

# IPsec

## Protocol security - where?

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**Application layer:** (+): easy access to user credentials, extend without waiting for OS vendor, understand data; (-): design again and again; e.g., PGP, ssh, Kerberos

**Transport layer:** (+): security mostly seamlessly, but difficult to get credentials; e.g., TLS

**Network layer:** (+): reduced key management, fewer application changes, fewer implementations, VPNs; (-) non-repudiation, multi-user machines, partial security in “middle boxes”

**Data link layer:** (+): speed; (-): hop-by-hop only

## Documents

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Document Roadmap	RFC 2411
Architecture	RFC 2401
IP Authentication Header (AH)	RFC 2402
IP Authentication Using Keyed MD5	RFC 1828
IP Encapsulating Security Payload (ESP)	RFC 2406
The Oakley Key Determination Protocol	RFC 2412
Internet Sec. Assoc. and Key Mmgt. P. (ISAKMP)	RFC 2408
The Internet Key Exchange (IKE)	RFC 2409
HMAC: Keyed-Hashing for Message AuthenticationA	RFC 2104

## IPSec services

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- IPv4 and IPv6 unicast
- access control
- connectionless integrity
- data origin authentication
- protection against replays (partial sequence integrity)
- confidentiality (encryption)
- limited traffic flow confidentiality.
- todo: NAT, multicast

## Architecture

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**Authentication header (AH):** access control, integrity, data origin authentication, replay protection

**Encapsulating Security Payload (ESP):** access control, confidentiality, traffic flow confidentiality.

**Key management protocols:** IKE = OAKLEY + ISAKMP, ...

- for any upper-layer protocol
- no effect on rest of Internet
- algorithm-independent, but default algorithms

## Architecture

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- between host and/or security gateways
- security gateway = router, firewall, ...
- security policy database (SPD)  $\Rightarrow$  IPsec, discarded, or bypass
- negotiate compression (why?)
- *tunnel mode or transport mode*
- granularity: single host-host tunnel vs. one per TCP connection

## Implementation

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- native IP implementation
- bump in the stack (BITS): beneath IP layer
- bump in the wire (BITW)

## Security Association (SA)

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- simplex
- AH *or* ESP
- identified by
  - Security Parameter Index (SPI),
  - IP destination address,
  - security protocol (AH or ESP) identifier.
- transport mode: two hosts
  - AH or ESP after IPv4 options, before UDP/TCP
  - IPv6: after base header and extensions, before/after destination options
  - mostly for higher-layer protocols (but: AH also some IP header parts)
- tunnel mode: one or two security gateways



- outer header → tunnel endpoint
- security header between outer and inner
- traffic hiding; ESP payload padding

## Nested Security Associations

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AH *and* ESP  $\Rightarrow$  two SAs (“SA bundle”):

- transport adjacency: AH, then ESP
- both tunnel endpoints the same
- one endpoint the same
- neither the same

## Security Policy Database

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- map to Security Association Database (per packet or per SPD entry)
- discard, bypass or apply to *inbound* or *outbound*
- ordered list of filters (stateless firewall)
- example: “use ESP in transport mode using 3DES-CBC with explicit IV, nested inside of AH in tunnel mode using HMAC-SHA-1.”
- selectors:
  - destination IP address: address, range, address + mask, wildcard
  - source IP address
  - name (for BITS/BITW hosts): user id, X.500 DN, system name, opaque, ...
  - data sensitivity label
  - transport layer protocol

- source/destination ports
- per socket setup or per packet (BITS, BITW, gateway)

## Security Association Database (SAD)

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- inbound: outer destination address
- IPsec protocol (AH or ESP)
- SPI (32-bit value)

## Examples of Implementations

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- end-to-end security ( $H1^* \text{ == } H2^*$ )
- VPN ( $H1 - SG1^* \text{ == } SG2^* - H2$ )
- e2e + VPN ( $H1^* - SG1^* \text{ == } SG2^* - H2^*$ )
- remote access ( $H1^* \text{ == } SG2^* - H2^*$ )

## Locating a Security Gateway

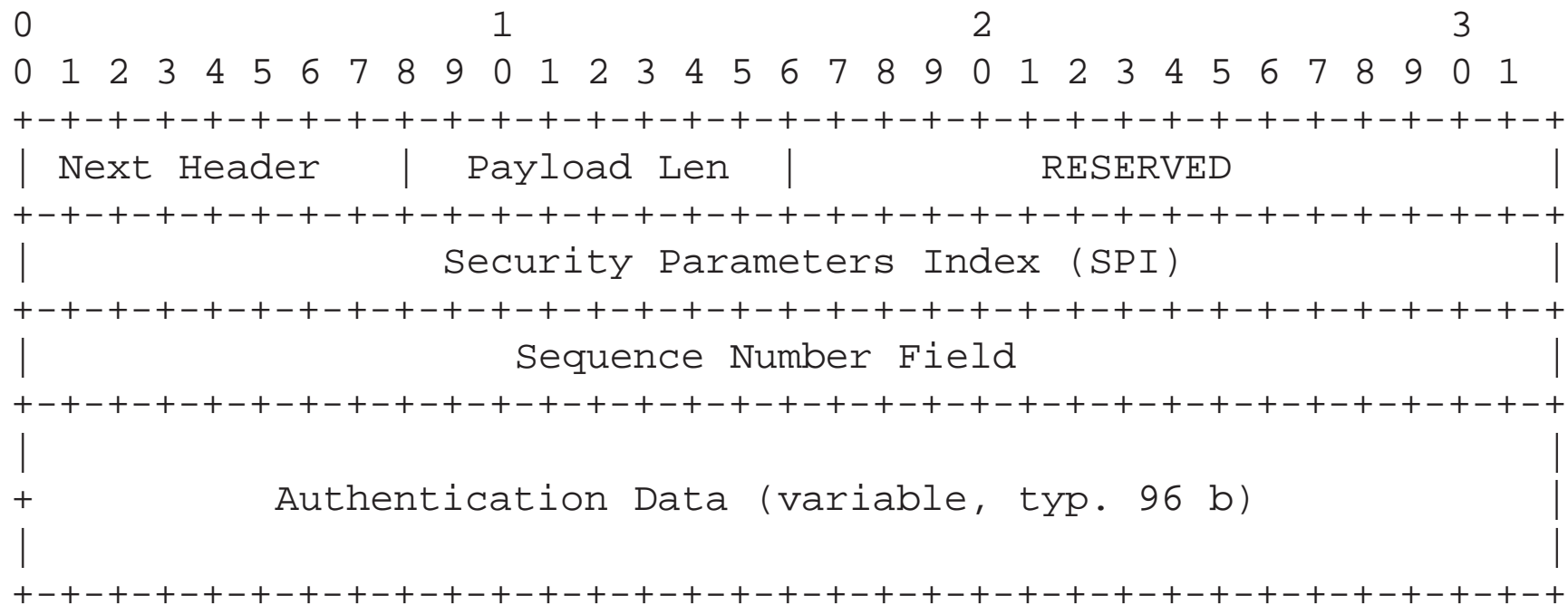
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- where's the gateway? authentication?
- currently done manually
- alternatives: SLP, multicast, DHCP, ...

## Authentication header (AH)

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protocol 51:





## Authentication Header: Transport Mode

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IPv4:

```
-----
|orig IP hdr  |      |      |      |
|(any options)| AH  | TCP  | Data |
-----
```

```
|<----- authenticated ----->|
except for mutable fields
```

IPv6:

```
-----
|             |hop-by-hop, dest*, |      | dest |      |      |
|orig IP hdr  |routing, fragment. | AH  | opt* | TCP  | Data |
-----
```

```
|<---- authenticated except for mutable fields ----->|
```

## Authentication Header: Tunnel Mode

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### IPv4:

```

-----
| new IP hdr* |      | orig IP hdr* |      |      |
| (any options) | AH | (any options) | TCP | Data |
-----
| <- authenticated except for mutable fields --> |
|           in the new IP hdr                    |

```

### IPv6:

```

-----
|           | ext hdrs* |      |           | ext hdrs* |      |
| new IP hdr* | if present | AH | orig IP hdr* | if present | TCP | Data |
-----
| <-- authenticated except for mutable fields in new IP hdr -> |

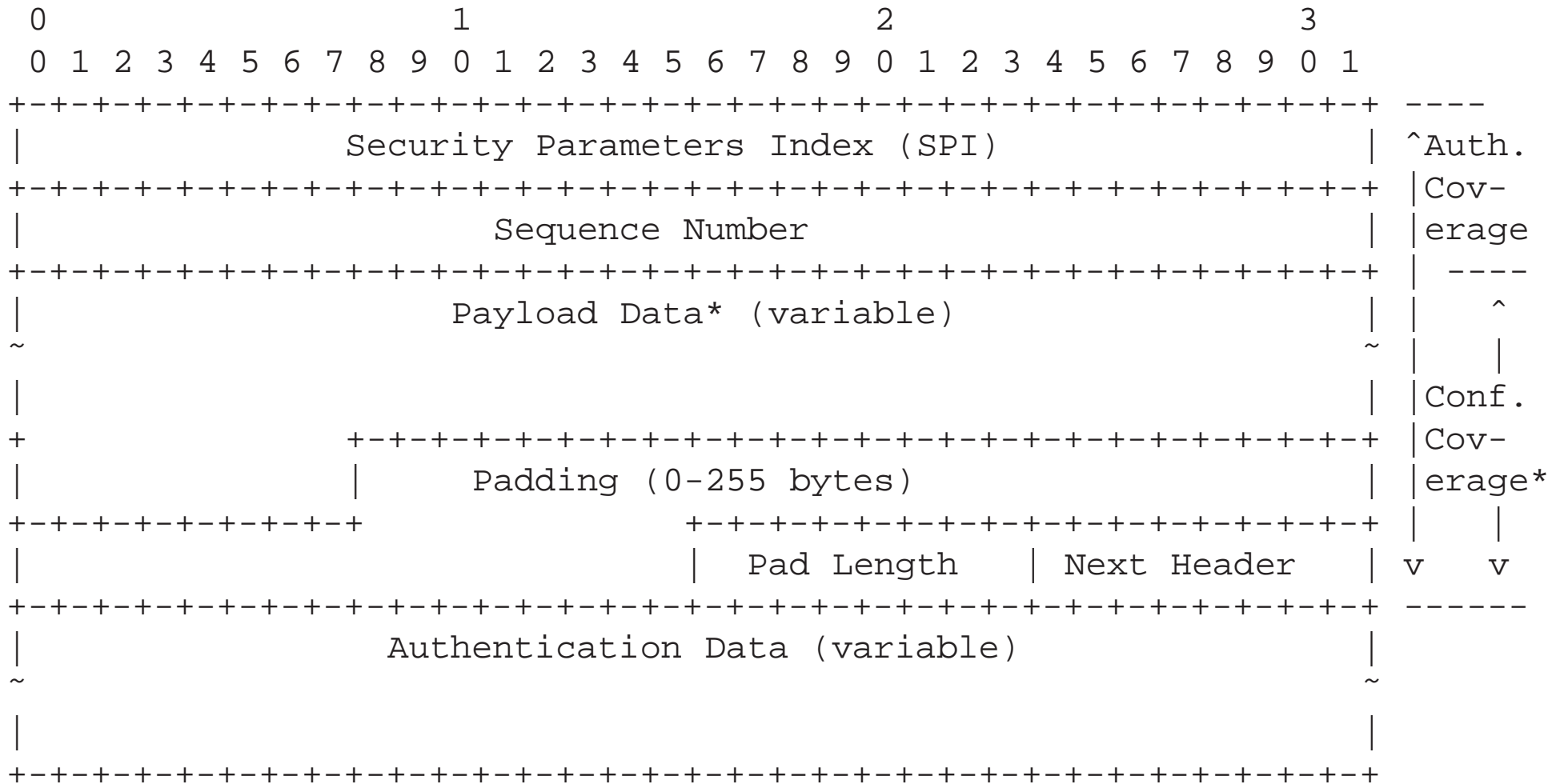
```

## Authentication

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- replay prevention: if seq. no. cycles, new SA; sliding window  $\Rightarrow$  reject lower than left window edge
- immutable or predictable IP header fields: version, IH length, total length, identification, protocol, source, destination (source route  $\Rightarrow$  predictable)
- set mutable fields to zero: TOS, flags, fragment, TTL, header checksum
- AH header, with zero ICV
- upper-layer data

# Encapsulating Security Payload (ESP)



## ESP for IPv4

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```

-----
|orig IP hdr | ESP |   |   |   ESP | ESP |
|(any options)| Hdr | TCP | Data | Trailer |Auth|
-----

```

```

                |<----- encrypted ----->|
|<----- authenticated ----->|

```

## ESP

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- DES in CBC mode [MD97]
- HMAC with MD5 (RFC 2104)
- HMAC with SHA-1
- NULL Authentication algorithm
- NULL Encryption algorithm

## Keyed Authentication (RFC 2104)

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- keyed MAC (message authentication codes)
- works with any iterated hash
- $\text{prf}(\text{key}, \text{msg}) = H((K \oplus \text{opad}) | H((K \oplus \text{ipad}) | \text{text}))$
- note: double hash, avoids continuation problem of  $H(K \text{---} m)$
- replace fixed IV of iterated hash by random (key) IV
- outer pad (opad) = 0x5c, ipad = 0x36 (Hamming distance!) to  $B = 64$  bytes
- may truncate hash – no less secure

## Internet Key Exchange (IKE)

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- IKE = ISAKMP + Oakley
- “negotiate and provide authenticated keying material for security associations in a protected manner”
- VPN, remote (“roaming”) user
- perfect forward secrecy (PFS): compromise of key  $\Rightarrow$  only single data item ( $\Rightarrow$  D-H)
- DOI = domain of interpretation  $\Rightarrow$  roughly, “name space” for algorithms (RFC 2407)
- ISAKMP phases, Oakley modes:

**Phase 1:** ISAKMP peers establish bidirectional secure channel using *main mode* or *aggressive mode*  $\rightarrow$  ISAKMP SA



**Phase 2:** negotiation of security services for IPsec (maybe for several SAs) using *quick mode*

- can have multiple Phase 2 exchanges, e.g., to change keys

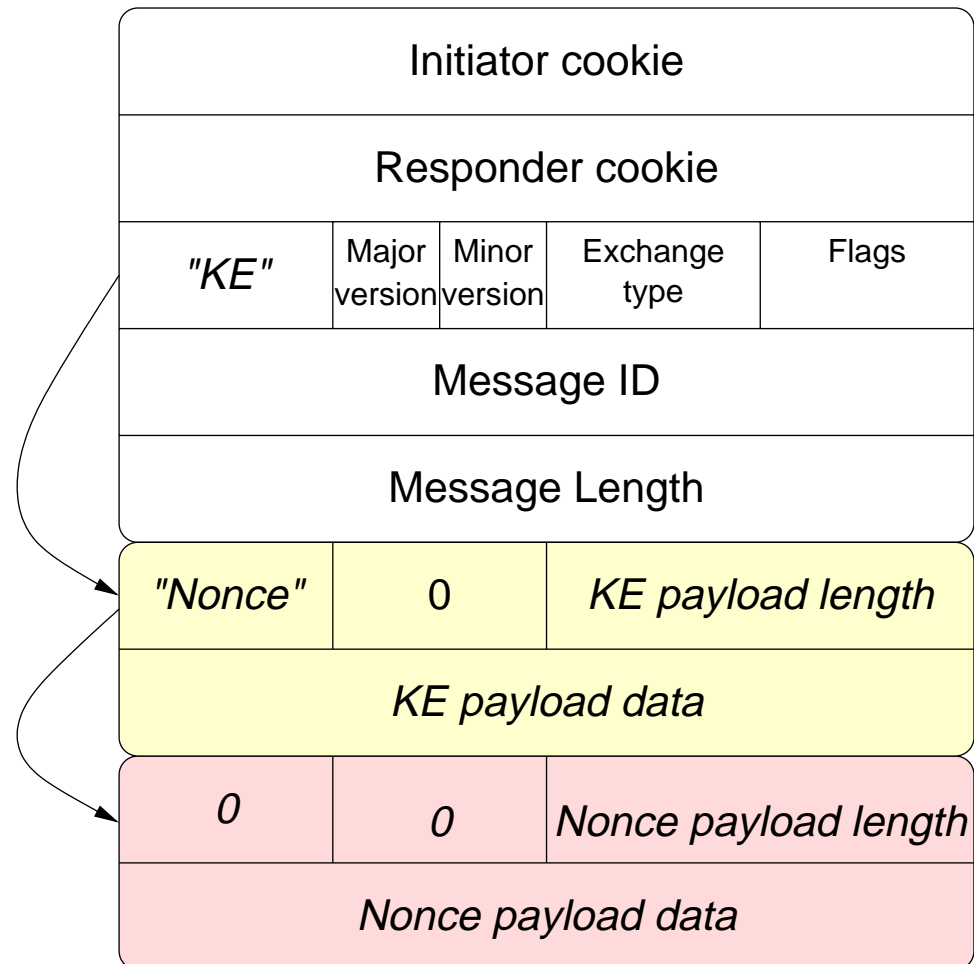
# ISAKMP

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Initiator cookie				
Responder cookie				
Next payload	Major version	Minor version	Exchange type	Flags
Message ID				
Message Length				
<i>generic header</i>	Next header	Reserved		Payload length

# ISAKMP example

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## Phase 1 ISAKMP exchange

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all based on ephemeral Diffie-Hellman exchange

**Main mode:** 6 messages = negotiate policy (2 msg.), D-H + nonces (2), authenticate D-H (2)

**Aggressive mode:** 3 messages = negotiate policy, exchange D-H public values, identities, authenticate responder (2 msg.), authenticate initiator

typically uses UDP (port 500), may use other protocols

## Policy proposals

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Allow AND (same number) and OR (different numbers); transforms are always OR

Proposal 1 AH

Transform 1: HMAC-SHA

Transform 2: HMAC-MD5

Proposal 2 ESP

Transform 1: 3DES with HMAC-SHA

Proposal 3 ESP

Transform 1: 3DES with HMAC-SHA

Proposal 3 PCP

Transform 1: LZS

Transform 2: Deflate

## ISAKMP Attacks

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**Connection hijacking:** linking authentication, key exchange, SA exchange

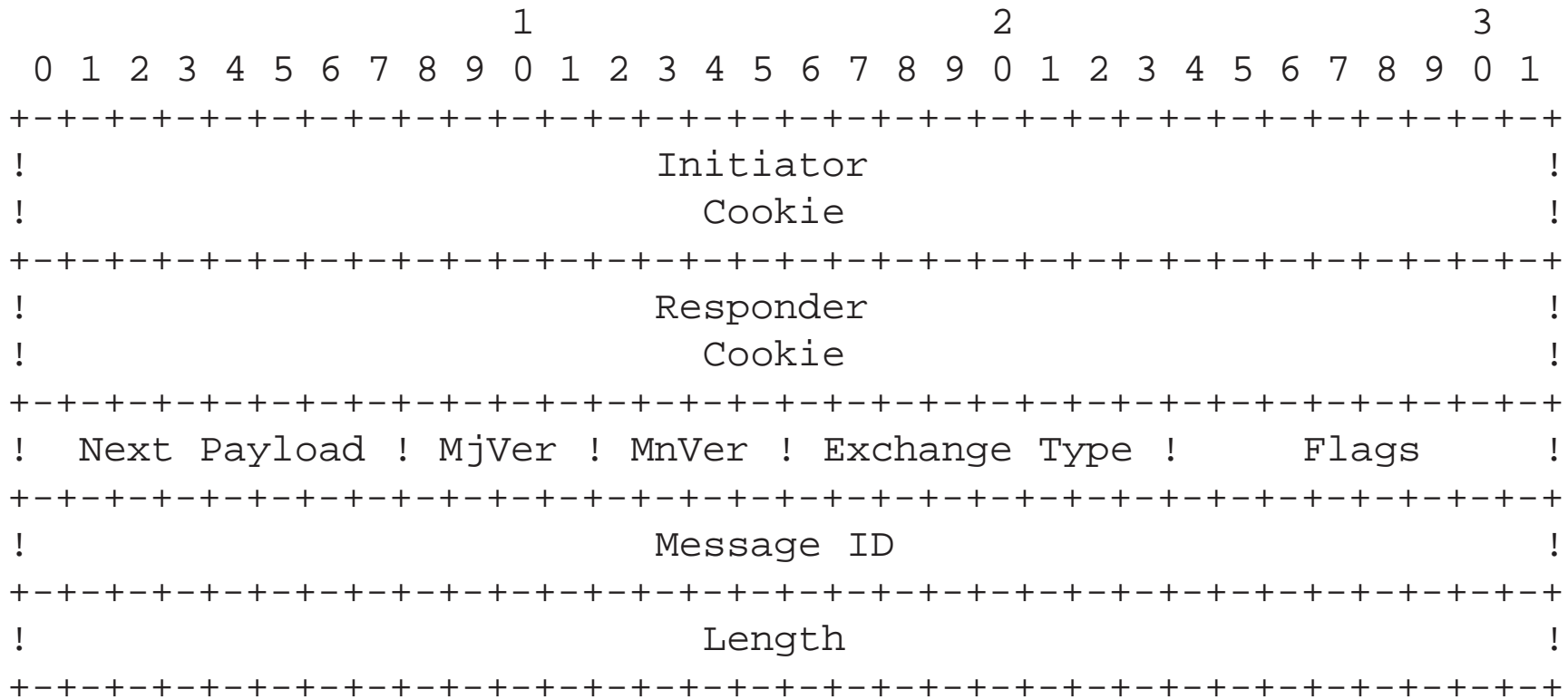
**Man-in-the-Middle:** linking  $\Rightarrow$  no insertion; deletion  $\Rightarrow$  no creation; reflection; modification

## ISAKMP Identification

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#	Operation	I-C.	R-C.	Message ID	SPI
1	Start ISAKMP SA negotiation	X	0	0	0
2	Respond ISAKMP SA negotiation	X	X	0	0
3	Init other SA negotiation	X	X	X	X
4	Respond other SA negotiation	X	X	X	X
5	Other (KE, ID, etc.)	X	X	X/0	NA
6	Security Protocol (ESP, AH)	NA	NA	NA	X

# ISAKMP Message





## ISAKMP Payloads

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NONE	0	Vendor ID (VID)	13
Security Association (SA)	1	RESERVED	14–127
Proposal (P)	2	Private Use	128–255
Transform (T)	3		
Key Exchange (KE)	4		
Identification (ID)	5		
Certificate (CERT)	6		
Certificate Request (CR)	7		
Hash (HASH)	8		
Signature (SIG)	9		
Nonce (NONCE)	10		
Notification (N)	11		
Delete (D)	12		

## Anti-Clogging Token ("Cookie") Creation

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- The cookie must depend on the specific parties;
- It must not be possible for anyone other than the issuing entity to generate cookies that will be accepted by that entity.
- The cookie generation function must be fast to thwart attacks intended to sabotage CPU resources.

▣ hash over the IP source and destination address, the UDP source and destination ports and a locally generated secret random value.

## ISAKMP

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- encrypted flag  $\Rightarrow$  SA(ic,rc)
- commit: done with phase, detect losses
- authentication

```

                                1                                2                                3
0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+
! Next Payload !   RESERVED   !           Payload Length           !
+-+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+--+

```

## IKE Keys

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SKEYID =

signatures  $\text{prf}(N_i|N_r, g^{xy})$

public key  $\text{prf}(h(N_i|N_r), C_i|C_r)$   $C_{i,r}$ : initiator or responder cookie

pre-shared  $\text{prf}(\text{shared key}, N_i|N_r)$