Outer Joins in Relational Algebra

- If a tuple in R doesn’t match any tuple in S, then it won’t be “represented” in $R \bowtie S$:

<table>
<thead>
<tr>
<th>R</th>
<th>A</th>
<th>B</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>x</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>y</td>
<td></td>
</tr>
<tr>
<td>2</td>
<td>z</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>S</th>
<th>B</th>
<th>C</th>
</tr>
</thead>
<tbody>
<tr>
<td>x</td>
<td>i</td>
<td></td>
</tr>
<tr>
<td>y</td>
<td>i</td>
<td></td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>$R \bowtie S$</th>
<th>A</th>
<th>B</th>
<th>C</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>x</td>
<td>i</td>
<td></td>
</tr>
<tr>
<td>1</td>
<td>y</td>
<td>i</td>
<td></td>
</tr>
</tbody>
</table>

“Lost” (2, z) from R!
Outer Joins in Relational Algebra

- If a tuple in R doesn’t match any tuple in S, then it won’t be “represented” in $R \bowtie S$:

  \[
  \begin{array}{ccc}
  R & & S \\
  1 & x & x & i \\
  1 & y & y & i \\
  2 & z & \text{null} & \text{null} \\
  \end{array}
  \]

  “Lost” (2, z) from R!

- Outer joins include such tuples in result, so natural left outer join of R and S is:

  \[
  \begin{array}{ccc}
  A & B & C \\
  1 & x & i \\
  1 & y & i \\
  2 & z & \text{NULL} \\
  \end{array}
  \]
Joins in SQL Revisited

SELECT [DISTINCT] target_list
FROM
  table1 [INNER | {LEFT|RIGHT|FULL} {OUTER}] JOIN table2
  ON qualification_list
WHERE ...

- INNER is default
- Difference in how to deal with NULL values

PostgreSQL documentation:
https://www.postgresql.org/docs/14/tutorial-join.html
Inner Join, Natural Join

SELECT S.sid, S.name, R.bid
FROM Sailors S, Reserves R
WHERE S.sid = R.sid

All equivalent!

SELECT S.sid, S.name, R.bid
FROM Sailors S INNER JOIN Reserves R
ON S.sid = R.sid

SELECT S.sid, S.name, R.bid
FROM Sailors S NATURAL JOIN Reserves R

Natural join means equijoin for each pair of attributes with same name
Find Sailor Names and Their Reserved bids

SELECT S.sid, S.name, R.bid
FROM Sailors S INNER JOIN Reserves R
ON S.sid = R.sid

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>rating</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Eugene</td>
<td>8</td>
<td>22</td>
</tr>
<tr>
<td>2</td>
<td>Luis</td>
<td>2</td>
<td>39</td>
</tr>
<tr>
<td>3</td>
<td>Ken</td>
<td>8</td>
<td>27</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>sid</th>
<th>bid</th>
<th>day</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>102</td>
<td>9/12</td>
</tr>
<tr>
<td>2</td>
<td>102</td>
<td>9/13</td>
</tr>
</tbody>
</table>

Query result (note no tuple for Ken!)

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>bid</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Eugene</td>
<td>102</td>
</tr>
<tr>
<td>2</td>
<td>Luis</td>
<td>102</td>
</tr>
</tbody>
</table>
Left Outer Join

- Returns all matched rows and all unmatched rows from table on left of join clause

So at least one row for each row in left table …
Left Outer Join

- Returns all matched rows and all unmatched rows from table on left of join clause

So at least one row for each row in left table ...

SELECT S.sid, S.name, R.bid
FROM Sailors S LEFT OUTER JOIN Reserves R
    ON S.sid = R.sid

- All sailors who have reserved boats appear with the corresponding bid’s
- All sailors who have reserved no boats also appear but with bid set to NULL
Left Outer Join

SELECT S.sid, S.name, R.bid
FROM Sailors S LEFT OUTER JOIN Reserves R
ON S.sid = R.sid

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>rating</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
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</tr>
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</tr>
</tbody>
</table>
Left Outer Join

```
SELECT S.sid, S.name, R.bid
FROM Sailors S LEFT OUTER JOIN Reserves R
ON S.sid = R.sid
```

<table>
<thead>
<tr>
<th>Sailors</th>
<th>Reserves</th>
</tr>
</thead>
<tbody>
<tr>
<td>sid</td>
<td>name</td>
</tr>
<tr>
<td>1</td>
<td>Eugene</td>
</tr>
<tr>
<td>2</td>
<td>Luis</td>
</tr>
<tr>
<td>3</td>
<td>Ken</td>
</tr>
</tbody>
</table>

Query result
(we now have a tuple for Ken!)
Right Outer Join

Same as LEFT OUTER JOIN, but guarantees results for rows in table on right side of JOIN

```
SELECT S.sid, S.name, R.bid
FROM Reserves R RIGHT OUTER JOIN Sailors S
ON R.sid = S.sid
```
Full Outer Join

Returns all matched or unmatched rows from both sides of JOIN

```
SELECT S.sid, S.name, R.bid
FROM
  Sailors S FULL OUTER JOIN Reserves R
ON S.sid = R.sid
```
Full Outer Join

SELECT S.sid, S.name, R.bid
FROM Sailors S FULL OUTER JOIN Reserves R
ON S.sid = R.sid

<table>
<thead>
<tr>
<th>sid</th>
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</thead>
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<tr>
<td>2</td>
<td>102</td>
<td>9/13</td>
</tr>
<tr>
<td>4</td>
<td>109</td>
<td>9/20</td>
</tr>
</tbody>
</table>

(ignore for this example the fact that Reserves violates referential integrity)
Full Outer Join

```
SELECT S.sid, S.name, R.bid
FROM Sailors S FULL OUTER JOIN Reserves R
ON S.sid = R.sid
```

<table>
<thead>
<tr>
<th>sid</th>
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<td>109</td>
<td>9/20</td>
</tr>
</tbody>
</table>

(ignore for this example the fact that Reserves violates referential integrity)

Why is sid NULL?
ORDER BY, LIMIT

SELECT [DISTINCT] target-list
FROM relation-list
WHERE qualification
GROUP BY grouping-list
HAVING group-qualification
ORDER BY order-list
LIMIT limit-expr [OFFSET offset-expr]
ORDER BY

SELECT S.name, S.rating, S.age
FROM Sailors S
ORDER BY S.rating ASC, S.age DESC

List of order-list expressions dictates ordering precedence:
- Sort primarily in ascending order by rating
- If there are ties on rating, sort them in descending order by age
- If there are ties on both rating and age, sort them arbitrarily
ORDER BY

:: means type “cast”

SELECT S.name, (S.rating/2)::INTEGER AS rat2, S.age
FROM Sailors S
ORDER BY (S.rating/2)::INTEGER ASC, S.age DESC

<table>
<thead>
<tr>
<th>sid</th>
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<th>age</th>
</tr>
</thead>
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</tr>
</tbody>
</table>
ORDER BY

(\(\::\) means type “cast”)

SELECT S.name, ((S.rating/2)::INTEGER) AS rat2, S.age
FROM Sailors S
ORDER BY (S.rating/2)::INTEGER ASC, S.age DESC

<table>
<thead>
<tr>
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</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>name</th>
<th>rat2</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>Luis</td>
<td>1</td>
<td>39</td>
</tr>
<tr>
<td>Ken</td>
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<td>27</td>
</tr>
<tr>
<td>Eugene</td>
<td>4</td>
<td>22</td>
</tr>
</tbody>
</table>

- Default is ASC
- Could add NULLS FIRST or NULLS LAST to indicate position of NULL in order for attribute in ORDER BY
- By default, NULL sorts as if it were larger than non-NULL values
SELECT S.name, (S.rating/2)::INTEGER AS rat2, S.age
FROM Sailors S
ORDER BY (S.rating/2)::INTEGER ASC, S.age ASC

<table>
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<thead>
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<th>name</th>
<th>rat2</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
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<td>39</td>
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<td>22</td>
</tr>
<tr>
<td>Ken</td>
<td>4</td>
<td>27</td>
</tr>
</tbody>
</table>
SELECT S.name, (S.rating/2)::INTEGER AS rat2, S.age
FROM Sailors S
ORDER BY (S.rating/2)::INTEGER ASC, S.age DESC
LIMIT 2

Sailors

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
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<th>age</th>
</tr>
</thead>
<tbody>
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</tr>
<tr>
<td>3</td>
<td>Ken</td>
<td>8</td>
<td>27</td>
</tr>
</tbody>
</table>

Query result

<table>
<thead>
<tr>
<th>name</th>
<th>rat2</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>Luis</td>
<td>1</td>
<td>39</td>
</tr>
<tr>
<td>Ken</td>
<td>4</td>
<td>27</td>
</tr>
</tbody>
</table>

Includes only the top-2 tuples (LIMIT 2)

- To have predictable results, LIMIT should always be used with ORDER BY
- LIMIT ALL and LIMIT NULL are equivalent to omitting LIMIT clause
LIMIT with OFFSET

SELECT S.name, (S.rating/2)::INTEGER AS rat2, S.age
FROM Sailors S
ORDER BY (S.rating/2)::INTEGER ASC, S.age DESC
LIMIT 2 OFFSET 1

Sailors

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>rating</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Eugene</td>
<td>8</td>
<td>22</td>
</tr>
<tr>
<td>2</td>
<td>Luis</td>
<td>2</td>
<td>39</td>
</tr>
<tr>
<td>3</td>
<td>Ken</td>
<td>8</td>
<td>27</td>
</tr>
</tbody>
</table>

Query result

<table>
<thead>
<tr>
<th>name</th>
<th>rat2</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>Ken</td>
<td>4</td>
<td>27</td>
</tr>
<tr>
<td>Eugene</td>
<td>4</td>
<td>22</td>
</tr>
</tbody>
</table>

Includes only the top-2 tuples (LIMIT 2) after skipping the top tuple (OFFSET 1)
**LIMIT with OFFSET**

```
SELECT S.name, (S.rating/2)::INTEGER AS rat2, S.age
FROM Sailors S
ORDER BY (S.rating/2)::INTEGER ASC, S.age DESC
LIMIT (SELECT COUNT(*) / 2
FROM Sailors AS S2)
```

→ can have expressions instead of constants

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>rating</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
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<td>8</td>
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</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>name</th>
<th>rat2</th>
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</tr>
</thead>
<tbody>
<tr>
<td>Luis</td>
<td>1</td>
<td>39</td>
</tr>
</tbody>
</table>
WITH: “Defining Tables” Just for a Query

WITH tablename(attr1, …) AS (select_query)
 [,tablename(attr1, …) AS (select_query)]
main_select_query
“Names and # of reservations of red boats that have been reserved no more than 3 times”
WITH

WITH RedBoats(bid, count) AS
 (SELECT B.bid, COUNT(*)
  FROM Boats B, Reserves R
  WHERE B.bid = R.bid AND B.color = 'red'
  GROUP BY B.bid)

SELECT B.bname, RB.count
FROM Boats AS B, RedBoats AS RB
WHERE B.bid = RB.bid AND RB.count < 4

Names of unpopular red boats, with their number of reservations
Views: Defining “Tables” in Terms of Other Tables

CREATE VIEW <view_name>
AS <select_statement>

- “Tables” defined as query results rather than through inserting base data
  - Helpful to make development simpler
  - Helpful for security
- At query time, references to view_name replaced with select_statement
- Similar to WITH but persistent, not associated with just one query
Defining a View for Popular Boats
Defining a View for Popular Boats

CREATE VIEW boat_counts
AS SELECT R.bid, COUNT(*)
FROM Reserves R
GROUP BY R.bid
HAVING COUNT(*) > 10

SELECT B.bname
FROM boat_counts BC,
Boats B
WHERE B.bid = BC.bid

Rewritten expanded query

SELECT B.bname
FROM (SELECT R.bid, COUNT(*)
FROM Reserves R
GROUP BY R.bid
HAVING COUNT(*) > 10) BC,
Boats B
WHERE B.bid = BC.bid

Query to find names of popular boats, expressed using view
Updates Over a View?

CREATE TABLE Students(
    uni VARCHAR(20),
    ssn CHAR(11),
    name VARCHAR(30),
    PRIMARY KEY(uni),
    UNIQUE(ssn));

<table>
<thead>
<tr>
<th>uni</th>
<th>ssn</th>
<th>name</th>
</tr>
</thead>
<tbody>
<tr>
<td>k101</td>
<td>111-11-1111</td>
<td>Ken</td>
</tr>
<tr>
<td>l102</td>
<td>321-33-3210</td>
<td>Luis</td>
</tr>
<tr>
<td>e105</td>
<td>123-45-6789</td>
<td>Eugene</td>
</tr>
</tbody>
</table>
Updates Over a View?

CREATE TABLE Students(
  uni VARCHAR(20),
  ssn CHAR(11),
  name VARCHAR(30),
  PRIMARY KEY(uni),
  UNIQUE(ssn));

CREATE VIEW StudentUNIs
AS SELECT S.uni, S.name
FROM Students S

→ to hide sensitive information (ssn) when necessary

Can we insert a tuple to StudentUNIs?

INSERT INTO StudentUNIs
VALUES (‘a103’, ‘Alex’)

→ expressed over view, but updates happen over “base” relation Students (view not “materialized” so its contents not stored)
Updates Over a View?

CREATE TABLE Students(
  uni VARCHAR(20),
  ssn CHAR(11),
  name VARCHAR(30),
  PRIMARY KEY(uni),
  UNIQUE(ssn));

CREATE VIEW StudentUNIs
AS SELECT S.uni, S.name
FROM Students S;

Can we insert a tuple to StudentUNIs?
INSERT INTO StudentUNIs
VALUES (‘a103’, ‘Alex’)
Yes, by padding missing attributes with NULL
CREATE TABLE Students(
    uni VARCHAR(20),
    ssn CHAR(11) NOT NULL,
    name VARCHAR(30),
    PRIMARY KEY(uni),
    UNIQUE(ssn));

CREATE VIEW StudentUNIs
AS SELECT S.uni, S.name
FROM Students S

Can we still insert a tuple to StudentUNIs?
INSERT INTO StudentUNIs
VALUES ('a103', 'Alex')
CREATE TABLE Students(
    uni VARCHAR(20),
    ssn CHAR(11) NOT NULL,
    name VARCHAR(30),
    PRIMARY KEY(uni),
    UNIQUE(ssn));

CREATE VIEW StudentUNIs
AS SELECT S.uni, S.name
FROM Students S

Can we still insert a tuple to StudentUNIs?
INSERT INTO StudentUNIs
VALUES ('a103', 'Alex')

No, because ssn cannot be assigned a NULL value
CREATE VIEW boat_counts
AS SELECT R.bid, COUNT(*)
    FROM Reserves R
    GROUP BY R.bid
    HAVING COUNT(*) > 10
Updates Over a View?

Would it make sense to add a tuple to boat_counts view, as follows?

```
INSERT INTO boat_counts
VALUES (103, 32)
```

CREATE VIEW boat_counts
AS SELECT R.bid, COUNT(*)
    FROM Reserves R
    GROUP BY R.bid
    HAVING COUNT(*) > 10

**No!** Reserves is where tuples “live,” and above insertion cannot be translated meaningfully to insertions over Reserves.
Updates Over a View?

Consider now a view with reserves for boat #103:
CREATE VIEW reserves_103
AS SELECT R.bid, R.sid, R.day
    FROM Reserves R
    WHERE R.bid=103

Would it make sense to add this tuple to reserves_103?
INSERT INTO reserves_103
VALUES (104, 22, 10/21/2021)
Updates Over a View?

Consider now a view with reserves for boat #103:

```
CREATE VIEW reserves_103
AS SELECT R.bid, R.sid, R.day
    FROM Reserves R
    WHERE R.bid=103
```

Would it make sense to add this tuple to reserves_103?

```
INSERT INTO reserves_103
VALUES (104, 22, 10/21/2021)
```

(Arguably) yes, but inserted tuple not in view!
Updatable Views

- Even trickier semantics when view is a join of multiple tables!
- Because of all this, views are generally not modifiable, except in limited cases

A view is updatable in SQL if all these conditions hold:
- FROM clause of view definition has only one relation
- SELECT clause contains only attribute names, without expressions, aggregates, or DISTINCT
- Any attribute not in SELECT clause can be set to NULL (i.e., not part of PRIMARY KEY and no NOT NULL constraint)
- No GROUP BY, HAVING clauses

PostgreSQL (slightly different) specifics:
https://www.postgresql.org/docs/14/sql-createview.html
CREATE TABLE with Query

CREATE TABLE <table_name> AS
<select_statement>
CREATE TABLE with Query

CREATE TABLE <table_name> AS <select_statement>

- Schema of table is inherited from SELECT but can be overridden:

```
CREATE TABLE boats_jane1 AS
  SELECT R.bid
  FROM Sailors S, Reserves R
  WHERE S.sid = R.sid AND
  S.sname='Jane'
```

```
CREATE TABLE boats_jane2 AS
  SELECT R.bid AS foo
  FROM Sailors S, Reserves R
  WHERE S.sid = R.sid AND
  S.sname='Jane'
```

boats_jane1(bid INTEGER)  boats_jane2(foo INTEGER)

- How is this different than views? (Hint: What if we insert a new tuple into Reserves?)
Modifying a Relation Schema

- ALTER TABLE Sailors ADD spouse_id INTEGER;
  spouse_id attribute initialized as NULL for existing Sailors
- ALTER TABLE Sailors ADD phone CHAR(12) DEFAULT 'unlisted';
  In contrast, phone attribute initialized with a default value ‘unlisted’
- ALTER TABLE Sailors DROP age;
  Modifications succeed only if they don’t conflict with the rest of the existing schema (e.g., cannot drop an attribute that is part of a primary key, for example)
Modifying a Relation Schema

- `ALTER TABLE Sailors DROP CONSTRAINT SidIsKey;`
  Can also drop constraints, but they have to have a name

- Can add constraints as well, but must be valid at the time they are added

- Can also “CASCADE CONSTRAINTS” (e.g., to drop all foreign keys that refer to the primary key attributes that you are dropping)
Event-Condition-Action Rules: Triggers

- Only awakened when a certain \textit{event} happens (e.g., insert, delete, update)
- A trigger tests a \textit{condition} when an event awakens it; if trigger condition is false, nothing happens
- If trigger condition is true, the associated \textit{action} is performed by the DBMS (e.g., prevent the event from happening or undo the effects of the event); action can be any sequence of database operations!
Triggers in SQL: Options

● A condition may be specified in WHEN clause: action executed only if the rule is triggered and the condition holds when triggering event occurs

● Action executed either before or after the triggering event

● Action can refer to both old and/or new values of tuples that were inserted, deleted, or updated in the event that triggered the action

● Action is performed either:
  ● Once for each modified tuple, or
  ● Once for all tuples changed in one database operation
Trigger to foil any attempt to lower the net worth of a movie executive

MovieExec(name, address, cert#, netWorth) relation
Trigger to foil any attempt to lower the net worth of a movie executive

MovieExec(name, address, cert#, netWorth) relation

CREATE TRIGGER NetWorthTrigger
AFTER UPDATE OF netWorth ON MovieExec
REFERENCING OLD ROW AS OldTuple,
NEW ROW AS NewTuple
FOR EACH ROW
WHEN (OldTuple.netWorth > NewTuple.netWorth)
UPDATE MovieExec
SET netWorth=OldTuple.netWorth
WHERE cert#=NewTuple.cert#;
Trigger to prevent average net worth of movie executives from dropping below $500K

MovieExec(name, address, cert#, netWorth) relation
Trigger to prevent average net worth of movie executives from dropping below $500K
MovieExec(name, address, cert#, netWorth) relation

CREATE TRIGGER AvgNetWorthTrigger
AFTER UPDATE OF netWorth ON MovieExec
REFERENCING OLD TABLE AS OldStuff,
   NEW TABLE AS NewStuff
FOR EACH STATEMENT
WHEN (500000 > (SELECT AVG(M.netWorth) FROM MovieExec M))
BEGIN
   DELETE FROM MovieExec
   WHERE (name, address, cert#, netWorth) IN NewStuff;
   INSERT INTO MovieExec
      (SELECT * FROM OldStuff)
END;
Modeling Total Participation, Finally
User-Defined Functions (UDFs)

- Custom functions that can be called in database
- Many languages: SQL, Python, C, Perl, …

```
CREATE FUNCTION function_name(p1 type, p2 type, ...) 
RETURNS type 
AS $$
-- logic
$$

$$ LANGUAGE language_name;
```
A Simple UDF, Written in SQL

CREATE FUNCTION mult1(v INTEGER) RETURNS INTEGER
AS $$
SELECT v*10;
$$ LANGUAGE SQL;

Last statement is returned

SELECT mult1(S.age) AS age10
FROM Sailors AS S

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>rating</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Eugene</td>
<td>8</td>
<td>22</td>
</tr>
<tr>
<td>2</td>
<td>Luis</td>
<td>2</td>
<td>39</td>
</tr>
<tr>
<td>3</td>
<td>Ken</td>
<td>8</td>
<td>27</td>
</tr>
</tbody>
</table>

Result

<table>
<thead>
<tr>
<th>age10</th>
</tr>
</thead>
<tbody>
<tr>
<td>220</td>
</tr>
<tr>
<td>390</td>
</tr>
<tr>
<td>270</td>
</tr>
</tbody>
</table>

https://www.postgresql.org/docs/14/xfunc-sql.html
Another UDF Written in SQL, with Tuple as Input

CREATE FUNCTION mult2(x Sailors) RETURNS REAL AS $$
SELECT (x.sid + x.age) / x.rating;
$$ LANGUAGE SQL;

SELECT mult2(*) AS age3
FROM Sailors AS S

<table>
<thead>
<tr>
<th>sid</th>
<th>name</th>
<th>rating</th>
<th>age</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Eugene</td>
<td>8</td>
<td>22</td>
</tr>
<tr>
<td>2</td>
<td>Luis</td>
<td>2</td>
<td>39</td>
</tr>
<tr>
<td>3</td>
<td>Ken</td>
<td>8</td>
<td>27</td>
</tr>
</tbody>
</table>

Result
<table>
<thead>
<tr>
<th>age3</th>
</tr>
</thead>
<tbody>
<tr>
<td>2.875</td>
</tr>
<tr>
<td>20.5</td>
</tr>
<tr>
<td>3.75</td>
</tr>
</tbody>
</table>
Procedural Language/SQL, or PL/SQL

Extension of SQL with instructions common in programming languages (IF ELSE statements, etc.)

```sql
CREATE FUNCTION proc(v INTEGER) RETURNS INTEGER AS $$
DECLARE
  -- define variables
BEGIN
  -- PL/SQL code
END;
$$ LANGUAGE plpgsql;
```

https://www.postgresql.org/docs/14/plpgsql.html
CREATE FUNCTION proc(v INTEGER) RETURNS INTEGER AS $$
DECLARE
    qty INTEGER = 10;
BEGIN
    qty = qty * v;
    INSERT INTO blah VALUES(qty);
    RETURN qty + 2;
END;
$$ LANGUAGE plpgsql;
Procedural Code in Python 2, or plpython2u (u=“untrusted”)

CREATE FUNCTION proc(v INTEGER) RETURNS INTEGER AS $$
import random
return random.randint(0, 100) * v
$$ LANGUAGE plpython2u;

- Very powerful: can do anything so must be careful; run in a Python interpreter with no security protection
- plpy Python module provides database access (e.g., plpy.execute("select 1"))

https://www.postgresql.org/docs/14/plpython.html
CREATE FUNCTION proc(v TEXT) RETURNS TEXT AS $$
import requests
resp = requests.get('http://google.com/search?q=%s' % v)
return resp.content.decode('unicode-escape')
$$ LANGUAGE plpython2u;