## Administrivia

- Textbooks delayed. ©
- Should I delay the homework?
- No class next Monday - Memorial Day


## Proofs

- Theorem: Statement that can be shown true via a proof.
- Conjecture is a statement whose truth value is unknown; turns into a theorem given a proof
- Asiom/postulate are underlying assumptions about mathematical structures, hypothesis, and previously proved theorems
- Rules of inference tie steps together
- Need to avoid fallacies
- Lemma: mini-proof used in other proofs; a corollary is a "side-effect" of a proof.
$4 \square$ Rules of inference
- Need these for proofs
- Modus ponens, or law of detachment
- Example: consider the tautology $(p \wedge(p \rightarrow q)) \rightarrow q$
- Either $p$ is true, in which case $p \rightarrow q$ depends on $q$, or $p$ is false, in which case $p \rightarrow q$ is always true
- Therefore, this is equivalent to $q$
- Other rules - see page 58 and 60
- Multiple ways of writing tautologies...

5
$\square$ Valid argument

- An argument is valid if all the hypotheses are true.
- Valid doesn't mean true!
- All the propositions must be true
- Scenario:
- It is not sunny this afternoon and it is colder than yesterday.
- We will go swimming only if it is sunny.
- If we do not go swimming, then we will take a canoe trip.
- If we take a canoe trip, then we will be home by sunset.
- Conclusion: we will be home by sunset.
$6 \square$ Fallacy
- An invalid argument
- Fallacy of affirming the conclusion: $[(p \rightarrow q) \wedge q] \rightarrow p$
- Just because q is true, doesn't mean p is
- "If you do homework, then you are smart"; "you are smart"; "therefore you did homework" doesn't fly, i.e., homework isn't the only criterion for becoming smart.
- Fallacy of denying the hypothesis: $[(p \rightarrow q) \wedge \neg p] \rightarrow \neg p$

7

## Rules of inference with quantifiers

- Universal instantiation: given $\forall x P(x)$, we can conclude $P(c)$.
- Universal generalization: given $P(c)$ true for all $c$, we can say $\forall x P(x)$ by selecting a truly arbitrary $c$.
- Existential instantiation: If $\exists x P(x)$, select an appropriate $c$ for which $P(c)$. We bind " $c$ " to it and use it through the argument.
- Existential generalization: If $P(c)$ is known for a $c$, we can state $\exists x P(x)$ is true.
- Example:
- "Everyone in this discrete mathematics class has taken a course in Computer Science
- "John is a student in this class"
- implies "John has taken a course in Computer Science."
- Mathematical theorems often omit the universal quantifier (i.e., for all real numbers, etc.) - it's all done implicitly


## Methods of proving theorems

- Direct proof: what we've been doing so far
- Assume $p$ is true and use rules of inference to show that $q$ must be true
- Example: If $n$ is an odd integer, then $n^{2}$ is an odd integer
- Definition 1 : $n$ is even if $k$ such that $n=2 k$ and odd if $k$ such that $n=2 k+1$
- Indirect proof: use contrapositive
- Show that if $q$ is false, $p$ must be false
- Example: If $3 n+2$ is odd, then $n$ is odd $\rightarrow$ assume $n$ is even
- Vacuous proof: if the hypothesis $p$ is false, then $p \rightarrow q$ is automatically true
- Example: $P(0)$ where $P(n)$ "If $n>1, n+1>1$."
$9 \square$ Proving theorems (II)
- Proof by contradiction: show that $\neg p \rightarrow q$ is true, i.e., $\neg p \rightarrow F$ or $q=F$. Therefore, $\neg p$ must be false and $p$ must be true.
- Example: Show at least 4 of any 22 days must fall on the same day of the week => assume this is false
- Proof by cases: decouple $\left(p_{1} \vee p_{2} \vee \ldots \vee p_{n}\right) \rightarrow q$ into $\left(p_{1} \rightarrow q\right) \wedge\left(p_{2} \rightarrow q\right) \wedge\left(p_{n} \rightarrow q\right)$.
- Proofs of equivalence: decouple $p \leftrightarrow q$ into $(p \rightarrow q) \wedge(q \rightarrow p)$


## Mistakes in proofs, techniques

- Theorem: If $\mathrm{n}^{2}$ is positive, then n is positive.
- "Proof:" Suppose $\mathrm{n}^{2}$ is positive. If n is positive, $\mathrm{n}^{2}$ is positive. Therefore $n$ is positive.
- Why: Let $P(n)$ be " $n$ is positive" and $Q(n)$ be " $n^{2}$ is positive". $\forall n(P(n) \rightarrow Q(n)), Q(n)$ doesn't mean $P(n)$
- How to choose right method?
- Black magic...
- Just a beginning
- We'll keep things simple in the course - I'll allow lots of leeway.


## Sets

- A set is an unordered collection of objects.
- Useful way of grouping discrete structures together.
- Everything builds on top of this abstract concept.
- The objects in a set are also called the elements or members of a set.
- Notation $\in$
- Duplicates make no difference, i.e., $\{1,3,5\}=\{1,1,3,3,5,5\}$
- How to describe?
- List all members $\mathrm{V}=\{\mathrm{a}, \mathrm{e}, \mathrm{i}, \mathrm{o}, \mathrm{u}\}$
- Set of integers less than $100=\{1,2,3, \ldots, 99\}$

12
$\square$ Common sets

- $\mathbf{N}=$ natural numbers $=\{0,1,2,3, \ldots\}-$ sometimes not zero
- $\mathbf{Z}=$ integers $=\{\ldots,-2,-1,0,1,2, \ldots\}$
- $\mathbf{Z}^{+}=$positive integers $=\{1,2, \ldots\}$
- $\mathbf{Q}=$ rational numbers $=\{p / q \mid p \in \mathbf{Z}, q \in \mathbf{Z}, q \neq 0\}$
- $\mathbf{R}=$ real numbers (incl. irrationsal)


## Other notations

- Set builder: State the property they must have to be members

$$
-\mathrm{O}=\{\mathrm{x} \mid \mathrm{x} \text { is an odd positive integer less than } 10\}
$$

- Venn diagram
- Remember Universal Set $U$ is the contents of the box
- Venn diagram showing vowels?
- Empty set: $\}$ or $\varnothing$
- Equal: Two sets are equal iff they have the same elements.
- Subset: $A \subseteq B-A$ is a subset of $B$ if and only if every element of $A$ is also an element of $B$. - For any set $\mathrm{S}, \varnothing \subseteq \mathrm{S}$ and $\mathrm{S} \subseteq \mathrm{S}$


## More set notation

- Proper subset, $\subset$
- Show equality by showing each set is a subset of the other (can't be proper)
- Can nest sets within sets
- Cardinality of a set is $|\mathrm{S}|$, number of distinct elements in a set, assuming S is finite.
- Power set of $S$ is the set of all subsets of $S$, or $P(S)$.
$-P(\{0,1,2\})=\{\varnothing,\{0\},\{1\}, \ldots,\{0,1,2\}\}$
$-P(\varnothing)=\{\varnothing\}$
- Power set of a set has $2^{n}$ elements.
$15 \square$ Tuples and Cartesian Product
- Generally ordered, as opposed to sets
- Ordered n-tuple ( $a_{1}, \ldots, a_{n}$ )
- Cartesian product of set $A$ and set $B, A \times B$, is the set of all ordered pairs $(a, b)$ where $a \in A$ and $b \in B$, i.e.,
$-A \times B=\{(a, b) \mid a \in A \wedge b \in B\}$
- Example: $A$ is the set of students, and $B$ the set of courses at a university - Cartesian product is the set of all possible enrollments of students in courses.
- N-way Cartesian product generates n-tuples (not nested tuples!)
$16 \square$ Set notation with quantifiers
- As I showed last time...
- $\forall x \in \mathbf{R}\left(x^{2} \geq 0\right)$
- Can also use set builder notation
$17 \square$ Set operations
- Union $(\cup)$ is the set that contains those elements that are in $A, B$, or both. - Generally don't include duplicates
- Intersection $(\cap)$ is the set containing elements in both A and B
- Illustrate using Venn diagrams
- Disjoint if intersection is the empty set.
- Difference, or $A-B$, is the set containing elements in $A$ but not in $B$.
- Complement, or $A$ with a bar on top, is the complement of $A$ with respect to $U$ (the universal set). - Difference is the intersection of $A$ and the complement of $U$


## Set identities

- Page 89, similar to logical equivalences
- Can use direct proof or membership table to demonstrate
- Example: prove that $\mathrm{A} \cap(\mathrm{B} \cup \mathrm{C})=(\mathrm{A} \cap \mathrm{B}) \cup(\mathrm{A} \cap \mathrm{C})$


## Generalized union/intersection, computer representation

- Concept remains same; notation is slightly different - Page 92/93
- How to represent in a computer?
- If finite, use bitstrings, assuming ordered.
- Can use NOT, AND, OR to do complement, intersection, and union.


## Functions

- A function $f$ from (set) $A$ to (set) $B$ is an assignment of exactly one element of $B$ to each element of $A$.
- $f(a)=b$ if $b$ is the unique element of $B$ assigned by the function $f$ to the element $a$ of $A$.
- $b$ is the image of $A$ and $a$ is a preimage of $b$.
- range of $f$ is the set of all images of elements of A.
- If $f$ is a function from $A$ to $B$, we can write $f: A \rightarrow B$.
- $A$ is the domain of $f$ and $B$ is the codomain of $f$.
- f "maps" A to B.
- Examples
- Page 97 for a visual representation
- $f: \mathbf{Z} \rightarrow \mathbf{Z}$ assigns the square of an integer to this integer. Then, $f(x)=x^{2}$.
- Note range and codomain may not be the same.

21

## Function Operations

- Real-valued functions can be added and/or multiplied - just "combine" the individual functions

$$
\text { - If } f_{1}(x)=x^{2} \text { and } f_{2}(x)=x-x^{2},\left(f_{1}+f_{2}\right)(x)=x \text { and }\left(f_{1} f_{2}\right)(x)=x^{3}-x^{4} \text {. }
$$

- If a subset of a domain is defined, you can define its image as well. $-f(S)=\{f(s) \mid s \in S\}$.


## One-to-one vs. onto

- Functions always map each preimage to a unique value.
- One-to-one suggests that every mapping maps to a unique image, i.e., $f(x)=f(y)$ implies that $x=y$.
- "Injection"
- $f(x)=x^{2}$ is not one-to-one, because of negative values.
- $f(x)=x+1$ is one-to-one.
- Onto suggests that each element of the codomain has a preimage.
- $f(x)=x^{2}$ is not onto, because of negative or skipped integers
- $f(x)=x+1$ is onto (infinite trick)
- "Surjection"
- One-to-one correspondence/bijection if it's both.
- See diagram on page 101.


## Inverse and composition

- The inverse of a function, $\mathrm{f}^{-1}$, assigns to an element $b$ in B the unique element $a$ in A such that $f(a)=b$.
- Must be one-to-one correspondence (i.e, one-to-one and onto).
- The composition ( f o g$)(\mathrm{a})=\mathrm{f}(\mathrm{g}(\mathrm{a}))$
- Not the same as ( $\mathrm{g} \circ \mathrm{f}$ )(a).
$24 \square$ Graphs, miscellaneous functions
- Exactly what you'd expect...
- Although not necessarily continuous
- Floor (or greatest integer) function $(\lfloor x\rfloor)$ returns the largest integer that is less than or equal to a real number x.
- Ceiling function $(\lceil x\rceil)$ returns the smallest integer that is greater than or equal to a real number $x$.
- Graphs of both on page 106
- Note open circles mean open intervals, e.g., floor has same value from $[\mathrm{n}, \mathrm{n}+1$ ) and ceiling has the same value from ( $n, n+1$ )
- Various useful properties on page 107
- Is $\lceil x+y\rceil=\lceil x\rceil+\lceil y\rceil$ ?


## Next time

- Algorithms, growth
- Integers and integer algorithms

