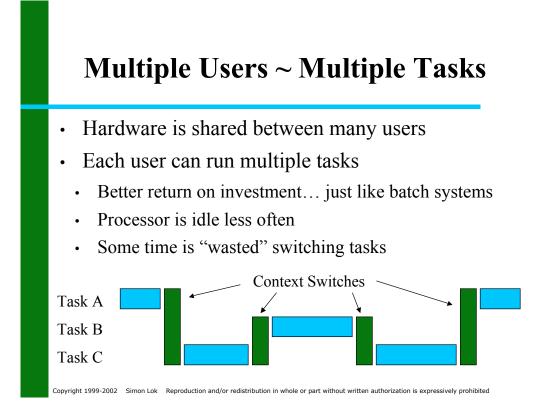


# In the Beginning...

- Computers ran a single task at a time...
  - Punch cards were placed into a feeder.
  - The cards were then read, compiled and run.
- Batch processing extended this...
  - Groups of punch cards could be run one after the next.
  - This increases return of investment in the hardware.

## Then Came the Operating System

- The OS sits between programs and the hardware.
- Contributions include:
  - Uniform interaction between hardware
    - I/O abstractions (e.g., filesystems)
    - Standardized interaction libraries (e.g., libc)
  - Multi-user Capabilities
    - Memory management and protection (virtual address space)
    - Scheduling and time sharing



# **Multi-Tasking**

- K users share the hardware running N tasks
- Tasks are time-sliced quickly to give the illusion that all tasks are running in parallel
- Each task thinks it's the only one on the machine
- Cooperative Multi-Tasking: (Win3.1, MacOS)
  - Each process yield the processor when it feels fit
- Pre-Emptive Multi-Tasking: (UNIX)
  - The OS scheduler decides who should run when

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# **Multi-Threading**

- If each user can run many tasks...
  - Why can't each task have many "sub-tasks"?
  - This is usually called multi-threading.
- Threads are like "lightweight" tasks...
  - Scheduling for execution is pretty much the same
- Differences include:
  - They share the same memory space.
  - They may not have as much OS overhead.



- Why multi-user / multi-tasking?
  - Processor is idle less, people can share a computer.
  - Better return on hardware investment
- We use threads for somewhat similar reasons:
  - Make sure processors are fully utilized
    - Don't block on I/O
    - · True parallel execution on multiprocessor hardware
  - Other cool things
    - · Games: intelligent user agents, animation

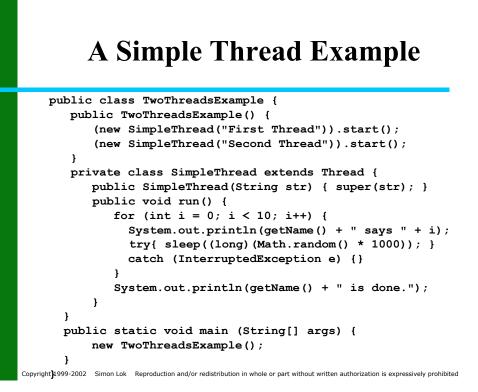
• Automatic garbage collection, hidden from the user

### How does the JVM affect this?

- Each JVM is a separate task on the native OS
- Most JVMs run a single JAVA program
- Each JVM (JAVA program) has many threads
  - In the simplest case, the GC and your main thread
  - · JVM threads to OS interaction depends on JVM
  - The newest JVMs (e.g., Sun HotSpot) will take advantage of physical multiprocessor hardware

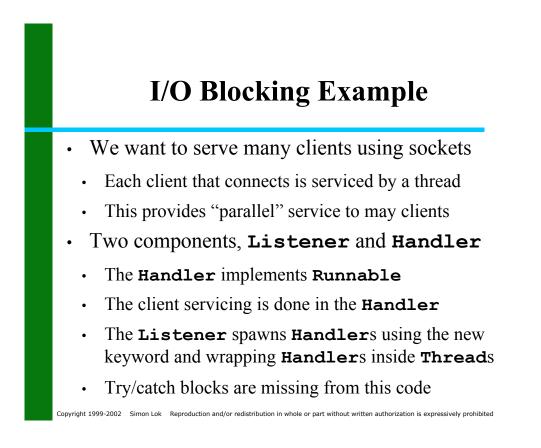
### **Programming Threads in JAVA**

- Two ways it can be done
  - Create a class, extend Thread
    - Override the run() method
    - Instantiate the class
    - Call start()
  - Create a class, implement Runnable
    - Implement the run() method
    - Instantiate your class
    - Instantiate a Thread class, pass your class in constructor
    - Call start()



## **TwoThreadsTest Output**

🚽 🚽 simon@sltpt20: /misc/ud/classes/java/codelib/threads 👘 🗆 🗙
[simon@sltpt20 threads]\$ java TwoThreadsExample
First Thread says 0
Second Thread says 0
First Thread says 1
First Thread says 2
Second Thread says 1
Second Thread says 2
First Thread says 3
First Thread says 4
Second Thread says 3
Second Thread says 4
First Thread says 5
First Thread says 6
First Thread says 7
First Thread says 8
Second Thread says 5
First Thread says 9
First Thread is done.
Second Thread says 6
Second Thread says 7
Second Thread says 8
Second Thread says 9
Second Thread is done.
💓 [simon@sltpt20 threads]\$
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#### **The Listener**

```
public class Listener {
  public static void main(String[] args) {
    ServerSocket srvSock = new ServerSocket(4567);
    while (keepRunning) {
        // when we get a connection, spawn off a
        // thread to handle it... this means we can
        // keep listening for other connections
        // while the first client is serviced
        Socket conn = srvSock.accept();
        (new Thread(new sockHandler(conn))).start();
    }
}
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```

#### **The Handler**

```
public class sockHandler implements Runnable {
    private Socket conn = null;
    public sockHandler(Socket conn) {
        this.conn = conn;
    }
    public void run() {
       InputStreamReader ISR = new
          InputStreamReader(conn.getInputStream());
      BufferedReader fromClient = new
          BufferedReader(ISR);
      OutputStreamReader OSR = new
          OutputStreamReader(conn.getOutputStream());
      PrintWriter toClient = new
          PrintWriter(OSR);
       // DO CLIENT SERVICING HERE
    }
 }
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```

### **Data Parallel Programming**

- We spawn off many threads to estimate PI
- As each thread completes, we update our estimate
- If we were running on MP hardware with a Hotspot JVM, these threads would run on separate processors and harness true parallelism
- Notice that the threads share a single memory space... that's why we can communicate between the sub-tasks and controller without RMI

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### **PI Estimation Task Thread**

```
public class PiEstimatorTask extends Thread {
    private EstimatePi Parent = null;
    private static final int iterations = 100000;
    public PiEstimatorTask(EstimatePi Parent) {
        this.Parent = Parent;
    }
    public void run() {
        int in = 0, out = 0;
        for (int i = 0; i < iterations; i++) {
            double x=2.0*Math.random()-1.0, y=2.0*Math.random()-1.0;
            if ((Math.sqrt(x*x+y*y) < 1.0)) { in++; } else { out++; }
        }
        double estimate = 4.0 * (double)in / (double)iterations;
        Parent.updateEstimate(estimate);
    }
}</pre>
```

### **PI Estimation Control Program**

```
public class EstimatePi {
    private double pi = 0.0;
    private final int numTasks = 12; // one for each processor
    private int allFinished = 0;
    private long starttime = 0;
    public synchronized void updateEstimate(double est) {
        long rt = System.currentTimeMillis() - starttime;
        System.out.println("Terminated at " + rt + " ms, est " + est);
        pi = (allFinished == 0) ? est : (pi + est) / 2;
        allFinished++;
    }
    public double getPi() { return pi; }
    // . . . continued on next slide . . .
```

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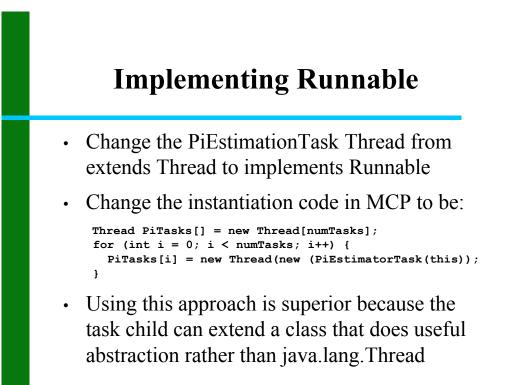
### **PI Estimation Control Program**

```
public void go() {
     PiEstimatorTask[] PiTasks = new PiEstimatorTask[numTasks];
     System.out.println("Instantiating " + numTasks + " threads");
     for (int i = 0; i < numTasks; i++) {
       PiTasks[i] = new PiEstimatorTask(this);
     ł
     starttime = System.currentTimeMillis();
     System.out.println("Starting threads, time = 0 ms");
     for (int i = 0; i < numTasks; i++)</pre>
       (PiTasks[i]).start();
     } // FIXME: try/catch InterruptedExeception below
     while(allFinished < numTasks) { Thread.sleep(1000); }</pre>
  }
  public static void main(String[] args) {
     EstimatePi MCP = new EstimatePi();
     MCP.go();
     System.out.println("Final estimate is: " + MCP.getPi());
  }
}
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```

# PI Estimation Output (Uniproessor Hardware)

File Sessions Options Help	
Instantiating 12 threads	
Starting 12 threads, time = 0 ms	
T1 terminated @ 965 ms, estimate 3.14752	
TO terminated @ 2001 ms, estimate 3.13968	
Master waiting for threads @ 2833 ms	
Master waiting for threads @ 4843 ms	
Master waiting for threads @ 6853 ms	
Master waiting for threads @ 8863 ms	
Master waiting for threads @ 10873 ms	
T2 terminated @ 11123 ms, estimate 3.15296	
T4 terminated @ 11308 ms, estimate 3,14236	
T3 terminated @ 11332 ms, estimate 3.14216	
T5 terminated @ 11680 ms, estimate 3.14432	
T9 terminated @ 11894 ms, estimate 3.14056	
T7 terminated @ 12089 ms, estimate 3,14456	
T8 terminated @ 12182 ms, estimate 3,136	
T11 terminated @ 12373 ms, estimate 3,13736	
T10 terminated @ 12389 ms, estimate 3,14248	
T6 terminated @ 12421 ms, estimate 3,14568	
Final estimate is: 3.14297 [simon@sltpt20 simon]\$	

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## **Did you notice?**

- The method updateEstimate in the control program is declared synchronized?
- This is JAVA's way of providing MutEx
  - MutEx is short for Mutual Exclusion
  - Only one person can go at a time
- When you have parallel processing, you will almost always run into MuTex problems

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### **Synchronized Methods**

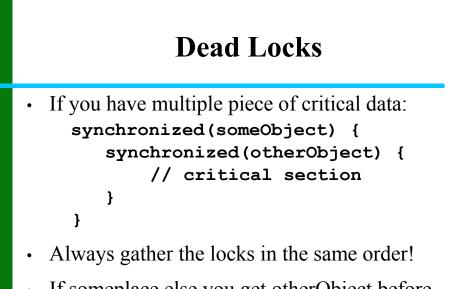
- public void synchronized doSomething() {
   // guaranteed that only one Thread
   // will be running this method
   // at any time
  }
- If one thread is running this method, any other thread calling it will wait until the first thread has returned from the method.
- Note that this can slow things down tremendously! Use with caution.

## **Synchronized Blocks**

```
• public void someMethod() {
    synchronized(someObject) {
        // do some stuff
    }
}
```

- Only one thread can be running the block at any time
- The someObject is the data that is critical
- Fine grain atomicity
  - Generally results in better performance
  - Harder to create code
  - Harder to read code later on

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• If someplace else you get otherObject before someObject, you might end up with deadlock.

### **In Summary**

- Multi-Threading enables use make better use of contemporary computers:
  - Prevents idle/busy waiting CPU
  - Non-blocking I/O
  - Parallel execution (on MP hardware)
  - Automatic garbage collection
  - Fun uses for games (animation / bad guys with AI)