Goals

Function is correct
Source code is concise, readable, maintainable
Time-critical sections of program run fast enough
Object code is small and efficient

Basically, optimize the use of three resources:

- Execution time
- Memory
- Development/maintenance time
You can say the same thing many different ways and mean the same thing.

There are many different ways to say the same thing.

The same thing may be said different ways.

There is more than one way to say it.

Many sentences are equivalent.

Be succinct.
Arithmetic

Integer Arithmetic

Fastest

Floating-point arithmetic in hardware

Slower

Floating-point arithmetic in software

Very slow

+, −

×

÷

sqrt, sin, log, etc.

slower
for (i = 0 ; i < 10000 ; ++i) 
   /* arithmetic operation */

On my desktop Pentium 4 with good hardware floating-point support,

<table>
<thead>
<tr>
<th>Operator</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>+ (int)</td>
<td>1</td>
</tr>
<tr>
<td>* (int)</td>
<td>5</td>
</tr>
<tr>
<td>/ (int)</td>
<td>12</td>
</tr>
<tr>
<td>&lt;&lt; (int)</td>
<td>2</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Operator</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>+ (double)</td>
<td>5</td>
</tr>
<tr>
<td>* (double)</td>
<td>5</td>
</tr>
<tr>
<td>/ (double)</td>
<td>10</td>
</tr>
<tr>
<td>sqrt</td>
<td>28</td>
</tr>
<tr>
<td>sin</td>
<td>48</td>
</tr>
<tr>
<td>pow</td>
<td>275</td>
</tr>
</tbody>
</table>
**Simple benchmarks**

On my Zaurus SL 5600, a 400 MHz Intel PXA250 Xscale (ARM) processor:

<table>
<thead>
<tr>
<th>Operator</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>+ (int)</td>
<td>1</td>
</tr>
<tr>
<td>* (int)</td>
<td>1</td>
</tr>
<tr>
<td>/ (int)</td>
<td>7</td>
</tr>
<tr>
<td>&lt;&lt; (int)</td>
<td>1</td>
</tr>
<tr>
<td>+ (double)</td>
<td>140</td>
</tr>
<tr>
<td>* (double)</td>
<td>110</td>
</tr>
<tr>
<td>/ (double)</td>
<td>220</td>
</tr>
<tr>
<td>sqrt</td>
<td>500</td>
</tr>
<tr>
<td>sin</td>
<td>3300</td>
</tr>
<tr>
<td>pow</td>
<td>820</td>
</tr>
</tbody>
</table>
Operations on `char`, `short`, `int`, and `long` probably run at the same speed (same ALU).

Same for unsigned variants

`int` or `long` slower when they exceed machine’s word size.

Operations on `float`s performed in double precision. `float` only useful for reducing memory.
Arithmetic Lessons

Try to use integer addition/subtraction
Avoid multiplication unless you have hardware
Avoid division
Avoid floating-point, unless you have hardware
Really avoid math library functions
C has many bit-manipulation operators.

- Bit-wise AND
- Bit-wise OR
- Bit-wise XOR
- Negate (one’s complement)
- Right-shift
- Left-shift

Plus assignment versions of each.
Bit-manipulation basics

\[
\begin{align*}
a &= 0x4; & \text{ /* Set bit 2 */} \\
b &= \sim 0x4; & \text{ /* Clear bit 2 */} \\
c &= \sim (1 \ll 3); & \text{ /* Clear bit 3 */} \\
d &= (1 \ll 5); & \text{ /* Toggle bit 5 */} \\
e &= 2; & \text{ /* Divide e by 4 */}
\end{align*}
\]
/* Set b to the rightmost 1 in a */
b = a & (a ^ (a - 1));

/* Set d to the number of 1’s in c */
char c, d;
d = (c & 0x55) + ((c & 0xaa) >> 1);
d = (d & 0x33) + ((d & 0xcc) >> 2);
d = (d & 0x0f) + ((d & 0xf0) >> 4);
Faking Multiplication

Addition, subtraction, and shifting are fast. Can sometimes supplant multiplication. Like floating-point, not all processors have a dedicated hardware multiplier. Recall the multiplication algorithm from elementary school, but think binary:

\[
\begin{array}{c}
101011 \\
\times \ 1101 \\
\hline
101011 \\
10101100 \\
+101011000 \\
\hline
1000101111
\end{array}
\]

\[= 43 + 43 \ll 2 + 43 \ll 3 = 559\]
Faking Multiplication

Even more clever if you include subtraction:

\[
\begin{array}{c}
101011 \\
\times 1110 \\
\hline
1010110 \\
10101100 \\
+101011000 \\
\hline
1001011010
\end{array}
\]

\[
= 43 \ll 1 + 43 \ll 2 + 43 \ll 3
= 43 \ll 4 - 43 \ll 2
= 602
\]

Only useful

- for multiplication by a constant
- for “simple” multiplicands
- when hardware multiplier not available
Division is a much more complicated algorithm that generally involves decisions. However, division by a power of two is just a shift:

\[ a \div 2 = a \gg 1 \]
\[ a \div 4 = a \gg 2 \]
\[ a \div 8 = a \gg 3 \]

There is no general shift-and-add replacement for division, but sometimes you can turn it into multiplication:

\[ a \div 1.33333333 = a \times 0.75 \]
\[ = a \times 0.5 + a \times 0.25 \]
\[ = a \gg 1 + a \gg 2 \]
Multi-way branches

```c
if (a == 1)
    foo();
else if (a == 2)
    bar();
else if (a == 3)
    baz();
else if (a == 4)
    qux();
else if (a == 5)
    quux();
else if (a == 6)
    corge();

switch (a) {
    case 1:
        foo(); break;
    case 2:
        bar(); break;
    case 3:
        baz(); break;
    case 4:
        qux(); break;
    case 5:
        quux(); break;
    case 6:
        corge(); break;
}
```
Nios code for if-then-else

```assembly
ldw r2, 0(fp)  # Fetch a from stack
cmpnei r2, r2, 1  # Compare with 1
bne r2, zero, .L2  # If not 1, jump to L2
call foo  # Call foo()
br .L3  # branch out

.L2:
ldw r2, 0(fp)  # Fetch a from stack (again!)
cmpnei r2, r2, 2  # Compare with 2
bne r2, zero, .L4  # If not 1, jump to L4
call bar  # Call bar()
br .L3  # branch out

.L4:
```
Nios code for switch (1)

```assembly
ldw    r2, 0(fp)       # Fetch a
cmpgeui r2, r2, 7     # Compare with 7
bne    r2, zero, .L2  # Branch if greater or equal
ldw    r2, 0(fp)       # Fetch a
muli   r3, r2, 4       # Multiply by 4
movhi  r2, %hiadj(.L9) # Load address .L9
addi   r2, r2, %lo(.L9)
add    r2, r3, r2      # = a * 4 + .L9
ldw    r2, 0(r2)       # Fetch from jump table
jmp    r2              # Jump to label

.section .rodata
.align 2

.L9:
.long  .L2             # Branch table
.long  .L3
.long  .L4
.long  .L5
.long  .L6
.long  .L7
.long  .L8
```

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Nios code for switch (2)

```
.section .text
.L3:
call foo
br .L2
.L4:
call bar
br .L2
.L5:
call baz
br .L2
.L6:
call qux
br .L2
.L7:
call quux
br .L2
.L8:
call corge
.L2:
```
Computing Discrete Functions

There are many ways to compute a “random” function of one variable:

/* OK, especially for sparse domain */
if (a == 0) x = 0;
else if (a == 1) x = 4;
else if (a == 2) x = 7;
else if (a == 3) x = 2;
else if (a == 4) x = 8;
else if (a == 5) x = 9;
/* Better for large, dense domains */
switch (a) {
case 0: x = 0; break;
case 1: x = 4; break;
case 2: x = 7; break;
case 3: x = 2; break;
case 4: x = 8; break;
case 5: x = 9; break;
}

/* Best: constant-time lookup table */
int f[] = {0, 4, 7, 2, 8, 9};
x = f[a]; /* assumes 0 <= a <= 5 */
Function calls

Modern processors, especially RISC, strive to make this cheap. Arguments passed through registers. Still has noticeable overhead.

Calling, entering, and returning:

```c
int foo(int a, int b) {
    int c = bar(b, a);
    return c;
}
```
foo:
    addi sp, sp, -20  # Allocate space on stack
    stw ra, 16(sp)   # Store return address
    stw fp, 12(sp)   # Store frame pointer
    mov fp, sp       # Frame pointer is new SP
    stw r4, 0(fp)    # Save a on stack
    stw r5, 4(fp)    # Save b on stack

    ldw r4, 4(fp)    # Fetch b
    ldw r5, 0(fp)    # Fetch a
    call bar         # Call bar()
    stw r2, 8(fp)    # Store result in c

    ldw r2, 8(fp)    # Return value in r2 = c
    ldw ra, 16(sp)   # Restore return address
    ldw fp, 12(sp)   # Restore frame pointer
    addi sp, sp, 20  # Release stack space
    ret              # Return from subroutine
Code for foo() (optimized)

foo:
    addi sp, sp, -4  # Allocate stack space
    stw ra, 0(sp)   # Store return address
    mov r2, r4      # Swap arguments (r4, r5)
    mov r4, r5      # using r2 as temporary
    mov r5, r2
    call bar        # Call bar() (return in r2)
    ldw ra, 0(sp)   # Restore return address
    addi sp, sp, 4  # Release stack space
    ret             # Return from subroutine
Strength Reduction

Why multiply when you can add?

```c
struct {
    int a;
    char b;
    int c;
} foo[10];
int i;
for (i=0 ; i<10 ; ++i) {
    foo[i].a = 77;
    foo[i].b = 88;
    foo[i].c = 99;
}

struct {
    int a;
    char b;
    int c;
} *fp, *fe, foo[10];
fe = foo + 10;
for (fp = foo ; fp != fe ; ++fp) {
    fp->a = 77;
    fp->b = 88;
    fp->c = 99;
}

Good optimizing compilers do this automatically.
```
.L2:
    ldw r2, 0(fp)           # Fetch i
    cmpgei r2, r2, 10      # i >= 10?
    bne r2, zero, .L1      # exit if true
    movhi r3, %hiadj(foo)  # Get address of foo array
    addi r3, r3, %lo(foo)
    ldw r2, 0(fp)           # Fetch i
    multi r2, r2, 12       # i * 12
    add r3, r2, r3         # foo[i]
    movi r2, 77
    stw r2, 0(r3)          # foo[i].a = 77
    movhi r3, %hiadj(foo)
    addi r3, r3, %lo(foo)
    ldw r2, 0(fp)
    multi r2, r2, 12
    add r2, r2, r3         # compute &foo[i]
    addi r3, r2, 4         # offset for b field
    movi r2, 88
    stb r2, 0(r3)          # foo[i].b = 88
.L2:
   lw   r3, 0(fp)   # fp
   lw   r2, 4(fp)   # fe
   beq  r3, r2, .L1 # fp == fe?
   lw   r3, 0(fp)
   movi r2, 77
   stw  r2, 0(r3)   # fp->a = 77
   lw   r3, 0(fp)
   movi r2, 88
   stb  r2, 4(r3)   # fp->b = 88
   lw   r3, 0(fp)
   movi r2, 99
   stw  r2, 8(r3)   # fp->c = 99
   lw   r2, 0(fp)
   addi r2, r2, 12
   stw  r2, 0(fp)   # ++fp
   br   .L2
Optimized (–O2) array code

```assembly
movi r6, 77          # Load constants
movi r5, 88
movi r4, 99
movhi r2, %hiadj(foo) # Load address of array
addi r2, r2, %lo(foo)
movi r3, 10          # iteration count

.L5:
    addi r3, r3, -1    # decrement iterations
    stw r6, 0(r2)      # foo[i].a = 77
    stb r5, 4(r2)      # foo[i].b = 88
    stw r4, 8(r2)      # foo[i].c = 99
    addi r2, r2, 12    # go to next array element
    bne r3, zero, .L5  # if there are more to do
    ret
```
movhi r6, %hiadj(foo+120)  # fe = foo + 10
addi  r6, r6, %lo(foo+120)
addi  r2, r6, -120        # fp = foo
movi  r5, 77              # Constants
movi  r4, 88
movi  r3, 99
.L5:
    stw r5, 0(r2)           # fp->a = 77
    stb r4, 4(r2)           # fp->b = 88
    stw r3, 8(r2)           # fp->c = 99
    addi r2, r2, 12         # ++fp
    bne r2, r6, .L5         # fp == fe?
    ret
How Rapid is Rapid?

How much time does the following loop take?

```c
for (i = 0; i < 1024; ++i) a += b[i];
```

<table>
<thead>
<tr>
<th>Operation</th>
<th>Cycles per iteration</th>
</tr>
</thead>
<tbody>
<tr>
<td>Memory read</td>
<td>2 or 7</td>
</tr>
<tr>
<td>Addition</td>
<td>1</td>
</tr>
<tr>
<td>Loop overhead</td>
<td>( \approx 4 )</td>
</tr>
<tr>
<td><strong>Total</strong></td>
<td>6–12</td>
</tr>
</tbody>
</table>

The Nios runs at 50 MHz, one instruction per cycle, so this takes

\[
6 \cdot 1024 \cdot \frac{1}{50\text{MHz}} = 0.12\mu s \quad \text{or} \quad 12 \cdot 1024 \cdot \frac{1}{50\text{MHz}} = 0.24\mu s
\]
Double-checking

GCC generates good code with -07:

```
movhi r4, %hiadj(b)  # Load &b[0]
addi r4, r4, %lo(b)
movi r3, 1024        # Iteration count

.L5:                  # cycles
    ldw r2, 0(r4)    # Fetch b[i]     2-7
    addi r3, r3, -1 # --i             1
    addi r4, r4, 4  # next b element  1
    add  r5, r5, r2 # a += b[i]       1
    bne r3, zero, .L5 # repeat if i > 0 3
    mov r2, r5       # result
    ret
```
Features in order of increasing cost

1. Integer arithmetic
2. Pointer access
3. Simple conditionals and loops
4. Static and automatic variable access
5. Array access
6. Floating-point with hardware support
7. Switch statements
8. Function calls
9. Floating-point emulation in software
10. Malloc() and free()
11. Library functions (sin, log, printf, etc.)
12. Operating system calls (open, sbrk, etc.)
/* fixed address: visible to other files */
int global_static;
/* fixed address: only visible within file */
static int file_static;

/* parameters always stacked */
int foo(int auto_param)
{
    /* fixed address: only visible to function */
    static int func_static;
    /* stacked: only visible to function */
    int auto_i, auto_a[10];
    /* array explicitly allocated on heap */
    double *auto_d = malloc(sizeof(double)*5);

    /* return value in register or stacked */
    return auto_i;
}
Dynamic Storage Allocation

↓ free()

↓ malloc(...)

↓ free()
Dynamic Storage Allocation

Rules:

Each allocated block contiguous (no holes)
Blocks stay fixed once allocated

`malloc()`
Find an area large enough for requested block
Mark memory as allocated

`free()`
Mark the block as unallocated
Simple Dynamic Storage Allocation

- Maintaining information about free memory
  - Simplest: Linked list
- The algorithm for locating a suitable block
  - Simplest: First-fit
- The algorithm for freeing an allocated block
  - Simplest: Coalesce adjacent free blocks
Dynamic Storage Allocation

\[ \text{malloc( } \text{\textcolor{gray}{gray\textcolor{black}{\downarrow}} } \text{ )} \]

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Simple Dynamic Storage Allocation

↓ free()

↓ free()
On most processors, access to automatic (stacked) data and globals is equally fast. Automatic usually preferable since the memory is reused when function terminates.

Danger of exhausting stack space with recursive algorithms. Not used in most embedded systems. The heap (malloc) should be avoided if possible:

- Allocation/deallocation is unpredictably slow
- Danger of exhausting memory
- Danger of fragmentation

Best used sparingly in embedded systems
“Magical” memory locations that, when written or read, send or receive data from hardware.

Hardware that looks like memory to the processor, i.e., addressable, bidirectional data transfer, read and write operations.

Does not always behave like memory:

- Act of reading or writing can be a trigger (data irrelevant)
- Often read- or write-only
- Read data often different than last written
```c
#define SWITCHES \ 
   ((volatile char *) 0x1800)
#define LEDS \ 
   ((volatile char *) 0x1810)

void main() {
   for (;;) {
      *LEDS = *SWITCHES;
   }
}
```
What's With the Volatile?

```c
#define ADDRESS ((char *) 0x1800)
#define VADDRESS ((volatile char *) 0x1800)

char foo() {
    char a = *ADDRESS;
    char b = *ADDRESS;
    return a + b;
}

char bar() {
    char a = *VADDRESS;
    char b = *VADDRESS;
    return a + b;
}
```

Compiled with optimization:

```assembly
foo:
    movi  r2, 6144
    ldbu  r2, 0(r2)
    add  r2, r2, r2
    andi  r2, r2, 0xff
    ret

bar:
    movi  r3, 6144
    ldbu  r2, 0(r3)
    ldbu  r3, 0(r3)
    add  r2, r2, r3
    andi  r2, r2, 0xff
    ret
```
/* Definitions of alt_u8, etc. */
#include "alt_types.h"

/* IORD_ALTERA_AVALON... for the ‘PIO’ device */
#include "altera_avalon_pio_regs.h"

/* Auto-generated addresses for all peripherals */
#include "system.h"

int main() {
    alt_u8 sw;
    for (;;) {
        sw = IORD_ALTERA_AVALON_PIO_DATA(SWITCHES_BASE);
        IOWR_ALTERA_AVALON_PIO_DATA(LEDS_BASE, sw);
    }
}

(From the Nios II Software Developer’s Handbook)
HW/SW Communication Styles

Memory-mapped I/O puts the processor in charge: only it may initiate communication.

Typical operation:

- Check hardware conditions by reading “status registers”
- When ready, send next “command” by writing control and data registers
- Check status registers for completion, waiting if necessary

Waiting for completion: “polling”

“Are we there yet?” “No.” “Are we there yet?” “No” “Are we there yet?” “No” “Are we there yet?” “No”
Idea: have hardware initiate communication when it wants attention.

Processor responds by immediately calling an interrupt handling routine, suspending the currently-running program.
The Unix environment provides “signals,” which behave like interrupts.

```c
#include <stdio.h>
#include <signal.h>

void handleint() {
    printf("Got an INT\n");
    /* some variants require this */
    signal(SIGINT, handleint);
}

int main() {
    /* Register signal handler */
    signal(SIGINT, handleint);
    /* Do nothing forever */
    for (;;) {
    }
    return 0;
}
```
#include "system.h"
#include "altera_avalon_pio_regs.h"
#include "alt_types.h"

static void button_isr(void* context, alt_u32 id)
{
    /* Read and store the edge capture register */
    *(volatile int *) context =
        IORD_ALTERA_AVALON_PIO_EDGE_CAP(BUTTON_PIO_BASE);

    /* Write to the edge capture register to reset it */
    IOWR_ALTERA_AVALON_PIO_EDGE_CAP(BUTTON_PIO_BASE, 0);

    /* Reset interrupt capability for the Button PIO */
    IOWR_ALTERA_AVALON_PIO_IRQ_MASK(BUTTON_PIO_BASE, 0xf);
}
interrupts under Altera (2)

#include "sys/alt_irq.h"
#include "system.h"

volatile int captured_edges;

static void init_button_pio()
{
    /* Enable all 4 button interrupts. */
    IOWR_ALTERA_AVALON_PIO_IRQ_MASK(BUTTON_PIO_BASE, 0xf);

    /* Reset the edge capture register. */
    IOWR_ALTERA_AVALON_PIO_EDGE_CAP(BUTTON_PIO_BASE, 0x0);

    /* Register the ISR. */
    alt_irq_register( BUTTON_PIO_IRQ,
                     (void *) &captured_edges,
                     button_isr );
}
Debugging Skills
The Edwards Way to Debug

1. Identify undesired behavior
2. Construct linear model for desired behavior
3. Pick a point along model
4. Form desired behavior hypothesis for point
5. Test
6. Move point toward failure if point working, away otherwise
7. Repeat #4–#6 until bug is found