CS1004: Intro to CS in Java, Spring 2005

Lecture \#20: Algorithms, cont'd.
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| Administrivia |
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| - HW\#\#4 due now |
| - Extra credits returned today |
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## Board examples

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- Finish Fibonacci numbers $\qquad$
- Array algorithms
- Search for a number (or an item in general) in a list $\qquad$
- Find the largest number in a list
- Sort numbers $\qquad$
- We'll do more in the homework and in the rest of the semester $\qquad$
$\qquad$
$\qquad$


## Algorithm correctness \& efficiency

- Define desirable characteristics in an algorithm:
- Correctness
- Does the algorithm solve the problem it is designed for?
- Does the algorithm solve the problem correctly?
- Ease of understanding
- How easy is it to understand or alter an algorithm?
- Important for program maintenance


## Attributes of Algorithms (continued)

- Elegance
- How clever or sophisticated is an algorithm?
- Sometimes elegance and ease of understanding work at cross-purposes
- Efficiency
- How much time and/or space does an algorithm require when executed?
- Perhaps the most important desirable attribute


## Measuring Efficiency

- Analysis of algorithms
- Study of the efficiency of various algorithms
- Efficiency measured as function relating size of input to time or space used
- For one input size, best case, worst case, and $\qquad$ average case behavior must be considered
- The $\Theta / O$ notation captures the order of magnitude of the efficiency function
- $\Theta$ ("big-Theta") vs. O ("big-Oh") notation
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## Order of Magnitude: Order n

- As n grows large, order of magnitude dominates running time, minimizing effect of coefficients and lower-order terms
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All functions that have a linear shape are considered equivalent

- Order of magnitude $n$
- Written $\mathrm{O}(\mathrm{n})$
- Functions vary as a constant times n


## Sequential Search, analyzed

- Comparison of the NAME being searched for against a $\qquad$ name in the list
- Central unit of work $\qquad$
- For lists with n entries:
- Best case
- NAME is the first name in the list, 1 comparison $\qquad$ - $\mathrm{O}(1)$
- Worst case
- NAME is the last name in the list, or not in list $\qquad$
- n comparisons, or $\mathrm{O}(\mathrm{n})$
- Average case
- Roughly $\mathrm{n} / 2$ comparisons, or $\mathrm{O}(\mathrm{n})$ $\qquad$
$\qquad$


## Sequential Search (continued)

- Space efficiency $\qquad$
■ Uses essentially no more memory storage than original input requires
- Very space-efficient
- But... is there a faster way to search through a $\qquad$ list?


## Binary Search

- Given ordered data,
- Search for $N A M E$ by comparing to middle element
- If not a match, restrict search to either lower or
$\qquad$ upper half only
- Each pass eliminates half the data
- Efficiency
- Best case
- 1 comparison: $\mathrm{O}(1)$
- Worst case
- $\lg \mathrm{n}$ comparisons: $\mathrm{O}(\lg \mathrm{n})$
- What's $\lg n$ ?


A Comparison of n and $\lg \mathrm{n}$ (S/G, pg. 109)

## Sorting

- What if we want to sort the numbers in a list? $\qquad$
- There are number of algorithms; book describes selection sort, but we'll also go over bubble sort
$\qquad$ very quickly.
- Let's begin!
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| Next time |
| :---: |
| - Finish working with algorithms (for now) |
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