### Sample complexity bounds for differentially private learning

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### Outline

- I. Learning and privacy model
- 2. Our results: sample complexity bounds for differentially-private learning
- 3. Recap & future work

### Part I. Learning and privacy model

# Data analytics with sensitive information

<u>eCommerce</u>: customers' browsing & purchase histories <u>Clinical studies</u>: patients' medical records & test results <u>Genomic studies</u>: subjects' genetic sequences

Patient I	age	34	
	test #l	1.76	
	test #2	86.6	
	has flu?	Ι	

Patient 2	age	31	
	test #l	1.62	
	test #2	67.5	
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Learn something useful about whole population from data about individuals.

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<u>This work</u>: learning a binary classifier from labeled examples, where each training example is an individual's sensitive information.

## Data analytics with sensitive information





Q: If a classifier is learned from some individuals' sensitive data, can releasing / deploying the classifier in public violate the privacy of individuals from the training data?



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A: Yes! Even after standard "anonymization", and even when just releasing aggregate statistics, because an adversary could have side-information.

## Example: genome-wide association studies

Has disease								Healthy													
1.00   .190 1.00   .216 .251   .186 .117   .154 .011   .190 .140   .270 .215   .101 .085   .239 .071   .471 .117   .179 .202	1.00 .047 .170 .102 .294 .170 .163 .243 .132	1.00 .083 .095 .248 .056 .111 .094 .094	1.00 .139 1 .140 . .234 . .161 . .144 . .087 .	1.00 141 099 093 123 159	1.00 .175 .199 .283 .207	1.00 .157 .216 .108	1.00 .274 .092	1.00	1.00		1.00 .141 .099 .093 .123 .159 .088 .046 .078 .045 .178	1.00 .175 .199 .283 .207 .152 .161 .392 .155 .135	1.00 .157 .216 .08 .075 .092 .122 .135 .102	1.00 .274 .092 .163 .072 .229 .139 .258	1.00 .294 .156 .157 .160 .110 .314	1.00 .220 .143 .172 .048 .165	1.00 .147 .145 .126 .147	1.00 .177 .104 .158	1.00 .169 .131	1.00	1.00

Correlations statistics

**Correlations statistics** 

Wang et al (2009): able to combine side-information and published correlation statistics to determine whether an individual from the study was in disease group or healthy group.



<u>Goals</u>: learn an accurate classifier from sensitive data while also preserving the privacy of the data.

<u>This work</u>: how many labeled examples are needed to achieve both of these goals simultaneously?

### Goal I: Differential privacy

What kind of privacy guarantee can a good learning algorithm provide?

<u>Differential privacy guarantee</u> [Dwork *et al*, 2006]: an individual's inclusion in the training data does not change (much) what an adversary could learn about that individual's sensitive information.

### Goal I: Differential privacy

(Definition from [Dwork, et al 2006], specialized to learning [Kasiviswanathan, et al 2008])

A learning algorithm  $\mathcal{A}: (\mathcal{X} \times \{0, 1\})^* \to \mathcal{H}$ is  $\alpha$ -differentially private if:

For all training sets S and S' differing in at most one example,

$$\forall \mathcal{G} \subseteq \mathcal{H}, \quad \frac{\Pr_{\mathcal{A}}[\mathcal{A}(S) \in \mathcal{G}]}{\Pr_{\mathcal{A}}[\mathcal{A}(S') \in \mathcal{G}]} \le e^{\alpha}.$$

- Probability is over internal randomness of the learning algorithm.
- Algorithm must behave similarly given similar training sets.
- Smaller  $\alpha \in [0,1]$  corresponds to stronger guarantee.

### Goal 2: Learning

#### Standard statistical learning guarantees:

If S is an i.i.d. sample from a distribution  $\mathcal{P}$  over  $\mathcal{X} \times \{0, 1\}$ , then  $\mathcal{A}(S)$  returns a hypothesis  $h \in \mathcal{H}$  such that w.p.  $\geq 1 - \delta$ (over random draw of S and randomness in  $\mathcal{A}$ )

$$\operatorname{err}_{\mathcal{P}}(h) \leq \min_{h' \in \mathcal{H}} \operatorname{err}_{\mathcal{P}}(h') + \epsilon$$

where  $\operatorname{err}_{\mathcal{P}}(\tilde{h}) = \operatorname{Pr}_{(x,y)\sim\mathcal{P}}[\tilde{h}(x)\neq y].$ 

### What was known (previous work)

• Sample complexity for finite hypothesis classes or VC classes over discrete data domains.

[Kasiviswanathan et al, 2008], [Blum et al, 2008], [Beimel et al, 2010]

$$C \cdot \left(\frac{1}{\alpha \epsilon} + \frac{1}{\epsilon^2}\right) \cdot \left(\min\{\log |\mathcal{H}|, \operatorname{VC}_{\mathcal{H}} \log |\mathcal{X}|\} + \log \frac{1}{\delta}\right)$$

• Related problems: (synthetic) data set release.

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• Related problems: (synthetic) data set release.

What about infinite classes & continuous data domains?

## Part 2. Sample complexity bounds for differentially-private learning

### Our results

- I. <u>Some bad news</u>: no distribution-independent sample complexity upper bound possible for differentially-private learning.
- 2. <u>Some hope</u>: differentially-private learning possible if
  - a. learner allowed some prior-knowledge, or
  - b. privacy requirement is relaxed.

## I. No distribution-independent sample complexity upper bound

Let  $\mathcal{H}$  be the class of threshold functions on the unit interval [0, 1], and pick any positive real number M.

For every  $\alpha$ -differentially private algorithm  $\mathcal{A}$ : ([0,1] ×  $\{0,1\}$ )\*  $\rightarrow \mathcal{H}$ , there is a distribution  $\mathcal{P}$  (with full support) over  $[0,1] \times \{0,1\}$  such that:

- 1. There exists a threshold  $h^* \in \mathcal{H}$  with  $\operatorname{err}_{\mathcal{P}}(h^*) = 0$ .
- 2. If S is an i.i.d. sample of size  $m \leq M$  from  $\mathcal{P}$ , then

$$\Pr_{S \sim \mathcal{P}^m, \mathcal{A}}\left[\operatorname{err}_{\mathcal{P}}(\mathcal{A}(S)) > \frac{1}{5}\right] \geq \frac{1}{2}.$$

I. No distribution-independent sample complexity upper bound

Implications:

 No direct analogue of VC theorem for differentially-private learning.



 Qualitative difference between finite hypothesis class / discrete data domains and infinite classes / continuous data domains.

VS

## I. No distribution-independent sample complexity upper bound

<u>Proof idea</u>: find data distributions P and P' such that a "successful" distribution over thresholds for P differs significantly from a "successful" distribution over thresholds for P'.



A differentially-private learner using just a small number of examples must behave similarly in both cases; therefore, it must fail for at least one of the cases.

#### 2. Some hope for differentiallyprivate learning

Possible ways around the lower-bound:

- a. Allow learner access to prior-knowledge (or prior belief) about unlabeled data distribution.
- b. Only guarantee the differential privacy of the labels in the training data.

- Allow learner access to a reference distribution U over unlabeled data X, chosen independently of the training data.
- Sample complexity upper bound depends on how close U is to D (true unlabeled data distribution).





U and D close

U and D far

Let  $\mathcal{P}$  be any distribution over  $\mathcal{X} \times \{0, 1\}$  with marginal  $\mathcal{D}$ over  $\mathcal{X}$ . There is a constant C > 0 and an  $\alpha$ -differentially private algorithm  $\mathcal{A}_1$  s.t. given an i.i.d. sample S of size

$$|S| \ge C \cdot \left(\frac{1}{\alpha\epsilon} + \frac{1}{\epsilon^2}\right) \cdot \left(d_{\mathcal{U}} \cdot \log \frac{\kappa(\mathcal{U}, \mathcal{D})}{\epsilon} + \log \frac{1}{\delta}\right),$$

w.p.  $\geq 1 - \delta$ ,  $\mathcal{A}_1(S)$  returns a hypothesis  $h \in \mathcal{H}$  with  $\operatorname{err}_{\mathcal{P}}(h) \leq \min_{h' \in \mathcal{H}} \operatorname{err}_{\mathcal{P}}(h') + \epsilon$ .

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Example:

- H = n-dimensional linear separators through the origin
- U = uniform distribution on unit sphere(so  $d_U = O(n)$ )
- Unlabeled data distribution D close to uniform:  $D(x) \le c \cdot U(x)$
- Sample complexity upper bound:

$$C \cdot \left(\frac{1}{\alpha \epsilon} + \frac{1}{\epsilon^2}\right) \cdot \left(n \cdot \log \frac{c}{\epsilon} + \log \frac{1}{\delta}\right)$$

#### 2(b). Label privacy

- Weaker privacy guarantee: only guarantee differential-privacy of the *labels*.
- Can still protect against some privacy attacks on training data.

A learning algorithm  $\mathcal{A}: (\mathcal{X} \times \{0, 1\})^* \to \mathcal{H}$ is  $\alpha$ -label private if:

For all training sets  $S, S' \subseteq \mathcal{X} \times \{0, 1\}$ differing in at most one *label*,  $\Pr_{\mathcal{A}}[\mathcal{A}(S) \in \mathcal{G}] \leq \Pr_{\mathcal{A}}[\mathcal{A}(S') \in \mathcal{G}] \cdot e^{\alpha} \quad (\forall \mathcal{G} \subseteq \mathcal{H})$ 

#### 2(b). Label privacy

- Label privacy avoids complications that arise with infinite hypothesis classes and continuous data domains
- Can obtain upper- and lower-bounds in terms of certain distribution-dependent complexity measures (covering number, doubling dimension).
- Bounds are (roughly) within  $1/\alpha$  factor of non-private sample complexity bounds.

### Recap & future work

- I. Differential-privacy requirement rules out distribution-independent proper learning.
- 2. Some ways out:
  - a. Data-dependent bounds based on priorknowledge.
  - b. Relaxed notion of privacy (label privacy).
- 3. Future directions:
  - a. Improper learning (some work in discrete settings by [Beimel et al, 2010]).
  - b. Other weaker notions of privacy.
  - c. More general statistical estimation tasks.

### Thanks!

#### I. Bad news: no distributionindependent sample complexity upper bound

<u>Idea</u>: Consider a set of distributions  $\{P_z\}$  for  $z \in [0,1]$ : the marginal of each  $P_z$  over X is an even mixture of

(I) uniform on [0,1], and

(2) uniform on  $[z-\eta, z+\eta]$  (where  $\eta = \Theta(\exp(-\alpha M))$ ); and labels are given by threshold  $h_z(x) = \mathbb{1}[x \ge z]$ .



<u>To show</u>: Every  $\alpha$ -differentially private learning algorithm using at most M training examples will fail on at least one distribution  $P_z$ .

Example:

- H = n-dimensional linear separators through the origin
- U = uniform distribution on unit sphere (so  $d_U = n$ )
- Unlabeled data distribution D uniform outside  $\Theta(1)$ -width band around equator.
- Sample complexity upper bound:

$$C \cdot \left(\frac{1}{\alpha \epsilon} + \frac{1}{\epsilon^2}\right) \cdot \left(n^2 + n \cdot \log \frac{1}{\epsilon} + \log \frac{1}{\delta}\right)$$

### Doubling dimension

• Hypothesis class H + unlabeled data distribution D→ disagreement metric space  $(\mathcal{H}, \rho_{\mathcal{D}})$ 

 $\rho_{\mathcal{D}}(h,h') = \Pr_{x \sim \mathcal{D}}[h(x) \neq h'(x)]$ 

- Doubling dimension is d if every ball of radius r can be covered by  $2^d$  balls of radius r/2 (and no fewer).
- (Non-private) sample complexity bound due to Bshouty et al (2009) for noiseless setting:

$$C \cdot \frac{1}{\epsilon} \left( d + \log \frac{1}{\delta} \right)$$

### Divergence K(U,D)

$$\kappa(\mathcal{U}, \mathcal{D}) = \inf \left\{ k > 0 \colon \Pr_{x \sim \mathcal{D}} [x \in A] \le k \cdot \Pr_{x \sim \mathcal{U}} [x \in A] \\ \forall \text{ measurable } A \right\}$$

(Quantifies absolute continuity of D w.r.t. U.)