Language Design Issues

**Syntax:** how programs look

- Names and reserved words
- Instruction formats
- Grouping

**Semantics:** what programs mean

- Model of computation: sequential, concurrent
- Control and data flow
- Types and data representation
C History

Developed between 1969 and 1973 along with Unix

Due mostly to Dennis Ritchie

Designed for systems programming
  - Operating systems
  - Utility programs
  - Compilers
  - Filters

Evolved from B, which evolved from BCPL
BCPL

Martin Richards, Cambridge, 1967

Typeless

• Everything a machine word (n-bit integer)
• Pointers (addresses) and integers identical

Memory: undifferentiated array of words

Natural model for word-addressed machines

Local variables depend on frame-pointer-relative addressing: no dynamically-sized automatic objects

Strings awkward: Routines expand and pack bytes to/from word arrays
C History

Original machine (DEC PDP-11) was very small:
24K bytes of memory, 12K used for operating system
Written when computers were big, capital equipment
Group would get one, develop new language, OS
C History

Many language features designed to reduce memory

- Forward declarations required for everything
- Designed to work in one pass: must know everything
- No function nesting

PDP-11 was byte-addressed

- Now standard
- Meant BCPL’s word-based model was insufficient
Euclid’s Algorithm in C

```c
int gcd(int m, int n ) {
    int r;
    while (((r = m % n) != 0) {
        m = n;
        n = r;
    }
    return n;
}
```

“New syle” function declaration lists number and type of arguments. Originally only listed return type. Generated code did not care how many arguments were actually passed, and everything was a word. Arguments are call-by-value.
Euclid’s Algorithm in C

```c
int gcd(int m, int n )
{
    int r;
    while ((r = m % n) != 0) {
        m = n;
        n = r;
    }
    return n;
}
```

Automatic variable
Allocated on stack
when function
entered, released
on return
Parameters &
automatic variables
accessed via frame
pointer
Other temporaries
also stacked
Euclid on the PDP-11

.globl _gcd
.text
_gcd:
    jsr r5, rsave
L2: mov 4(r5), r1
    sxt r0
    div 6(r5), r0
    mov r1, -10(r5)
    jeq L3
    mov 6(r5), 4(r5)
    mov -10(r5), 6(r5)
    jbr L2
L3: mov 6(r5), r0
    jbr L1
L1: jmp rretrn

GPRs: r0–r7
r7=PC, r6=SP, r5=FP
Save SP in FP
r1 = n
sign extend
r0, r1 = m ÷ n
r = r1 (m % n)
if r == 0 goto L3
m = n
n = r
non-optimizing compiler
return r0 (n)
Euclid on the PDP-11

.globl _gcd
.text
_gcd:
  jsr r5, rsave
L2:  mov 4(r5), r1
     sxt r0
     div 6(r5), r0
     mov r1, -10(r5)
     jeq L3
     mov 6(r5), 4(r5)
     mov -10(r5), 6(r5)
     jbr L2
L3:  mov 6(r5), r0
     jbr L1
     jbr L1
L1:  jmp rretrn

Very natural mapping from C into PDP-11 instructions.
Complex addressing modes make frame-pointer-relative accesses easy.
Another idiosyncrasy: registers were memory-mapped, so taking address of a variable in a register is straightforward.
The Design of C

Taken from Dennis Ritchie’s *C Reference Manual* (Appendix A of Kernighan & Ritchie)
Lexical Conventions

Identifiers (words, e.g., foo, printf)

Sequence of letters, digits, and underscores, starting with
a letter or underscore

Keywords (special words, e.g., if, return)

C has fairly few: only 23 keywords. Deliberate: leaves
more room for users’ names

Comments (between / * and */)

Most fall into two basic styles: start/end sequences as in
C, or until end-of-line as in Java’s //
Lexical Conventions

C is a *free-form* language where whitespace mostly serves to separate tokens. Which of these are the same?

```
1+2  return this
1 + 2  returnthis
foo bar
foobar
```

Space is significant in some language. Python uses indentation for grouping, thus these are different:

```
if x < 3:
    y = 2
    z = 3
if x < 3:
    y = 2
    z = 3
```
Constants/Literals

Integers (e.g., 10)

Should a leading – be part of an integer or not?

Characters (e.g., ‘a’)

How do you represent non-printable or ‘ characters?

Floating-point numbers (e.g., 3.5e–10)

Usually fairly complex syntax, easy to get wrong.

Strings (e.g., "Hello")

How do you include a " in a string?
# What’s in a Name?

In C, each name has a **storage class** (where it is) and a **type** (what it is).

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<tr>
<th>Storage classes:</th>
<th>Fundamental types:</th>
<th>Derived types:</th>
</tr>
</thead>
<tbody>
<tr>
<td>automatic</td>
<td>char</td>
<td>arrays</td>
</tr>
<tr>
<td>static</td>
<td>int</td>
<td>functions</td>
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<td>float</td>
<td>pointers</td>
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<tr>
<td>register</td>
<td>double</td>
<td>structures</td>
</tr>
</tbody>
</table>
Objects and lvalues

Object: area of memory

lvalue: refers to an object

An lvalue may appear on the left side of an assignment

```c
a = 3; /* OK: a is an lvalue */
3 = a; /* 3 is not an lvalue */
```
Conversions

C defines certain automatic conversions:

- A `char` can be used as an `int`
- `int` and `char` may be converted to `float` or `double` and back. Result is undefined if it could overflow.
- Adding an integer to a pointer gives a pointer
- Subtracting two pointers to objects of the same type produces an integer
Expressions

Expressions are built from identifiers (foo), constants (3), parenthesis, and unary and binary operators.

Each operator has a **precedence** and an **associativity**

**Precedence tells us**

\[
1 \times 2 + 3 \times 4 \text{ means } (1 \times 2) + (3 \times 4)
\]

**Associativity tells us**

\[
1 + 2 + 3 + 4 \text{ means } ((1 + 2) + 3) + 4
\]
# C's Operators in Precedence Order

<table>
<thead>
<tr>
<th>Operator</th>
<th>Description</th>
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<td><code>f(r, r, ...)</code></td>
<td>Function call</td>
</tr>
<tr>
<td><code>a[i]</code></td>
<td>Array element access</td>
</tr>
<tr>
<td><code>p-&gt;m</code></td>
<td>Pointer dereference</td>
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<tr>
<td><code>s.m</code></td>
<td>Structure member access</td>
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<tr>
<td><code>!b</code></td>
<td>Logical negation</td>
</tr>
<tr>
<td><code>~i</code></td>
<td>Bitwise negation</td>
</tr>
<tr>
<td><code>-i</code></td>
<td>Unary negation</td>
</tr>
<tr>
<td><code>++l</code></td>
<td>Pre-increment</td>
</tr>
<tr>
<td><code>--l</code></td>
<td>Pre-decrement</td>
</tr>
<tr>
<td><code>l++</code></td>
<td>Post-increment</td>
</tr>
<tr>
<td><code>l--</code></td>
<td>Post-decrement</td>
</tr>
<tr>
<td><code>*p</code></td>
<td>Dereference</td>
</tr>
<tr>
<td><code>&amp;l</code></td>
<td>Address of</td>
</tr>
<tr>
<td><code>(type) r</code></td>
<td>Type cast</td>
</tr>
<tr>
<td><code>sizeof(t)</code></td>
<td>Size of</td>
</tr>
<tr>
<td><code>n * o</code></td>
<td>Multiplication</td>
</tr>
<tr>
<td><code>n / o</code></td>
<td>Division</td>
</tr>
<tr>
<td><code>i % j</code></td>
<td>Modulus</td>
</tr>
<tr>
<td><code>n + o</code></td>
<td>Addition</td>
</tr>
<tr>
<td><code>n - o</code></td>
<td>Subtraction</td>
</tr>
<tr>
<td><code>i &lt;&lt; j</code></td>
<td>Bitwise left shift</td>
</tr>
<tr>
<td><code>i &gt;&gt; j</code></td>
<td>Bitwise right shift</td>
</tr>
<tr>
<td><code>n &lt; o</code></td>
<td>Less than</td>
</tr>
<tr>
<td><code>n &gt; o</code></td>
<td>Greater than</td>
</tr>
<tr>
<td><code>n &lt;= o</code></td>
<td>Less than or equal to</td>
</tr>
<tr>
<td><code>n &gt;= o</code></td>
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</tr>
<tr>
<td><code>r == r</code></td>
<td>Equality</td>
</tr>
<tr>
<td><code>r != r</code></td>
<td>Inequality</td>
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<tr>
<td><code>i &amp; j</code></td>
<td>Bitwise AND</td>
</tr>
<tr>
<td><code>i ^ j</code></td>
<td>Bitwise XOR</td>
</tr>
<tr>
<td>`i</td>
<td>j`</td>
</tr>
<tr>
<td><code>b &amp;&amp; c</code></td>
<td>Logical AND</td>
</tr>
<tr>
<td>`b</td>
<td></td>
</tr>
<tr>
<td><code>b ? r : r</code></td>
<td>Conditional operator</td>
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<tr>
<td><code>l = r</code></td>
<td>Assignment</td>
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<tr>
<td><code>l += n</code></td>
<td>Additive assignment</td>
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<tr>
<td><code>l -= n</code></td>
<td>Subtractive assignment</td>
</tr>
<tr>
<td><code>l *= n</code></td>
<td>Multiplicative assign</td>
</tr>
<tr>
<td><code>l /= n</code></td>
<td>Divisive assignment</td>
</tr>
<tr>
<td><code>l %= i</code></td>
<td>Modulo assignment</td>
</tr>
<tr>
<td><code>l &amp;= i</code></td>
<td>Bitwise AND assignment</td>
</tr>
<tr>
<td>`l</td>
<td>=- i`</td>
</tr>
<tr>
<td><code>l &gt;&gt;= i</code></td>
<td>Bitwise right shift assignment</td>
</tr>
<tr>
<td><code>r1, r2</code></td>
<td>Comma operator</td>
</tr>
</tbody>
</table>
Declarators

Declaration: string of specifiers followed by a declarator

\[
\text{static unsigned int } (\star \text{f[10]}) (\text{int, char*});
\]

Declarator’s notation matches that of an expression: use it to return the basic type.

Largely regarded as the worst syntactic aspect of C: both pre- (pointers) and post-fix operators (arrays, functions).
Storage-Class Specifiers

**auto** Automatic (stacked), default

**static** Statically allocated

**extern** Look for a declaration elsewhere

**register** Kept in a register, not memory

C trivia: Originally, a function could only have at most three `register` variables, may only be `int` or `char`, can’t use address-of operator `&`.

Today, `register` simply ignored. Compilers try to put most automatic variables in registers.
Type Specifiers

```c
int
cchar
float
double
struct { declarations }
struct identifier { declarations }
struct identifier
```
Declarators

identifier

( declarator )

Grouping

declarator ()

Function

declarator [ optional-constant ]

Array

* declarator

Pointer

C trivia: Originally, number and type of arguments to a function wasn’t part of its type, thus declarator just contained ().

Today, ANSI C allows function and argument types, making an even bigger mess of declarators.
Declarator syntax

Is `int *f()` a pointer to a function returning an `int`, or a function that returns a pointer to an `int`?

Hint: precedence rules for declarators match those for expressions.

Parentheses resolve such ambiguities:

- `int *(f())` Function returning pointer to `int`
- `int (*f)()` Pointer to function returning `int`
Statements

expression ;
{ statement-list }
if ( expression ) statement else statement
while ( expression ) statement
do statement while ( expression );
for ( expression ; expression ; expression ) statement
switch ( expression ) statement
case constant-expression :
default :
break;
continue;
return expression ;
goto label ;
label :
External Definitions

“A C program consists of a sequence of external definitions”

Functions, simple variables, and arrays may be defined.

“An external definition declares an identifier to have storage class extern and a specified type”
Function definitions

type-specifier declarator ( parameter-list )
type-decl-list
{
  declaration-list
  statement-list
}

Example:
int max(a, b, c)
int a, b, c;
{
  int m;
  m = (a > b) ? a : b ;
  return m > c ? m : c ;
}
More C trivia

The first C compilers did not check the number and type of function arguments.

The biggest change made when C was standardized was to require the type of function arguments to be defined:

Old-style

```c
int f();
int f(a, b, c)
int a, b;
double c;
{
}
```

New-style

```c
int f(int, int, double);
int f(int a, int b, double c)
int a, b;
{
}
```
Data Definitions

type-specifier init-declarator-list ;
declarator optional-initializer

Initializers may be constants or brace-enclosed, comma-separated constant expressions. Examples:

```
int a;

struct { int x; int y; } b = { 1, 2 };

float a, *b, c;
```
Scope Rules

Two types of scope in C:

1. Lexical scope
   Essentially, place where you don’t get “undeclared identifier” errors

2. Scope of external identifiers
   When two identifiers in different files refer to the same object. E.g., a function defined in one file called from another.
Lexical Scope

Extends from declaration to terminating } or end-of-file.

```c
int a;

int foo()
{
    int b;
    if (a == 0) {
        printf("A was 0");
        a = 1;
    }
    b = a; /* OK */
}

int bar()
{
    a = 3; /* OK */
    b = 2; /* Error: b out of scope */
}
External Scope

file1.c:
int foo()
{
    return 0;
}

int bar()
{
    foo();  /* OK */
}

file2.c:
int baz()
{
    foo();  /* Error */
}

extern int foo();

int baff()
{
    foo();  /* OK */
}
The Preprocessor

Violates the free-form nature of C: preprocessor lines \textit{must} begin with \#.

Program text is passed through the preprocessor before entering the compiler proper.

Define replacement text:

\# \texttt{define} \textit{identifier} \texttt{token-string}

Replace a line with the contents of a file:

\# \texttt{include} "\textit{filename} "

## C’s Standard Libraries

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<th>Description</th>
<th>Example</th>
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<td>Generate runtime errors</td>
<td><code>assert(a &gt; 0)</code></td>
</tr>
<tr>
<td><code>&lt;ctype.h&gt;</code></td>
<td>Character classes</td>
<td><code>isalpha(c)</code></td>
</tr>
<tr>
<td><code>&lt;errno.h&gt;</code></td>
<td>System error numbers</td>
<td><code>errno</code></td>
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<td><code>&lt;float.h&gt;</code></td>
<td>Floating-point constants</td>
<td><code>FLT_MAX</code></td>
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<tr>
<td><code>&lt;limits.h&gt;</code></td>
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<td><code>&lt;locale.h&gt;</code></td>
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<td>Math functions</td>
<td><code>sin(x)</code></td>
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<td>Signal handling</td>
<td><code>signal(SIGINT,&amp;f)</code></td>
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<tr>
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<td>Variable-length arguments</td>
<td><code>va_start(ap, st)</code></td>
</tr>
<tr>
<td><code>&lt;stddef.h&gt;</code></td>
<td>Some standard types</td>
<td><code>size_t</code></td>
</tr>
<tr>
<td><code>&lt;stdio.h&gt;</code></td>
<td>File I/O, printing</td>
<td><code>printf(&quot;%d&quot;, i)</code></td>
</tr>
<tr>
<td><code>&lt;stdlib.h&gt;</code></td>
<td>Miscellaneous functions</td>
<td><code>malloc(1024)</code></td>
</tr>
<tr>
<td><code>&lt;string.h&gt;</code></td>
<td>String manipulation</td>
<td><code>strcmp(s1, s2)</code></td>
</tr>
<tr>
<td><code>&lt;time.h&gt;</code></td>
<td>Time, date calculations</td>
<td><code>localtime(tm)</code></td>
</tr>
</tbody>
</table>
Language design

Language design is library design.
— Bjarne Stroustrup

Programs consist of pieces connected together.

Big challenge in language design: making it easy to put pieces together correctly. C examples:

- The function abstraction (local variables, etc.)
- Type checking of function arguments
- The \texttt{#include} directive