

Lecture 6: Relationship to Linear Programming

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Relationship to Linear Programming (LP)

Relaxations

Our original optimization problem:

$$\arg \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} f(y) + g(z)$$

such that

$$y(i, t) = z(i, t)$$

for all i, t .

The Dual

- ▶ The Lagrangian:

$$L(u, y, z) = f(y) + g(z) + \sum_{i,t} u(i,t)(y(i,t) - z(i,t))$$

- ▶ The dual:

$$L(u) = \max_{y \in \mathcal{Y}, z \in \mathcal{Z}} L(u, y, z)$$

- ▶ The dual problem:

$$\min_u L(u)$$

- ▶ Theorem: for any value of u ,

$$L(u) \geq f(y^*) + g(z^*)$$

A New Optimization Problem

$$\arg \max_{\alpha \in \Delta_y, \beta \in \Delta_z} \sum_y \alpha_y f(y) + \sum_z \beta_z g(z)$$

such that

$$\sum_y \alpha_y y(i, t) = \sum_z \beta_z z(i, t)$$

for all i, t .

Here

$$\Delta_y = \{ \alpha \in \mathbb{R}^{|\mathcal{Y}|} : \forall y, \alpha_y \geq 0, \sum_y \alpha_y = 1 \}$$

$$\Delta_z = \{ \beta \in \mathbb{R}^{|\mathcal{Z}|} : \forall z, \beta_z \geq 0, \sum_z \beta_z = 1 \}$$

The New Dual

- ▶ The Lagrangian:

$$\begin{aligned} M(u, \alpha, \beta) = & \sum_y \alpha_y f(y) + \sum_z \beta_z g(z) \\ & + \sum_{i,t} u(i,t) \left(\sum_y \alpha_y y(i,t) - \sum_z \beta_z z(i,t) \right) \end{aligned}$$

- ▶ The dual:

$$M(u) = \max_{\alpha \in \Delta_y, \beta \in \Delta_z} M(u, \alpha, \beta)$$

- ▶ The dual problem:

$$\min_u M(u)$$

Theorems for the New Problem

Theorem 1:

$$\min_u M(u) = \sum_y \alpha_y^* f(y) + \sum_z \beta_z^* g(z)$$

Theorem 2: For any value of u ,

$$M(u) = L(u)$$