T
  | Atk -> String_T) t, c)

let (return_type, _, formals_types) = try
  List.find (fun (_, s, _) -> s = fname) funcs.functions
with Not_found ->
  raise(Failure("" ^ fname ^ " was not declared in this scope") in
let new_actuals = List.map (expr scope funcs) actuals in
let actuals_types = List.map (fun (_, t, _) -> t) new_actuals in
if formals_types = actuals_types then
    (Sast.Call(fname, new_actuals), return_type, false)
else
    raise(Failure("type mismatch between formals and actuals of function \"\" ^ fname ^ \"\")
| Ast.NoExpr -> (Sast.NoExpr, Void_T, false)
let init_transform scope funcs = function
    (* Ast.expr -> Sast.expr for initialization of variable declaration *)
    Some(e) -> Some(expr scope funcs e)
    None -> None
let rec count_var_break scope num =
    if scope.in_loop || scope.in_switch then
        num
    else
        match scope.parent with
            Some(parent) -> count_var_break parent (num + List.length scope.variables)
            | None -> raise(Failure("break statement not within loop or switch"))
let rec count_var_continue scope num =
    if scope.in_loop then
        num
    else
        match scope.parent with
            Some(parent) -> count_var_continue parent (num + List.length scope.variables)
            | None -> raise(Failure("continue statement not within a loop"))
let var_decl_transform decl =
    (* Ast.var_decl -> Sast.var_decl, needs init_transform to transform init *)
    let t = fst decl in
    try
        List.rev (List.fold_left (fun l -> function
            (name, Ast.NoExpr) -> (t, name, None):l
            | (name, e) -> (t, name, Some(e)):l) [] (snd decl))
    in raise(Failure("redeclaration of \"\" ^ string_of_type t ^ \"\" ^ name ^ \"\")

let new_init = init_transform scope funcs e in
let init_t =  
match new_init with
  | Some(a, b, c) -> b
| None -> t in  (* No need to check types, thus t = init_t is trivial *)
let new_var_decl =  
if t = init_t then
  (t, name, new_init)
else
  raise(Failure ("type mismatch of\"^string_of_type t
^"^"^string_of_type init_t\" in the assignment of\"^name^"))
in (new_scope, funcs, new_var_decl::new_var_list)

let get_inherit_list = function
  | [] -> []
  | l -> List.map (fun (t, s, _) -> (t, s)) l

let rec stmt scope funcs inherit_list = function
  | Ast.Expr(e) -> Sast.Expr(expr scope funcs e)
  | Ast.Block(block) ->
    let block_scope = { parent = Some(scope) ; variables = get_inherit_list inherit_list ;
in_loop = false; in_switch = false } in
    let check_stmt (scope', st1) s =  
     match s with
      | Ast.VarDecl(var_decl) -> (* variable declarations can update the scope *)
        let var_list = var_decl_transform var_decl in
        let (new_scope, _, new_var_list) = List.fold_left add_var (scope',
            funcs, []) var_list
        in (new_scope,
            (List.map (fun x -> Sast.VarDecl(x)) new_var_list) @ st1)
      | _ -> (scope', (stmt scope' funcs [])::st1) in
        let (new_block_scope, st1) = List.fold_left check_stmt (block_scope, [])
        in Sast.Block(List.rev st1, List.rev new_block_scope.variables)
  | Ast.Return(e) -> Sast.Return(expr scope funcs e)
let exprList1 = (List.rev (List.fold_left (fun l e ->
    let (_, t, _) = expr1 in
    if t = Boolean_T then
        expr1::l
    else raise(Failure("predicates of if statements must be type \"boolean\"")))
[l exprlist)) in
let stmtlist1 = List.rev (List.fold_left (fun l s ->
    let st1 = stmt scope funcs [] s in
    let stmt1 =
        match st1 with
        Ast.Block(_, _) -> st1
        | _ -> raise(Failure("unexpected error"))
        in stmt1::l) [] stmtlist)
in Sast.If(exprList1, stmtlist1)
let stmtlist1 = List.rev (List.fold_left (fun l s ->
    let st1 = stmt scope funcs [] s in
    let stmt1 =
        match st1 with
        Ast.Block(_, _) -> st1
        | _ -> raise(Failure("unexpected error"))
        in stmt1::l) [] stmtlist)
in Sast.If(exprList1, stmtlist1)
let stmtlist1 = List.rev (List.fold_left (fun l s ->
    let st1 = stmt scope funcs [] s in
    let stmt1 =
        match st1 with
        Ast.Block(_, _) -> st1
        | _ -> raise(Failure("unexpected error"))
        in stmt1::l) [] stmtlist)
in Sast.If(exprList1, stmtlist1)
let block_for =
match stmt1 with
  Ast. Block(_) -> stmt1
| _ -> raise(Failure("unexpected error")) in
let e2 = expr new_scope funcs expr2 in
let new_stmt = stmt for_scope funcs new_var_list block_for in
Sast.ForWithDecl(List.rev new_var_list, e1, e2, new_stmt)
else
raise(Failure("the second expression in for statements must be type 'boolean'"))

let block_for =
match stmt1 with
  Ast. Block(_) -> stmt1
| _ -> raise(Failure("unexpected error")) in
let e2 = expr new_scope funcs expr2 in
let new_stmt = stmt for_scope funcs new_var_list block_for in
Sast.ForWithDecl(List.rev new_var_list, e1, e2, new_stmt)
else
raise(Failure("the second expression in for statements must be type 'boolean'"))

let exprList1 = (List.rev (List.fold_left (fun l e ->
  let e1 = expr scope(funcs e in
  let (_, t1, c1) = e1 in
  if c1 = false then
raise(Failure("case label must be a constant-expression"))
  else if t1 = Float_T then
顺应("case label must have the same type as switch quantity")
else
raise(Failure("case label must have the same type as switch quantity"))
}) exprlist)) in
let stmtlist1 = List.map (fun case_stmt 

let case_block = match case_stmt with
  Ast.Block(_) -> case_stmt
  | _ -> raise(Failure("unexpected error"))
in stmt new_scope funcs [] case_block
  in stmtlist
else raise(Failure("switch quantity can not be a float"))

let rec check_return fdecl s_list = function
  Sast.Expr(e) -> false
  Sast.Block(s_list, _) -> List.fold_left (fun a b -> a || b) false (List.map (check_return fdecl) s_list)
  Sast.Return(e) -> let (_, t, _) = e in
    if t = fdecl.Ast.rtype then true else raise(Failure("return expression has the type " ^
      string_of_type t ^
      but a return expression was expected of type " ^
      string_of_type fdecl.Ast.rtype ^
      in function " ^
      fdecl.Ast.fname ^
      "(" ^
      (List.map string_of_type (List.map fst l))) fdecl.Ast.formals)
    ^ ")")
  Sast.If(e_list, s_list) -> List.fold_left (fun a b -> a || b) false (List.map (check_return fdecl) s_list)
  Sast.For(e1, e2, e3, s) -> check_return fdecl s
  Sast.ForWithDecl(vdecl_list, e1, e2, s) -> check_return fdecl s
  Sast.While(e, s) -> check_return fdecl s
  Sast.Switch(e, e_list, s_list) -> List.fold_left (fun a b -> a || b) false (List.map (check_return fdecl) s_list)
  Sast.Continue(i) -> false
  Sast.Break(i) -> false
  Sast.VarDecl(vdecl) -> false

let semantically_check program = (* AST.program -> SAST.program *)
let init_funcs = { functions = [
  (Void_T, "__leftClickPoint", [Point_T]);
  (Void_T, "__rightClickPoint", [Point_T]);
]
(Void_T, "__doubleClick", [Point_T]);
(Void_T, "__movePoint", [Point_T]);
(Void_T, "__leftDown", []);
(Void_T, "__rightDown", []);
(Void_T, "__leftUp", []);
(Void_T, "__rightUp", []);
(Void_T, "__keyDown", [Key_T]);
(Void_T, "__keyUp", [Key_T]);
(Void_T, "__keyStroke", [Key_T]);
(Void_T, "__comboStroke", [Key_T; Key_T; Key_T; Key_T; Key_T; Key_T; Key_T; Key_T]);
(Point_T, "__getPoint", []);
(Color_T, "__getcolor", [Point_T]);
(Void_T, "__print", [String_T]);
(String_T, "__stringOfInt", [Int_T]);
(String_T, "__stringOfFloat", [Float_T]);
(String_T, "__stringOfChar", [Char_T]);
(String_T, "__stringOfBoolean", [Boolean_T]);
(String_T, "__stringOfPoint", [Point_T]);
(String_T, "__stringOfKey", [Key_T]);
(String_T, "__stringOfColor", [Color_T]);
(Int_T, "__getASCI", [Char_T]);
(Char_T, "__getChar", [Int_T]);
(Int_T, "__getStringLength", [String_T]);
(Char_T, "__getCharOfString", [String_T; Int_T]);
(String_T, "__setCharOfString", [String_T; Int_T; Char_T]);
(String_T, "__subOfString", [Int_T; Int_T]);
(Void_T, "__delay", [Int_T]);
let init_scope = {
  parent = None;
  variables = [];
  in_loop = false;
  in_switch = false } in
let (new_var_decl, new_func_decl) =
  let rec check var_decl func_decl scope funcs = function
  | [] -> (var_decl, List.rev func_decl)
  | Ast.VarDecl(vdecl)::tl ->
    let new_vdecl = var_decl_transform vdecl in
    let sast_vdecl = List.map (fun (t, name, init) ->
      try
    in

let _ = List.find (fun (_, s) -> s = name) scope.variables

in raise(Failure("redeclaration of \"" ^ string_of_type t ^ " \\
^ name ^ "}))

with Not_found ->

let new_init = init_transform scope funcs init in

let init_t =

match new_init with

Some(a, b, c) -> b

| None -> t in (* No need to check types, thus t = init_t is trivial *)

if t = init_t then

(t, name, new_init)

else raise(Failure("type mismatch of \"" ^ string_of_type t

^ \\
^ string_of_type init_t ^ "} in the

assignment of \"" ^ name ^ "")

) new_vdecl in

let updated_scope = { scope with

variables = scope.variables @ List.map (fun (t, name, _) -> (t, name)) new_vdecl

}

in check (var_decl @ sast_vdecl) func_decl updated_scope funcs tl

try

let _ = List.find (fun (_, s, _) -> s = fdecl.Ast.fname) funcs.functions

in raise(Failure("redeclaration of function \"" ^ string_of_type

fdecl.Ast.fname ^ "(" ^ ((fun l -> String.concat ", "

(List.map string_of_type

(List.map fst l)) fdecl.Ast.formals) ^ "}")")

) fdecl.Astype ^ " \\

fdecl.Ast.fname ^ "(" ^ ((fun l -> String.concat ", 

(List.map string_of_type

(List.map fst l)) fdecl.Ast.formals) ^ "}")

) fdecl.Astype ^ " \\

with Not_found ->

let rec formals_name_test = function

[] -> (false, ")")

| hd::tl -> let name = snd hd in

if List.exists (fun (_, s) -> s = name) tl then

(true, name)

else formals_name_test tl in

let (same_name, name) = formals_name_test fdecl.Ast.formals in

if same_name then

raise(Failure("\"" ^ name ^ "} has a previous declaration")

else
let formals_types = List.map fst fdecl.Ast.formals in
let updated_funcs = {
    functions = (fdecl.Ast.rtype, fdecl.Ast.fname, formals_types)::func
}
in
let func_scope = { scope with
    parent = Some(scope);
    variables = fdecl.Ast.formals
}
in
let sast_block = stmt func_scope updated_funcs []

let valid_return = check_return fdecl sast_block in

if(not valid_return) then
    raise(Failure("no return statement in function " ^
        string_of_type fdecl.Ast.rtype ^
        fdecl.Ast.fname ^ '(' ^ (fun l -> String.concat '
', (List.map string_of_type (List.map fst l))) fdecl.Ast.formals) ^
        ')")
else
    let (sast_body, sast_locals) = match sast_block with
        Sast.Block(stmts_list, locals_list) ->
        List.iter (fun (_, name) ->
            try
                let _ = List.find (fun (_, s) -> s = name)
            in raise(Failure("declaration of " ^ name ^ ",' shadows a parameter")
        ) fdecl.Ast.formals; (stmts_list, locals_list)
    in
    let sast_fdecl = {
        rtype = fdecl.Ast.rtype;
        fname = fdecl.Ast.fname;
        formals = fdecl.Ast.formals;
        locals = sast_locals;
        body = sast_body
    } in
    check var_decl (sast_fdecl::func_decl) scope updated_funcs tl

in
check [] [] init_scope init_funcs (List.rev program)
in (new_var_decl, new_func_decl)

(* 10:45 Dec 16, 2012 wzh *)
open Type

type expr =
(*) Expressions *)

  Id of string
  (* foo *)

  Literal of constant
  (* 12 $PgDn$ '@' *)

  Assign of string * expr
  (* foo = 823 *)

  Binop of expr * binop * expr
  (* x=a+b *)

  Unop of unop * expr
  (* !a -t p@x *)

  Call of string * expr list
  (* foo(a,23) *)

  NoExpr
  (* for(;;) *)

(* Variable initialization *)

  WithInit of string * expr
  (* int a = 3; *)

  WithoutInit of string
  (* float b; *)

(* Variable declaration *)

t * (string * expr) list
(* int a, b = 3, c; *)

(* Statements *)

  Expr of expr
  (* foo = bar + 3; *)

  Block of stmt list
  (* while(i<10) ... end while *)

  Return of expr
  (* return 42; *)

  If of expr list * stmt list
  (* if(foo==42) ... elseif(foo<42) ... else ... end if *)

  For of expr * expr * expr * stmt
  (* for(i=0;i<10;i=i+1) ... end for *)

  ForWithDecl of var_decl * expr * expr * stmt
  (* for(int i=0;i<10;i=i+1) ... end for *)

  While of expr * stmt
  (* while(i<10) ... end while *)

  Switch of expr * expr list * stmt list
  (* switch(x+y) case 1: a=8; b=23; case 2: c=8; d=23; end switch *)

  Continue
  (* continue; *)

  Break
  (* break; *)

(* Function declaration *)

rtype : t;
(* return type of the function *)

fname : string;
(* name of the function *)

formals : (t * string) list;
(* types and names of formal arguments *)

body : stmt list;
(* body of the function *)
/ast.mli:type decls =
/ast.mli: Var_decl of var_decl
/ast.mli: Func_decl of func_decl
/ast.mli:
/ast.mli:type program = decls list (* global vars, funcs *)

/bytecode.ml:(* Nov 29, @Jinqi*)
/bytecode.ml:
/bytecode.ml:open Type
/bytecode.ml:
/bytecode.ml:type bstmt =
/bytecode.ml: Glb of int (* Indicate the number of gloabal variables *)
/bytecode.ml: Pop of int (* Pop out several elements from the stack *)
/bytecode.ml: Bin of Type.binop (* Perform arithmetic on the top two elements of stack *)
/bytecode.ml: Uni of Type.unop (* Perform unary operation on the top of stack *)
/bytecode.ml: Psi of int (* Push an integer *)
/bytecode.ml: Psp of int * int (* Push a pointer *)
/bytecode.ml: Psco of int * int * int (* push a color *)
/bytecode.ml: Psf of float (* Push a floating number *)
/bytecode.ml: Pse of bool (* Push a boolean number *)
/bytecode.ml: Ps of char (* Push a char *)
/bytecode.ml: Pss of string (* Push a string *)
/bytecode.ml: Lod of int (* Fetch global variable *)
/bytecode.ml: Str of int (* Store global variable *)
/bytecode.ml: Lfp of int (* Load frame pointer relative *)
/bytecode.ml: Sfp of int (* Store frame pointer relative *)
/bytecode.ml: Jsr of int (* Call function by absolute address *)
/bytecode.ml: Ent of int (* Push FP, FP -> SP, SP += i *)
/bytecode.ml: Rts of int (* Restore FP, SP, consume formals, push result *)
/bytecode.ml: Beq of int (* Branch relative if top-of-stack is zero *)
/bytecode.ml: Bne of int (* Branch relative if top-of-stack is non-zero *)
/bytecode.ml: Bra of int (* Branch relative *)
/bytecode.ml: Hlt (* Terminate *)

/bytecode.ml:let string_of_stmt = function
/bytecode.ml: Glb(i) -> "glb " ^ string_of_int i
/bytecode.ml:Psi(i) -> "psi " ^ string_of_int i
/bytecode.ml:Psp(i,j) -> "psp " ^ string_of_int i ^ " " ^ string_of_int j
/bytecode.ml:Psc(i,j,k) -> "psco " ^ string_of_int i ^ " " ^ string_of_int j ^ " " ^ string_of_int k
/bytecode.ml:Psf(i) -> "psf " ^ string_of_float i
/bytecode.ml:Psb(i) -> if i then "psb " ^ "1" else "psb " ^ "0"
/bytecode.ml:Psc(i) -> "psc " ^ Char.escaped i
/bytecode.ml:Pss(i) -> "pss " ^ i ^ "\0"
/bytecode.ml:Psk(k) -> "psk " ^ k
/bytecode.ml:Pop(i) -> "pop " ^ string_of_int i
/bytecode.ml:Bin(Type.Add) -> "add"
/bytecode.ml:Bin(Type.Sub) -> "sub"
/bytecode.ml:Bin(Type.Mul) -> "mul"
/bytecode.ml:Bin(Type.Div) -> "div"
/bytecode.ml:Bin(Type.Mod) -> "mod"
/bytecode.ml:Bin(Type.And) -> "and"
/bytecode.ml:Bin(Type.Or) -> "or"
/bytecode.ml:Bin(Type.Eq) -> "eql"
/bytecode.ml:Bin(Type.Neq) -> "neq"
/bytecode.ml:Bin(Type.Lt) -> "lt"
/bytecode.ml:Bin(Type.Leq) -> "leq"
/bytecode.ml:Bin(Type.Gt) -> "gt"
/bytecode.ml:Bin(Type.Geq) -> "geq"
/bytecode.ml:Uni(Type.Neg) -> "neg"
/bytecode.ml:Uni(Type.Not) -> "not"
/bytecode.ml:Uni(Type.Atx) -> "atx"
/bytecode.ml:Uni(Type.Aty) -> "aty"
/bytecode.ml:Uni(Type.Atr) -> "atr"
/bytecode.ml:Uni(Type.Atg) -> "atg"
/bytecode.ml:Uni(Type.Ab) -> "ab"
/bytecode.ml:Uni(Type.Atb) -> "atb"
/bytecode.ml:Uni(Type.Atk) -> "atk"
/bytecode.ml:Lod(i) -> "lod " ^ string_of_int i
/bytecode.ml:Str(i) -> "str " ^ string_of_int i
/bytecode.ml:Lfp(i) -> "lfp " ^ string_of_int i
/bytecode.ml:Sfp(i) -> "sfp " ^ string_of_int i
/bytecode.ml:Jsri) -> "jsr " ^ string_of_int i
/bytecode.ml:Ent(i) -> "ent " ^ string_of_int i
/bytecode.ml:Rts(i) -> "rts " ^ string_of_int i
/bytecode.ml:Bne(i) -> "bne " ^ string_of_int i
/bytecode.ml:Beq(i) -> "beq " ^ string_of_int i
/bytecode.ml:Bra(i) -> "bra " ^ string_of_int i
/bytecode.ml:Hlt -> "hlt"
let string_of_prog p = "" ^ let funca = Array.mapi (fun i s -> "" ^ string_of_stmt s) p.text in String.concat "\n" (Array.to_list funca)

open Bytecode
open Type
open Sast

module StringMap = Map.Make(String)

(*test region*)
(*Keqiu Hu & Lizhong Zhang*)
(*Reference: Microc--Stephen Edwards & TML--Yan Zou*)
(*Symbol table as MICROC*)

type env={
    function_index:int StringMap.t;
    global_index:int StringMap.t;
    (* local_index:int StringMap.t;*)
}

(*function to merge var_decl into a Block t*string list*)

(*let merge (t,str,_)=function*)

(*add names in list to the map (referto zou)*)

let add_to_map list map start_index = snd (List.fold_left (fun (i, map) name -> (*print_endline (name ^": " ^ (string_of_int i)); (* debug *)*) (i + 1, StringMap.add name i map)) (start_index, map) list)

(*function to perform preglobal expression handling*)

let add_to_map list map start_index = snd (List.fold_left (fun (i, map) name -> (*print_endline (name ^": " ^ (string_of_int i)); (* debug *)*) (i + 1, StringMap.add name i map)) (start_index, map) list)

(**)

let exprId =function

let takestr (a,b,c)=a

let expr2 env =function

Assign(s,e)->[Str (StringMap.find s env.global_index)]*)
let rec enum stride n = function
  | [] -> []
  | hd::tl -> (n, hd)::enum stride (n+stride) tl

(*add pairs into a map as MICORC*)
let string_map_pairs map pairs =
  List.fold_left (fun m (i, n) -> StringMap.add n i m) map pairs

(*main translate function, modified from MICROC*)
let res = function
  | WithInit (a, b) -> a
  | WithoutInit (c) -> c

let initos = function
  (t, v) -> List.map res v

(*take the t*string list out from a BLOCK*)
let takelistout = function
  | Block (a, e) -> e
  | d -> [Int_T, "noinit"]

(*turn variable list into t*string list*)
let var_list t =
  List.map (fun (a, b, c) -> (a, b)) t

let init_value t = function
  | Some (e) -> e
  | None ->

let lit =
  match t with
  | Void_T -> Integer(0)
  | Point_T -> Point(0, 0)
  | Int_T -> Integer(0)
  | Key_T -> String("Null")
  | Color_T -> Color(0, 0, 0)
  | Float_T -> Float(0.0)
  | Char_T -> Char("000")
  | String_T -> String("0")
  | Boolean_T -> Boolean(false)

in Literal(lit), t, true
let translate (globals,functions)=
  let glbexpr=(List.map (fun (a,b,c)->c) globals) in
  let precube=List.map (fun (a,b,c)->b) globals in
  (* Allocate "addresses" for each global variable *)
  let global_indexes = string_map_pairs StringMap.empty (enum 1 0 precute) in
  (* Assign indexes to function names; built-in "print" is special *)
  let built_in_functions = StringMap.add "print" (-1) StringMap.empty in

  (*buildin function manipulation*)
  let built_in_functions = string_map_pairs StringMap.empty (enum (1) (1) (let builtin =
    let builtin =
      let mutable index = 0;
      let rec add_builtin (name,expr)=
        match expr with
        | Integer(i) -> StringMap.add name (index) expr
        | Char(c)    -> StringMap.add name (index) expr
        | String(s)  -> StringMap.add name (index) expr
        | Float(f)   -> StringMap.add name (index) expr
        | Boolean(b) -> StringMap.add name (index) expr
        | Point(x,y) -> StringMap.add name (index) expr
        | Key(k)     -> StringMap.add name (index) expr
        | Color(r,g,b) -> StringMap.add name (index) expr
        | StringOfInt()    -> StringMap.add name (index) expr
        | StringOfFloat()  -> StringMap.add name (index) expr
        | StringOfChar()   -> StringMap.add name (index) expr
        | StringOfBoolean() -> StringMap.add name (index) expr
        | StringOfPoint()  -> StringMap.add name (index) expr
        | StringOfKey()    -> StringMap.add name (index) expr
        | StringOfColor()  -> StringMap.add name (index) expr
        | StringOfAscii()  -> StringMap.add name (index) expr
        | StringOfChar()   -> StringMap.add name (index) expr
        | StringOfStringLength() -> StringMap.add name (index) expr
        | StringOfCharOfString() -> StringMap.add name (index) expr
        | StringOfSubOfString() -> StringMap.add name (index) expr
        | StringOfDelay()  -> StringMap.add name (index) expr
        | StringOfCompareColor() -> StringMap.add name (index) expr
      in
      add_builtin (name,expr)
    in
    builtin
  in

  let function_indexes = string_map_pairs built_in_functions (enum 1 0 (List.map (fun f->f.fname) functions)) in

  (* Translate a function in AST form into a list of bytecode statements *)
  let rec block local_index next_index num_params sl=
    let rec expr (exprdetail,t,int)=match exprdetail with
    | Literal i->begin match i with
      | Integer(j)->[Psi j]
      | Char(c)->[Psc c]
      | String(s)->[Pss s]
      | Float(f)->[Psf f]
      | Boolean(b)->[Psb b]
      | Point(x,y)->[Psp(x,y)]
      | Key(k)->[Pss k]
      | Color(r,g,b)->[Psco (r,g,b)](*later change*)
    end
    | Id(s)->[Pf]
    | (try [Lfp (StringMap.find s local_index)]
      with Not_found -> try [Lod (StringMap.find s global_indexes)]
  in

  let rec expr detail t=int=match exprdetail with
  | Literal i->begin match i with
    | Integer(j)->[Psi j]
    | Char(c)->[Psc c]
    | String(s)->[Pss s]
    | Float(f)->[Psf f]
    | Boolean(b)->[Psb b]
    | Point(x,y)->[Psp(x,y)]
    | Key(k)->[Pss k]
    | Color(r,g,b)->[Psco (r,g,b)](*later change*)
  end
  | Id(s)->[Pf]
  | (try [Lfp (StringMap.find s local_index)]
    with Not_found -> try [Lod (StringMap.find s global_indexes)]

with Not_found -> raise (Failure ("undeclared variable"^s)))

|Binop(e1,op,e2) -> expr e1 @ expr e2 @[Bin op]

|Unop (op, e) -> (expr e) @ [Uni op]

|Assign (s,e)->expr e @

(try [Sfp (StringMap.find s local_index)]
with Not_found->try [Str (StringMap.find s
global_indexes)]
with Not_found->raise (Failure("Undeclared variable"^s)))

|Call(fname,actuals) ->

if(fname="__leftClickPoint") then
if(List.length actuals>0) then
(List.concat (List.map (fun e->
(expr e) @ [Jsr (-1);Pop 1]) actuals))@ [Psi 0]
else []
else if(fname="__rightClickPoint") then
List.concat (List.map (fun e->(expr e)) actuals) @[Jsr (-2);Pop 1]@[Psi 0]
else if(fname="__doubleClick") then
List.concat (List.map (fun e->(expr e)) actuals) @[Jsr (-3);Pop 1]@[Psi 0]
else if(fname="__movePoint") then
List.concat (List.map (fun e->(expr e)) actuals) @[Jsr (-4)]@[Psi 0]
else if(fname="__leftDown") then
List.concat (List.map (fun e->(expr e)) actuals) @[Jsr (-5)]@[Psi 0]
else if(fname="__rightDown") then
List.concat (List.map (fun e->(expr e)) actuals)@[Jsr (-6)]@[Psi 0]
else if(fname="__leftUp") then
List.concat (List.map (fun e->(expr e)) actuals)@[Jsr (-7);Pop 1;Pop 1]@[Psi 0]
else if(fname="__rightUp") then
List.concat (List.map (fun e->(expr e)) actuals)@[Jsr (-8);Pop 1;Pop 1]@[Psi 0]
else if(fname="__keyDown") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-9); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__keyUp") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-10); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__keyStroke") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-11); Pop 1]@[Psi 0]
else if (fname = "__comboStroke") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-12); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__getPoint") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-13)]
else if (fname = "__getline") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-14)]
else if (fname = "__print") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-15); Pop 1]@[Psi 0]
else if (fname = "__stringOfInt") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-16)]
else if (fname = "__stringOfFloat") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-17); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__stringOfChar") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-18); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__stringOfBoolean") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-19); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__stringOfPoint") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-20); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__stringOfKey") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr (-21); Pop 1; Pop 1]@[Psi 0]
else if (fname = "__stringOfColor") then
List.concat (List.map (fun e -> (expr e)) actuals) @[Jsr
else if (fname = "__getASCII") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-22); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__getChar") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-23); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__getStringLength") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-24); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__getCharOfString") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-25); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__setCharOfString") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-26); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__subOfString") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-27); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__delay") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-28); Pop 1; Pop 1] @[Psi 0]
else if (fname = "__compareColor") then
    List.concat (List.map (fun e -> (expr e)) actuals) @ [Json
      (-29); Pop 1; Pop 1] @[Psi 0]
else
  (try
    (List.concat (List.map expr (List.rev actuals))) @ [Json
      (StringMap.find fname function_indexes)]
    with Not_found -> raise (Failure("undefined function " ^ fname)))
| NoExpr -> []
in
(* modified in order to create new variables inside a block(trial)*)
let rec stmt = function
  Block (sl, vl) ->
    let vname_list = List.map (fun (_, n) -> n) vl in
    let new_local_index =
      add_to_map vname_list local_index next_index
    in
    let num_locals = List.length vl in
    (block new_local_index (next_index + num_locals) num_params
sl) @(if num_locals>0 then
        [Pop num_locals]
    else[])

Expr e -> expr e @ [Pop 1]
Return e -> expr e @ [Rts num_params]

(*no idea how to change this if*)

If (p, t) -> begin match (p, t) with
    |([],[])->[Bra 1]
    |([],t)->let t'=(List.hd t) in stmt t'
    |(p,t)->let p'=List.hd p and t'= stmt (List.hd t) and f'= stmt (If(List.tl p,List.tl t)) in expr p' @ [Beq(2 + List.length t')] @t' @ [Bra(1 + List.length f')] @ f'
end

VarDecl(t,name,init)->expr(init_value t init)(*begin match e1 with
(a,b,Some(c))->expr c
    |(a,b,None)-> [Bra 1] end*)
For (e1, e2, e3, b) ->
    stmt (Block([Expr(e1); While(e2,;
        Block([b;Expr(e3)],takelistout b)],takelistout b))
ForWithDecl (v1, e2, e3, b) ->
    let li=(takelistout b)@(var_list v1) in
    stmt (Block([While(e2,Block([b;Expr(e3)],li))],li))
Switch (e, el, sl) -> begin match (e, el, sl) with
    |(_, [],[])->[Bra 1]
    |(_, [], t) -> let t'=(List.hd t) in stmt t'
    |(e, el, sl) -> let p = List.hd el and t' = stmt (List.hd sl) and f' = stmt (Switch(e, List.tl el, List.tl sl)) in expr e @ expr p @ [Bin(Type.Eq)] @ [Bne(2)] @ [Bra(2 + List.length t')] @ t' @ [Bra(1+List.length f')] @ f'
end
While (e, b) ->
let b' = stmt b and e' = expr e in
[Bra (1+ List.length b')] @ b' @ e' @
[Bne ((List.length b' + List.length e'))]

in
List.concat (List.map stmt sl)
in
let data=
let assign_globals=
List.map (fun (t,name,init)->
Expr(Assign(name,init_value t init),t,true)) globals
in
[Glb (List.length assign_globals)]@
block StringMap.empty 0 0 assign_globals @
try
[Jsr (StringMap.find "main" function_indexes);Hlt]
with Not_found->
raise (Failure "main function not found")
in
let text=
let trans_func f=
let param_names=List.map snd f.formals in
let num_params=List.length param_names in
let param_index=
add_to_map param_names StringMap.empty (-1 -num_params)
in
let local_names=List.map snd f.locals in
let num_locals=List.length local_names in
let local_index=
add_to_map local_names param_index 1
in
let bodycodes=
block local_index (num_locals +1) num_params f.body
in
List.iter (function b->()) bodycodes;(*not quite understand*)
[Ent num_locals]@bodycodes@
(*
Void_T->Integer(0)
Point_T -> Point(0,0)
Int_T -> Integer(0)
Key_T -> String("")
Color_T -> Color(0,0,0)
Float_T -> Float(0.0)
Char_T -> Char('000')
String_T -> String("")
Boolean_T -> Boolean(false)

(match f.rtype with
|Int_T -> [Psi 0]
|Key_T -> [Psk ""]
|Color_T -> [Psco (0,0,0)]
|Point_T -> [Psp (0,0)]
|Float_T -> [Psf 0.0]
|Char_T -> [Psc '000']
|String_T -> [Pss ""]
|Boolean_T -> [Psb false]
|Void_T -> [Psi 0]) @
[Rts num_params]

let func_bodies=
(List.map trans_func functions)
in
let (_func_offsets)=
(List.fold_left (fun (i,offsets) body->
((List.length body) +i,i::offsets))
((List.length data),[])func_bodies
in
let array_offsets=Array.of_list (List.rev func_offsets) in
List.map (function
|Jsr i when i>=0->Jsr array_offsets.(i)
|_ as s->s)
(List.concat (data::func_bodies))
in
Bytecode.string_of_prog {num Globals=10;text=Array.of_list text}
(Open Bytecode*)
(Open Bytecode*)
if (Array.length Sys.argv<2) then
    print_endline "Usage: easkey <file_name>"
else
    let in_filename=Sys.argv.(1) in
    let fin =open_in in_filename in
    let lexbuf=Lexing.from_channel fin in
    let program=Analyzer.semantically_check (Parser.program Scanner.token lexbuf) in
    let listing=Compile.translate program in
    print_endline listing

#include <memory>
#include <string>
#include "tyname.h"

enum Type
{
  glb, psi, psf, pss, psb, psc, pop, uop, bin, lod, str, lfr,
  alc, fld, sfd, pst, scd, next, add, sub, mul, divs, mod, eq, neq, lt, leq, gt, geq, val, chd,
  neg, num, at, cln, fat, null
};

enum subT
using namespace std;

class instructions
{
public:

instructions(string type, shared_ptr<void> operand);
~instructions(void);

//Type getType();
//subT getSubType();

shared_ptr<void> getOperand();

string getType();

private:

string type;
//Type type;

subT subType;

shared_ptr<void> operand;

};
/* 10:45 Dec 16, 2012 wzh */

/* open Type */

/* open Ast */

/*
token LPAREN RPAREN LBRACKET RBRACKET COLON SEMI COMMA PLUS MINUS TIMES DIVIDE MOD ASSIGN EQ NEQ LT
*/

/*
token LEQ GT GEQ NOT AND OR ATX ATY ATR ATG ATB ATK IF ELSEIF ELSE SWITCH CASE DEFAULT FOR WHILE
*/

/*
token CONTINUE BREAK FUNCTION RETURN END VOID INT FLOAT CHAR STRING BOOLEAN POINT COLOR KEY EOF
*/

/*
token <int> INT_LITERAL
*/

/*
token <char> CHAR_LITERAL
*/

/*
token <float> FLOAT_LITERAL
*/

/*
token <string> STRING_LITERAL
*/

/*
token <bool> BOOLEAN_LITERAL
*/

/*
token <int*int> POINT_LITERAL
*/

/*
token <int*int*int> COLOR_LITERAL
*/

/*
token <string> KEY_LITERAL
*/

/*
token <string> ID
*/
%left ASSIGN
%left OR
%left AND
%left EQ NEQ
%left LT GT LEQ GEQ
%left PLUS MINUS
%left TIMES DIVIDE MOD
%nonassoc UMINUS NOT ATX ATY ATR ATG ATB ATK

%start program
%type <Ast.program> program

program:
  /* nothing */ { [] }
| program vdecl { Var_decl($2) :: $1 }
| program fdecl { Func_decl($2) :: $1 } /* in reverse order */
// parser.mly: FUNCTION return_type_specifier ID LPAREN formals_opt RPAREN stmt_list END

// parser.mly:   { { rtype= $2; 

// parser.mly:     fname   = $3; 

// parser.mly:     formals = $5; 

// parser.mly:     body= List.rev $7 } } 

// parser.mly:

// parser.mly:formals_opt:

// parser.mly:   /* nothing */   { [] } 

// parser.mly:| formal_list       { List.rev $1 } 

// parser.mly:

// parser.mly:formal_list:

// parser.mly:   type_specifier ID       { [(1, 2)] } 

// parser.mly:| formal_list COMMA type_specifier ID   { (1, 3)::1 } /* in reverse order */

// parser.mly:

// parser.mly:stmt_list:

// parser.mly:   /* nothing */   { [] } 

// parser.mly:| stmt_list stmt       { 2 :: 1 } /* in reverse order */

// parser.mly:

// parser.mly:stmt:

// parser.mly:   expr SEMI         { Expr($1) }
RETURN expr SEMI  { Return($2) }

RETURN SEMI  { Return(NoExpr) }

IF LPAREN expr RPAREN stmt_list END IF  { If([$3], [Block(List.rev $5)]) }

IF LPAREN expr RPAREN stmt_list ELSE stmt_list END IF  { If([$3], [Block(List.rev $5); Block(List.rev $7)]) }

IF LPAREN expr RPAREN stmt_list elseif_stmt END IF  { If($3::fst $6, Block(List.rev $5)::snd $6) }

FOR LPAREN expr_opt SEMI expr_opt SEMI expr_opt RPAREN stmt_list END FOR  { For($3, $5, $7, Block(List.rev $9)) }

FOR LPAREN vdecl expr_opt RPAREN stmt_list END FOR  { ForWithDecl($3, $4, $6, Block(List.rev $8)) }

WHILE LPAREN expr RPAREN stmt_list END WHILE  { While($3, Block(List.rev $5)) }

SWITCH LPAREN expr RPAREN stmt_list case_stmt END SWITCH  { Switch($3, fst $6, snd $6) }

SWITCH LPAREN expr RPAREN stmt_list END SWITCH  { Switch($3, [], []) } /* stmts before cases will not be executed */

CONTINUE SEMI  { Continue }

BREAK SEMI  { Break }

vdecl  { VarDecl((fst $1, snd $1)) }

/* in reverse order */

init  { [ $1 ] }

init_list COMMA init  { $3 :: $1 }
```plaintext
./parser.mly:

./parser.mly:init:

./parser.mly: ID ASSIGN expr   { ($1, $3) }

./parser.mly: ID               { ($1, NoExpr) }

./parser.mly:vdecl:

./parser.mly: typeSpecifier init_list SEMI   { ($1, List.rev $2) }

./parser.mly:.expr:

./parser.mly: ID               { Id($1) }

./parser.mly: constant         { Literal($1) }

./parser.mly: ID ASSIGN expr   { Assign($1, $3) }

./parser.mly: expr PLUS expr   { Binop($1, Add, $3) }

./parser.mly: expr MINUS expr  { Binop($1, Sub, $3) }

./parser.mly: expr TIMES expr  { Binop($1, Mul, $3) }

./parser.mly: expr DIVIDE expr { Binop($1, Div, $3) }

./parser.mly: expr MOD expr    { Binop($1, Mod, $3) }

./parser.mly: expr GT expr     { Binop($1, Gt, $3) }

./parser.mly: expr LT expr     { Binop($1, Lt, $3) }

./parser.mly: expr GEQ expr    { Binop($1, Geq, $3) }
```

./parser.mly: expr LEQ expr { Binop($1, Leq, $3) }

./parser.mly: expr EQ expr { Binop($1, Eq, $3) }

./parser.mly: expr NEQ expr { Binop($1, Neq, $3) }

./parser.mly: expr AND expr { Binop($1, And, $3) }

./parser.mly: expr OR expr { Binop($1, Or, $3) }

./parser.mly: MINUS expr %prec UMINUS { Unop(Neg, $2) }

./parser.mly: NOT expr { Unop(Not, $2) }

./parser.mly: expr ATX { Unop(Atx, $1) }

./parser.mly: expr ATY { Unop(Aty, $1) }

./parser.mly: expr ATR { Unop(Atr, $1) }

./parser.mly: expr ATG { Unop(Atg, $1) }

./parser.mly: expr ATB { Unop(Atb, $1) }

./parser.mly: expr ATK { Unop(Atk, $1) }

./parser.mly: ID LPAREN actuals_opt RPAREN { Call($1, $3) }

./parser.mly: LPAREN expr RPAREN { $2 }

./parser.mly:

./parser.mly:constant:

./parser.mly: INT_LITERAL { Integer($1) }

./parser.mly: CHAR_LITERAL { Char($1) }

./parser.mly: STRING_LITERAL { String($1) }

./parser.mly: FLOAT_LITERAL { Float($1) }
Boolean($1) 

Point(fst $1, snd $1) 

Key($1) 

Color((let fst (a,_,_) = a in fst $1), (let snd (_,b,_) = b in snd $1), (let thd (_,_,c) = c in thd $1))) 

ELSEIF LPAREN expr RPAREN stmt_list 

$3::fst $6, Block(List.rev $5)::snd $6 

ELSEIF LPAREN expr RPAREN stmt_list ELSE stmt_list 

$3, [Block(List.rev $5); Block(List.rev $7)]) 

CASE expr COLON stmt_list 

$2, [Block(List.rev $4)]) 

CASE expr COLON stmt_list case_stmt 

$2::fst $5, Block(List.rev $4)::snd $5 

DEFAULT COLON stmt_list 

[], [Block(List.rev $3)]) 

expr_opt: 

/* nothing */ 

NoExpr 

e $1
/* nothing */
{ [] }

actuals_list
{ List.rev $1 }

actuals_list:

expr
{ [$1] }

actuals_list COMMA expr
{ $3 :: $1 } /* in reverse order */

typeSpecifier:

INT
{ Int_T }

FLOAT
{ Float_T }

CHAR
{ Char_T }

STRING
{ String_T }

BOOLEAN
{ Boolean_T }

POINT
{ Point_T }

COLOR
{ Color_T }

KEY
{ Key_T }

return_typeSpecifier:
#include "point.h"

/*Author: Keqiu Hu

the point class
*/

point::point(int a,int b)
{
x=a;
y=b;
}

point::point()
{
x=0;
y=0;
}

point::~point(void)
{
@Author: Keqiu Hu

the point class declaration

*/

class point
{
    public:
        int x;
        int y;
        point(int,int);
        point();
        ~point(void);
};
# program.h

```cpp
#include <vector>
#include <stack>

using namespace std;

class program {
public:
  int inpoint;
  vector<tyname> progs;
  //stack< tyname > progs;
  vector<instructions> instru;
  void dump();
  vector< tyname > globals;
  program(vector <instructions> ins);
  ~program();
  void execute();
};
```

# rgb.cpp

```cpp
#include "rgb.h"

rgb::rgb(int a, int d, int c) {
```

```cpp```
rgb.cpp: r=a;
rgb.cpp: g=d;
rgb.cpp: b=c;
rgb.cpp:
rgb.cpp:}
rgb.cpp:rgb::rgb(){
rgb.cpp: r=0;
rgb.cpp: g=0;
rgb.cpp: b=0;
rgb.cpp:}
rgb.cpp:rgb::~rgb(void){
rgb.cpp:}
rgb.cpp:
rgb.cpp:
rgb.cpp:rgb.h:#pragma once
rgb.h:class rgb
rgb.h:{
rgb.h:public:
rgb.h:    int r;
./rgb.h:  int g;
./rgb.h:  int b;

./rgb.h:  rgb(int,int,int);
./rgb.h:  rgb();
./rgb.h:  ~rgb(void);
./rgb.h:  ;

./sast.mli:(* 11:37 Dec 16, 2012 hxy, wzh *)

./sast.mli:

./sast.mli:open Type

./sast.mli:

./sast.mli:type expr_detail =         (* Expressions *)

./sast.mli:  Id of string              (* foo *)
./sast.mli:  Literal of constant       (* 12 $PgDn$ '@' *)
./sast.mli:  Assign of string * expr   (* foo = 823 *)
./sast.mli:  Binop of expr * binop * expr (* x=a+b *)
./sast.mli:  Unop of unop * expr       (* !a -t p@x *)
./sast.mli:  Call of string * expr list (* foo(a,23) *)
./sast.mli:  NoExpr                     (* for(;;) *)
Expressions with type and integer that indicates constant expression or not

Variable declaration

Statements

foo = bar + 3

statement list and var list

return 42

if()...

expression list

statement list

for(int i=0; i<10; i++)

end for

continue statement with num of locals that need to be popped

break statement num of locals that need to be popped

int a, b = 3, c;

Function declaration
rtypen : t;

fname : string ;

formals : (t * string) list;

locals : (t * string) list ;

body : stmt list ;

program =

var_decl list * func_decl list

let get_point str = (* string -> int * int = <fun>, e.g. "{12,345}" -> (12,345) *)

let index = String.index str ',' in

(int_of_string (String.sub str 1 (index-1)), int_of_string (String.sub str (index+1) (String.length str-index-2)))

let get_color str = (* string -> int * int * int = <fun>, e.g. "#12,23,34#" -> (12,23,34) *)

let index1 = String.index_from str 0 ',' in

let index2 = String.index_from str (index1+1) ',' in

(int_of_string (String.sub str 1 (index1-1)), int_of_string (String.sub str (index2+1) (String.length str-index2-2)))

let check_color_boundary (r, g, b) =

let list = [r; g; b] in

let checked_list = List.map (fun c -> if c >= 0 && c <= 255 then c else raise(Failure("color #" ^ (string_of_int r) ^ "," ^ (string_of_int g) ^ "," ^ (string_of_int b) ^ "# out of boundary"))) list

in (fun [r; g; b] -> (r, g, b)) checked_list
Char.chr (int_of_string (get_ascii lxm))
raise (Failure("Character " ^ lxm ^ " out of boundary")))

{ FLOAT_LITERAL(float_of_string lxm) } (* float *)

{ STRING_LITERAL(get_str (get_key lxm) "") } (* string *)

{ BOOLEAN_LITERAL(bool_of_string lxm) } (* boolean *)

{ POINT_LITERAL(get_point lxm) } (* point *)

{ COLOR_LITERAL(check_color_boundary(get_color lxm)) } (* color *)

{ ID(lxm) } (* key *)

(* End-of-comment *)

(* Eat everything else *)

#include "tyname.h"


/tyname.cpp:

/tyname.cpp: tyname::~tyname(void)

/tyname.cpp: {

/tyname.cpp: }

/tyname.h:#pragma once

/tyname.h:#include<string>

/tyname.h:#include<memory>

/tyname.h:using namespace std;

/tyname.h://to implement type name specification

/tyname.h:class tyname

/tyname.h:

/tyname.h:public:

/tyname.h: tyname(string,shared_ptr<void>);

/tyname.h: ~tyname(void);

/tyname.h: string name;

/tyname.h: shared_ptr<void> ptr;

/tyname.h:;