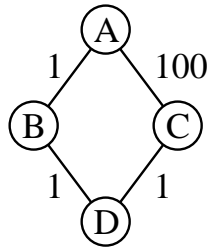


HW #5

ELEN E4710 - Intro to Network Engineering
Spring 2002

Due 4/2/2002
Prof. Rubenstein

Homework must be turned in at the beginning of class on the due date indicated above. CVN students have one additional day. Late assignments will not be accepted.



1. Consider the small 4-node network above with weighted edges. Assume that the nodes communicate on a round-by-round basis, where a node can communicate once to each of its neighbors within a single round, but only one node can send a message in a given round.
 - (a) Run the Belman-Ford Algorithm (i.e., Distance Vector Protocol) without Poison Reverse/Split Horizon implemented to compute the shortest paths trees to all nodes.
 - (b) Assume that the weight of edge (A, B) suddenly changes to 200. Perform the first five rounds of the algorithm after the change. Indicate how many rounds take place in total until convergence.
 - (c) Run the Belman-Ford Algorithm on the original graph (where (A, B) has weight 1) with Poison Reverse/Split Horizon implemented to compute the shortest paths trees to all nodes.
 - (d) Assume again that the weight of edge (A, B) suddenly changes to 200. Perform all rounds of the algorithm until messages cease to be sent.
 - (e) Assume an additional edge (B, C) with weight 3 is added to the original graph. Modify the routing tables from part 1c to include this edge when Poison Reverse/Split Horizon is implemented (you need not draw all the steps, just the final tables).
 - (f) Suppose the weight of (A, B) once again shifts to 200. Perform the first 5 iterations of the algorithm. Explain why Poison Reverse did not prevent the count to infinity problem.
 - (g) Did Poison Reverse help at all in part 1f? Explain your answer.
2. Draw the trie structure that maps addresses to the interface corresponding to the longest matching prefix in the table below:

prefix	interface
0	2
1	2
01	1
111	1
011	3
0110	1
1011	3