

Questions for Flipped Classroom Session of COMS 4705 Week 3, Fall 2014. (Michael Collins)

Question 1 Consider a trigram HMM tagger with:

- The set \mathcal{K} of possible tags equal to $\{D, N, V\}$
- The set \mathcal{V} of possible words equal to $\{\text{the, dog, barks}\}$
- The following parameters:

$$\begin{aligned}q(D|*, *) &= 1 \\q(N|*, D) &= 1 \\q(V|D, N) &= 1 \\q(\text{STOP}|N, V) &= 1 \\e(\text{the}|D) &= 1 \\e(\text{dog}|N) &= 0.4 \\e(\text{barks}|N) &= 0.6 \\e(\text{dog}|V) &= 0.1 \\e(\text{barks}|V) &= 0.9\end{aligned}$$

with all other parameter values equal to 0.

Question: Write down the set of all pairs of sequences $x_1 \dots x_n, y_1 \dots y_{n+1}$ such that the following properties hold:

- $p(x_1 \dots x_n, y_1 \dots y_{n+1}) > 0$
- $x_i \in \mathcal{V}$ for all $i \in 1 \dots n$
- $y_i \in \mathcal{K}$ for all $i \in 1 \dots n$, and $y_{n+1} = \text{STOP}$

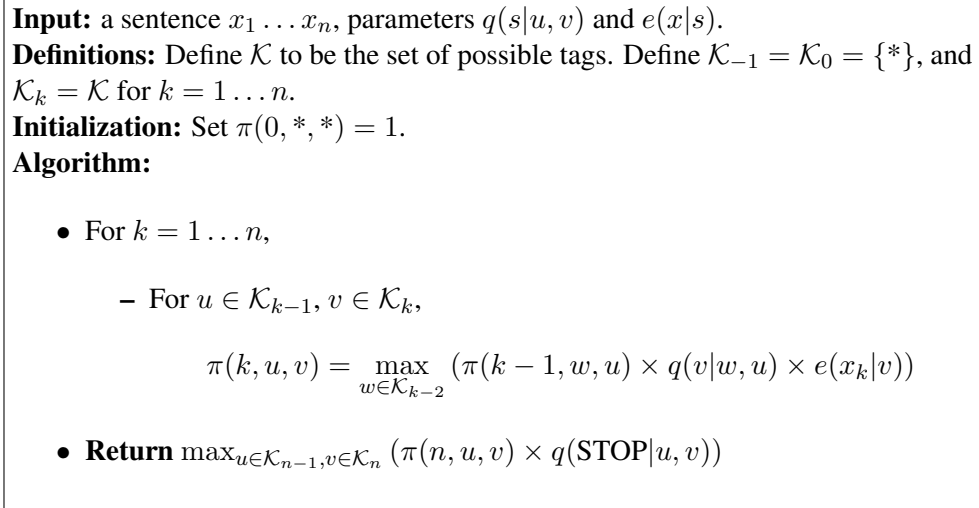


Figure 1: The basic Viterbi Algorithm.

Question 2 Consider a trigram HMM, as introduced in class. We saw that the Viterbi algorithm could be used to find

$$\max_{y_1 \dots y_{n+1}} p(x_1 \dots x_n, y_1 \dots y_{n+1})$$

where the max is taken over all sequences $y_1 \dots y_{n+1}$ such that $y_i \in \mathcal{K}$ for $i = 1 \dots n$, and $y_{n+1} = \text{STOP}$. (Recall that \mathcal{K} is the set of possible tags in the HMM.) In a trigram tagger we assume that p takes the form

$$p(x_1 \dots x_n, y_1 \dots y_{n+1}) = \prod_{i=1}^{n+1} q(y_i|y_{i-2}, y_{i-1}) \prod_{i=1}^n e(x_i|y_i) \quad (1)$$

Recall that we have assumed in this definition that $y_0 = y_{-1} = *$, and $y_{n+1} = \text{STOP}$. The Viterbi algorithm is shown in figure 1.

Now consider a four-gram tagger, where p takes the form

$$p(x_1 \dots x_n, y_1 \dots y_{n+1}) = \prod_{i=1}^{n+1} q(y_i|y_{i-3}, y_{i-2}, y_{i-1}) \prod_{i=1}^n e(x_i|y_i) \quad (2)$$

We have assumed in this definition that $y_0 = y_{-1} = y_{-2} = *$, and $y_{n+1} = \text{STOP}$.

Question: In the box below, give a version of the Viterbi algorithm that takes as input a sentence $x_1 \dots x_n$, and finds

$$\max_{y_1 \dots y_{n+1}} p(x_1 \dots x_n, y_1 \dots y_{n+1})$$

for a four-gram tagger, as defined in Eq. 4.

Input: a sentence $x_1 \dots x_n$, parameters $q(w|t, u, v)$ and $e(x|s)$.

Definitions: Define \mathcal{K} to be the set of possible tags. Define $\mathcal{K}_{-2} = \mathcal{K}_{-1} = \mathcal{K}_0 = \{*\}$, and $\mathcal{K}_k = \mathcal{K}$ for $k = 1 \dots n$.

Initialization:

Algorithm:

Return:

Question: In the box below, give a version of the Viterbi algorithm that takes as input an integer n , and finds

$$\max_{y_1 \dots y_{n+1}, x_1 \dots x_n} p(x_1 \dots x_n, y_1 \dots y_{n+1})$$

for a trigram tagger, as defined in Eq. 3. **Hence the input to the algorithm is an integer n , and the output from the algorithm is the highest scoring pair of sequences $x_1 \dots x_n, y_1 \dots y_{n+1}$ under the model.**

Input: an integer n , parameters $q(w|u, v)$ and $e(x|s)$.

Definitions: Define \mathcal{K} to be the set of possible tags. Define $\mathcal{K}_{-1} = \mathcal{K}_0 = \{*\}$, and $\mathcal{K}_k = \mathcal{K}$ for $k = 1 \dots n$. Define \mathcal{V} to be the set of possible words.

Initialization:

Algorithm:

Return:

Question 3 Consider a trigram HMM, as introduced in class. We saw that the Viterbi algorithm could be used to find

$$\max_{y_1 \dots y_{n+1}} p(x_1 \dots x_n, y_1 \dots y_{n+1})$$

where the max is taken over all sequences $y_1 \dots y_{n+1}$ such that $y_i \in \mathcal{K}$ for $i = 1 \dots n$, and $y_{n+1} = \text{STOP}$. (Recall that \mathcal{K} is the set of possible tags in the HMM.) In a trigram tagger we assume that p takes the form

$$p(x_1 \dots x_n, y_1 \dots y_{n+1}) = \prod_{i=1}^{n+1} q(y_i | y_{i-2}, y_{i-1}) \prod_{i=1}^n e(x_i | y_i) \quad (3)$$

Recall that we have assumed in this definition that $y_0 = y_{-1} = *$, and $y_{n+1} = \text{STOP}$. The Viterbi algorithm is shown in figure 1.

Now consider a “skip” tagger, where p takes the form

$$p(x_1 \dots x_n, y_1 \dots y_{n+1}) = \prod_{i=1}^{n+1} q(y_i | y_{i-2}) \prod_{i=1}^n e(x_i | y_i) \quad (4)$$

We have assumed in this definition that $y_0 = y_{-1} = *$, and $y_{n+1} = \text{STOP}$. Note that a “skip” tagger replaces the term $q(y_i | y_{i-2}, y_{i-1})$ in a regular trigram tagger with

$$q(y_i | y_{i-2})$$

We call it a skip tagger because y_{i-1} is now omitted from the conditioning information.

Question: In the box below, give a version of the Viterbi algorithm that takes as input a sentence $x_1 \dots x_n$, and finds

$$\max_{y_1 \dots y_{n+1}} p(x_1 \dots x_n, y_1 \dots y_{n+1})$$

for a skip tagger, as defined in Eq. 4. (Note: it is fine if the runtime of your algorithm is $O(n|\mathcal{K}|^3)$.)

Input: a sentence $x_1 \dots x_n$, parameters $q(w|v)$ and $e(x|s)$.

Definitions: Define \mathcal{K} to be the set of possible tags. Define $\mathcal{K}_{-1} = \mathcal{K}_0 = \{*\}$, and $\mathcal{K}_k = \mathcal{K}$ for $k = 1 \dots n$.

Initialization:

Algorithm:

Return:

Question 4 Say we have a training set consisting of two tagged sentences:

the/DT can/NN is/VB in/IN the/DT shed/NN

the/DT dog/NN can/VB see/VB the/DT cat/NN

We train a bigram tagger of the form

$$p(x_1 \dots x_n, y_1 \dots y_{n+1}) = \prod_{i=1}^{n+1} q(y_i | y_{i-1}) \prod_{i=1}^n e(x_i | y_i)$$

using simple maximum-likelihood estimates for the q and e parameters.

If we then use the Viterbi algorithm to find the maximum probability tag sequence for each of the training sentences, show that the tagger tags both sentences correctly.

Question 5 Now come up with a training set such that when we train a bigram tagger using maximum likelihood estimates, the resulting model makes at least one mistake on the training set.