

A Domain-Specific Language for Generating Dataflow Analyzers

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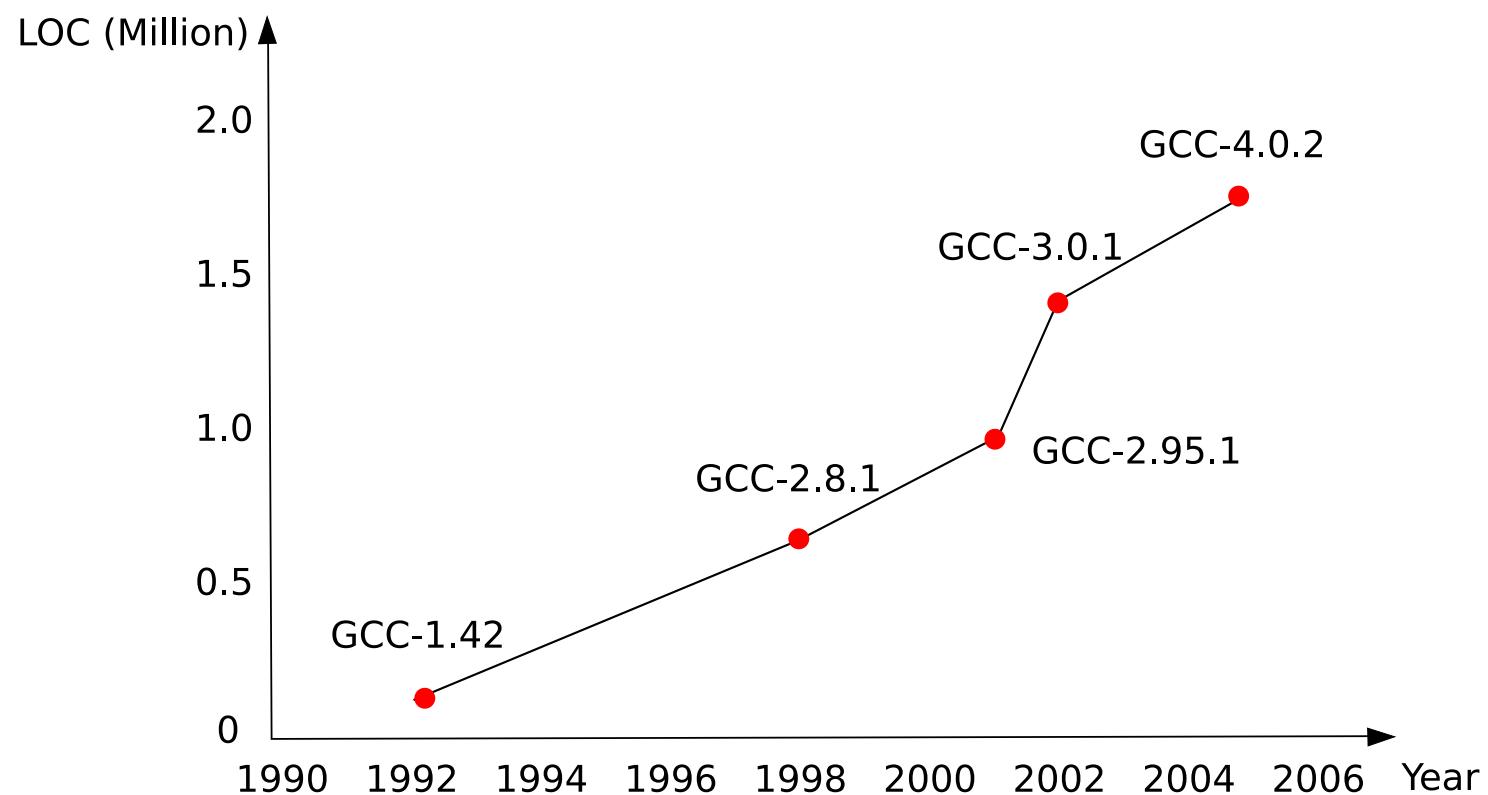
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Motivation



GCC Code Sizes

Motivation

- Algorithm: Live variable analysis in Dragon Book

Input: A flow graph with def and use

Output: $out[B]$

for each block B **do** $in[B] = \emptyset$

while changes to any of the in 's occur **do**

for each block B **do begin**

$out[B] = \cup_{S \text{ a successor of } B} in[S]$

$in[B] = use[B] \cup (out[B] - def[B])$

end

Motivation: AG vs. Phoenix C++

Input: A flow graph with *def* and *use*

Output: $out[B]$

for each block B do $in[B] = \emptyset$

while changes to any of the in 's occur do

for each block B do begin

$$out[B] = \bigcup_S \text{a successor of } B \text{ } in[S]$$

$$in[B] = use[B] \cup (out[B] - def[B])$$

end

```

Phase Liveness(Func func){
    extend class Cond{
        SetAlias< Def;
        SetAlias< Use;
    }

    void init(){
        Opnd opnd = this;

        if(opnd->isDef()){
            foreach(Cond alias_of_tag(alias_tag, opnd->AliasTag,AliasInfo)
                opnd->def += alias_tag;
        }
        else{//uses
            foreach(MayPartialAliasOfTag alias_tag,opnd->AliasTag,AliasInfo)
                opnd->use += alias_tag;
        }
    }
}

extend class Instr {
    SetAlias< Def; //set will be translate to sparse bit vector on opnd->ID
    SetAlias< Use;
}

void init() {
    Inst *inst = this;
    foreach(Opnd opnd in inst where (dataflow && src)){
        inst->use += opnd->use;
    }
    foreach(Opnd opnd in inst where (dataflow && dst){
        inst->def += opnd->def;
    }
}

extend class Block{
    SetAlias< Def;
    SetAlias< Use;
}

void init(){
    Block block = this;

    foreach(Instr inst in block)
        block->use = (block->use + inst->use) - block->def;
    block->def = (block->def + inst->def) - block->use;
}

SetAlias< TransFunc(Backward,Iterative,Block){
    Compose(N){
        In = (Out - N->Def) + N->Use;
    }

    Meet(N){
        Out += N->In;
    }

    Closure(N){
        Result(N){
            Print(Out);
            Print(Out);
        }
    }
}

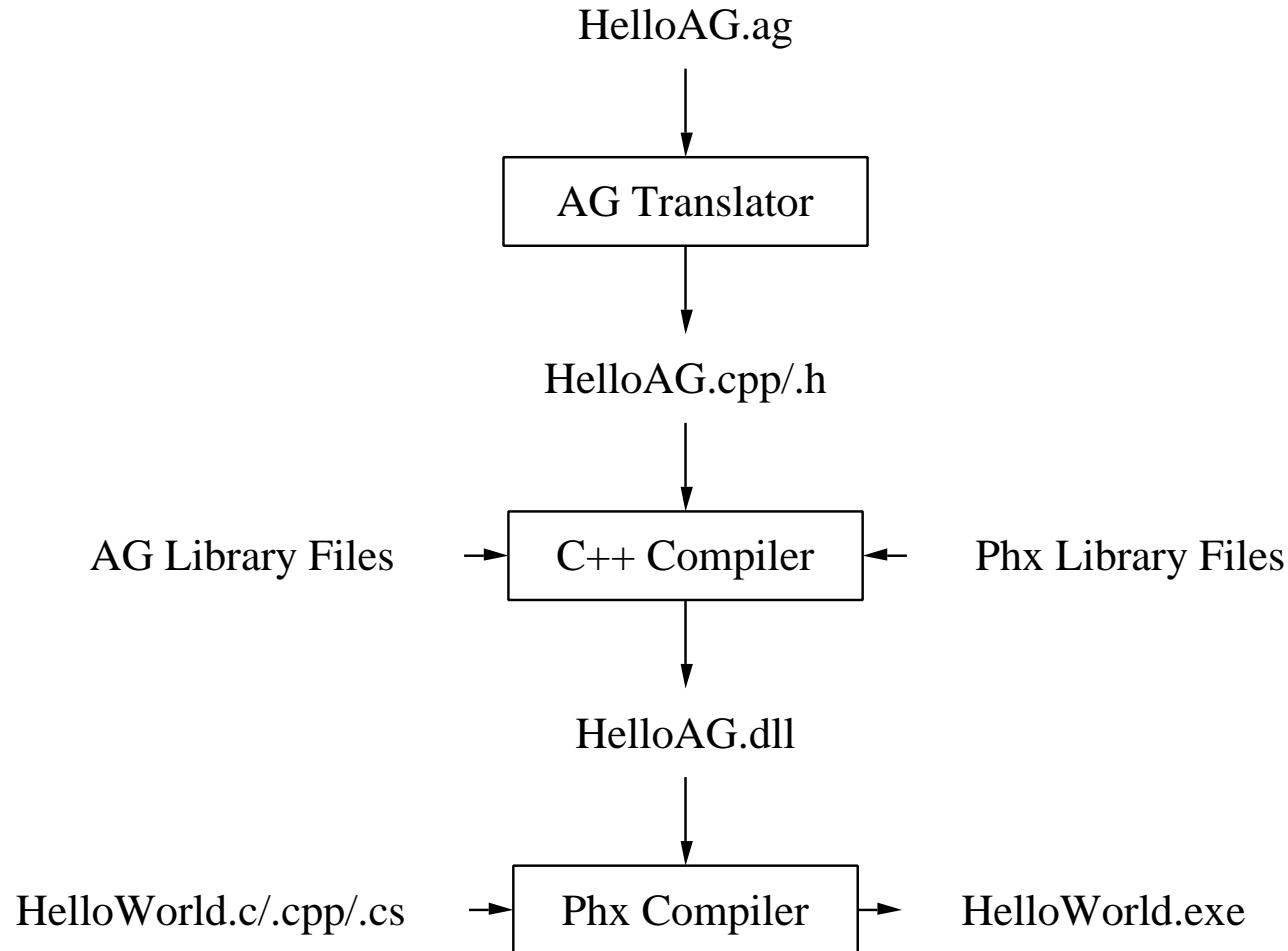
```

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Previous Work

- Theory of dataflow analysis
 - Gary A. Kildall [1973]: unified lattice-based framework
 - Kam and Ullman [1976]: iterative approach
- Applicable tools
 - Tjiang's Sharlit [1992]: efficiency > simplicity
 - Alt and Martin's PAG [1995]: lattice specification
 - Yi & Harrison's work [1993]: interprocedural analysis

Overview of AG Tool



An AG Program

```
Phase name {
    extend class name {
        //field declarations...
        //method declarations...
        void Init() { ... }
    }
    ...
    type TransFunc(direction) {
        Meet(P) { ... }
        Compose(N) { ... }
        Result(N) { ... }
    }
}
```

Example: Reaching Definitions

A *definition* of a variable is the operand in an instruction that may assign a value to the variable.

```
y = 0;  
if (x > 0) {  
    y = x * 2;  
    x --;  
}  
else  
    x ++;  
z = y * x;
```

Example: Reaching Definitions

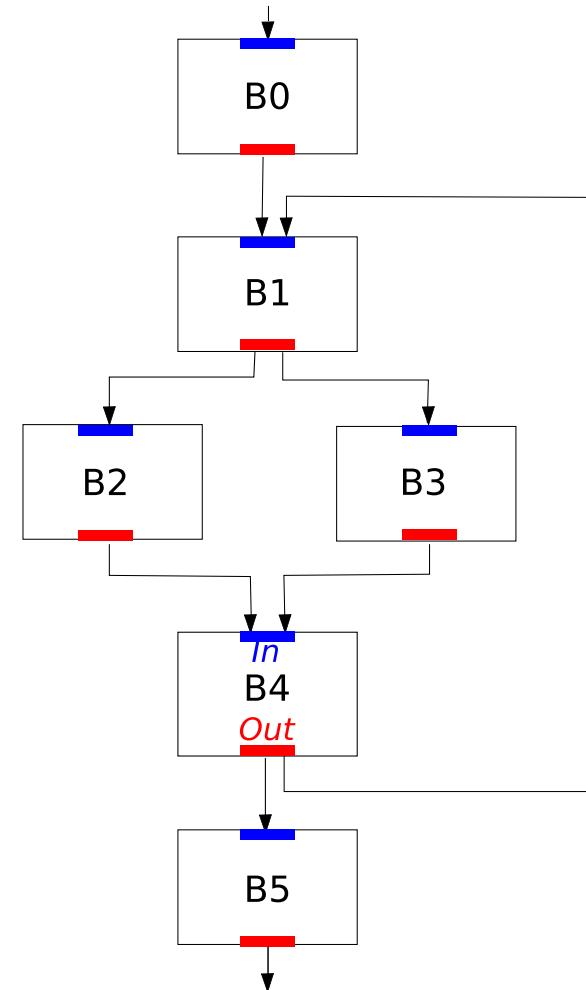
A *definition* of a variable is the operand in an instruction that may assign a value to the variable.

```
y = 0;  
if (x > 0) {  
    y = x * 2;  
    x --;  
}  
else  
    x ++;  
z = y * x;
```

The diagram illustrates the flow of definitions for variables y, x, and z. A red curved arrow starts at the assignment `y = 0;` and points to the first occurrence of `y` in the expression `y = x * 2;`. Another red curved arrow starts at the assignment `x ++;` and points to the first occurrence of `x` in the same expression. A third red curved arrow starts at the assignment `z = y * x;` and points to the first occurrence of `y` in that expression.

Example: Reaching Definitions

```
for each block  $B$  do begin
    Compute  $gen[B]$  and  $kill[B]$ ;
end
for each block  $B$  do  $in[B] = \emptyset$ ;
 $change = \text{true}$ ;
while  $change$  do begin
     $change = \text{false}$ ;
    for each block  $B$  do begin
         $oldout = out[B]$ ;
         $in[B] = \cup_P$  a predecessor of  $B$   $out[P]$ ;
         $out[B] = gen[B] \cup (in[B] - kill[B])$ ;
        if  $out[B] \neq oldout$  then  $change = \text{true}$ ;
    end
end
```



Example: Reaching Definitions in AG

for each block B do begin

 Compute $gen[B]$ and $kill[B]$;

end

for each block B do $in[B] = \emptyset$;

$change = \text{true}$;

while $change$ do begin

$change = \text{false}$;

 for each block B do begin

$oldout = out[B]$;

$in[B] = \bigcup_P$ a predecessor of B $out[P]$;

$out[B] = gen[B] \cup (in[B] - kill[B])$;

 if $out[B] \neq oldout$ then $change = \text{true}$;

 end

end

Phase $name$ {
 extend class $name$ {
 //field declarations...
 //method declarations...
 void Init() { ... }
 }
 ...
 type TransFunc(direction)
 Meet(P) { ... }
 Compose(N) { ... }
 Result(N) { ... }
 }
}

Example: Reaching Definitions in AG

```
Phase ReachingDefs {  
    extend class Opnd {  
        ...  
    }  
    extend class Instr {  
        ...  
    }  
    extend class Block {  
        ...  
    }  
    Set<Opnd> TransFunc(Forward) {  
        ...  
    }  
}
```

Open Classes

```
class Instr {
    int id;
    list<Opnd> srcOpnds;
    list<Opnd> dstOpnds;
}

class MyInstr : Instr {
    list<Opnd> gen;
    list<Opnd> kill;
    void Init() {...}
}
```

Open Classes

```
class Instr {  
    int id;  
    list<Opnd> srcOpnds;  
    list<Opnd> dstOpnds;  
}  
  
class MyInstr : Instr {  
    list<Opnd> gen;  
    list<Opnd> kill;  
    void Init() {...}  
}
```

```
class Instr {  
    int id;  
    list<Opnd> srcOpnds;  
    list<Opnd> dstOpnds;  
  
    //added fields  
    list<Opnd> gen;  
    list<Opnd> kill;  
  
    //added methods  
    void Init() {...}  
}
```

Previous Work

	AspectJ	MultiJava	AG
Usage	General Purpose	General Purpose	Dataflow Analysis
Methods	+	+	+
Fields	+	-	+
Recompiling	Yes	No	No
Dynamic	-	+	+

Example AG: Extend Class Block

```
extend class Block {
    Set<Opnd> Gen;
    Set<Opnd> Kill;

    void Init() {
        Block block = this;

        foreach (Instr instr in block) {
            block->Gen = instr->Gen +
                (block->Gen - instr->Kill);
            block->Kill = block->Kill + instr->Kill
                - instr->Gen;
        }
    }
}
```

Example AG: Extend Class Instr

```
extend class Instr {
    Set<Opnd> Gen;
    Set<Opnd> Kill;

    void Init() {
        Instr instr = this;
        foreach (Opnd dstOpnd in instr
                    where (dataflow && dst)) {
            instr->Gen += dstOpnd->Gen;
            instr->Kill += dstOpnd->Kill;
        }
    }
}
```

Example AG: Extend Class Opnd

```
extend class Opnd {
    Set<Opnd> Gen;
    Set<Opnd> Kill;

    void Init() {
        Opnd opnd = this;
        if (opnd->IsDef) {
            opnd->Gen += opnd;
            foreach_must_total_alias_of_tag(alias_tag,
                opnd->AliasTag, AliasInfo) {
                opnd->Kill += DstAliasTable(alias_tag);
            }
            opnd->Kill -= opnd;
        }
    }
}
```

Example AG: C++ Code Generated

```
void OpndExtensionObject::Init( Phx::FuncUnit *func_unit,
                               Phx::BitVector::Sparse *PHX_ARRAY(dst_alias_table))
{
    Phx::IR::Opnd *opnd = _this;
    if(opnd->IsDef) {
        this->Gen->SetBit(this->uid);
        foreach_must_total_alias_of_tag(alias_tag,
                                         opnd->AliasTag, func_unit->AliasInfo) {
            this->Kill->Or(dst_alias_table[alias_tag]);
        }
        next_must_total_alias_of_tag;
        this->Kill->ClearBit(this->uid);
    }
}
```

Example AG: Transfer Functions

```
Set<Opnd> TransFunc(Forward) {  
    Meet(P) {  
        In += P->Out;  
    }  
  
    Compose(N) {  
        Out = In - N->Kill + N->Gen;  
    }  
}
```

AG Syntax - Highlights

data types	Set Map int bool void
special variables	In Out
operators	+ - * = += -= *= && ! !=
built-in classes	Opnd Instr Block Alias Expr Func Region
built-in constants	Forward Backward
statements	<code>foreach (type var in range where cond. direction)</code> <code>{ ... }</code> <i>phoenix-iterator (...) { ... }</i> <i>lvalue = expression;</i> <code>if (expression) { ... } else { ... }</code> <i>/% arbitrary C++ code %/</i>
built-in functions	DstAliasTable SrcAliasTable Print

Experimental Results

	Reaching Definitions	Live Variables	Uninitialized Variables
C++ LOC (manual)	791	303	108
C++ Style	dynamic	global var	SSA
AG LOC (manual)	64	55	94
C++ LOC (generated)	626	519	682
C++ runtime	7.3s	0.8s	
AG runtime	7.4s	3.1s	13.6s

- Benchmark:

Phoenix Microsoft Intermediate Language reader
(> 500,000 lines)

Conclusions

- Strengths
 - Very compact code: < 1/10 of manually-written C++ program
 - Similar performance: speed penalty < 4 times
 - Mechanism for extending existing IR classes
- Weaknesses
 - Fixed framework: iterative, MIR-based
 - Debug environment

Acknowledgments

- ☺ Thanks to Al Aho!
- ☺ Thanks to the Microsoft Phoenix group!