

- 1 **CS3134 #22**
 - 11/20/03
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- 2 **Administrivia**
 - Minor typo on HW#5 (points)
 - Scheduling final exam?
- 3 **Agenda**
 - Minimum spanning tree
 - Directed graph algorithms
- 4 **Complexity of BFS and DFS?**
 - Optimally, $O(V+E)$ – we visit every vertex a constant number of times and potentially travel every edge a constant number of times
 - But this is only for an adjacency list; in an adjacency matrix version, it's $O(V^2)$ – we scan every row and every column in the adjacency matrix once
 - Admittedly inefficient, but we knew that
- 5 **Minimum spanning trees**
 - A (minimum) spanning tree is a subgraph with no cycles
 - Different in weighted graphs
 - Remove graph redundancy
 - Useful for many applications
 - Ex: minimize wiring
 - In a minimum spanning tree, $\#E = \#V - 1$
 - Simple algorithm (p. 644): DFS and record the edges traveled
 - Don't worry about backtracking
 - Can also use BFS...
- 6 **Directed graphs**
 - As earlier mentioned, useful for situations where we need to model “one-way” information
 - Streets
 - Trees are a subclass of directed graphs
 - Book: course prerequisites
 - Topological sorting: come up with a legitimate ordering of processing the nodes
 - Often useful for *partial ordering* problems, such as aforementioned course prerequisites
 - Result: a order where no vertex y comes before a vertex x where $x \rightarrow y$
 - There can be multiple correct answers!
- 7 **Topological sort**
 - Find a vertex that has no successors, i.e., arrows that point to *it*
 - Look at columns of the adjacency matrix
 - Delete that vertex and print it out
 - Repeat
 - What kinds of graphs doesn't this work for?
 - Cycles – what happens?
 - “Catch-22” in real life

- In other words, works on generalized trees (multiple roots, etc.) – *DAG*
- Complexity again $O(V+E)/O(V^2)$

8 **Topological sort (II)**

- How to find node with no successors?
- How do you delete a node?

9 **Next time**

- Warshall's Algorithm
- Start weighted graphs