CS1001

Lecture 24

Overview

- Encryption
- Artificial Intelligence
- Homework 4

Reading

- Brookshear, 11.6
- Brookshear, 10

Homework 4

- Check Courseworks
- Problems based on
 - Handout (Smullyan, Natural Deduction)
 - Question from Ch. 10
 - Question from Ch. 11

Some Trivial Schemes

- Caesar cipher: substitution cipher:
 - $-A \rightarrow D, B \rightarrow E$
- Captain Midnight Secret Decoder rings:
 - shift variable by n: IBM → HAL, or :
 - (letter + offset) mod 26
 - only 26 possible ways of secret coding.
- Monoalphabetic cipher:
 - generalization, arbitrary mapping of one letter to another
 - 26!, approximately 4×10^{26}
 - statistical analysis of letter frequencies
- One-time pad
 - A random sequence of 0's and 1's XORed to plaintext

Definitions

- Process data into unintelligible form, reversible, without data loss
- Usually one-to-one (not compression)
- Other services:
 - Integrity checking: no tampering
 - Authentication: not an imposter
- Plaintext encryption → ciphertext decryption → plaintext

Computational Difficulty

- Algorithm needs to be efficient.
 - Otherwise only short keys can be used.
- Most schemes can be broken: depends on \$\$\$.
 - E.G. Try all possible keys.
- Longer key is often more secure:
 - Encryption O(N+1).
 - Brute-force cryptanalysis: $O(2^{N+1})$, twice as hard with each additional bit.
- Cryptanalysis tools:
 - Special-purpose hardware.
 - Parallel machines.
 - Internet coarse-grain parallelism.

Secret Key vs. Secret Algorithm

- Secret algorithm: additional hurdle
- Hard to keep secret if used widely:
 - Reverse engineering, social engineering
- Commercial: published
 - Wide review, trust
- Military: avoid giving enemy good ideas

Cryptanalysis: Breaking an Encryption Scheme

- Ciphertext only:
 - Exhaustive search until "recognizable plaintext"
 - Need enough ciphertext
- Known plaintext:
 - Secret may be revealed (by spy, time), thus<ciphertext, plaintext> pair is obtained
 - Great for monoalphabetic ciphers
- Chosen plaintext:
 - Choose text, get encrypted
 - Useful if limited set of messages

Models for Evaluating Security

- Unconditional security (perfect secrecy)
 - Observation of ciphertext provides no information
 - Uncertainty/entropy H(p) = H(p|c)
- Complexity-theoretic security
- Provable security
 - As difficult to break as solving well-known and supposedly difficult problem
- Computational security
- Ad hoc security

Brute Force Attacks

- Number of encryption/sec: 1 million to 1 billion/sec
- 56-bit key broken in 1 week with 120,000 processors (\$6.7m)
- 56-bit key broken in 1 month with 28,000 processors (\$1.6m)
- 64-bit key broken in 1 week with 3.1 × 10⁷ processors (\$1.7b)
- 128-bit key broken in 1 week with 5.6 × 10²⁶ processors

Types of Cryptography

- Hash functions: no key
- Secret key cryptography: one key
- Public key cryptography: two keys public, private

Secret Key Cryptography

- Same key is used for encryption and decryption
 - Symmetric cryptography
- Ciphertext approximately the same length as plaintext
- Substitution codes, DES, IDEA
- Message transmission:
 - Agree on key (but how?)
 - Communicate over insecure channel
- Secure storage: crypt

Secret Key Cryptography (Cont'd)

- Strong authentication: prove knowledge of key without revealing it:
 - Send challenge r, verify the returned encrypted {r}
 - Fred can obtain chosen plaintext, cihpertext pairs
 - Challenge should chosen from a large pool
- Integrity check: fixed-length checksum for message
 - Send MIC along with the message

Public Key Cryptography

- Asymmetric cryptography
- Invented/published in 1975
- Two keys: private (*d*), public (*e*)
 - Encryption: public key; Decryption: private key
 - Signing: private key; Verification: public key
- Much slower than secret key cryptography

Public Key Cryptography (Cont'd)

- Data transmission:
 - Alice encrypts m_a using e_B , Bob decrypts to m_a using d_b .

Storage:

 Can create a safety copy: using public key of trusted person.

Authentication:

- No need to store secrets, only need public keys.
- Secret key cryptography: need to share secret key for every person to

Public Key Cryptography (Cont'd)

- Digital signatures
 - Encrypt hash h(m) with private key
 - Authorship
 - Integrity
 - Non-repudiation: can't do with secret key cryptography

Hash Algorithms

- Message digests, one-way transformations
- Length of h(m) much shorter then length of m
- Usually fixed lengths: 48-128 bits
- Easy to compute *h*(*m*)
- Given h(m), no easy way to find m
- Computationally infeasible to find m_1 , m_2 s.t. $h(m_1) = h(m_2)$
- Example: $(m+c)^2$, take middle *n* digits

Hash Algorithms (Cont'd)

Password hashing

- Doesn't need to know password to verify it
- Store h(p+s), s (salt), and compare it with the user-entered p
- Salt makes dictionary attack less convenient
- Message integrity
 - Agree on a password p
 - Compute h(p|m) and send with m
 - Doesn't require encryption algorithm, so the technology is exportable

Public Key Crypto's Trick

- Consider the Knapsack problem (p482)
- We have a knapsack filled with numbers
- The goal is to select out a series of numbers that adds to some desired number
- How do we do this efficiently?

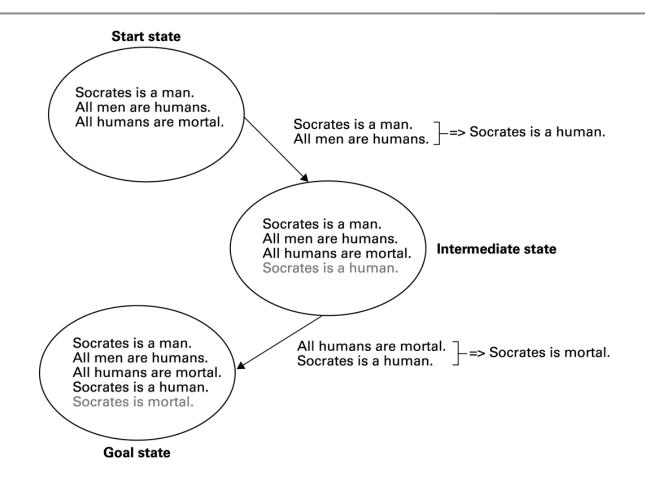
The Trick

- The trick is factoring
- You know ahead of time what certain properties of the number will be. This allows you to reduce the problem to a computable one
- Otherwise, you are dealing with a nonpolynomial problem

Artificial Intelligence

- Reasoning (Production Systems
 - Goal is to derive a solution given facts and rules
- Searching
 - You are given facts and rules and search all possible combinations to find some desired solution (usually minimum/max)
- Heuristics
 - Operate based on guidelines you know to be true about a problem

A Production System



Artificial Intelligence

- Neural Networks
- Genetic Algorithms
- Machine Learning