

Anatomy of a Small Compiler

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The Scanner

C-style numeric constants

```
Number
: ('0'..'9')+
  ( '.' ('0'..'9')* (Exponent)?
    ( ('f'|'F') { $setType(FloatConst); }
    | /* empty */
      { $setType(DoubleConst); }
    )
  | /* empty */ { $setType(Integer); }
;
```

CEC

CEC is the Columbia Esterel Compiler that my group is currently developing.

You can find the source code (well-documented C++) off the "software" link on my homepage.

Compiles the Esterel language into hardware and software.

A concurrent language: uses a concurrent control-flow graph as an intermediate representation.

Options

```
class EsterelLexer extends Lexer;

options {
  // Lookahead to distinguish, e.g., : and :=
  k = 2;
  // Handle all 8-bit characters
  charVocabulary = '\3'..\377';
  // Export these token types for tree walkers
  exportVocab = Esterel;
  // Disable checking every rule against keywords
  testLiterals = false;
}
```

C-style numeric constants contd.

```
FractionalNumber
: '.' ('0'..'9')+ (Exponent)?
  ( ('f'|'F') { $setType(FloatConst); }
  | /* empty */
    { $setType(DoubleConst); }
  )
;

protected
Exponent
: ('e'|'E') ('+'|'-')? ('0'..'9')+
;
```

Esterel Syntax

Standard free-form style:

```
module test_present2:
input A;
output B, C;

present A then
  emit B
else
  emit C
end present

end module
```

Punctuation and Identifiers

```
PERIOD : '.' ;
POUND : '#' ;
PLUS : '+' ;
DASH : '-' ;
SLASH : '/' ;
STAR : '*' ;
PARALLEL : "||" ;
/* etc. */

ID options { testLiterals = true; }
: ('a'..'z' | 'A'..'Z')
  ('a'..'z' | 'A'..'Z' | '_' | '0'..'9')*
;
```

Strings, whitespace, newlines

```
StringConstant
: '"' ( ~('"' | '\n') | ('"'! '"') )* '"'
;

Whitespace
: (' ' | '\t' | '\f')+
  { $setType(antlr::Token::SKIP); }
;

Newline
: ('\n' | "\r\n" | '\r')
  { $setType(antlr::Token::SKIP);
  newline(); }
;
```

The Parser

Options

```
class EsterelParser extends Parser;
options {
  // Lookahead
  k = 2;
  // Construct an AST during parsing
  buildAST = true;
  // Export these token types for the tree walker
  exportVocab = Esterel;
  // Create AST nodes with line numbers
  ASTLabelType = "RefLineAST";
  // Don't automatically catch every exception
  defaultErrorHandler = false;
}
```

Tokens

Extra token types; don't correspond to keywords. Used to build additional structure into the AST.

```
tokens {
  SIGS;
  VARS;
  TYPES;
  DECLS;
  TRAPS;
  SEQUENCE;
  ARGS;
  /* etc. */
}
```

File and module

```
file
: (module)+ EOF!
;

module
: "module" ^ moduleIdentifier COLON!
  declarations
  statement
(
  "end"! "module"!
| PERIOD! // Deprecated syntax
)
;
```

Declarations

```
declarations
: (interfaceDecls)*
  { #declarations = #([DECLS, "decls"],
    #declarations); }
;

interfaceDecls :
  typeDecls
| constantDecls
| functionDecls
| procedureDecls
| taskDecls
| interfacesignalDecls
| sensorDecls
| relationDecls
;
```

Various Declarations

```
typeDecls
: "type" ^ typeIdentifier
  (COMMA! typeIdentifier)* SEMICOLON!
;

constantDecls
: "constant" ^ constantDecl
  (COMMA! constantDecl)* SEMICOLON!
;
```

Expressions

```
expression
: orexpr
;

orexpr
: andexpr ("or" ^ andexpr)*
;

andexpr
: notexpr ("and" ^ notexpr)*
;

notexpr
: "not" ^ cmpexpr
| cmpexpr
;
```

Expressions

```
mulexpr
: unaryexpr
  ( (STAR ^ | SLASH ^ | "mod" ^) unaryexpr)*
;

unaryexpr
: DASH ^ unaryexpr
| LPAREN! expression RPAREN!
| QUESTION ^ signalIdentifier
| "pre" ^
  LPAREN! QUESTION ^ signalIdentifier RPAREN!
| DQUESTION ^ trapIdentifier
| functionCall
| constant
;
```

Statements in Parallel

```
statement
: sequence (PARALLEL! sequence)*
  { if (#statement &&
    #statement->getNextSibling()) {
    #statement = #([PARALLEL, "||"],
    #statement);
  }
}
;
```

Statements in Sequence

```
sequence
: atomicStatement
  (options {greedy=true;} :
   SEMICOLON! atomicStatement)*
  (SEMICOLON!)?
  { if (#sequence &&
     #sequence->getNextSibling()) {
     #sequence = #([SEQUENCE, ";"],
                  #sequence);
   }
 }
;
```

The Present (if) Statement

Two forms:

```
present S then      present
  nothing           case C do nothing
else                case D
  nothing           else pause
end                 end present

present
: "present"^
  (presentThenPart | (presentCase)+
  (elsePart)? "end"! ("present"!)?
;
```

The Present (if) Statement

```
presentThenPart
: presentEvent ("then"! statement)?
  { #presentThenPart = #([CASE,"case"],
                       presentThenPart); }
;
elsePart
: "else"^ statement
;
presentCase
: "case"! presentEvent ("do"! statement)?
  { #presentCase = #([CASE,"case"],
                    presentCase); }
;
```

The AST Classes

My AST Classes

ANTLR, by default, builds its AST out of one type of object, an AST node with numeric type, a string, a first child, and a next sibling.

It has a facility for building heterogeneous ASTs (one class per token type), but I chose not to use it.

Instead, I created a new set of AST classes and translated the homogeneous AST into these classes during static semantics.

AST Classes

- Symbols (modules, signals, variables, functions)
Name and usually a type
- Symbol table
Holds symbols, points to a parent symbol table
- Expressions (literals, variables, operators)
Each has a type
- Modules (like a function declaration)
Has many symbol tables and a body
- Statement sequences and parallel groups
- Statements, one class per statement type

Example AST class

```
class Assignment : Statement {
  VariableSymbol *variable;
  Expression *value;
};
```

Example AST Classes

```
class CaseStatement : Statement {
  vector<PredicatedStatement *> cases;
  Statement *default;
};

class BodyStatement : Statement {
  Statement *body;
};

class PredicatedStatement : BodyStatement {
  Expression *predicate;
};
```

The Symbol Table Class

```
class SymbolTable : public ASTNode {
public:
  SymbolTable *parent;
  typedef map<string, Symbol*> stmap;
  stmap symbols;

  SymbolTable() : parent(NULL) {}

  bool local_contains(const string) const;
  bool contains(const string) const;
  void enter(Symbol *);
  Symbol* get(const string);
};
```

SymbolTable contains tests

```
bool SymbolTable::
local_contains(const string s) const {
    return symbols.find(s) != symbols.end();
}

bool SymbolTable::
contains(const string s) const {
    for ( const SymbolTable *st = this ; st ;
          st = st->parent )
        if (st->symbols.find(s) !=
            st->symbols.end()) return true;
    return false;
}
```

SymbolTable::enter

```
void SymbolTable::enter(Symbol *sym) {
    assert(sym);
    assert(symbols.find(sym->name) ==
           symbols.end());
    symbols.insert(
        std::make_pair(sym->name, sym)
    );
}
```

SymbolTable::get

```
Symbol* SymbolTable::get(const string s) {
    map<string, Symbol*>::const_iterator i;
    for ( SymbolTable *st = this; st ;
          st = st->parent ) {
        i = st->symbols.find(s);
        if (i != st->symbols.end()) {
            assert((*i).second);
            assert((*i).second->name == s);
            return (*i).second;
        }
    }
    assert(0);
}
```

Static Semantics

Static Semantics

Checks that every symbol is defined

Checks types (simple in Esterel)

Translates the ANTLR-generated AST into my own specialized version.

Written as a tree walker

The Tree Walker

```
class EsterelTreeParser extends TreeParser;

options {
    // Get the Esterel token types
    importVocab = Esterel;
    // Expect AST nodes with line numbers
    ASTLabelType = "RefLineAST";
}

file [Modules *ms, string filename]
: { assert(ms); }
  ( module[ms] )+
;
```

The Module Rule

```
module [Modules* modules]
: #( "module" moduleName:ID
{
    assert(modules);
    string name = moduleName->getText();
    if (modules->
        module_symbols.local_contains(name))
        throw LineError(moduleName,
            "Duplicate module " + name);
    ModuleSymbol *ms = new ModuleSymbol(name);
    Module *m = new Module(ms);
    ms->module = m;
    modules->add(m);
}
```

The notion of a Context

When you're translating, say, an expression, you need to know in which symbol table to look for symbols and other useful things.

I implemented a class called "Context" to hold this information.

Encountering a scope-generating statement creates a new context.

Translation routines pass the context to whatever they call.

Contexts are not part of the AST and are discarded after a scope closes.

Context

```
struct Context {
    Module *module;
    SymbolTable *variables;
    SymbolTable *traps;
    SymbolTable *signals;
    BuiltinTypeSymbol *boolean_type;
    BuiltinTypeSymbol *integer_type;
    BuiltinTypeSymbol *float_type;
    BuiltinTypeSymbol *double_type;
    BuiltinConstantSymbol *true_constant;
    BuiltinConstantSymbol *false_constant;
    Context(Module *m) :
        module(m), variables(m->constants),
        traps(0), signals(m->signals) {}
};
```

The Module Rule

```
Context c(m);

m->types->enter(c.boolean_type =
    new BuiltinTypeSymbol("boolean"));
m->constants->enter(c.false_constant =
    new BuiltinConstantSymbol("false", c.boolean_type, 0);
m->functions->enter(new BuiltinFunctionSymbol("and"));
/* ... */

VariableSymbol *vs =
    new VariableSymbol("tick", c.boolean_type, 0);
m->variables->enter(vs);
m->signals->enter(
    new BuiltinSignalSymbol("tick", 0,
        "input", 0, vs, 0));
```

Signal Declarations

```
signalDecl [Context *c, string direction,
    bool isGlobal]
: #( SDECL signalName:ID
  {
    string name = signalName->getText();
    if (c->signals->local_contains(name))
        throw LineError(signalName,
            "Redeclaration of " + name);
    Expression *e = 0;
  }
  ( #(COLEQUALS e=expr:expression[c] )?
  { TypeSymbol *t = 0; FunctionSymbol *fs = 0; }
```

Signal Declarations

```
{
    new_signal(c, name, t, direction, fs, e);
    if (e && (e->type != t))
        throw LineError(signalName,
            "initializer does not "
            "match type of signal");
}
);
```

The Module Rule

```
Statement *s; /* Local variable in module rule */
}

declarations[&c]
s=statement[&c] { m->body = s; }
) /* matches #("module" ... */
;
```

Signal Declarations

```
(t=typeToken:type[c]
  ( func:ID
    {
      string name = func->getText();
      if (!c->module->functions
          ->local_contains(name))
          throw LineError(func,
              "Undeclared function " + name);
      Symbol *sym = c->module->functions->get(name);
      fs = dynamic_cast<FunctionSymbol*>(sym);
      assert(fs);
    }
  )?
);
```

Signal Expressions

```
sigExpression [Context *c] returns [Expression *e]
: { Expression *e1, *e2; }
  ( #( "and" e1=sigExpression[c] e2=sigExpression[c] )
  { e = new BinaryOp(c->boolean_type,"and",e1,e2); }
  | sig:ID
  {
    string name = sig->getText();
    if (!c->signals->contains(name))
        throw LineError(sig,
            "unrecognized signal " + name);
    SignalSymbol *ss = dynamic_cast<SignalSymbol*>(
        c->signals->get(name));
    e = new LoadSignalExpression(ss);
  }
  );
```

Signal Declarations

```
input s1,
    s2 : boolean,
    s3 : combine integer with +,
    s8 := 3 : integer,
    s9 := 5 : combine integer with +;
```

Signal Declarations

```
| pcf:predefinedCombineFunction
  {
    string name = pcf->getText();
    assert(c->module->functions->contains(name));
    Symbol *sym = c->module->functions->get(name);
    fs = dynamic_cast<BuiltinFunctionSymbol*>(sym);
    assert(fs);
  }
  )?
);
```

Local Signal Statements

```
signal ls2,
    ls3 : boolean,
    ls4 := 3 + v1 : integer,
    ls5 := v3 or true :
        combine boolean with or in
    emit ls2;
    emit ls4(10);
    emit ls5(false)
end
```

Local Signal Statement

```
| #( "signal"  
  {  
    Signal *sig = new Signal();  
    Context nc = *c;  
    nc.signals = sig->symbols = new SymbolTable();  
    sig->symbols->parent = c->signals;  
  }  
  #( SIGS ( signalDecl[&nc, "local", false] )+ )  
  { Statement *s; }  
  s=statement[&nc]  
  {  
    sig->body = s;  
    st = sig;  
  }  
)
```

Dismantling

Some Statistics

File	Role	# lines
esterel.g	Parser/Scanner	850
staticsemantics.g	AST builder	1025
AST.nw	AST class source	1488
IR.nw	XML Serialization	827
Dismantle.nw	Dismantling	1571
AST.hpp*	AST classes	1828
AST.cpp*	AST classes	1525

* auto-generated

Type checking expressions

```
expression [Context *c] returns [Expression *e]  
:  
{  
  Expression *e1 = 0, *e2 = 0; // for safety  
  e = 0; // for safety  
}  
#( PLUS e1=expression[c] e2=expression[c] )  
{ e = numeric_binop(#expression,  
                    c, "+", e1, e2); }  
| #( STAR e1=expression[c] e2=expression[c] )  
{ e = numeric_binop(#expression,  
                    c, "*", e1, e2); }
```

Dismantling

Many more complicated Esterel statement are equivalent to multiple simple statements, e.g.,

```
present                                if (p1) s1  
  case p1 do s1                        else if (p2) s2  
  case p2 do s2                        else s3  
  else s3  
end
```

Type checking expressions

```
static Expression*  
numeric_binop(RefLineAST l, Context *c, string op,  
              Expression *e1, Expression *e2)  
{  
  assert(c); assert(e1); assert(e2);  
  
  if (e1->type != e2->type ||  
      !(e1->type == c->integer_type ||  
        e1->type == c->float_type ||  
        e1->type == c->double_type ))  
    throw LineError(l,  
                    "arguments of " + op + " must be numeric");  
  return new BinaryOp(e1->type, op, e1, e2);  
}
```

Dismantling Case Statements

```
IfThenElse *dismantle_case(CaseStatement &c) {  
  IfThenElse *result = 0; IfThenElse *lastif = 0;  
  for (vector<PredicatedStatement*>::iterator i =  
        c.cases.begin(); i != c.cases.end(); i++) {  
    IfThenElse *thisif =  
      new IfThenElse((*i)->predicate);  
    thisif->then_part = transform((*i)->body);  
    if (result) lastif->else_part = thisif;  
    else result = thisif;  
    lastif = thisif;  
  }  
  lastif->else_part = c.default_stmt;  
  return transform(result); // Recurse  
}
```