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² Administrivia

- Textbooks delayed. ☺
 - Should I delay the homework?
- No class next Monday Memorial Day

³ Proofs

- Theorem: Statement that can be shown true via a proof.
 - Conjecture is a statement whose truth value is unknown; turns into a theorem given a proof
- Asiom/postulate are underlying assumptions about mathematical structures, hypothesis, and previously proved theorems
- Rules of inference tie steps together
- Need to avoid fallacies
- Lemma: mini-proof used in other proofs; a corollary is a "side-effect" of a proof.

⁴ Bules of inference

- Need these for proofs
- · Modus ponens, or law of detachment
 - Example: consider the tautology $(p \land (p \rightarrow q)) \rightarrow q$
 - Either p is true, in which case $p \rightarrow q$ depends on q, or p is false, in which case $p \rightarrow q$ is always true
 - Therefore, this is equivalent to q
 - Other rules see page 58 and 60
- Multiple ways of writing tautologies...

5 🔲 Valid argument

- An argument is valid if all the hypotheses are true.
- Valid doesn't mean true!
 - All the propositions must be true
- · Scenario:
 - It is not sunny this afternoon and it is colder than yesterday.
 - We will go swimming only if it is sunny.
 - $-\,$ If we do not go swimming, then we will take a canoe trip.
 - If we take a canoe trip, then we will be home by sunset.
 - Conclusion: we will be home by sunset.

6 🔲 Fallacy

- · An invalid argument
- Fallacy of affirming the conclusion: $[(p \rightarrow q) \land q] \rightarrow p$
 - Just because q is true, doesn't mean p is
 - "If you do homework, then you are smart"; "you are smart"; "therefore you did homework" doesn't fly, i.e., homework isn't the only criterion for becoming smart.
- Fallacy of denying the hypothesis: [(p→q) ∧ ¬p] → ¬p

⁷ Rules of inference with quantifiers

- Universal instantiation: given $\forall x P(x)$, we can conclude P(c).
- Universal generalization: given P(c) true for all c, we can say $\forall x P(x)$ by selecting a truly arbitrary c.
- Existential instantiation: If $\exists xP(x)$, select an appropriate c for which P(c). We bind "c" to it and use it through the argument.
- Existential generalization: If P(c) is known for a c, we can state $\exists x P(x)$ is true.

- Example:
 - "Everyone in this discrete mathematics class has taken a course in Computer Science"
 - "John is a student in this class"
 - implies "John has taken a course in Computer Science."
- Mathematical theorems often omit the universal quantifier (i.e., for all real numbers, etc.) it's all done implicitly.

8 B Methods of proving theorems

- Direct proof: what we've been doing so far
 - Assume p is true and use rules of inference to show that q must be true
 - Example: If n is an odd integer, then n^2 is an odd integer
 - Definition 1: *n* is even if *k* such that n = 2k and odd if *k* such that n = 2k+1
- Indirect proof: use contrapositive
 - Show that if q is false, p must be false
 - Example: If 3n+2 is odd, then n is odd \rightarrow assume n is even.
- Vacuous proof: if the hypothesis p is false, then p → q is automatically true
 Example: P(0) where P(n) "If n > 1, n+1 > 1."

Proving theorems (II)

- Proof by contradiction: show that ¬p → q is true, i.e., ¬p → F or q = F. Therefore, ¬p must be false and p must be true.
 - Example: Show at least 4 of any 22 days must fall on the same day of the week => assume this is false
- Proof by cases: decouple $(p_1 \lor p_2 \lor ... \lor p_n) \rightarrow q$ into $(p_1 \rightarrow q) \land (p_2 \rightarrow q) \land (p_n \rightarrow q)$.
- Proofs of equivalence: decouple $p \leftrightarrow q$ into $(p \rightarrow q) \land (q \rightarrow p)$

10 Mistakes in proofs, techniques

- Theorem: If n² is positive, then n is positive.
 - "Proof:" Suppose n^2 is positive. If n is positive, n^2 is positive. Therefore *n* is positive.
 - − Why: Let P(n) be "n is positive" and Q(n) be "n² is positive". $\forall n(P(n) \rightarrow Q(n)), Q(n)$ doesn't mean P(n)
- How to choose right method?
 Black magic...
- Just a beginning
 - We'll keep things simple in the course I'll allow lots of leeway.

11 🔲 Sets

- A set is an unordered collection of objects.
- Useful way of grouping discrete structures together.
- · Everything builds on top of this abstract concept.
- The objects in a set are also called the elements or members of a set.
 Notation ∈
 - Duplicates make no difference, i.e., $\{1, 3, 5\} = \{1, 1, 3, 3, 5, 5\}$
- How to describe?
 - List all members V = {a, e, i, o, u}
 - Set of integers less than $100 = \{1, 2, 3, ..., 99\}$

¹² Common sets

- N = natural numbers = {0, 1, 2, 3, ...} sometimes not zero
- **Z** = integers = {..., -2, -1, 0, 1, 2, ...}
- **Z**⁺ = positive integers = {1, 2, ...}
- **Q** = rational numbers = { $p/q | p \in \mathbf{Z}, q \in \mathbf{Z}, q \neq 0$ }
- **R** = real numbers (incl. irrationsal)

¹³ Other notations

· Set builder: State the property they must have to be members

- $O = \{x \mid x \text{ is an odd positive integer less than 10} \}$
- Venn diagram
 - Remember Universal Set U is the contents of the box
 - Venn diagram showing vowels?
- Empty set: { } or \varnothing
- Equal: Two sets are equal iff they have the same elements.
- Subset: A ⊆ B -- A is a subset of B if and only if every element of A is also an element of B.
 For any set S, Ø ⊆ S and S ⊆ S

¹⁴ More set notation

- Proper subset, ⊂
- Show equality by showing each set is a subset of the other (can't be proper)
- · Can nest sets within sets
- Cardinality of a set is |S|, number of *distinct* elements in a set, assuming S is *finite*.
- Power set of S is the set of all subsets of S, or P(S).
 - $\ \mathsf{P}(\{0,\,1,\,2\}) = \{ \varnothing,\,\{0\},\,\{1\},\,\ldots,\,\{0,\,1,\,2\} \}$
 - $\mathsf{P}(\emptyset) = \{\emptyset\}$
 - Power set of a set has 2ⁿ elements.

¹⁵ Tuples and Cartesian Product

- · Generally ordered, as opposed to sets
 - Ordered n-tuple (a1, ..., an)
 - Cartesian product of set A and set B, A x B, is the set of all ordered pairs (a,b) where $a \in A$ and $b \in B$, i.e.,
 - $\hspace{0.1cm} A \hspace{0.1cm} x \hspace{0.1cm} B = \{(a,b) \hspace{0.1cm} | \hspace{0.1cm} a \hspace{0.1cm} \in \hspace{0.1cm} A \hspace{0.1cm} \wedge \hspace{0.1cm} b \hspace{0.1cm} \in \hspace{0.1cm} B\}$
 - Example: A is the set of students, and B the set of courses at a university Cartesian product is the set
 of all possible enrollments of students in courses.
- N-way Cartesian product generates n-tuples (not nested tuples!)

¹⁶ Set notation with quantifiers

- As I showed last time...
- $\forall x \in \mathbf{R} (x^2 \ge 0)$
- · Can also use set builder notation

¹⁷ Set operations

- Union (\cup) is the set that contains those elements that are in A, B, or both. – Generally don't include duplicates
- Intersection (\cap) is the set containing elements in both A and B
- Illustrate using Venn diagrams
- Disjoint if intersection is the empty set.
 Difference, or A-B, is the set containing elements in A but not in B.
- Complement, or A with a bar on top, is the complement of A with respect to U (the universal set).
 Difference is the intersection of A and the complement of U

18 Set identities

- Page 89, similar to logical equivalences
- Can use direct proof or *membership table* to demonstrate
 Example: prove that A ∩ (B ∪ C) = (A ∩ B) ∪ (A ∩ C)
- ¹⁹ Generalized union/intersection, computer representation
 - Concept remains same; notation is slightly different – Page 92/93
 - How to represent in a computer?

- If finite, use bitstrings, assuming ordered.
- Can use NOT, AND, OR to do complement, intersection, and union.

²⁰ E Functions

- A function f from (set) A to (set) B is an assignment of exactly one element of B to each element of A.
 - f(a) = b if b is the unique element of B assigned by the function f to the element a of A.
 b is the *image* of A and a is a *preimage* of b.
 - range of f is the set of all images of elements of A.
 - If f is a function from A to B, we can write f: $A \rightarrow B$.
 - A is the *domain* of f and B is the *codomain* of f.
 - f "maps" A to B.
- Examples
 - Page 97 for a visual representation
 - f: Z → Z assigns the square of an integer to this integer. Then, f(x) = x².
 Note range and codomain may not be the same.
- ²¹ Function Operations
 - Real-valued functions can be added and/or multiplied just "combine" the individual functions

- If $f_1(x) = x^2$ and $f_2(x) = x - x^2$, $(f_1 + f_2)(x) = x$ and $(f_1f_2)(x) = x^3 - x^4$.

If a subset of a domain is defined, you can define its image as well.
 - f(S) = {f(s) | s ∈ S}.

22 One-to-one vs. onto

- Functions always map each preimage to a unique value.
- One-to-one suggests that every mapping maps to a unique image, i.e., f(x) = f(y) implies that x = y.
 "Injection"
 - $f(x) = x^2$ is *not* one-to-one, because of negative values.
 - f(x) = x+1 is one-to-one.
- Onto suggests that each element of the codomain has a preimage.
 - $f(x) = x^2$ is not onto, because of negative or skipped integers
 - f(x) = x+1 is onto (infinite trick)
 - "Surjection"
- · One-to-one correspondence/bijection if it's both.
- See diagram on page 101.

²³ Inverse and composition

- The inverse of a function, f⁻¹, assigns to an element b in B the unique element a in A such that f(a) = b.
 - Must be one-to-one correspondence (i.e, one-to-one and onto).
- The composition $(f \circ g)(a) = f(g(a))$
 - Not the same as (g o f)(a).

²⁴ Graphs, miscellaneous functions

- Exactly what you'd expect...
 - Although not necessarily continuous
- Floor (or greatest integer) function ([x]) returns the largest integer that is less than or equal to a real number x.
- Ceiling function (\bar{x}) returns the smallest integer that is greater than or equal to a real number x.
- Graphs of both on page 106
- Note open circles mean open intervals, e.g., floor has same value from [n, n+1) and ceiling has the same value from (n, n+1)
 Various useful properties on page 107
 - Is[x+y] = [x] + [y]?

²⁵ Next time

- · Algorithms, growth
- · Integers and integer algorithms