

A Domain-Specific Language for Generating Dataflow Analyzers

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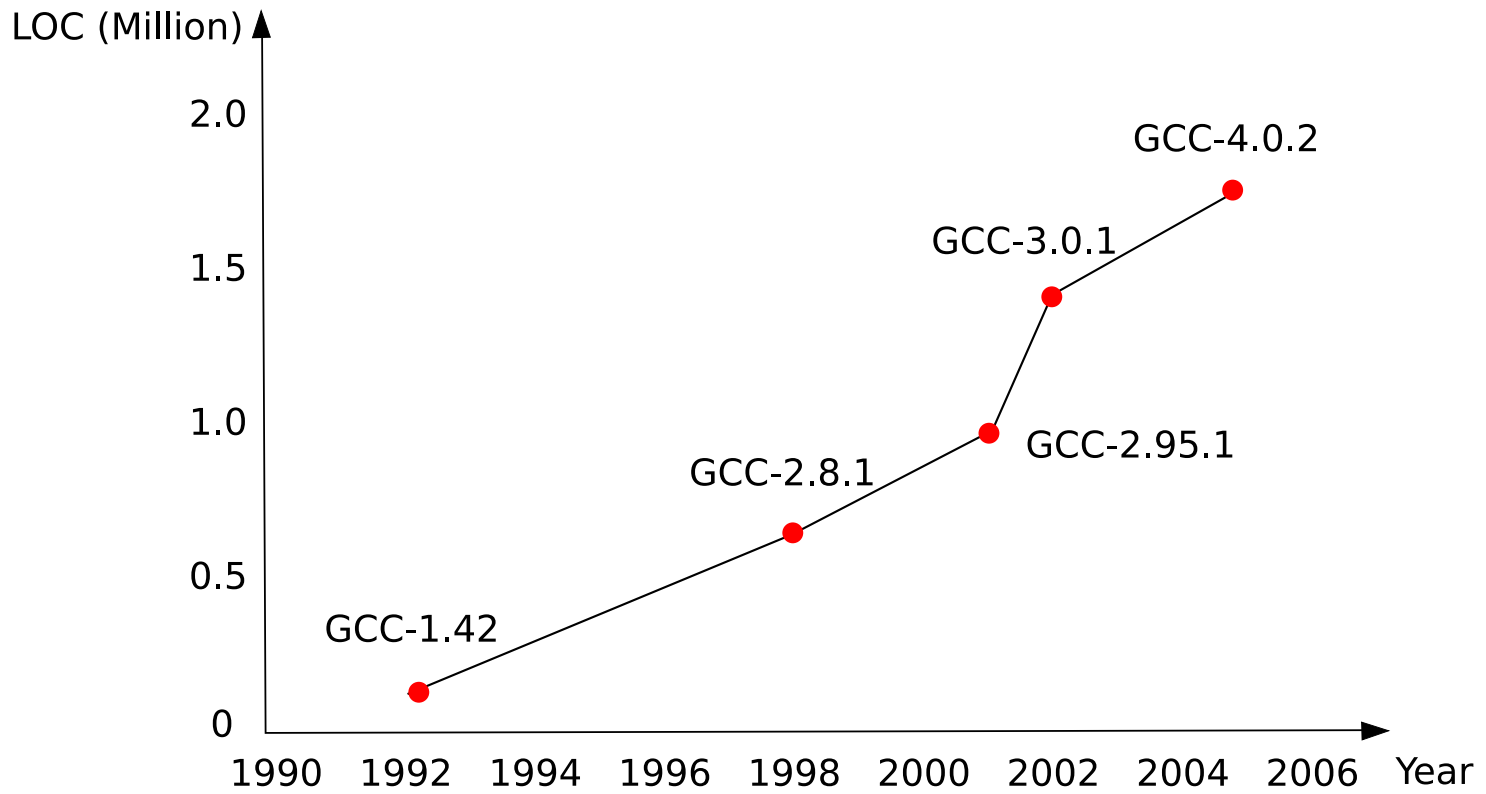
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Motivation



GCC Code Sizes

Motivation

- Algorithm: Live variable analysis in Dragon Book

Input: A flow graph with *def* and *use*

Output: $out[B]$

for each block B **do** $in[B] = \emptyset$

while changes to any of the *in*'s occur **do**

for each block B **do begin**

$out[B] = \cup_{S \text{ a successor of } B} in[S]$

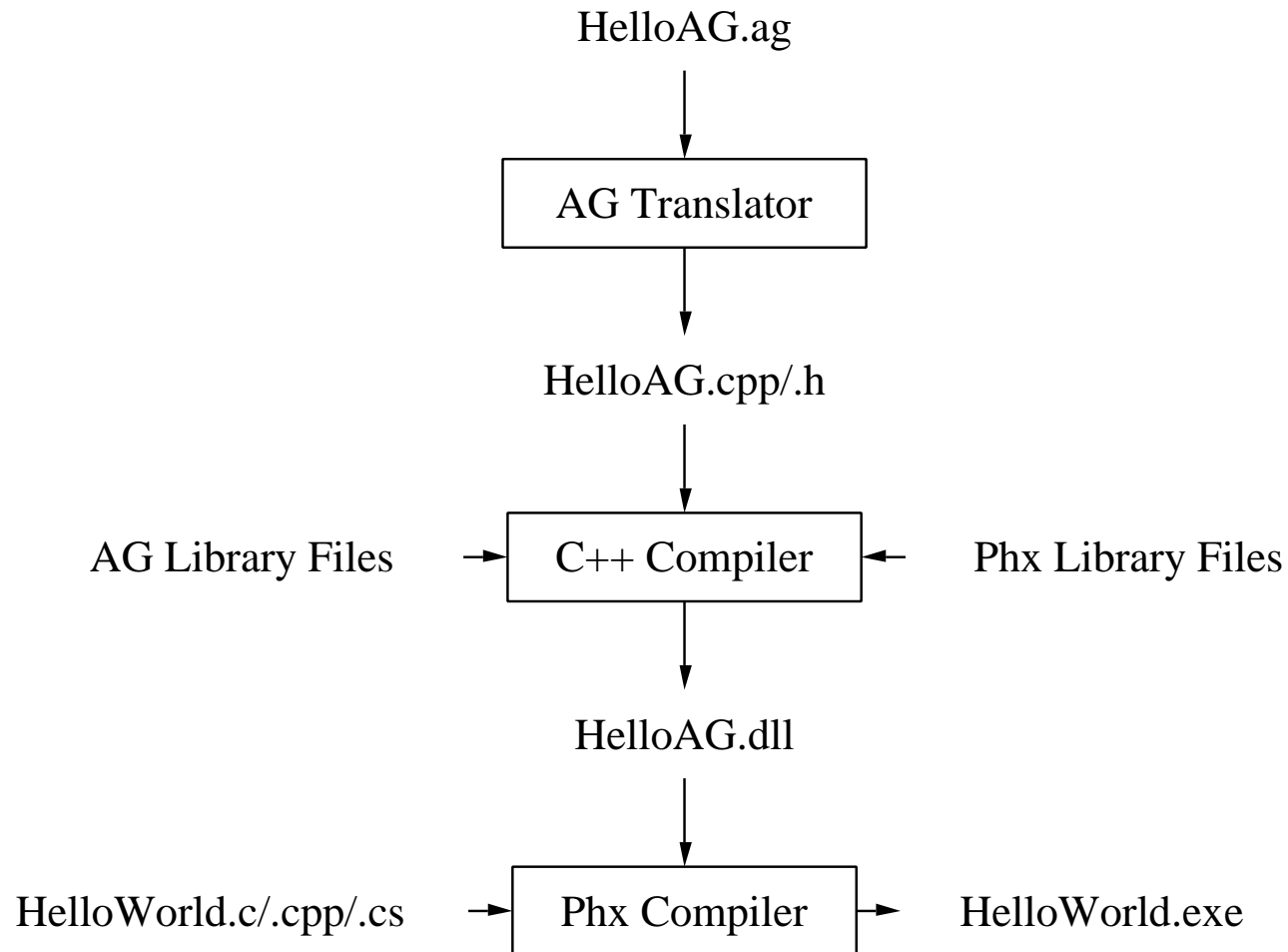
$in[B] = use[B] \cup (out[B] - def[B])$

end

Previous Work

- Theory of dataflow analysis
 - Gary A. Kildall [1973]: unified lattice-based framework
 - Kam and Ullman [1976]: iterative approach
- Applicable tools
 - Tjiang's Sharlit [1992]: efficiency > simplicity
 - Alt and Martin's PAG [1995]: lattice specification
 - Yi & Harrison's work [1993]: interprocedural analysis

Overview of AG Tool



An AG Program

```
Phase name {
  extend class name {
    //field declarations...
    //method declarations...
    void Init() { ... }
  }
  ...
  type TransFunc(direction) {
    Meet(P) { ... }
    Compose(N) { ... }
    Result(N) { ... }
  }
}
```

Example: Reaching Definitions

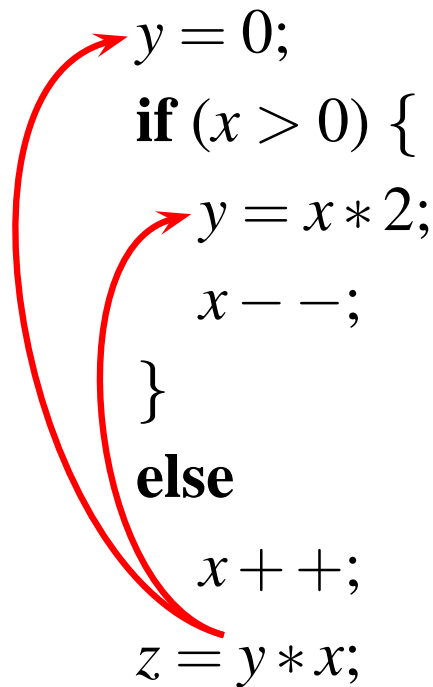
A *definition* of a variable is the operand in an instruction that may assign a value to the variable.

```
y = 0;  
if (x > 0) {  
    y = x * 2;  
    x --;  
}  
else  
    x ++;  
z = y * x;
```


Example: Reaching Definitions

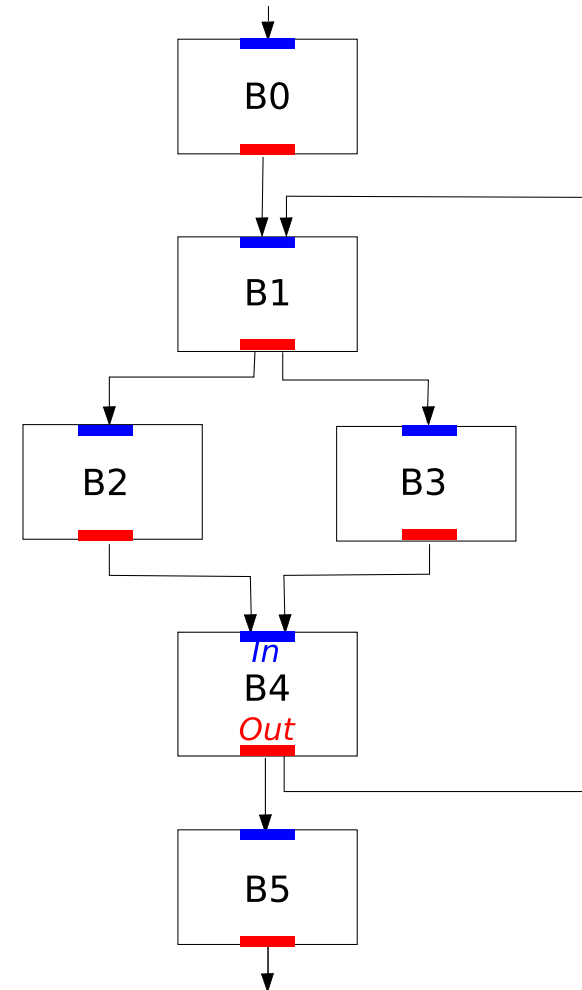
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```
    y = 0;  
    if (x > 0) {  
        y = x * 2;  
        x --;  
    }  
    else  
        x ++;  
    z = y * x;
```



Example: Reaching Definitions

```
for each block  $B$  do begin
    Compute  $gen[B]$  and  $kill[B]$ ;
end
for each block  $B$  do  $in[B] = \emptyset$ ;
change = true;
while change do begin
    change = false;
    for each block  $B$  do begin
         $oldout = out[B]$ ;
         $in[B] = \cup_{P \text{ a predecessor of } B} out[P]$ ;
         $out[B] = gen[B] \cup (in[B] - kill[B])$ ;
        if  $out[B] \neq oldout$  then change = true;
    end
end
end
```



Example: Reaching Definitions in AG

for each block B **do begin**

 Compute $gen[B]$ and $kill[B]$;

end

for each block B **do** $in[B] = \emptyset$;

$change = \mathbf{true}$;

while $change$ **do begin**

$change = \mathbf{false}$;

for each block B **do begin**

$oldout = out[B]$;

$in[B] = \cup_{P \text{ a predecessor of } B} out[P]$;

$out[B] = gen[B] \cup (in[B] - kill[B])$;

if $out[B] \neq oldout$ **then** $change = \mathbf{true}$;

end

end

Phase $name$ {

extend class $name$ {

 //field declarations...

 //method declarations...

 void Init() { ... }

}

...

type **TransFunc**($direction$)

 Meet(P) { ... }

 Compose(N) { ... }

 Result(N) { ... }

}

}

Example: Reaching Definitions in AG

```
Phase ReachingDefs {
  extend class Opnd {
    ...
  }
  extend class Instr {
    ...
  }
  extend class Block {
    ...
  }
  Set<Opnd> TransFunc(Forward) {
    ...
  }
}
```

Open Classes

```
class Instr {  
    int id;  
    list<Opnd> srcOpnds;  
    list<Opnd> dstOpnds;  
}
```

```
class MyInstr : Instr {  
    list<Opnd> gen;  
    list<Opnd> kill;  
    void Init() {...}  
}
```

Open Classes

```
class Instr {  
    int id;  
    list<Opnd> srcOpnds;  
    list<Opnd> dstOpnds;  
}  
  
class MyInstr : Instr {  
    list<Opnd> gen;  
    list<Opnd> kill;  
    void Init() {...}  
}
```

```
class Instr {  
    int id;  
    list<Opnd> srcOpnds;  
    list<Opnd> dstOpnds;  
  
    //added fields  
    list<Opnd> gen;  
    list<Opnd> kill;  
  
    //added methods  
    void Init() {...}  
}
```

Previous Work

	AspectJ	MultiJava	AG
Usage	General Purpose	General Purpose	Dataflow Analysis
Methods	+	+	+
Fields	+	—	+
Recompiling	Yes	No	No
Dynamic	—	+	+

Example AG: Extend Class Block

```
extend class Block {  
    Set<Opnd> Gen;  
    Set<Opnd> Kill;  
  
    void Init() {  
        Block block = this;  
  
        foreach (Instr instr in block) {  
            block->Gen = instr->Gen +  
                (block->Gen - instr->Kill);  
            block->Kill = block->Kill + instr->Kill  
                - instr->Gen;  
        }  
    }  
}
```


Example AG: Extend Class Instr

```
extend class Instr {  
    Set<Opnd> Gen;  
    Set<Opnd> Kill;  
  
    void Init() {  
        Instr instr = this;  
        foreach (Opnd dstOpnd in instr  
            where (dataflow && dst)) {  
            instr->Gen += dstOpnd->Gen;  
            instr->Kill += dstOpnd->Kill;  
        }  
    }  
}
```

Example AG: Extend Class Opnd

```
extend class Opnd {  
    Set<Opnd> Gen;  
    Set<Opnd> Kill;  
  
    void Init() {  
        Opnd opnd = this;  
        if (opnd->IsDef) {  
            opnd->Gen += opnd;  
            foreach_must_total_alias_of_tag(alias_tag,  
                opnd->AliasTag, AliasInfo) {  
                opnd->Kill += DstAliasTable(alias_tag);  
            }  
            opnd->Kill -= opnd;  
        }  
    }  
}
```

Example AG: C++ Code Generated

```
void OpndExtensionObject::Init( Phx::FuncUnit *func_unit,
                               Phx::BitVector::Sparse *PHX_ARRAY(dst_alias_table))
{
    Phx::IR::Opnd *opnd = _this;
    if(opnd->IsDef) {
        this->Gen->SetBit(this->uid);
        foreach_must_total_alias_of_tag(alias_tag,
                                         opnd->AliasTag, func_unit->AliasInfo) {
            this->Kill->Or(dst_alias_table[alias_tag]);
        }
        next_must_total_alias_of_tag;
        this->Kill->ClearBit(this->uid);
    }
}
```

Example AG: Transfer Functions

```
Set<Opnd> TransFunc(Forward) {  
    Meet(P) {  
        In += P->Out;  
    }  
  
    Compose(N) {  
        Out = In - N->Kill + N->Gen;  
    }  
}
```

AG Syntax - Highlights

data types Set Map int bool void

special variables In Out

operators + - * = += -= *= && || ! !=

built-in classes Opnd Instr Block Alias Expr Func Region

built-in constants Forward Backward

statements *foreach* (*type var in range where cond. direction*)
{ ... }

phoenix-iterator (...) { ... }

lvalue = expression;

if (*expression*) { ... } *else* { ... }

!% arbitrary C++ code %!

built-in functions DstAliasTable SrcAliasTable Print

Experimental Results

	Reaching Definitions	Live Variables	Uninitialized Variables
C++ LOC (manual)	791	303	108
C++ Style	dynamic	global var	SSA
AG LOC (manual)	64	55	94
C++ LOC (generated)	626	519	682
C++ runtime	7.3s	0.8s	
AG runtime	7.4s	3.1s	13.6s

- Benchmark:

Phoenix Microsoft Intermediate Language reader
(> 500,000 lines)

Conclusions

- Strengths
 - Very compact code: $< 1/10$ of manually-written C++ program
 - Similar performance: speed penalty < 4 times
 - Mechanism for extending existing IR classes
- Weaknesses
 - Fixed framework: iterative, MIR-based
 - Debug environment

Acknowledgments

☺ Thanks to Al Aho!

☺ Thanks to the Microsoft Phoenix group!