

Questions for Flipped Classroom Session of COMS 4705 Week 5, Fall 2014. (Michael Collins)

Question 1 In this question our goal is to design an algorithm that takes a sentence s and a context-free grammar in Chomsky normal form as input, and as its output returns *the number of parse trees for the sentence s* as its output.

For example, if s is the sentence $a a a$, and the context-free grammar is

$X \rightarrow X X$
 $X \rightarrow a$

with start symbol X , the algorithm should return the value 2, because there are two parses for the sentence under this grammar:



Question: Complete the following algorithm so that it returns the number of possible parse trees for the input sentence s .

Input: a sentence $s = x_1 \dots x_n$, a context-free grammar $G = (N, \Sigma, S, R)$.

Initialization:
 For all $i \in \{1 \dots n\}$, for all $X \in N$,

$$\pi(i, i, X) = \begin{cases} 1 & \text{if } X \rightarrow x_i \in R \\ 0 & \text{otherwise} \end{cases}$$

Algorithm:

- For $l = 1 \dots (n - 1)$
 - For $i = 1 \dots (n - l)$
 - * Set $j = i + l$
 - * For all $X \in N$, calculate

$$\pi(i, j, X) = \sum_{\substack{X \rightarrow YZ \in R, \\ s \in \{i \dots (j-1)\}}} \underbrace{\hspace{10em}}_{\text{COMPLETE THE DEFINITION HERE}}$$

Output: Return $\pi(1, n, S)$

Question 2 Consider the CKY algorithm for finding the maximum probability for any tree when given as input a sequence of words x_1, x_2, \dots, x_n . As usual, we use N to denote the set of non-terminals in the grammar, and S to denote the start symbol.

The base case in the recursive definition is as follows: for all $i = 1 \dots n$, for all $X \in N$,

$$\pi(i, i, X) = \begin{cases} q(X \rightarrow x_i) & \text{if } X \rightarrow x_i \in R \\ 0 & \text{otherwise} \end{cases}$$

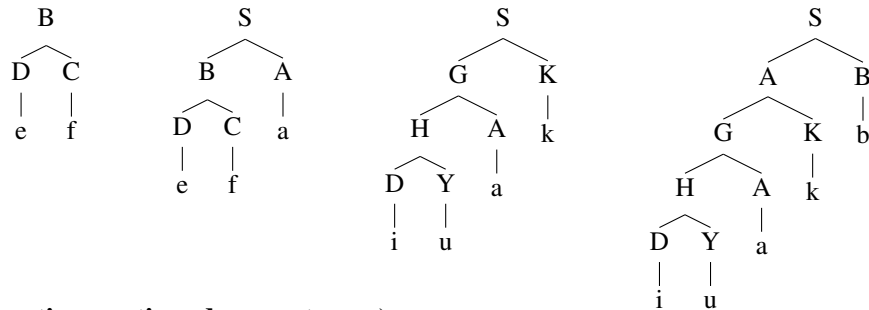
and the recursive definition is as follows: for all (i, j) such that $1 \leq i < j \leq n$, for all $X \in N$,

$$\pi(i, j, X) = \max_{\substack{X \rightarrow YZ \in R, \\ s \in \{i \dots (j-1)\}}} (q(X \rightarrow YZ) \times \pi(i, s, Y) \times \pi(s+1, j, Z))$$

Finally, we return

$$\pi(1, n, S) = \max_{t \in \mathcal{T}_G(s)} p(t)$$

Now assume that we want to find the maximum probability for any *left-branching* tree for a sentence. Here are some example left-branching trees:



(Question continued on next page)

It can be seen that in left-branching trees, whenever a rule of the form $X \rightarrow Y Z$ is seen in the tree, then the non-terminal Z must directly dominate a terminal symbol.

Question: Complete the recursive definition below, so that the algorithm returns the maximum probability for any **left-branching** tree underlying a sentence x_1, x_2, \dots, x_n .

Base case: for all $i = 1 \dots n$, for all $X \in N$,

$$\pi(i, i, X) = \begin{cases} q(X \rightarrow x_i) & \text{if } X \rightarrow x_i \in R \\ 0 & \text{otherwise} \end{cases}$$

Recursive case: (Complete below)

Return:

$$\pi(1, n, S) = \max_{t \in \mathcal{T}_G(s)} p(t)$$

Question 3 Consider the following PCFG (probabilities for each rule are shown after the rule):

$S \rightarrow NP VP$	1.0
$NP \rightarrow DT NBAR$	1.0
$NBAR \rightarrow NN$	0.7
$NBAR \rightarrow NBAR NBAR$	0.3
$VP \rightarrow sleeps$	1.0
$DT \rightarrow the$	1.0
$NN \rightarrow mechanic$	0.1
$NN \rightarrow car$	0.2
$NN \rightarrow metal$	0.7

Now consider a PCFG based on this context-free grammar. What parse tree will be returned as the highest probability tree for *the metal car mechanic sleeps*?

Question 4 Consider the following HMM:

- States in the HMM are $\{A, B\}$.
- q parameters of the HMM are $q(y|x)$ for $x \in \{A, B, *\}$ and $y \in \{A, B, \text{STOP}\}$.
- Vocabulary in the HMM is $\{s, t\}$.
- e parameters in the HMM of the form $e(y|x)$ for $x \in \{A, B, *\}$ and $y \in \{s, t\}$.

Our aim in this question is to write down a PCFG such that for any sentence $x_1 \dots x_n$ and tag sequence $y_1 \dots y_{n+1}$ with probability

$$p = p(x_1 \dots x_n, y_1 \dots y_{n+1})$$

under the HMM, there is a parse tree for the sentence $x_1 \dots x_n$ with the same probability p under the PCFG.

Complete the probabilities for the following rules in the PCFG (hint: try writing down parse trees for simple sentence/tag sequences such as $s/A, s/A \ t/B$ etc.):

$S \rightarrow A \text{ FA}$

$S \rightarrow B \text{ FB}$

$S \rightarrow A$

$S \rightarrow B$

$\text{FA} \rightarrow A \text{ FA}$

$\text{FA} \rightarrow A$

$\text{FA} \rightarrow B \text{ FB}$

$\text{FA} \rightarrow B$

$\text{FB} \rightarrow A \text{ FA}$

$\text{FB} \rightarrow A$

$\text{FB} \rightarrow B \text{ FB}$

$\text{FB} \rightarrow B$

$A \rightarrow s$

$A \rightarrow t$

$B \rightarrow s$

$B \rightarrow t$