

# IPsec

## Slide 1

### Protocol security - where?

---

**Application layer:** (+): easy access to user credentials, extend without waiting for OS vendor, understand data; (-): design again and again; e.g., PGP, ssh, Kerberos

**Transport layer:** (+): security mostly seamlessly, but difficult to get credentials; e.g., TLS

**Network layer:** (+): reduced key management, fewer application changes, fewer implementations, VPNs; (-) non-repudiation, multi-user machines, partial security in “middle boxes”

**Data link layer:** (+): speed; (-): hop-by-hop only

## Slide 2

## Documents

---

Document Roadmap	RFC 2411
Architecture	RFC 2401
IP Authentication Header (AH)	RFC 2402
IP Authentication Using Keyed MD5	RFC 1828
IP Encapsulating Security Payload (ESP)	RFC 2406
The Oakley Key Determination Protocol	RFC 2412
Internet Sec. Assoc. and Key Mgmt. P. (ISAKMP)	RFC 2408
The Internet Key Exchange (IKE)	RFC 2409
HMAC: Keyed-Hashing for Message AuthenticationA	RFC 2104

### Slide 3

## IPSec services

---

- IPv4 and IPv6 unicast
- access control
- connectionless integrity
- data origin authentication
- protection against replays (partial sequence integrity)
- confidentiality (encryption)
- limited traffic flow confidentiality.
- todo: NAT, multicast

### Slide 4

## Architecture

---

**Authentication header (AH):** access control, integrity, data origin authentication, replay protection

**Encapsulating Security Payload (ESP):** access control, confidentiality, traffic flow confidentiality.

**Key management protocols:** IKE = OAKLEY + ISAKMP, ...

- for any upper-layer protocol
- no effect on rest of Internet
- algorithm-independent, but default algorithms

### Slide 5

## Architecture

---

- between host and/or security gateways
- security gateway = router, firewall, ...
- security policy database (SPD)  $\rightsquigarrow$  IPsec, discarded, or bypass
- negotiate compression (why?)
- *tunnel mode* or *transport mode*
- granularity: single host-host tunnel vs. one per TCP connection

### Slide 6

## Implementation

---

- native IP implementation
- bump in the stack (BITS): beneath IP layer
- bump in the wire (BITW)

### Slide 7

## Security Association (SA)

---

- simplex
- AH *or* ESP
- identified by
  - Security Parameter Index (SPI),
  - IP destination address,
  - security protocol (AH or ESP) identifier.
- transport mode: two hosts
  - AH or ESP after IPv4 options, before UDP/TCP
  - IPv6: after base header and extensions, before/after destination options
  - mostly for higher-layer protocols (but: AH also some IP header parts)
- tunnel mode: one or two security gateways

### Slide 8

- outer header  $\Rightarrow$  tunnel endpoint
- security header between outer and inner
- traffic hiding; ESP payload padding

## Slide 9

### Nested Security Associations

---

AH and ESP  $\Rightarrow$  two SAs (“SA bundle”):

- transport adjacency: AH, then ESP
- both tunnel endpoints the same
- one endpoint the same
- neither the same

## Slide 10

## Security Policy Database

---

- map to Security Association Database (per packet or per SPD entry)
- discard, bypass or apply to *inbound* or *outbound*
- ordered list of filters (stateless firewall)
- example: “use ESP in transport mode using 3DES-CBC with explicit IV, nested inside of AH in tunnel mode using HMAC-SHA-1.”
- selectors:
  - destination IP address: address, range, address + mask, wildcard
  - source IP address
  - name (for BITS/BITW hosts): user id, X.500 DN, system name, opaque, ...
  - data sensitivity label
  - transport layer protocol

### Slide 11

- source/destination ports
- per socket setup or per packet (BITS, BITW, gateway)

### Slide 12

## Security Association Database (SAD)

---

- inbound: outer destination address
- IPsec protocol (AH or ESP)
- SPI (32-bit value)

**Slide 13**

## Examples of Implementations

---

- end-to-end security ( $H1^* == H2^*$ )
- VPN ( $H1 - SG1^* == SG2^* - H2$ )
- e2e + VPN ( $H1^* - SG1^* == SG2^* - H2^*$ )
- remote access ( $H1^* == SG2^* - H2^*$ )

**Slide 14**

## Locating a Security Gateway

---

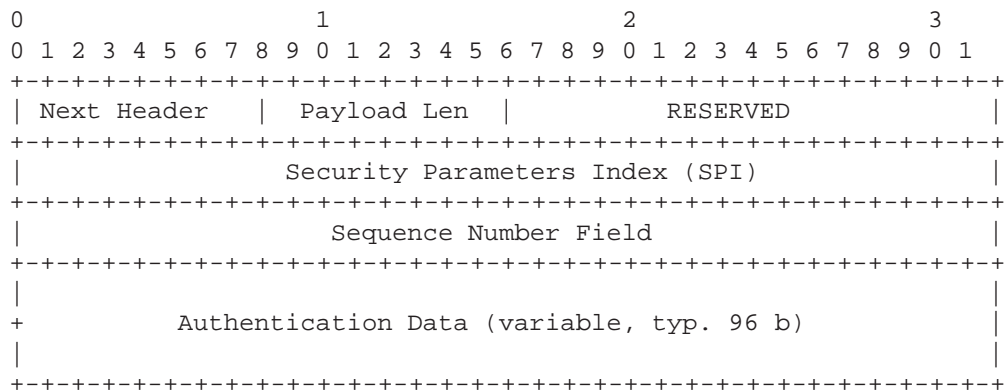
- where's the gateway? authentication?
- currently done manually
- alternatives: SLP, multicast, DHCP, ...

### Slide 15

## Authentication header (AH)

---

protocol 51:



### Slide 16



## Authentication Header: Transport Mode

---

IPv4:

```
-----
|orig IP hdr  |      |      |      |
|(any options)| AH  | TCP | Data |
-----
|<----- authenticated ----->|
except for mutable fields
```

IPv6:

```
-----
|      | hop-by-hop, dest*, |      | dest |      |
|orig IP hdr | routing, fragment. | AH | opt* | TCP | Data |
-----
|<---- authenticated except for mutable fields ----->|
```

### Slide 17

## Authentication Header: Tunnel Mode

---

IPv4:

```
-----
| new IP hdr* |      | orig IP hdr* |      |      |
|(any options)| AH  | (any options) | TCP | Data |
-----
|<- authenticated except for mutable fields -->|
|           in the new IP hdr                |
```

IPv6:

```
-----
|      | ext hdrs* |      |      | ext hdrs* |      |
|new IP hdr*|if present| AH |orig IP hdr*|if present|TCP|Data|
-----
|<-- authenticated except for mutable fields in new IP hdr ->|
```

### Slide 18



## ESP for IPv4

---

```

-----
|orig IP hdr | ESP |   |   |   | ESP | ESP |
|(any options)| Hdr | TCP | Data | Trailer | Auth|
-----
                |<----- encrypted ----->|
                |<----- authenticated ----->|

```

### Slide 21

## ESP

---

- DES in CBC mode [MD97]
- HMAC with MD5 (RFC 2104)
- HMAC with SHA-1
- NULL Authentication algorithm
- NULL Encryption algorithm

### Slide 22

## Keyed Authentication (RFC 2104)

---

- keyed MAC (message authentication codes)
- works with any iterated hash
- $\text{prf}(\text{key}, \text{msg}) = H((K \oplus \text{opad}) | H((K \oplus \text{ipad}) | \text{text}))$
- note: double hash, avoids continuation problem of  $H(K || m)$
- replace fixed IV of iterated hash by random (key) IV
- outer pad (opad) = 0x5c, ipad = 0x36 (Hamming distance!) to  $B = 64$  bytes
- may truncate hash – no less secure

### Slide 23

## Internet Key Exchange (IKE)

---

- IKE = ISAKMP + Oakley
- “negotiate and provide authenticated keying material for security associations in a protected manner”
- VPN, remote (“roaming”) user
- perfect forward secrecy (PFS): compromise of key  $\Rightarrow$  only single data item ( $\Rightarrow$  D-H)
- DOI = domain of interpretation  $\Rightarrow$  roughly, “name space” for algorithms (RFC 2407)
- ISAKMP phases, Oakley modes:

**Phase 1:** ISAKMP peers establish bidirectional secure channel using *main mode* or *aggressive mode*  $\rightarrow$  ISAKMP SA

### Slide 24

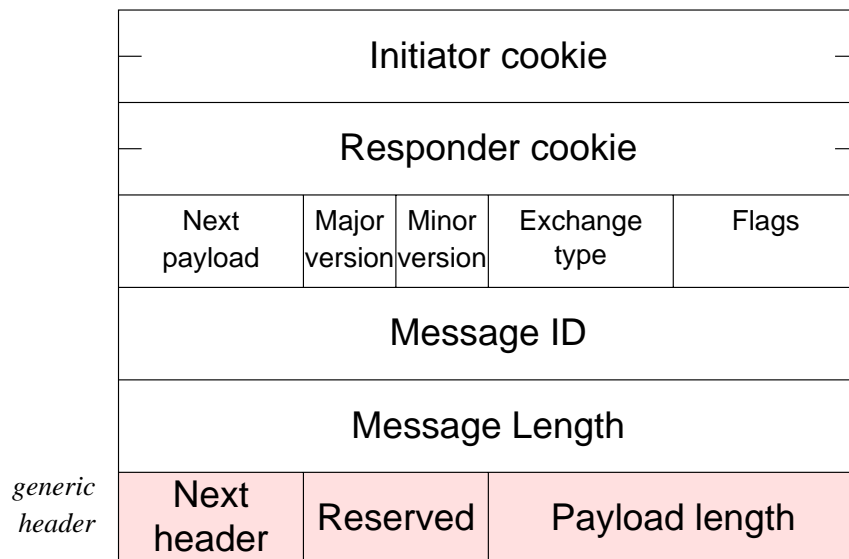
**Phase 2:** negotiation of security services for IPsec (maybe for several SAs) using *quick mode*

- can have multiple Phase 2 exchanges, e.g., to change keys

### Slide 25

## ISAKMP

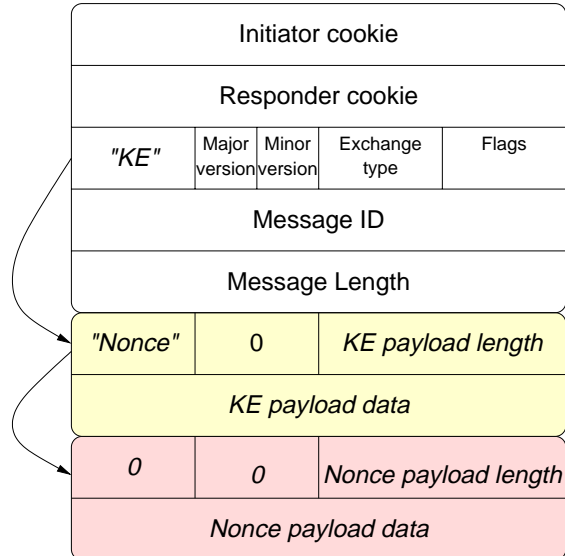
---



### Slide 26

## ISAKMP example

---



### Slide 27

## Phase 1 ISAKMP exchange

---

all based on ephemeral Diffie-Hellman exchange

**Main mode:** 6 messages = negotiate policy (2 msg.), D-H + nonces (2), authenticate D-H (2)

**Aggressive mode:** 3 messages = negotiate policy, exchange D-H public values, identities, authenticate responder (2 msg.), authenticate initiator

typically uses UDP (port 500), may use other protocols

### Slide 28

## Policy proposals

---

Allow AND (same number) and OR (different numbers); transforms are always OR

- Proposal 1 AH
  - Transform 1: HMAC-SHA
  - Transform 2: HMAC-MD5
- Proposal 2 ESP
  - Transform 1: 3DES with HMAC-SHA
- Proposal 3 ESP
  - Transform 1: 3DES with HMAC-SHA
- Proposal 3 PCP
  - Transform 1: LZS
  - Transform 2: Deflate

### Slide 29

## ISAKMP Attacks

---

**Connection hijacking:** linking authentication, key exchange, SA exchange

**Man-in-the-Middle:** linking  $\Rightarrow$  no insertion; deletion  $\Rightarrow$  no creation; reflection; modification

### Slide 30

## ISAKMP Identification

#	Operation	I-C.	R-C.	Message ID	SPI
1	Start ISAKMP SA negotiation	X	0	0	0
2	Respond ISAKMP SA negotiation	X	X	0	0
3	Init other SA negotiation	X	X	X	X
4	Respond other SA negotiation	X	X	X	X
5	Other (KE, ID, etc.)	X	X	X/0	NA
6	Security Protocol (ESP, AH)	NA	NA	NA	X

### Slide 31

## ISAKMP Message

```

                                1                               2                               3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-----+-----+-----+-----+-----+-----+-----+-----+-----+
!                               Initiator                               !
!                               Cookie                                 !
+-----+-----+-----+-----+-----+-----+-----+-----+-----+
!                               Responder                             !
!                               Cookie                                 !
+-----+-----+-----+-----+-----+-----+-----+-----+-----+
! Next Payload ! MjVer ! MnVer ! Exchange Type !           Flags      !
+-----+-----+-----+-----+-----+-----+-----+-----+-----+
!                               Message ID                             !
+-----+-----+-----+-----+-----+-----+-----+-----+-----+
!                               Length                                 !
+-----+-----+-----+-----+-----+-----+-----+-----+-----+

```

### Slide 32



## ISAKMP Payloads

---

NONE	0	Vendor ID (VID)	13
Security Association (SA)	1	RESERVED	14–127
Proposal (P)	2	Private Use	128–255
Transform (T)	3		
Key Exchange (KE)	4		
Identification (ID)	5		
Certificate (CERT)	6		
Certificate Request (CR)	7		
Hash (HASH)	8		
Signature (SIG)	9		
Nonce (NONCE)	10		
Notification (N)	11		
Delete (D)	12		

### Slide 33

## Anti-Clogging Token ("Cookie") Creation

---

- The cookie must depend on the specific parties;
  - It must not be possible for anyone other than the issuing entity to generate cookies that will be accepted by that entity.
  - The cookie generation function must be fast to thwart attacks intended to sabotage CPU resources.
- ➡ hash over the IP source and destination address, the UDP source and destination ports and a locally generated secret random value.

### Slide 34

## ISAKMP

---

- encrypted flag  $\Rightarrow$  SA(ic,rc)
- commit: done with phase, detect losses
- authentication

```

                                1                2                3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-----+-----+-----+-----+-----+-----+-----+-----+
! Next Payload !  RESERVED  !           Payload Length           !
+-----+-----+-----+-----+-----+-----+-----+-----+

```

### Slide 35

## IKE Keys

---

SKEYID =

signatures	$\text{prf}(N_i N_r, g^{xy})$	
public key	$\text{prf}(h(N_i N_r), C_i C_r)$	$C_{i,r}$ : initiator or responder cookie
pre-shared	$\text{prf}(\text{shared key}, N_i N_r)$	

### Slide 36