



A Snapshot on MPLS Reliability Features

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Outline

- ◆ ***Introduction***
- ◆ Fast Reroute
- ◆ Graceful Restart
- ◆ Summary



MPLS in a Nutshell

◆ *Tunnels*

- ❖ Drop a packet in, and out it comes at the other end

◆ *Explicit (aka source) routing*

◆ *Label stack*

- ❖ e.g., 2-label stack: "outer" label defines the tunnel; "inner" label demultiplexes

◆ *Layer 2 independence*

- ❖ Just like IP



Why tunnels...

- ◆ ***Transfer Non-IP (or private addressed IP) packets over the backbones e.g.:***
 - ❖ **Layer 3 VPN (BGP/MPLS VPN)**
 - ❖ **Layer 2 VPN (draft-kompella-ppvpn-vpn)**
 - ❖ **Virtual Private LAN Service (VPLS)**
 - ❖ ***This is potentially a huge market!***
- ◆ ***Map user traffic according to your plan.***
 - ❖ **Guarantee bandwidth to user “flows”**
 - ❖ **More efficient use of network resources**



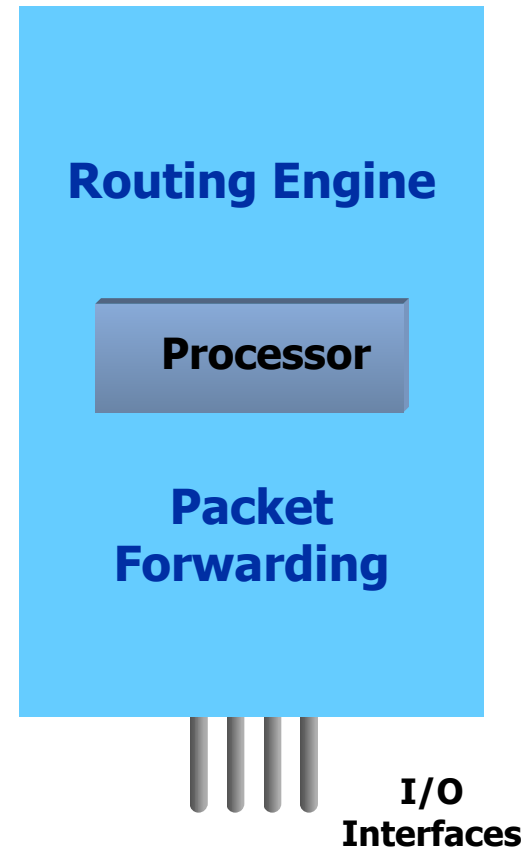
Challenges

- ◆ *What if my MPLS tunnels break...*
- ◆ *Hold on...*
 - ❖ *Let's first take a look at router's internal structure.*



Legacy Router Architecture

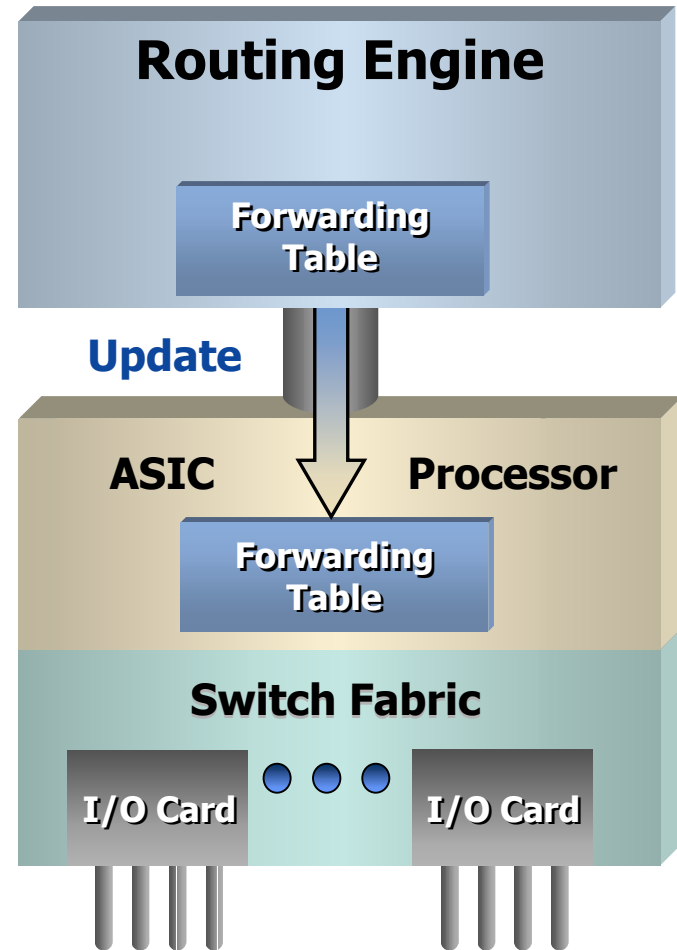
- ◆ *Data plane and control plane are **together**.*
- ◆ *If either data or control plane fails, the entire router will get effected, which, in turn, can disrupt the data traffic.*





New-generation Router Architecture

- ◆ *The **separation** of data and control planes*
- ◆ *Either data or control processor failure will not effect the entire router.*





Comparison

◆ Legacy routers:

- ❖ *Control and data plane live and die together!*

◆ New-generation routers:

- ❖ *Control and data plane can be managed separately.*

◆ *Observation:*

- ❖ *For various reasons (e.g., software upgrade, control software crash), the control plane needs to be restarted more frequent than the data plane.*



Ask me again...

◆ *What if my MPLS tunnels break...*

❖ **Link outage:**

◆ **Solution: reroute at data plane**

❖ **Control plane up/down, e.g.:**

◆ **Solution: sustain the data plane, while recovering the control plane**

◆ ***The bottom line: we need to have high availability at data plane for MPLS tunnels!***



A Snapshot on MPLS Redundancy

- ◆ ***Redundant Hardware and Software***
 - ❖ ... but what if it's the adjacent links and nodes are in trouble?
- ◆ ***Backup Tunnels from ingress***
 - ❖ ... but this may not be fast enough.
- ◆ ***Fast Reroute***
 - ❖ At data forwarding level, redirect user traffic on the fly.
- ◆ ***Graceful Restart***
 - ❖ At control plane, recover the control information on the "down" nodes without disturbing data traffic.



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Fast Reroute

- ◆ **Reroute around link or node failure... *fast***
 - ❖ **~10s of msec reroute time**
- ◆ **Reroute paths immediately available**
 - ❖ **Make-Before-Break**
- ◆ **Crank back to the node closest to the failure, not ingress router**
 - ❖ **Local repair is the key.**
- ◆ **Short term solution for traffic protection**
 - ❖ **The ingress should re-compute alternative routes eventually.**



Fast Reroute (signaling protocol)

◆ History:

- ❖ Juniper and Cisco both have working solutions.
- ❖ Due to customer demand, we merged our ideas:
 - ◆ draft-ietf-mpls-rsvp-lsp-fastreroute-00.txt

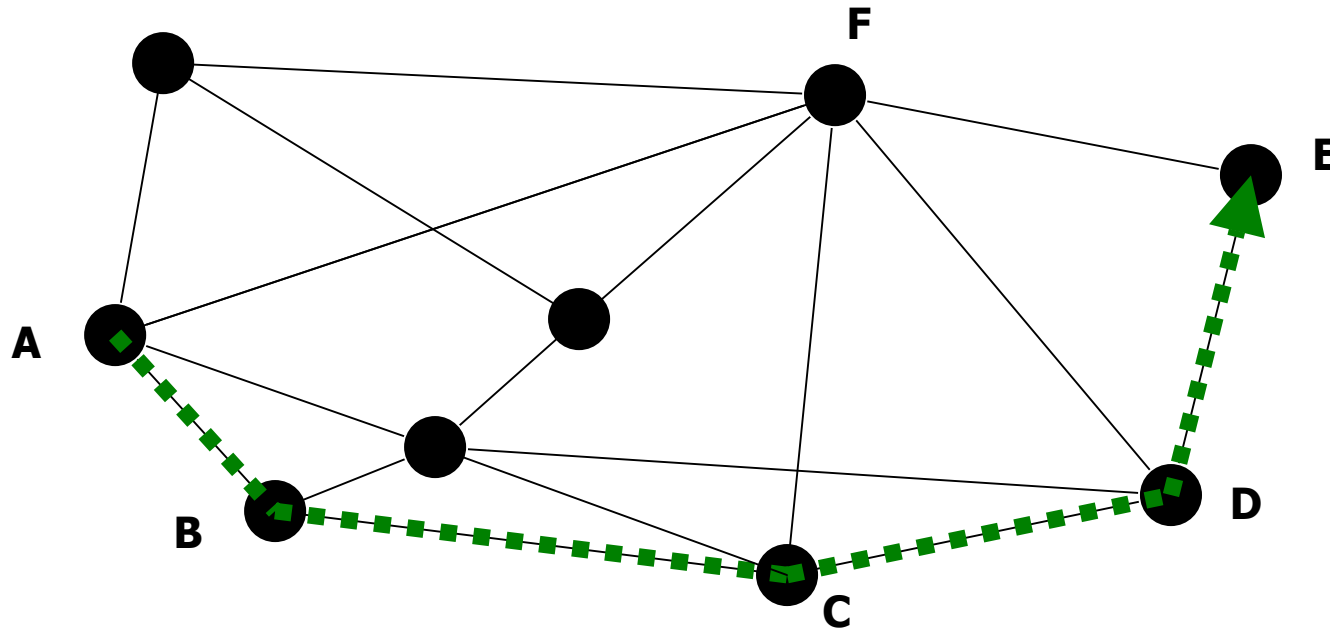
◆ RSVP Protocol Extensions:

- ❖ One-to-one backup
 - ◆ Backup each LSP separately.
 - ◆ More flexible
 - ◆ Simple to configure
- ❖ Many-to-one backup
 - ◆ Backup a bunch of LSPs with one LSP
 - ◆ Less states with label stacking
 - ◆ Requires configuring backup LSPs
- ❖ Use common set of RSVP mechanisms



One-to-one backup: example

◆ A LSP from A to E

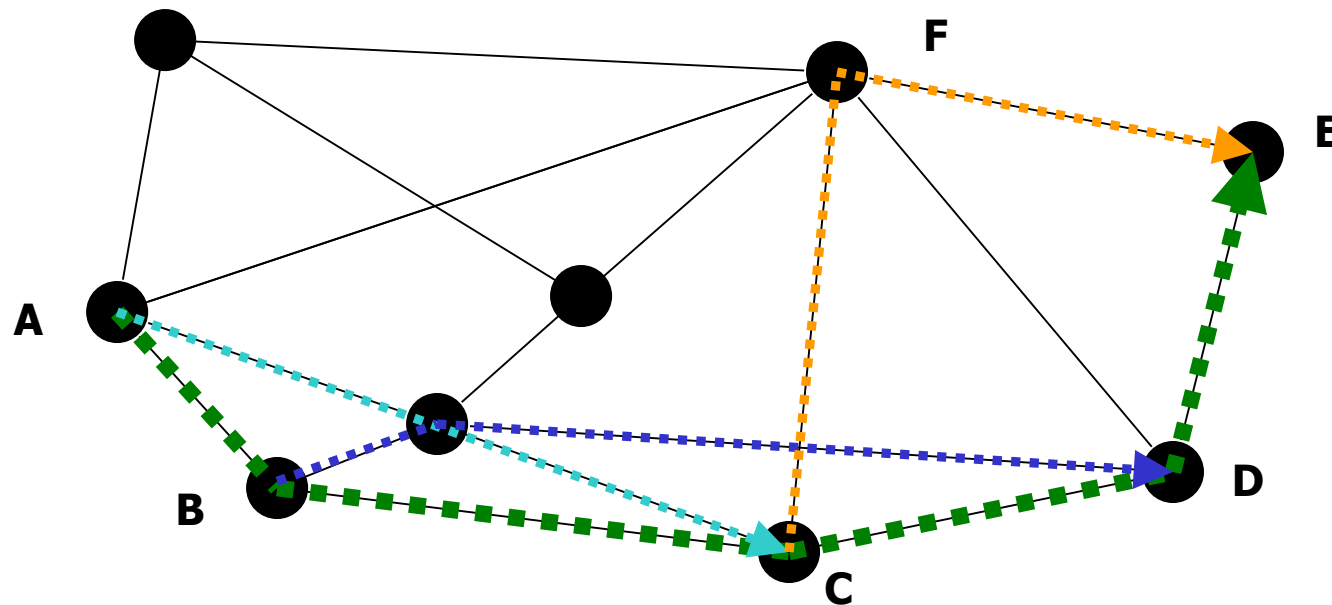




One-to-one backup: example

◆ Enable fast reroute on ingress

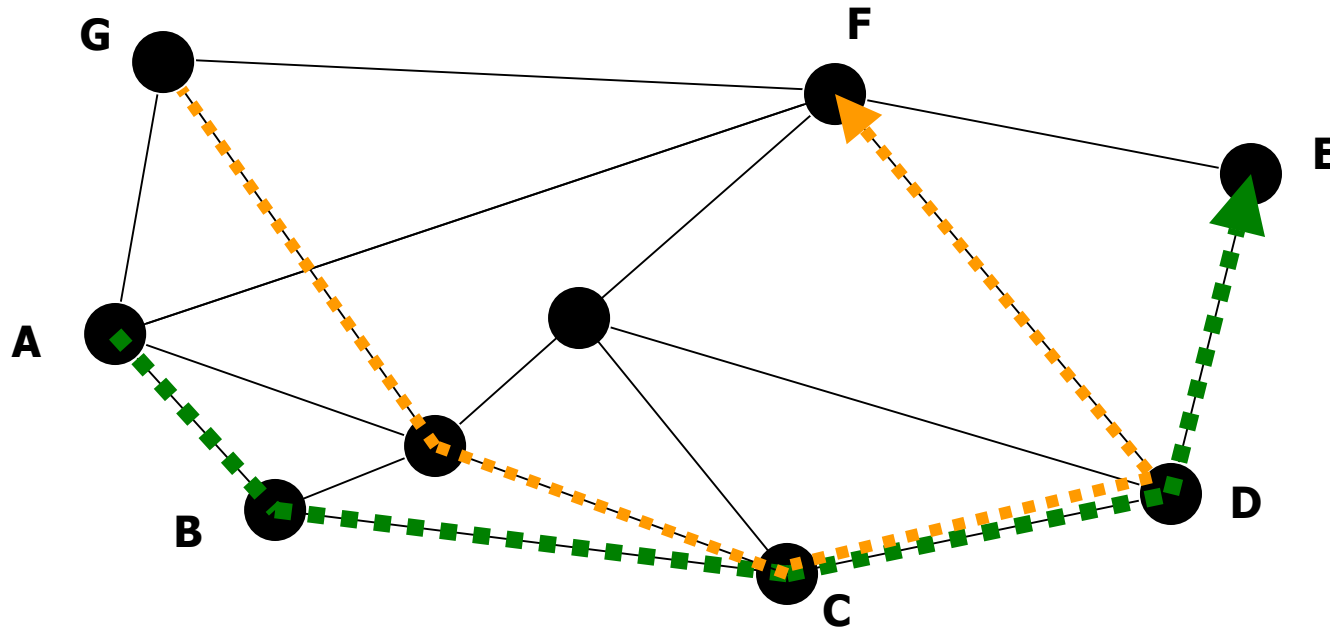
- ❖ A creates detour around B
- ❖ B creates detour around C
- ❖ C creates detour around D
- ❖ No additional configuration required on B, C, D, etc...





Many-to-one backup: example

- ◆ Two User LSPs going over link C-D.

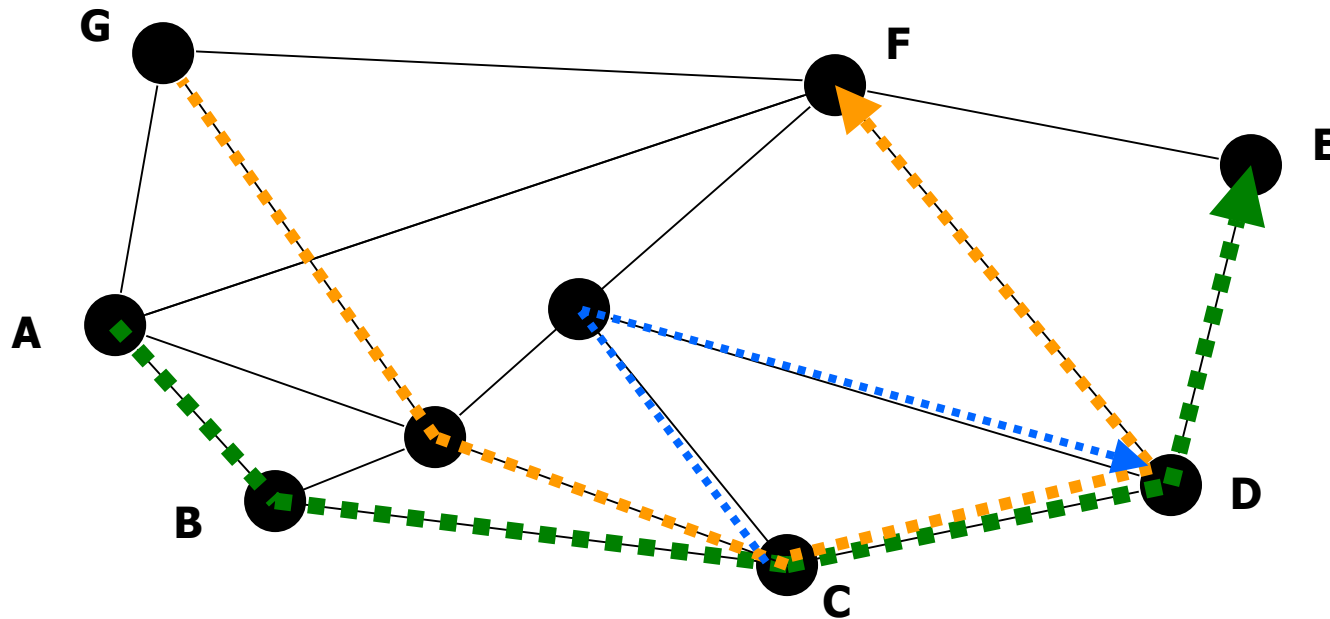




Many-to-one backup: example

◆ Enable link-protection

- ❖ Each LSP that uses link protection has to be identified as such at the ingress (via configuration)
- ❖ Requires configuration for every link that has to be protected
- ❖ **C creates a LSP that will bypass C-D.**

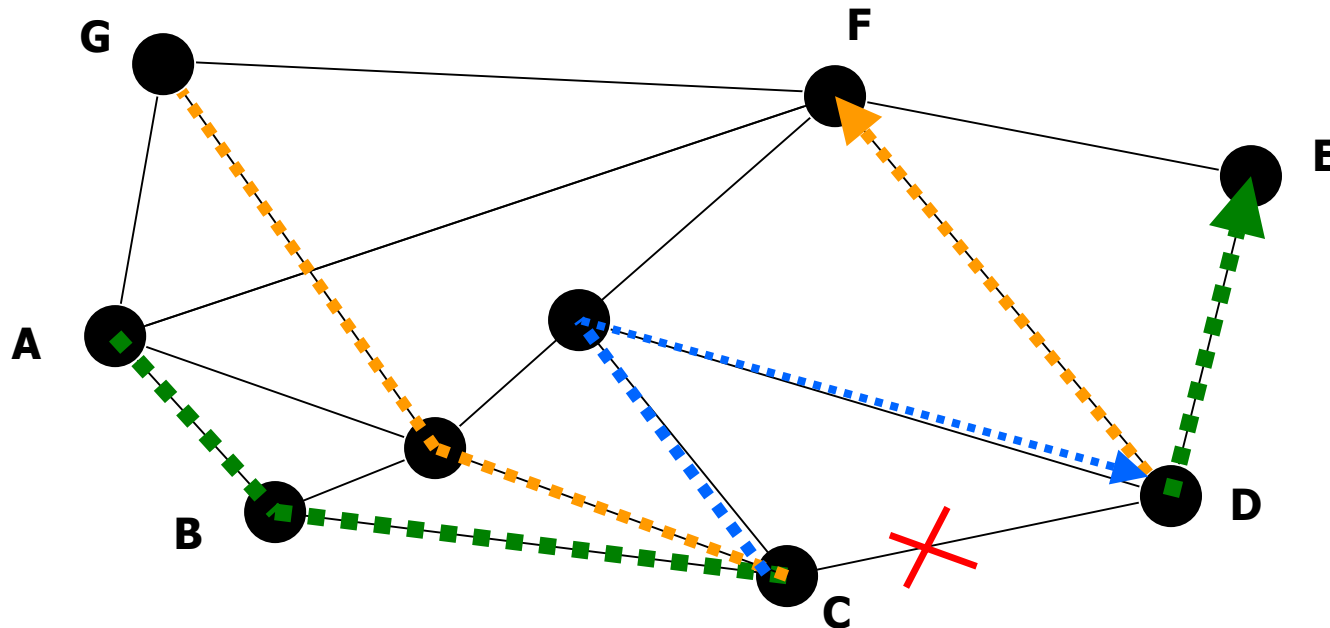




Many-to-one backup: example

◆ Link C-D fails

- ❖ C reroutes user traffic with label-stacking ("outer" label + "inner-1" or "inner-2" labels)
- ❖ C signals to A and G that failure occurred





Fast Reroute Issues

◆ Network Operation:

- ❖ **Having too many configuration parameters complicates the usage**

- ◆ **One-to-one backup: only ingress routers initiate fast reroute.**

- ◆ **Many-to-one backup: both ingress and transit routers need to configure.**

◆ Performance:

- ❖ **On Juniper routers, for both one-to-one and many-to-one backups, the data-plane reroute time after the detection of a failure:**

- ◆ **An OC12 link is protected via an OC48 link.**

- ◆ **100 packet sources, 20,000 pps, load balancing.**

- ◆ **~0 for 1 LSP**

- ◆ **~40 msec for 10 LSP's**



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Graceful Restart

- ◆ **A generic solution to**
 - ❖ **BGP**
 - ❖ **ISIS**
 - ❖ **OSPF**
 - ❖ **LDP**
 - ❖ **RSVP-TE**
 - ❖ **Various MPLS VPN solutions**
- ◆ **RSVP-TE graceful restart:**
 - ❖ `draft-ietf-mpls-generalized-rsvp-te`



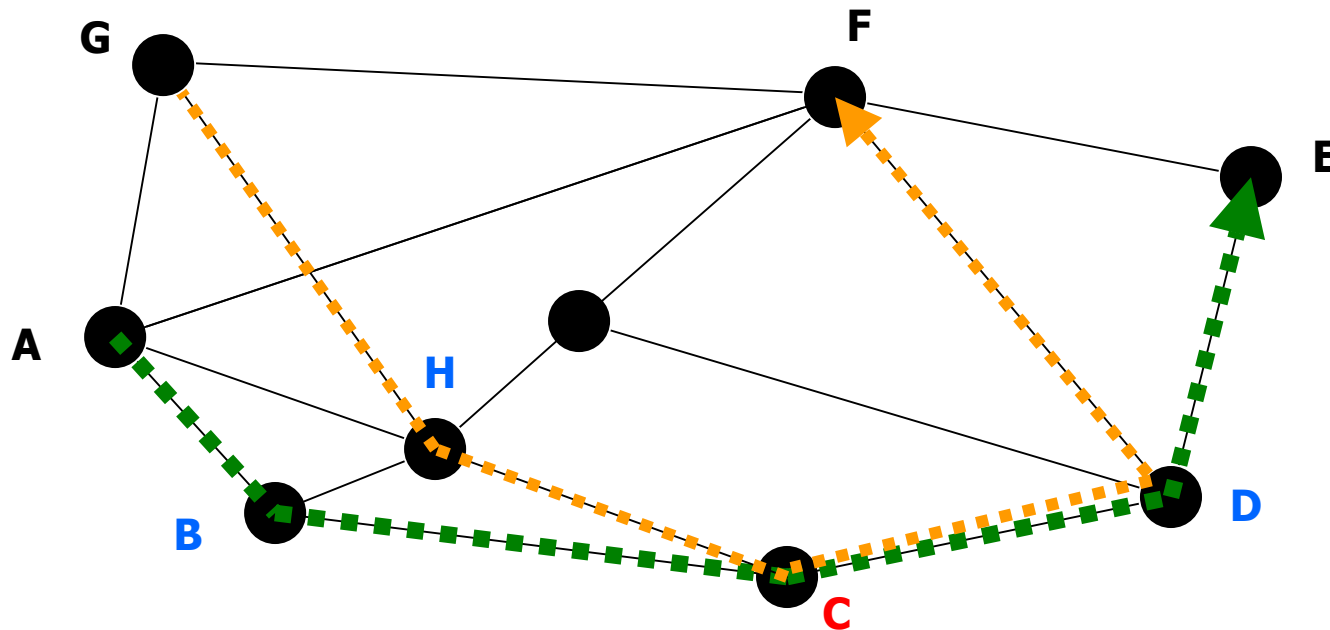
Graceful Restart...

- ◆ **Currently, while data forwarding is OK,**
 - ❖ ***IF....***
 - ◆ the router control plane restarts (due to crash or s/w upgrade)
 - ◆ the control channel between a pair of routers restarts
 - ❖ ***Then...***
 - ◆ **All LSP's traversing the router are terminated.**
 - ◆ **Major traffic disruption inside the network**
- ◆ ***With Graceful Restart,***
 - ❖ **the control plane can be recovered,**
 - ❖ **... without disturbing the data plane**
 - ◆ **no disruption to data/user traffic**



Graceful Restart (example)

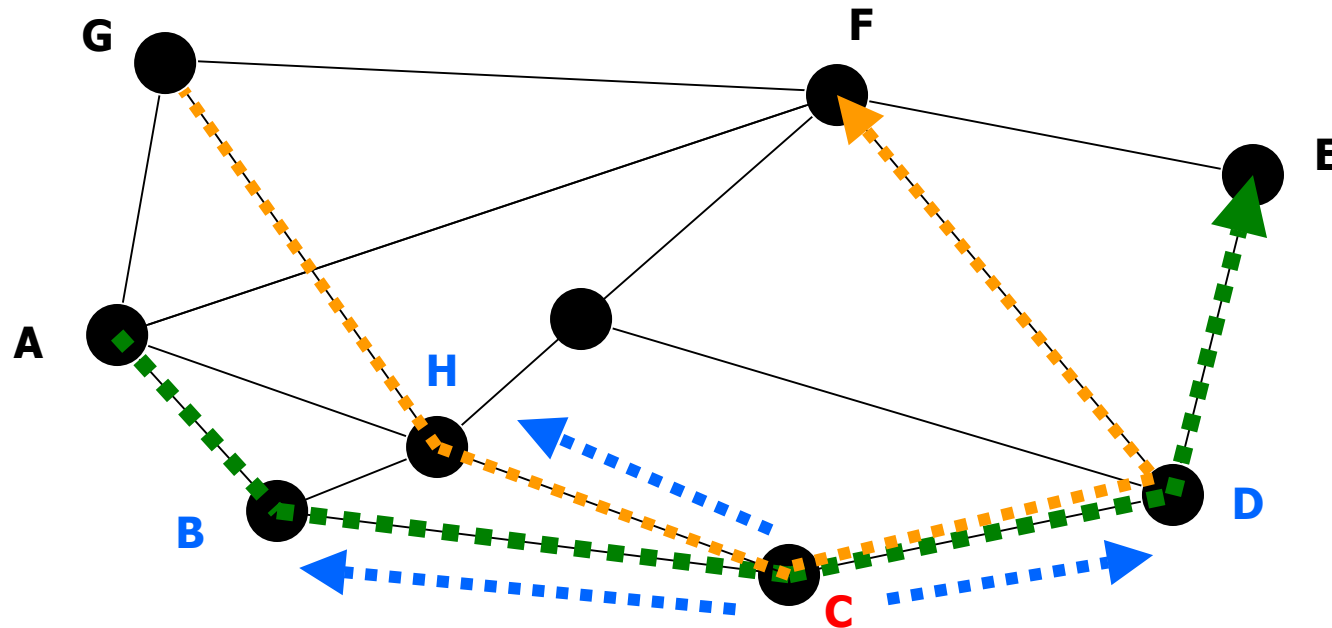
- ◆ Two LSPs going through C.
- ◆ B, D and H have the knowledge about the labels that are used for data forwarding on C.





Graceful Restart (example)

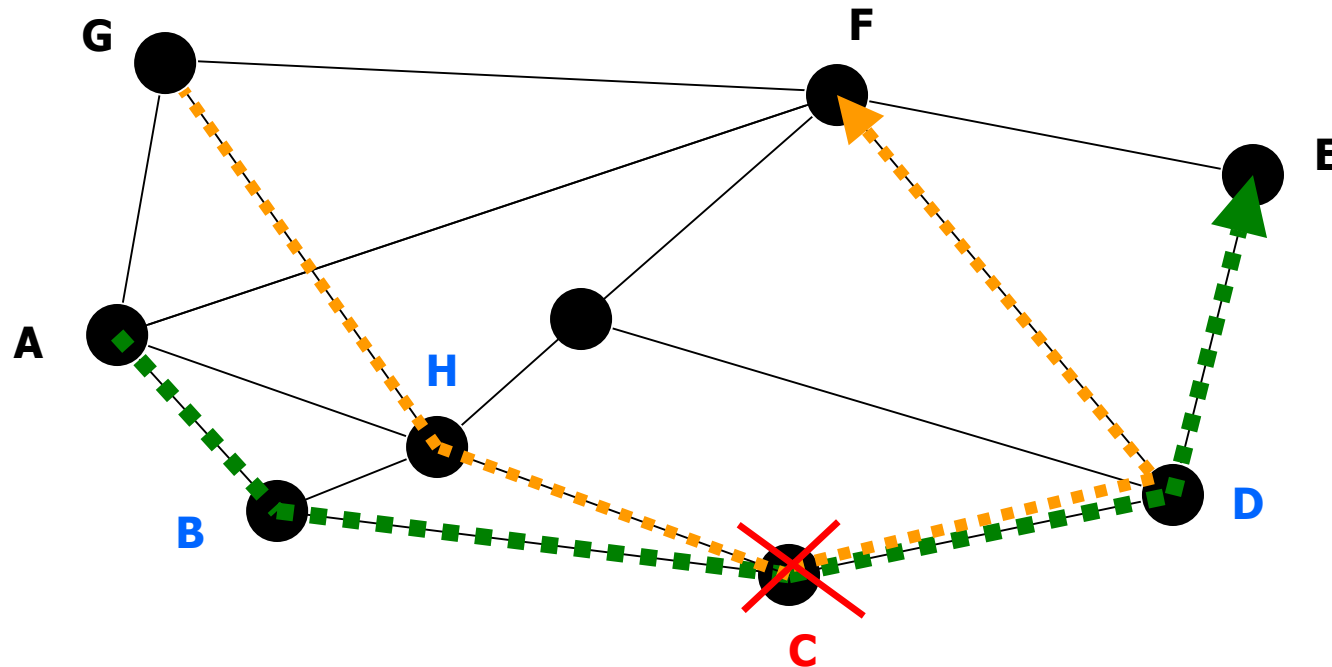
- ◆ C advertises the Graceful Restart capability to neighbors, B, H, D.





Graceful Restart (example)

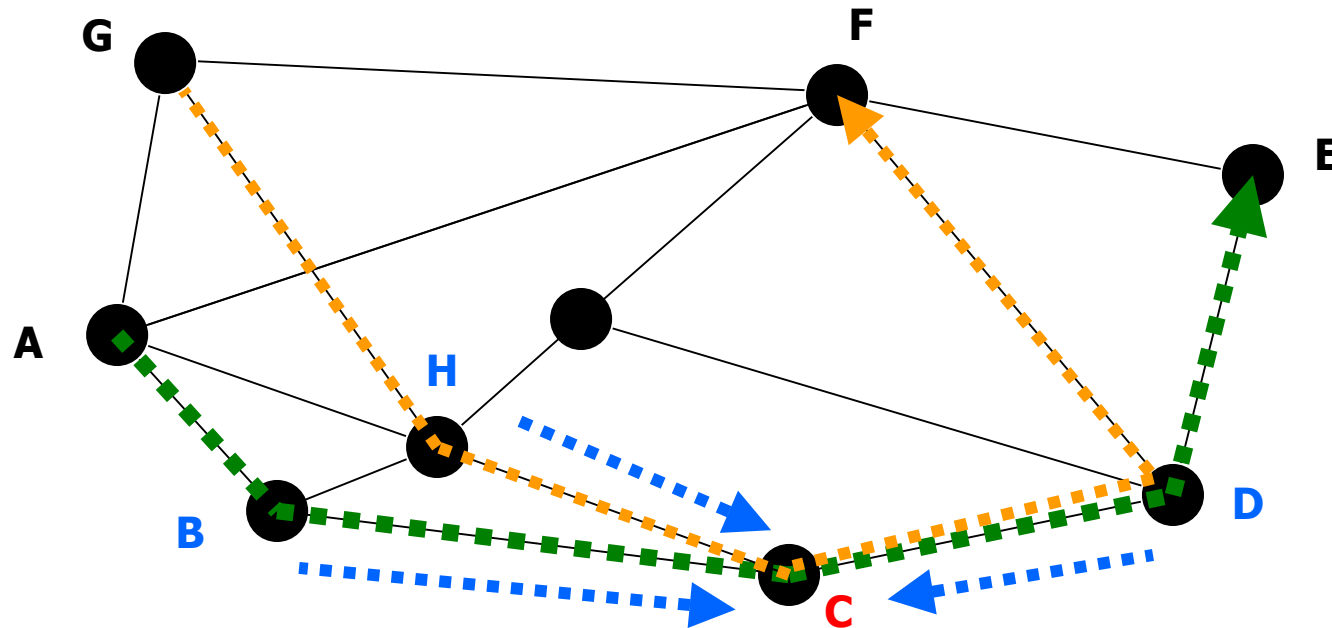
- ◆ The control plane on C has crashed.
- ◆ If data forwarding is OK, B, H and D won't over-react, and keep the LSPs intact.





Graceful Restart (example)

- ◆ After detecting C is up again, B, D and H sends labels information to C to help its recovery.





Graceful Restart Issues

- ◆ Only applicable on **new-generation** routers:
 - ❖ Requires the separation of data & control plane
- ◆ This is perceived to be especially important in the context of **GMPLS**



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Summary

- ◆ Both **Fast Reroute** and **Graceful Restart** are designed to improve data plane availability in the face of network failures.
- ◆ From our measurement, the reroute timing on MPLS **Fast Reroute** is as good as SONET APS.
- ◆ MPLS **Graceful Restart** can help to prevent traffic disruption in today's network.
- ◆ Requires **new-generation** routers.



Fast Reroute and Graceful Restart Comparison (1)

◆ Fast Reroute:

- ❖ Backup tunnels may consume network resources (e.g. bandwidth in case of SONET/SDH or OXCs).
 - ◆ Can become a serious constraint in optical networks
- ❖ Many-to-one backups rely on label-stack
 - ◆ Not available in environments such as optical networks
- ❖ Configuration can be a problem.
- ❖ Cannot protect user traffic at ingress routers
 - ◆ Works very well on transit routers.



Fast Reroute and Graceful Restart Comparison (2)

◆ Graceful Restart:

- ❖ Does **not** consume any network resource
 - ◆ Very desirable for optical networks
- ❖ Configuration is simple
 - ◆ Thanks to the capability advertisement
- ❖ Can protect ingress routers
 - ◆ As well as transit and egress routers
- ❖ Require new-generation routers

Thank you!

