

Fundamentals of Computer Systems

Thinking Digitally

Stephen A. Edwards and Martha Kim

Columbia University

Spring 2012

The Subject of this Class

0

The Subjects of this Class

0

1

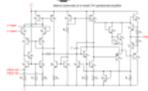
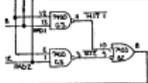
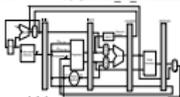
But let your communication be, Yea, yea; Nay, nay: for whatsoever is more than these cometh of evil.

— Matthew 5:37

Engineering Works Because of Abstraction



```
;; voice 1 wave select
ld a, (#CHI_W_NUM)
and a
ld a, (#CHI_W_SEL)
jr nz, #00b4
ld a, (#CHI_E_TABLE0)
```



Application Software

Operating Systems

Architecture

Micro-Architecture

Logic

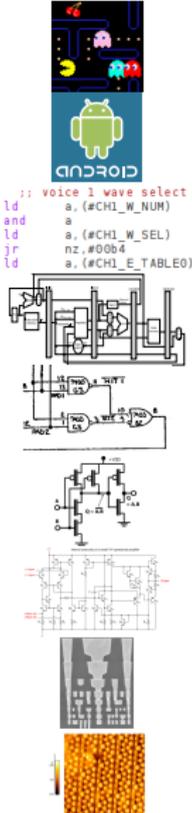
Digital Circuits

Analog Circuits

Devices

Physics

Engineering Works Because of Abstraction



Application Software COMS 3157, 4156, et al.

Operating Systems COMS W4118

Architecture Second Half of 3827

Micro-Architecture Second Half of 3827

Logic First Half of 3827

Digital Circuits First Half of 3827

Analog Circuits ELEN 3331

Devices ELEN 3106

Physics ELEN 3106 et al.

Boring Stuff

Mailing list: csee3827-staff@lists.cs.columbia.edu

<http://www.cs.columbia.edu/~sedwards/classes/2012/3827-spring/>

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Lectures 1:10–2:25 PM, Mon, Wed, 614 Schermerhorn
Jan 18–Apr 30
Holidays: Mar 12–16 (Spring Break)

Assignments and Grading

Weight	What	When
40%	Six homeworks	See Webpage
30%	Midterm exam	March 7th
30%	Final exam	During Finals Week (May 4–11)

Homework is due at the beginning of lecture.

We will drop the lowest of your six homework scores;

you can { skip
omit
forget
ignore
blow off
screw up
feed to dog
flake out on
sleep through } one with no penalty.

Rules and Regulations

You may collaborate with classmates on homework.

Each paper turned in must be unique; work must ultimately be your own.

List your collaborators on your homework.

Don't cheat: if you're stupid enough to try, we're smart enough to catch you.

Tests will be closed-book with a one-page “cheat sheet” of your own devising.

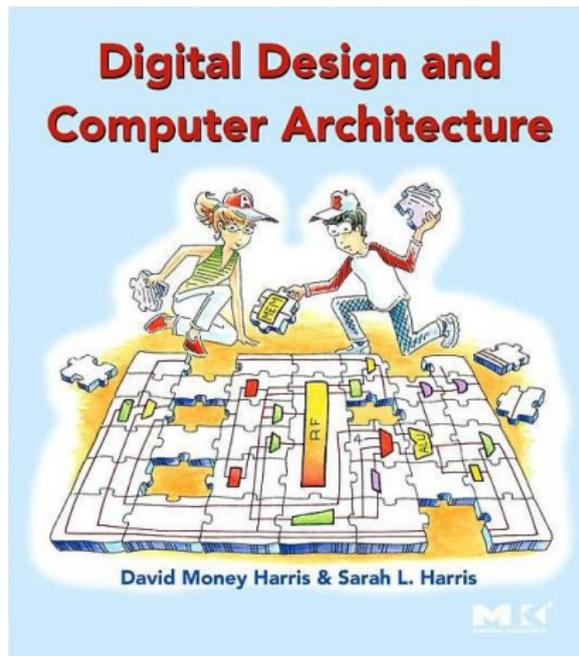
The Text

David Harris and Sarah Harris.

Digital Design and Computer Architecture.

Morgan-Kaufmann, 2007.

Almost precisely right for the scope of this class: digital logic and computer architecture



GILDAN
ULTRA
COTTON

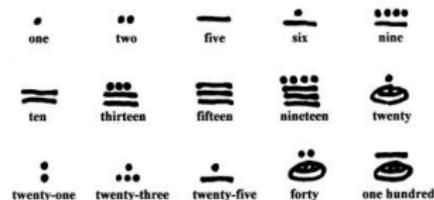
There are only 10 types
of people in the world:
Those who understand binary
and those who don't.

Which Numbering System Should We Use?

Some Older Choices:



Roman: I II III IV V VI VII VIII IX X



Mayan: base 20, Shell = 0

1	𐎠	11	𐎠𐎡	21	𐎠𐎡𐎢	31	𐎠𐎡𐎢𐎣	41	𐎠𐎡𐎢𐎣𐎤	51	𐎠𐎡𐎢𐎣𐎤𐎥
2	𐎠𐎠	12	𐎠𐎡𐎢	22	𐎠𐎡𐎢𐎣	32	𐎠𐎡𐎢𐎣𐎤	42	𐎠𐎡𐎢𐎣𐎤𐎥	52	𐎠𐎡𐎢𐎣𐎤𐎥𐎦
3	𐎠𐎠𐎠	13	𐎠𐎡𐎢𐎣	23	𐎠𐎡𐎢𐎣𐎤	33	𐎠𐎡𐎢𐎣𐎤𐎥	43	𐎠𐎡𐎢𐎣𐎤𐎥𐎦	53	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧
4	𐎠𐎠𐎠𐎠	14	𐎠𐎡𐎢𐎣𐎤	24	𐎠𐎡𐎢𐎣𐎤𐎥	34	𐎠𐎡𐎢𐎣𐎤𐎥𐎦	44	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧	54	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨
5	𐎠𐎠𐎠𐎠𐎠	15	𐎠𐎡𐎢𐎣𐎤𐎥	25	𐎠𐎡𐎢𐎣𐎤𐎥𐎦	35	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧	45	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨	55	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩
6	𐎠𐎠𐎠𐎠𐎠𐎠	16	𐎠𐎡𐎢𐎣𐎤𐎥𐎦	26	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧	36	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨	46	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩	56	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪
7	𐎠𐎠𐎠𐎠𐎠𐎠𐎠	17	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧	27	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨	37	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩	47	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪	57	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫
8	𐎠𐎠𐎠𐎠𐎠𐎠𐎠𐎠	18	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩	28	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪	38	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫	48	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫𐎬	58	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫𐎬𐎭
9	𐎠𐎠𐎠𐎠𐎠𐎠𐎠𐎠𐎠	19	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫	29	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫𐎬	39	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫𐎬𐎭	49	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫𐎬𐎭𐎮	59	𐎠𐎡𐎢𐎣𐎤𐎥𐎦𐎧𐎨𐎩𐎪𐎫𐎬𐎭𐎮𐎯
10	𐎡	20	𐎡𐎢	30	𐎡𐎢𐎣	40	𐎡𐎢𐎣𐎤	50	𐎡𐎢𐎣𐎤𐎥		

Babylonian: base 60

The Decimal Positional Numbering System



Ten figures: 0 1 2 3 4 5 6 7 8 9

$$7 \times 10^2 + 3 \times 10^1 + 0 \times 10^0 = 730_{10}$$

$$9 \times 10^2 + 9 \times 10^1 + 0 \times 10^0 = 990_{10}$$

Why base ten?



Hexadecimal, Decimal, Octal, and Binary

Hex	Dec	Oct	Bin
0	0	0	0
1	1	1	1
2	2	2	10
3	3	3	11
4	4	4	100
5	5	5	101
6	6	6	110
7	7	7	111
8	8	10	1000
9	9	11	1001
A	10	12	1010
B	11	13	1011
C	12	14	1100
D	13	15	1101
E	14	16	1110
F	15	17	1111

Binary and Octal



DEC PDP-8/I, c. 1968

Oct	Bin
0	0
1	1
2	10
3	11
4	100
5	101
6	110
7	111

$$\begin{aligned} \text{PC} &= 0 \times 2^{11} + 1 \times 2^{10} + 0 \times 2^9 + 1 \times 2^8 + 1 \times 2^7 + 0 \times 2^6 + \\ & 1 \times 2^5 + 1 \times 2^4 + 1 \times 2^3 + 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 \\ &= 2 \times 8^3 + 6 \times 8^2 + 7 \times 8^1 + 5 \times 8^0 \\ &= 1469_{10} \end{aligned}$$

Hexadecimal Numbers

Base 16: 0 1 2 3 4 5 6 7 8 9 A B C D E F

Instead of groups of 3 bits (octal), Hex uses groups of 4.

$$\begin{aligned} \text{CAFEF00D}_{16} &= 12 \times 16^7 + 10 \times 16^6 + 15 \times 16^5 + 14 \times 16^4 + \\ &\quad 15 \times 16^3 + 0 \times 16^2 + 0 \times 16^1 + 13 \times 16^0 \\ &= 3,405,705,229_{10} \end{aligned}$$

C	A	F	E	F	0	0	D		Hex		
11001010111111101111000000001101									Binary		
3	1	2	7	7	5	7	0	0	1	5	Octal

Computers Rarely Manipulate True Numbers

Infinite memory still very expensive

Finite-precision numbers typical

32-bit processor: naturally manipulates 32-bit numbers

64-bit processor: naturally manipulates 64-bit numbers

How many different numbers can you

represent with 5

binary	
octal	
decimal	digits?
hexadecimal	

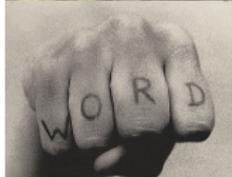
Jargon



Bit Binary digit: 0 or 1



Byte Eight bits



Word Natural number of bits for the processor, e.g., 16, 32, 64



LSB Least Significant Bit (“rightmost”)



MSB Most Significant Bit (“leftmost”)

Decimal Addition Algorithm

$$\begin{array}{r} 434 \\ +628 \\ \hline \end{array}$$

$$4 + 8 = 12$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1 \\ 434 \\ + 628 \\ \hline 2 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1 \\ 434 \\ + 628 \\ \hline 62 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} \text{1 1} \\ 434 \\ + 628 \\ \hline 062 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Decimal Addition Algorithm

$$\begin{array}{r} 1\ 1 \\ 434 \\ +628 \\ \hline 1062 \end{array}$$

$$4 + 8 = 12$$

$$1 + 3 + 2 = 6$$

$$4 + 6 = 10$$

+	0	1	2	3	4	5	6	7	8	9
0	0	1	2	3	4	5	6	7	8	9
1	1	2	3	4	5	6	7	8	9	10
2	2	3	4	5	6	7	8	9	10	11
3	3	4	5	6	7	8	9	10	11	12
4	4	5	6	7	8	9	10	11	12	13
5	5	6	7	8	9	10	11	12	13	14
6	6	7	8	9	10	11	12	13	14	15
7	7	8	9	10	11	12	13	14	15	16
8	8	9	10	11	12	13	14	15	16	17
9	9	10	11	12	13	14	15	16	17	18
10	10	11	12	13	14	15	16	17	18	19

Binary Addition Algorithm

$$\begin{array}{r} 10011 \\ +11001 \\ \hline \end{array}$$

$$1 + 1 = 10$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 1 \\ 10011 \\ +11001 \\ \hline 0 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 11 \\ 10011 \\ +11001 \\ \hline 00 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 011 \\ 10011 \\ +11001 \\ \hline 100 \end{array}$$

$$1 + 1 = 10$$

$$1 + 1 + 0 = 10$$

$$1 + 0 + 0 = 01$$

$$0 + 0 + 1 = 01$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 0011 \\ 10011 \\ +11001 \\ \hline 1100 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \\ 0 + 0 + 1 = 01 \\ 0 + 1 + 1 = 10 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Binary Addition Algorithm

$$\begin{array}{r} 10011 \\ 10011 \\ +11001 \\ \hline 101100 \end{array}$$

$$\begin{array}{l} 1 + 1 = 10 \\ 1 + 1 + 0 = 10 \\ 1 + 0 + 0 = 01 \\ 0 + 0 + 1 = 01 \\ 0 + 1 + 1 = 10 \end{array}$$

+	0	1
0	00	01
1	01	10
10	10	11

Signed Numbers: Dealing with Negativity

A rectangular image showing a handwritten signature in cursive script. The signature reads "John Hancock" and is written in dark ink on a light-colored, aged paper background. The signature is highly stylized, with long, sweeping flourishes, particularly under the "H" and the final "k".

How should both positive and negative numbers be represented?

Signed Magnitude Numbers

You are most familiar with this: negative numbers have a leading –

In binary, a leading 1 means negative:

$$0000_2 = 0$$

$$0010_2 = 2$$

$$1010_2 = -2$$

$$1111_2 = -7$$

$$1000_2 = -0?$$

Can be made to work, but addition is annoying:

If the signs match, add the magnitudes and use the same sign.

If the signs differ, subtract the smaller number from the larger; return the sign of the larger.

One's Complement Numbers

Like Signed Magnitude, a leading 1 indicates a negative One's Complement number.

To negate a number, complement (flip) each bit.

$$0000_2 = 0$$

$$0010_2 = 2$$

$$1101_2 = -2$$

$$1000_2 = -7$$

$$1111_2 = -0?$$

Addition is nicer: just add the one's complement numbers as if they were normal binary.

Really annoying having a -0 : two numbers are equal if their bits are the same or if one is 0 and the other is -0 .



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ZEROS
ARE CREATED
EQUAL**

ZERO CALORIES. MAXIMUM PEPSI™ TASTE.



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Two's Complement Numbers



Really neat trick: make the most significant bit represent a *negative* number instead of positive:

$$1101_2 = -8 + 4 + 1 = -3$$

$$1111_2 = -8 + 4 + 2 + 1 = -1$$

$$0111_2 = 4 + 2 + 1 = 7$$

$$1000_2 = -8$$

Easy addition: just add in binary and discard any carry.

Negation: complement each bit (as in one's complement) then add 1.

Very good property: no -0

Two's complement numbers are equal if all their bits are the same.

Number Representations Compared

Bits	Binary	Signed Mag.	One's Comp.	Two's Comp.
0000	0	0	0	0
0001	1	1	1	1
⋮				
0111	7	7	7	7
1000	8	-0	-7	-8
1001	9	-1	-6	-7
⋮				
1110	14	-6	-1	-2
1111	15	-7	-0	-1

Smallest number

Largest number

Fixed-point Numbers



How to represent fractional numbers? In decimal, we continue with negative powers of 10:

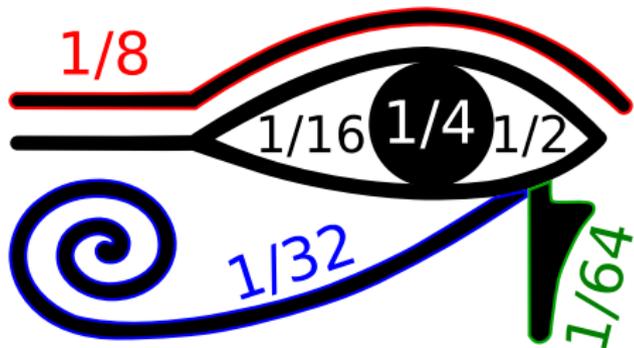
$$31.4159 = 3 \times 10^1 + 1 \times 10^0 + 4 \times 10^{-1} + 1 \times 10^{-2} + 5 \times 10^{-3} + 9 \times 10^{-4}$$

The same trick works in binary:

$$\begin{aligned} 1011.0110_2 &= 1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0 + \\ &\quad 0 \times 2^{-1} + 1 \times 2^{-2} + 1 \times 2^{-3} + 0 \times 2^{-4} \\ &= 8 + 2 + 1 + 0.25 + 0.125 \\ &= 11.375 \end{aligned}$$

F a
u c
Interesting

The ancient Egyptians used binary fractions:



The Eye of Horus

Floating-point Numbers

How can we represent very large and small numbers with few bits?

Floating-point numbers: a kind of scientific notation

IEEE-754 floating-point numbers:

$$\begin{array}{l} \underbrace{1}_{\text{sign}} \underbrace{10000001}_{\text{exponent}} \underbrace{011000000000000000000000}_{\text{significand}} \\ = -1.011_2 \times 2^{129-127} \\ = -1.375 \times 4 \\ = -5.5 \end{array}$$

Binary-Coded Decimal



thinkgeek.com

Humans prefer reading decimal numbers; computers prefer binary.

BCD is a compromise: every four bits represents a decimal digit.

Dec	BCD
0	0000 0000
1	0000 0001
2	0000 0010
⋮	⋮
8	0000 1000
9	0000 1001
10	0001 0000
11	0001 0001
⋮	⋮
18	0001 1000
19	0001 1001
20	0010 0000
⋮	⋮

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 158 \\ +242 \\ \hline \end{array}$$

$$\begin{array}{r} 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \end{array} \text{ First group}$$



BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 158 \\ +242 \\ \hline \end{array}$$

$$\begin{array}{r} 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +0110 \\ \hline \end{array} \begin{array}{l} \text{First group} \\ \text{Correction} \end{array}$$

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 1 \\ 158 \\ +242 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 1 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +\ 0110 \\ \hline 1010\ 0000 \end{array} \begin{array}{l} \text{First group} \\ \text{Correction} \\ \text{Second group} \end{array}$$

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 1 \\ 158 \\ +242 \\ \hline 0 \end{array}$$

$$\begin{array}{r} 1 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +\ 0110 \\ \hline 1010\ 0000 \\ +\ 0110 \\ \hline \hline \end{array}$$

First group
Correction
Second group
Correction

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 11 \\ 158 \\ +242 \\ \hline 00 \end{array}$$

$$\begin{array}{r} 11 \\ 0001\ 0101\ 1000 \\ +0010\ 0100\ 0010 \\ \hline 1010 \\ +0110 \\ \hline 1010\ 0000 \\ +0110 \\ \hline 0100\ 0000 \\ \hline \hline \end{array}$$

First group
Correction
Second group
Correction
Third group

BCD Addition

Binary addition followed by a possible correction.

Any four-bit group greater than 9 must have 6 added to it.

Example:

$$\begin{array}{r} 11 \\ 158 \\ +242 \\ \hline 400 \end{array}$$

	1	1		
	0001	0101	1000	
	+0010	0100	0010	
	<hr/>			
		1010		First group
		+ 0110		Correction
		<hr/>		
		1010	0000	Second group
		+ 0110		Correction
		<hr/>		
		0100	0000	Third group
				(No correction)
		<hr/>		
		0100	0000 0000	Result