

Internet Telephony: More than just re-inventing the telephone

Henning Schulzrinne
Dept. of Computer Science
Columbia University
New York, New York
schulzrinne@cs.columbia.edu

University of Maryland

December 2, 1999

(Joint work with Ellen Hahne, Jonathan Lennox, Gautam Nair, Maria Papadopouli, Ping Pan, Lili Qiu, Jonathan Rosenberg, Xin Wang, Elin Wedlund, and Jianqi Yin)

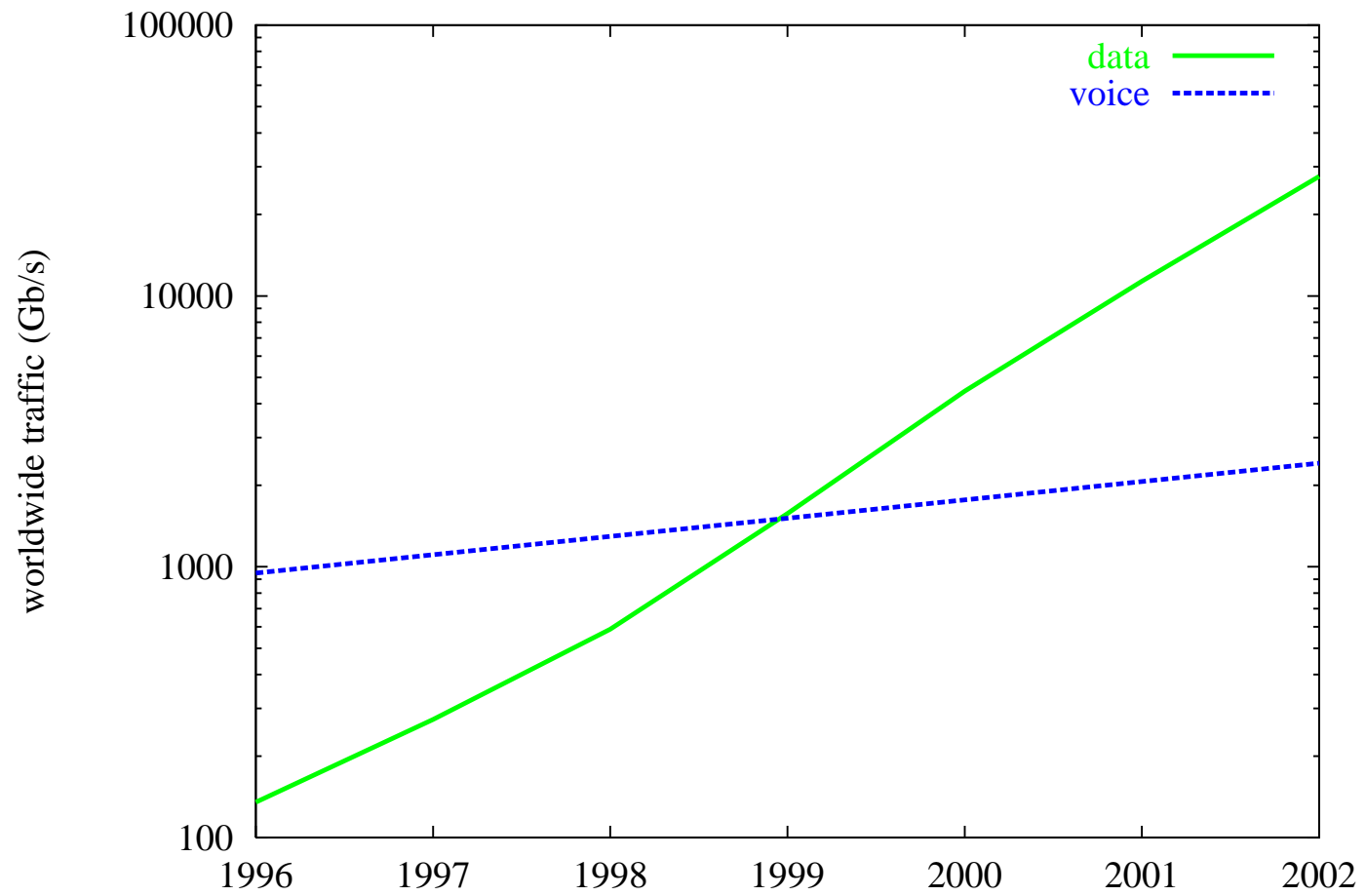
Overview

- Internet telephony: motivation and problems
- protocol architecture
- quality of service:
 - light-weight resource reservation
 - forward error control
- services \Rightarrow signaling
- the “programmable” telephone
- Internet telephony “appliances”
- mobile services

Historical perspective

- 1876 invention of telephone
- 1915 first transcontinental telephone (NY–SF)
- 1920's first automatic switches
- 1956 TAT-1 transatlantic cable (35 lines)
- 1962 digital transmission (T1)
- 1965 1ESS analog switch
- 1977 4ESS digital switch
- 1980s Signaling System #7 (out-of-band)
- 1990s Advanced Intelligent Network (AIN) services deployed

Data vs. Voice Traffic



The phone works — why bother with VoIP?

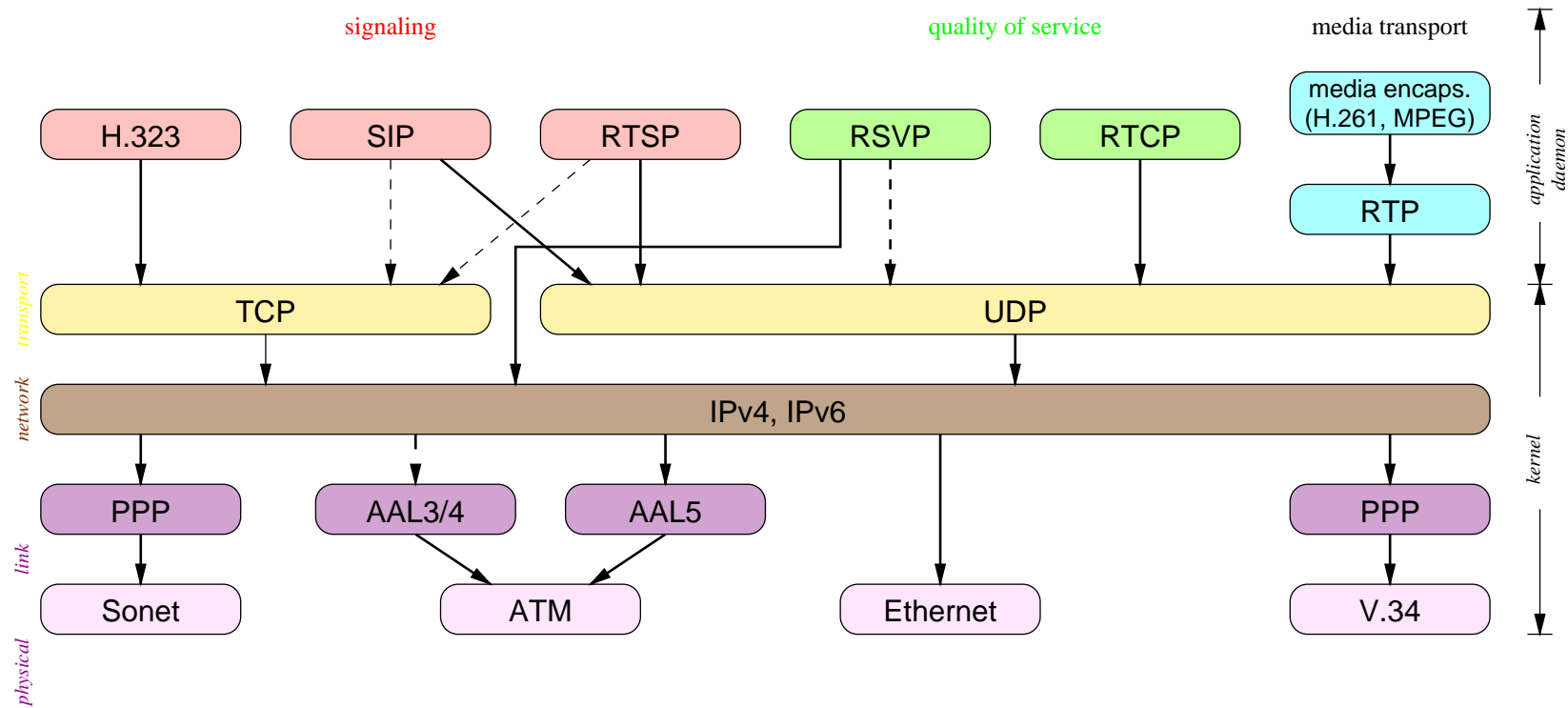
user perspective

- variable compression: tin can to broadcast quality
- security through encryption
- caller, talker identification
- better user interface (browser)
- internat. calls: TAT transatlantic cable = \$0.03/hr
- no local access fees (3.4c)
- easy: video, whiteboard, ...

carrier perspective

- silence suppression \Rightarrow traffic \downarrow
- shared facilities \Rightarrow management, redundancy
- advanced services (email/web integration)
- cheaper switching (\$0.005 vs. \$5/kb/s)
- fax uses 9.6 kb/s of 64 kb/s line

Internet multimedia protocol stack

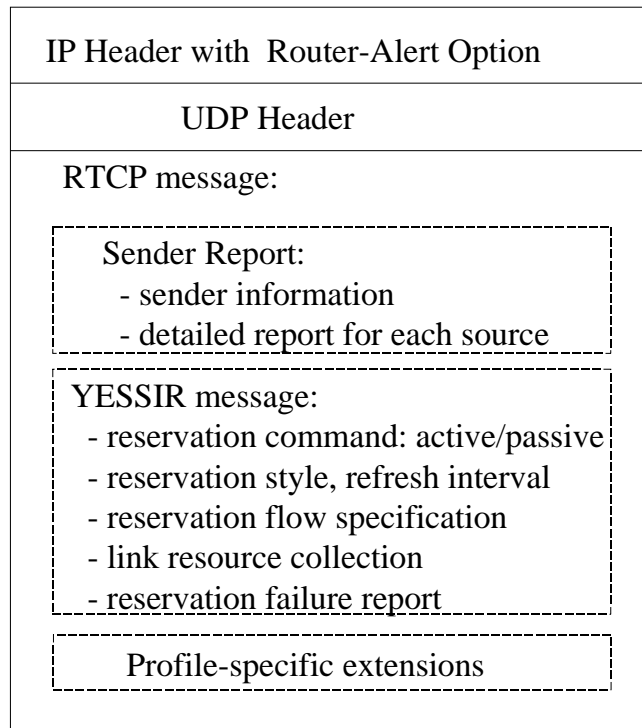


YESSIR: Yet another Sender Session Internet Reservation

- RSVP: separate daemon, API
- \Rightarrow integrate into application that needs it (embedded systems!)
- in-band \Rightarrow easier firewall
- RTP: common data transport protocol for audio/video
- router alert option in RTCP packets
- resource demands: payload type, measurement, flow specs, ...
- soft-state + RTCP BYE
- partial reservations: add links as session ages \leftrightarrow fragmentation

YESSIR

plain RTCP, SRs or additional information:



end-to-end refresh (vs. hop-by-hop)

RSVP and YESSIR performance

	setup	refresh
	<i>μs</i>	
RSVP	1,105	624
YESSIR	356	344

Charging for Multimedia Services

- service models:

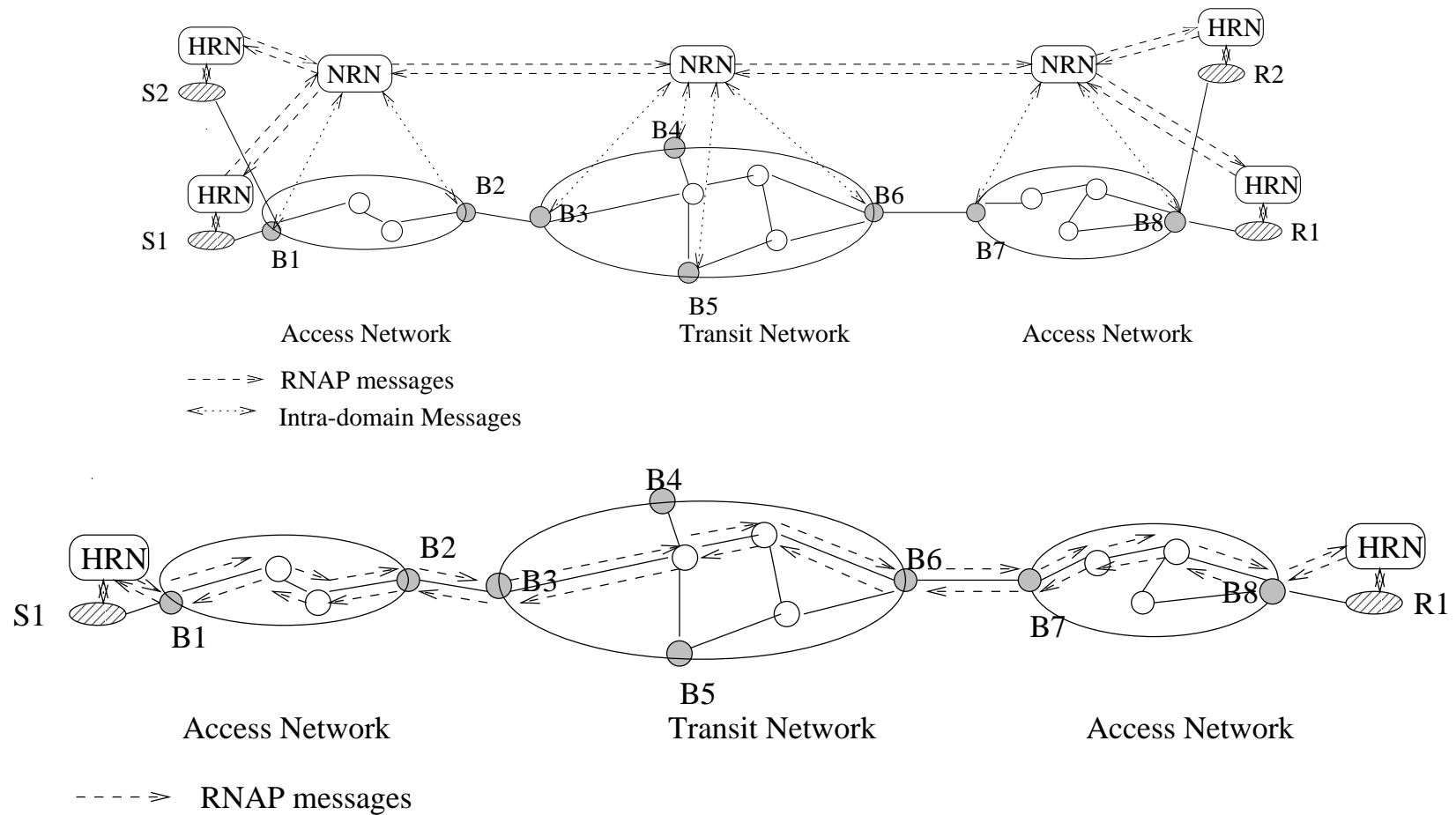
best-effort: not predictable, all drop below threshold

adaptive, “TCP-friendly”: no incentive to adjust

reserved: long-lived connections \longrightarrow blocking \uparrow

thus, we define *adaptive reserved services* with pricing incentives

RNAP: Architecture



RNAP: Pricing

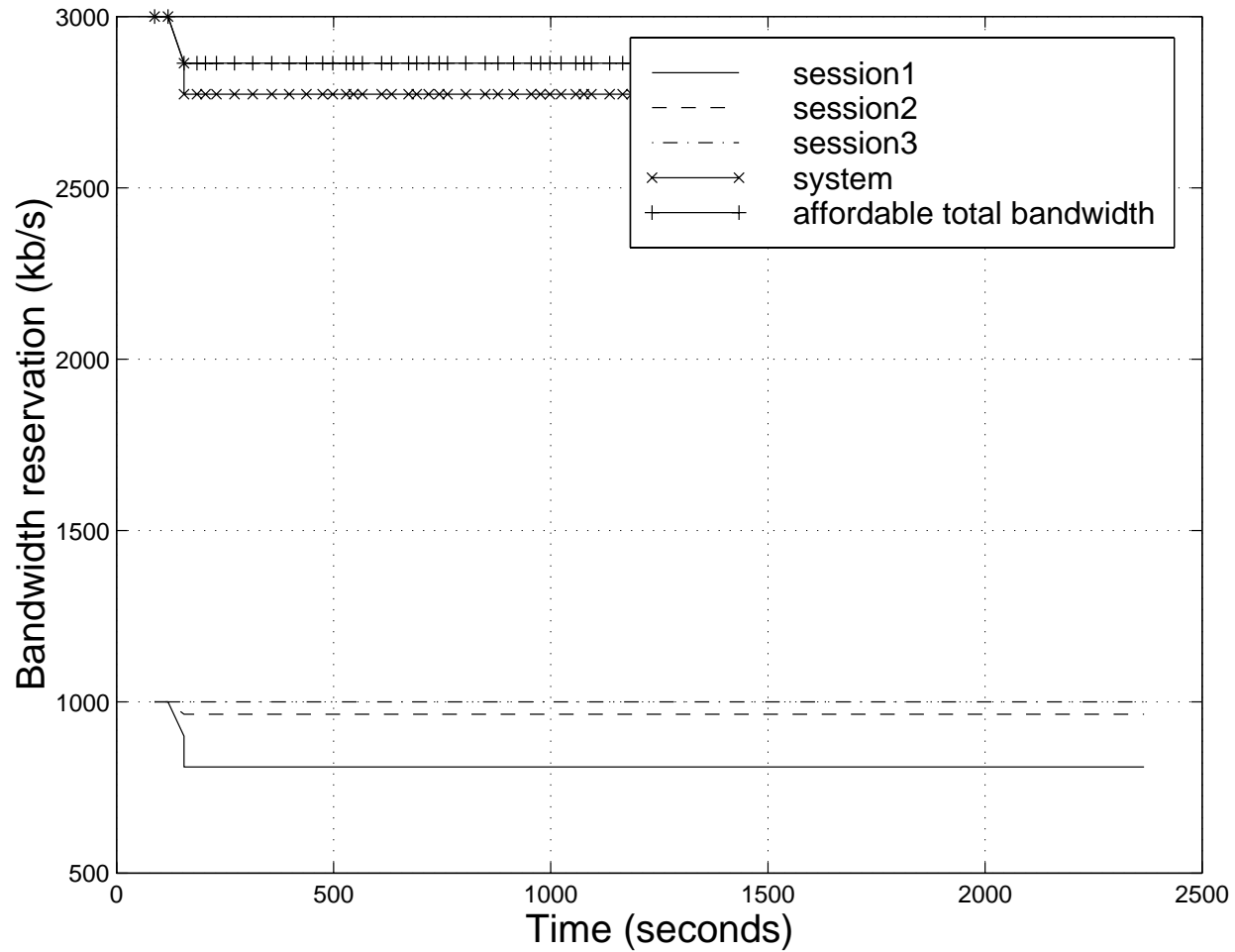
holding cost: opportunity cost for holding resources (= price of next-lower quality level)

usage cost: infrastructure cost amortized over usage

congestion cost: discouragement mechanism

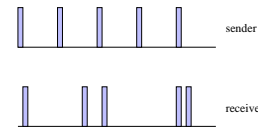
RNAP

Adjustment with different utility functions:



Integrating packet FEC into adaptive voice playout buffers

- playout buffer: trade loss (2...20%) for delay (50...500 ms)

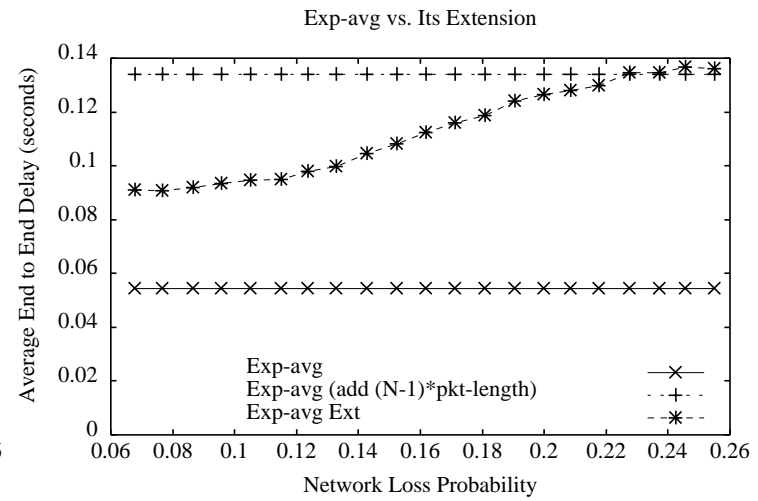
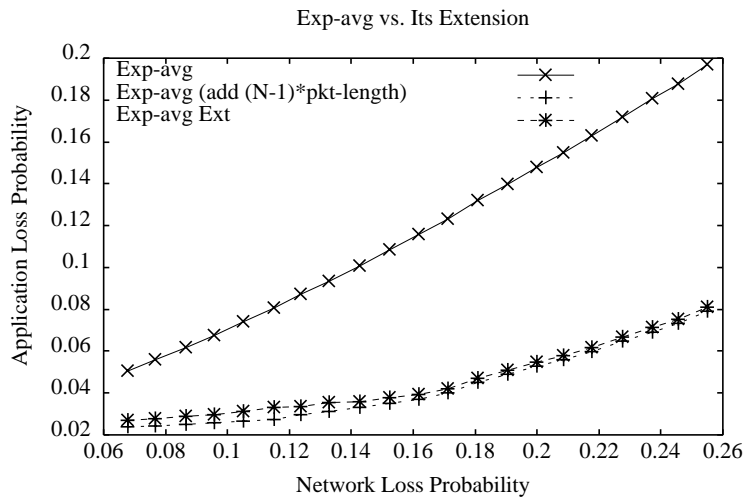


- (n, k) FEC: add $n - k$ additional packets for total of n
- algorithms:
 - exponential average and fast exp-avg
 - minimum delay
 - spike delays
 - window: spike mode + q^{th} quantile

FEC: virtual delay

- virtual delay = $\min(\text{arrival}, \text{recovery}) - \text{departure}$
- playout delay $\approx \alpha \cdot \sigma$
- if loss $<$ target loss, $\alpha \leftarrow \alpha + \delta$
- recover lost and **late** packets
- 20% loss: application loss/5, delay * 2

FEC loss and delay



New playout algorithm: delayed optimal

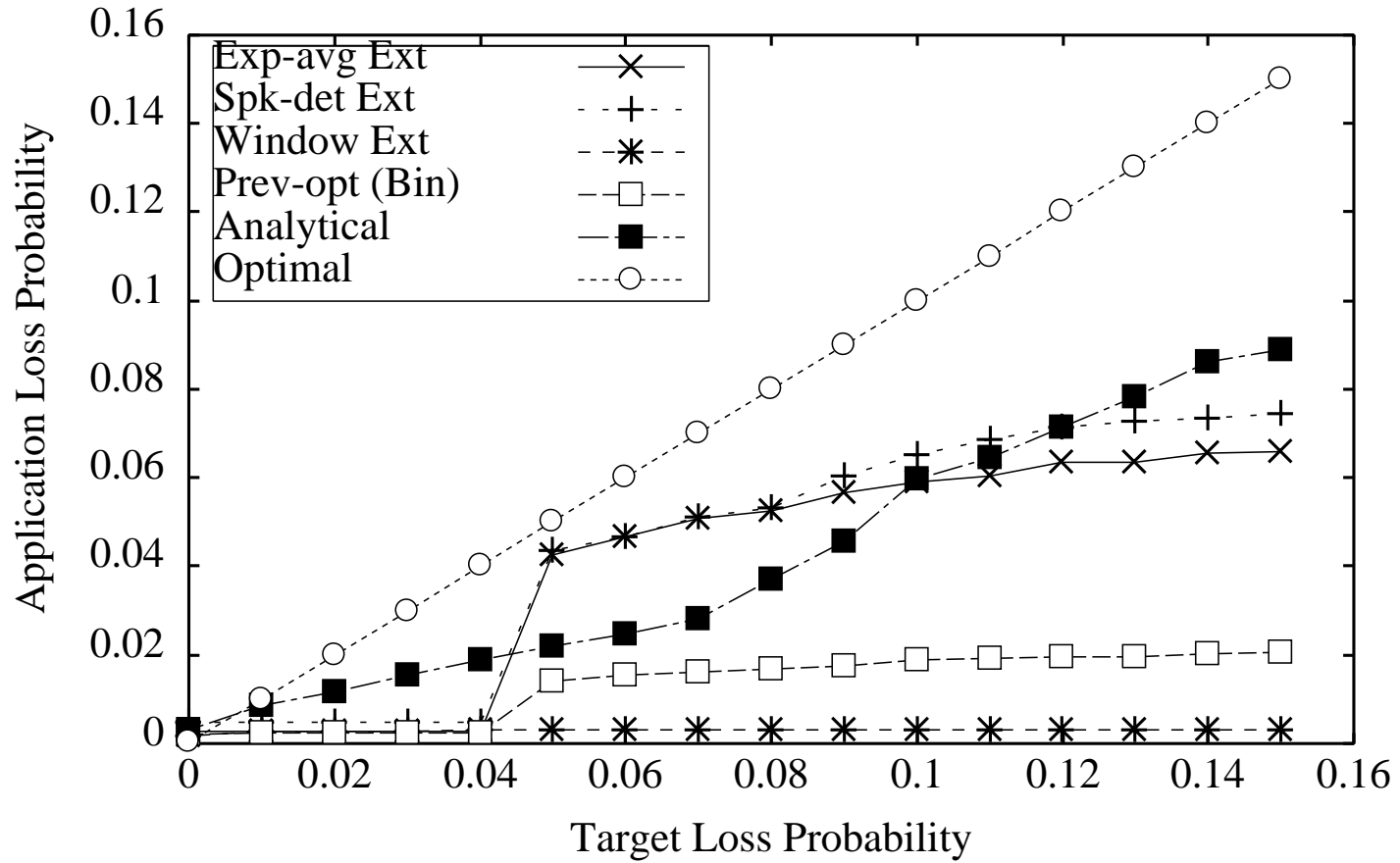
- after talkspurt, one knows optimal delay
- use optimal combination for next talkspurt
- $D_i = \alpha D_{i-1} + (1 - \alpha) D_{\text{opt}}$
- user perception function: minimal delay that achieves loss target

New playout algorithm: Analytical

- assume independent loss p , delayed randomly d
- compute playout probability
- tabulate delay distribution histogram
- at end of talkspurt, find matching playout delay for target loss rate

Playout algorithm comparison

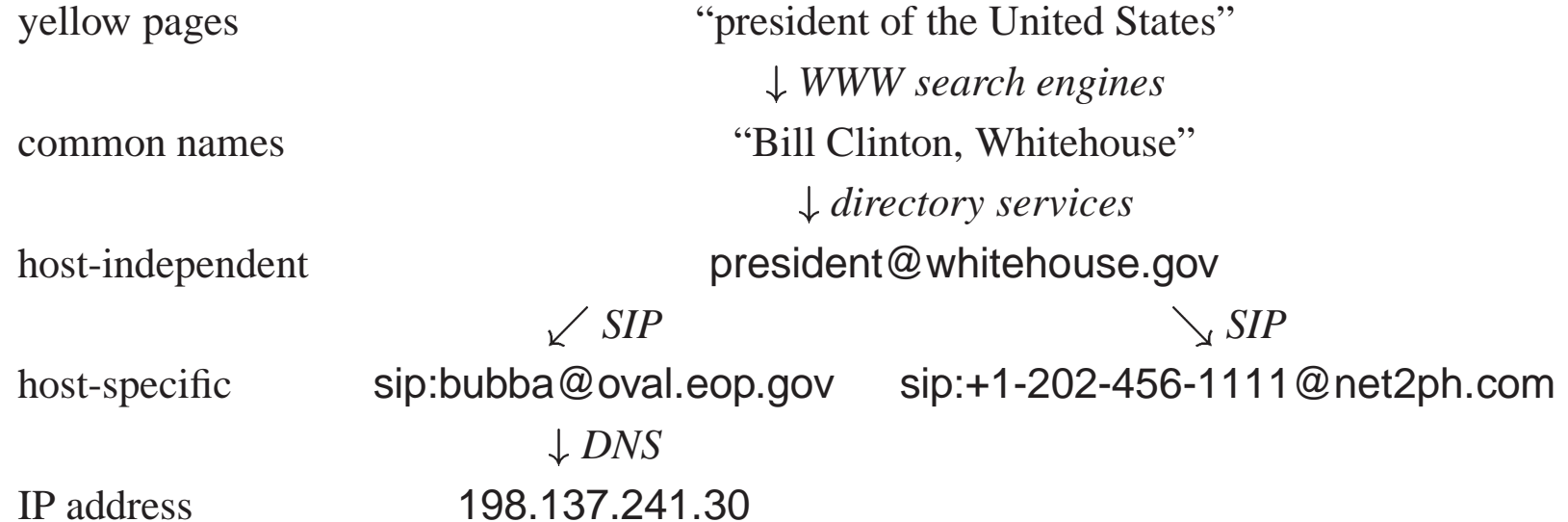
Trace 1



SIP: Session Initiation Protocol

- call user
- re-negotiate call parameters
- manual and automatic forwarding
- call center: reach first (load distribution) or reach all (department conference)
- *personal mobility* (complements data link/IP mobility) ⇒ change of terminal (PC, digital cordless, palmtop), location
- “forking” of calls: one person, multiple locations
- terminate and transfer calls
- web security, cookies

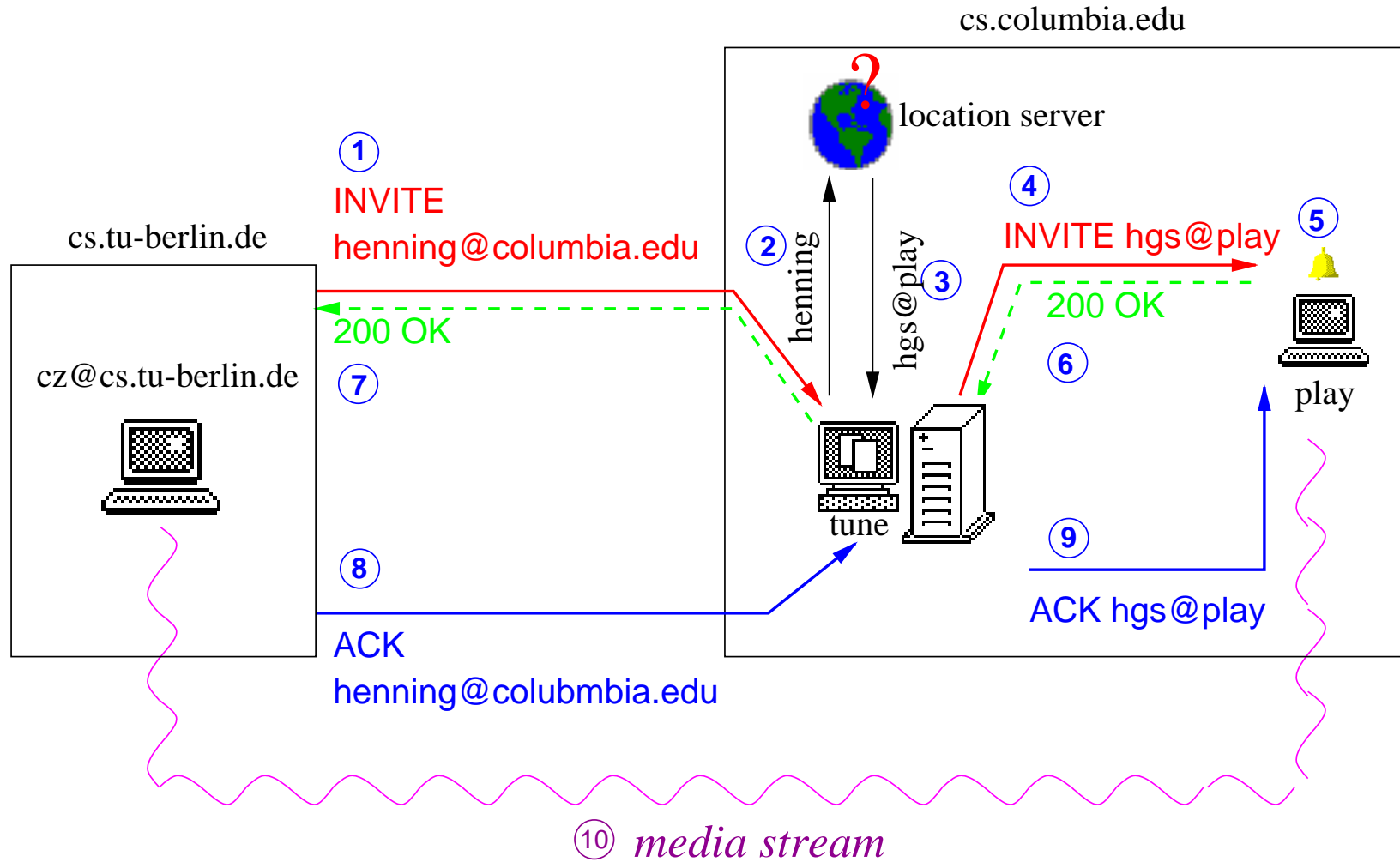
SIP addresses food chain



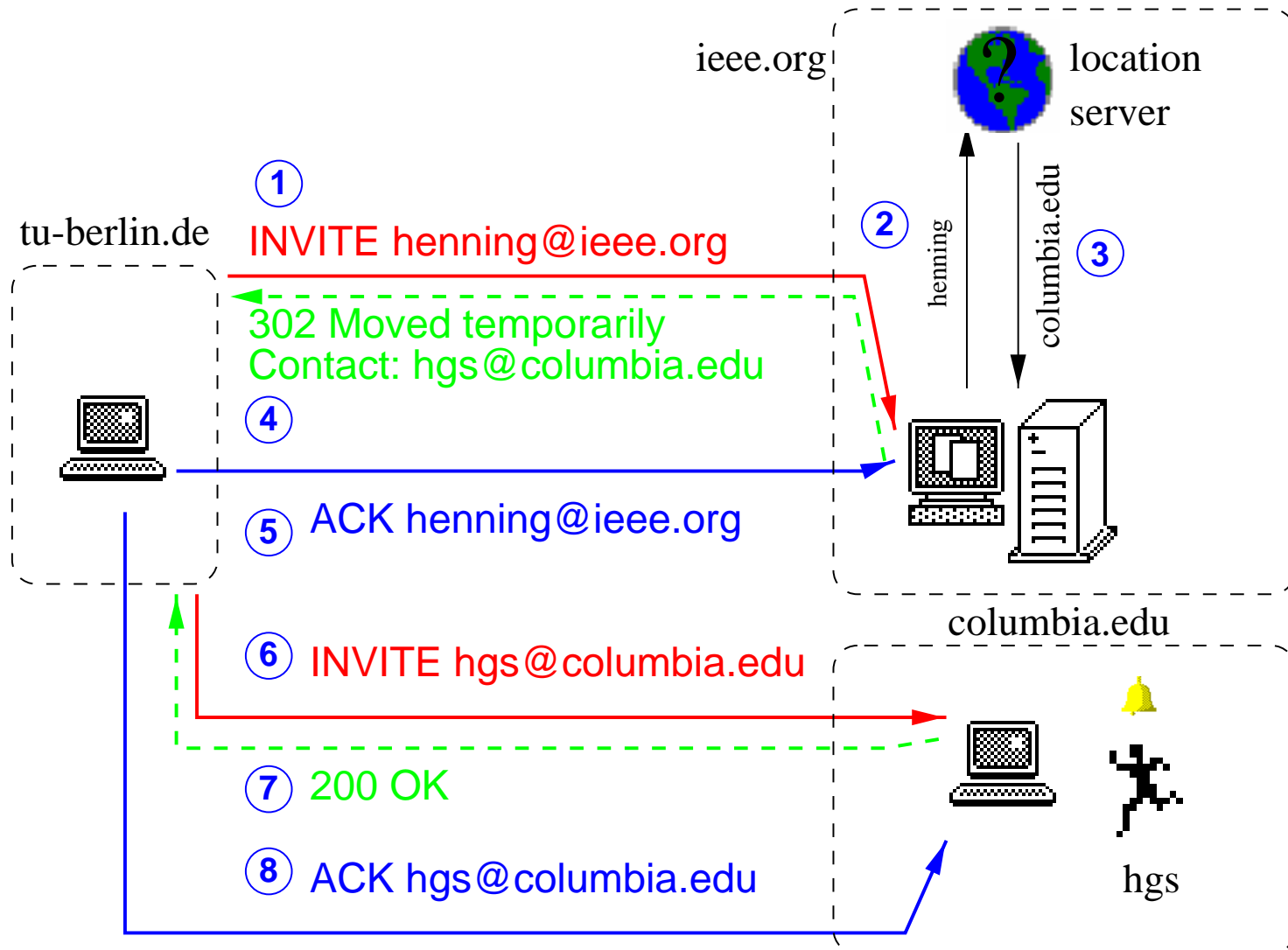
SIP: basic operation

1. use directory service (e.g., LDAP) to map name to *user@domain*
2. locate SIP servers using DNS SRV, CNAME, A
3. called server may map name to *user@host*
4. callee accepts, rejects, forward (→ new address)
5. if new address, go to step 2
6. if accept, caller confirms
7. ... conversation ...
8. caller or callee sends BYE

SIP operation in proxy mode



SIP operation in redirect mode



SIP protocol design: robustness

SIP is designed to be robust against server failures:

- no state in proxy servers during call
- responses are “self-routing”
- subsequent requests and retransmissions can take different path (backup server)
- proxy servers can “lose memory” any time \Rightarrow still function
- UDP \Rightarrow less state than TCP, no time-wait

Invitation modes

invitation		conference
	unicast	multicast
<hr/>		
unicast	telephony	Internet TV session
multicast	reach first	dept. conference

SIP user location

- SIP is independent of mechanism to locate user
- examples:
 - local multicast of invitation
 - login-based via NFS
 - recursive “finger”-traversal
 - name translation: *Alexander.G.Bell* \mapsto *agb*
 - active badges
 - SIP:
 - * REGISTER announces location, with time limit
 - * REGISTER + Location sets new location
 - * forwarding within host (\neq standard port)

Implementations

- 33 vendors at December 1999 SIP “bake-off”
- Columbia sipd:
 - registration via unicast and multicast
 - handles mailing lists (ug-students@cs), ambiguous names (lee@cs)
 - maps names (b.clinton@whitehouse)
 - Apache (httpd)-style configuration and logging
 - “basic” authentication
 - how many servers for 2300 requests/second?

Signaling ← event notification

- call queueing ... buddy lists ... event notification
- SUBSCRIBE to events
- server NOTIFY
- can use forking
- handle subscriptions using CPL
- transition to multicast if large group of subscribers

Programmable phone service

- “caller proposes, callee disposes, administrator decides”
- web = static pages → cgi-bin → Java
- “if somebody is trying to call for the 3rd time, allow mobile”
- “try office and lab in parallel, if that fails, try home”
- “allow call to mobile if I’ve talked to person before”
- “if on telemarketing list, forward to dial-a-joke”
- phone: CTI = complex, not generally for end users
- “cgi-bin” for Internet telephones: generate requests, proxy, responses
 - ▣▶ sip-cgi, complete control

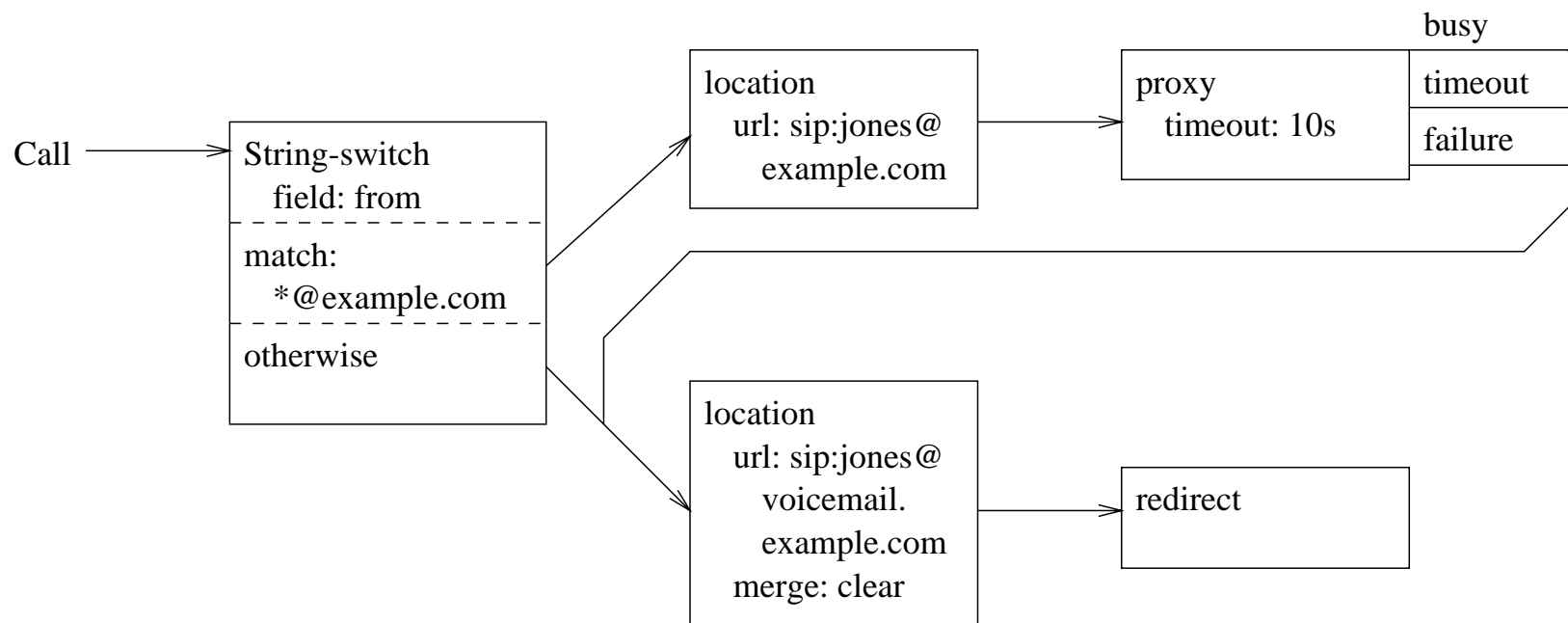
“Active Phone Networks”

language:

- don't want Turing-complete language
- fail safe: make phone calls even if crashes
- predictable resource consumption
- hide parallelism (searches)
- hide timers
- execute in callee's proxy server or end system (or phone button)

⇒ CPL, an XML-based language

CPL example



CPL example


```
<call>
  <location url="sip:jones@jonespc.example.com">
    <proxy timeout="8s">
      <busy>
        <location url="sip:jones@voicemail.example.com" merge="clear"
          id="voicemail" >
          <proxy />
        </location>
      </busy>
      <noanswer>
        <link ref="voicemail" />
      </noanswer>
    </proxy>
  </location>
</call>
```

e*phone, an Internet phone “appliance”

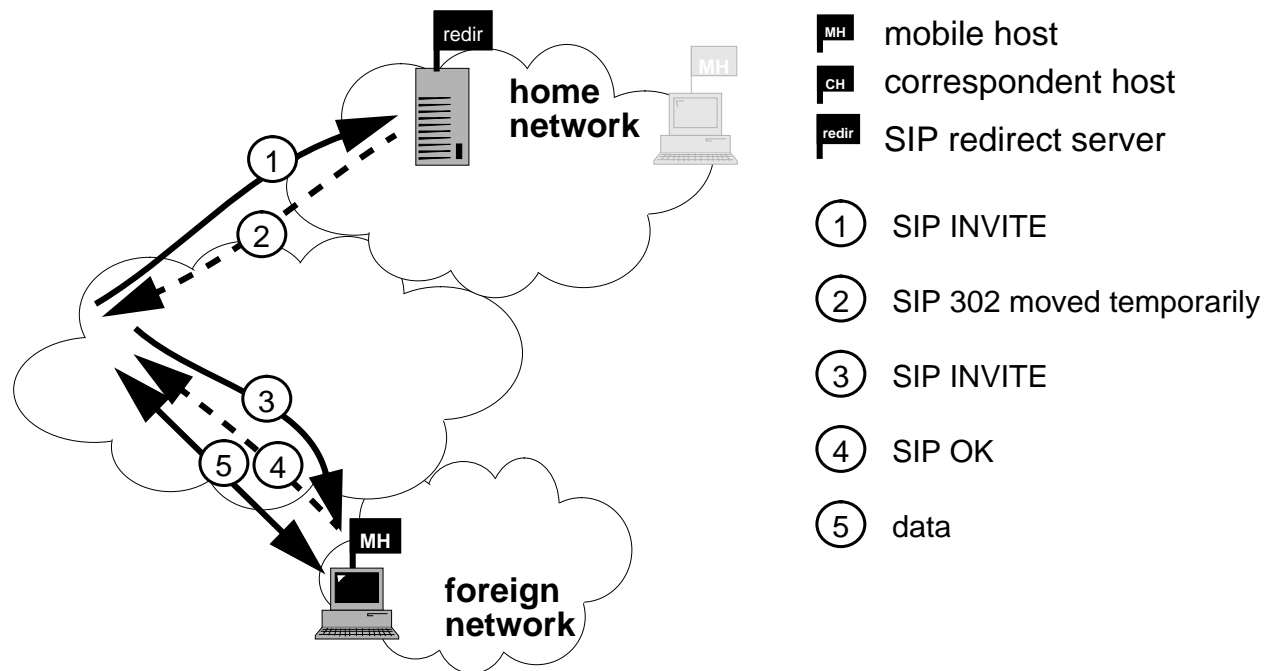
- phone = \$49.99; PC > \$600 (GPF included)
- *Ethernet phone* ⇒ no PBX for switching
- minimal operating system: threads and event flags (CRTX, 2 kB)
- DSP for voice coding ⇒ limited memory (128 kB SRAM!)
- implemented minimal IP stack (IP/UDP/RTP, DHCP, SIP)
- DNS and TCP not absolutely needed
- MP3 radio
- interface to the analog world: sensors, X10, ...



Internet cellular phone

- mobile IP: mask mobility to TCP connections
- imposes overhead:
 - all registrations to home agent
 - triangle routing (mostly)
 - encapsulation, address filtering problems
-  use SIP and RTP for mobility management

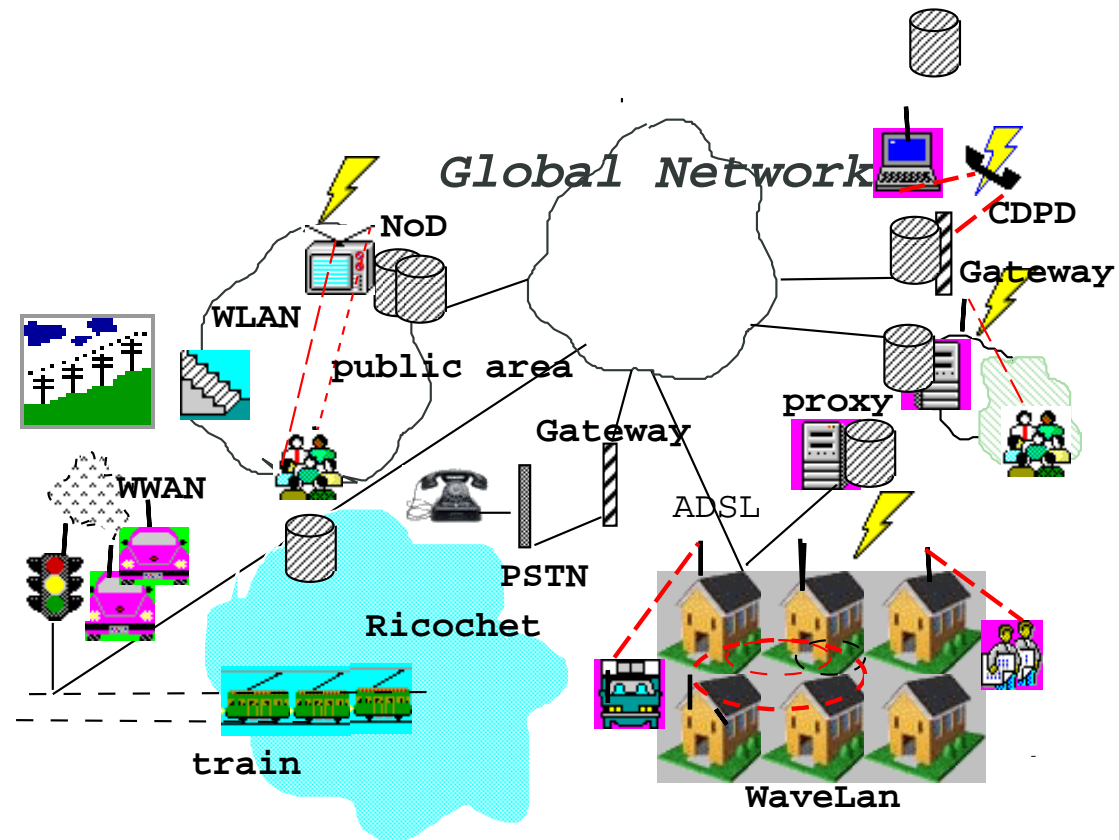
SIP mobility



Other work: signaling

- touch-tone transmission
- interoperation of SIP with SS7, ISDN and POTS
- large-scale IPtel gateways
- locating IPtel gateways (and other wide-area resources)
- charging for (adaptive) services and resources
- Internet voice mail

Internet mobile services



Internet mobile services: “social” ad-hoc networks

- connection sharing
 - multiple wireless networks: 2–10 Mb/s (WL Ethernet, IR) to 28 kb/s (Ricochet, CDPD)
 - share wide-area connections with neighbors
 - load sharing with mobile or in-home gateways
- social caching
 - subway model: in-car high-speed receiver updated in stations
 - socially optimal retrieval
 - anticipatory caching of streaming media
 - “leave the newspaper behind”

Conclusion

- Internet telephony = first new service since web
- last new/old service?
- touches QOS, signaling, programming
- deployment inside out or outside in?
- operational issues: billing, 911, CALEA, ...

More information

Internet and telecom statistics:

<http://www.cs.columbia.edu/~hgs/internet>

Papers: <http://www.cs.columbia.edu/IRT>

RTP: <http://www.cs.columbia.edu/~hgs/rtp>

SIP: <http://www.cs.columbia.edu/sip>