CSEE W3827 Fundamentals of Computer Systems Homework Assignment 6

Prof. Martha A. Kim Columbia University Due December 8, 2015 at **10:10**.

Write your name and UNI on your solutions

Show your work for each problem; we are more interested in how you get the answer than whether you get the right answer.

The remaining problems on this assignment require analysis of the two functions, adapted from the HW#4 solutions. Except for sll all of the instructions are supported by the pipelined processor from the lecture slides. The sll instruction will flow through the pipeline in the same manner as an addi.

The first is manhattan_dist:

i0: slt \$t0. \$a2. \$a0 i1: beq \$t0, \$0, i4 i2: sub \$t0, \$a0, \$a2 i3: beq \$0, \$0, i5 i4: sub \$t0, \$a2, \$a0 i5: slt \$t0, \$a3, \$a1 i6: beq \$t0, \$0, i9 sub \$t1, \$a1, \$a3 i7: i8: beq \$0, \$0, i10 i9: sub \$t1, \$a3, \$a1 i10: add \$v0. \$t0. \$t1

The second is edge_count_iter:

0_start: addi \$t0, \$0, 0 i1: addi \$t1. \$0. -1 i2: addi \$v0, \$0, 0 i3_ntop: beq \$t0, \$a1, i17_ndone i4: sll \$t2, \$t0, 2 i5: add \$t2, \$a0, \$t2 i6: lw \$t2, 0(\$t2) i7: beg \$t2, \$t1, i15_edone i8: sll \$t2, \$t2, 2 i9: add \$t2, \$a0, \$t2 i10_etop: lw \$t3, 0(\$t2) i11: beq \$t3, \$t1, i15_edone i12: addi \$v0, \$v0, 1 i13: addi \$t2, \$t2, 4 i14: beg \$0, \$0, i10_etop i15_edone: addi \$t0, \$t0, 1 i16: beq \$0, \$0, i3_ntop i17 ndone:

 (20 pts.) Give a cycle-by-cycle execution trace for the following implementation of manhattan_dist, when called on (0,0) and (5,5) (i.e., \$a0=0, \$a1=0, \$a2=5, \$a3=5) on a five-stage pipeline with *no forwarding*. The only option to resolve hazards is to stall (i.e., insert bubbles). 2. (20 pts.) Provide a second trace of manhattan_dist on the fully bypassed pipeline (i.e., branches resolved in D, forwarding from W-E, M-E, and M-D).

3. (20 pts.) Imagine how edge_count_iter would execute on the fully bypassed pipeline. List all pairs of instructions between which one or more bubbles would occur. If a bubble occurs between i3 and i4, then you should write i3 → i4. (HINT: Think systematically through all scenarios that result in an empty slot in the pipeline.)

4. (20 pts.) Now list where in the program (still edge_count_iter on the fully bypassed 5-stage pipe) data operands would be forwarded, and which forwarding path would be used. If i3 forwards the future value of \$a0 to i4 using the M-E forwarding path, write \$a0, i3 → i4, M-E, per the example below. (HINT: Think through the values consumed by each instruction systematically.)

5. (20 pts.) Assuming a very large graph with one million nodes and an average of 500 edges per node, what is the CPI of the edge_count_iter code? Assume that the beq in i7 is taken 2% of the time.