

Machine Learning

4771

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Lecture 12: Large Margin & Optimal Hyperplane

- Structural Risk Minimization (SRM)
- Large Margin, Optimal Hyperplane (Burges Tutorial)
- Optimization
- Support Vector Machines (Bishop 7.1, Burges Tutorial)

Constructive Bound

- With probability $(1-\eta)$, for the function that minimizes empirical risk, the inequality below holds true

$$R(\alpha_\ell) < R_{emp}(\alpha_\ell) + \frac{E(\ell)}{2} \left(1 + \sqrt{1 + \frac{4R_{emp}(\alpha_\ell)}{E(\ell)}} \right)$$

where

$$E(\ell) = 4 \frac{h(1 + \ln(2\ell/h)) - \ln(\eta/4)}{\ell}$$

Large Sample Size

- Suppose we have a *large sample size* (ℓ/h is large)
 - The value of actual risk is determined by value of empirical risk
 - The principle of ERM gives good results in practice
- Justification (we drop constants and show what the bound is proportional to):

$$E(\ell) = 4 \frac{h(1 + \ln(2\ell/h)) - \ln(\eta/4)}{\ell} \approx \frac{h}{\ell} + \frac{\ln(2\ell/h)}{(\ell/h)} \approx \delta$$

$$R_{emp}(\alpha_\ell) + \frac{E(\ell)}{2} \left(1 + \sqrt{1 + \frac{4R_{emp}(\alpha_\ell)}{E(\ell)}} \right) \approx R_{emp}(\alpha_\ell) + \delta \left(1 + \sqrt{1 + \frac{R_{emp}(\alpha_\ell)}{\delta}} \right)$$

$$\approx R_{emp}(\alpha_\ell) + \delta \left(\sqrt{\frac{R_{emp}(\alpha_\ell)}{\delta}} \right) \approx R_{emp}(\alpha_\ell) + \sqrt{\delta R_{emp}(\alpha_\ell)}$$

Large Sample Size

- Suppose we have a large sample size (ℓ/h is large)
 - The value of actual risk is determined by value of empirical risk
 - The principle of ERM gives good results
- Justification (we drop constants and show what the bound is proportional to):

$$R(\alpha_\ell) \leq \left\{ R_{emp}(\alpha_\ell) + \sqrt{\delta R_{emp}(\alpha_\ell)} \right\}$$

Small Sample Size

- Suppose we have a *small sample size* ($\ell/h < 20$)
 - Small empirical risk doesn't guarantee small actual risk anymore
 - Need to minimize bound over both terms simultaneously
 - To do this, we make the VC dimension (capacity) a *controlling variable*
- This observation motivates a new *induction principle*: **Structural Risk Minimization**
- What do we mean by a controlling variable?
- How do we justify this new induction principle?

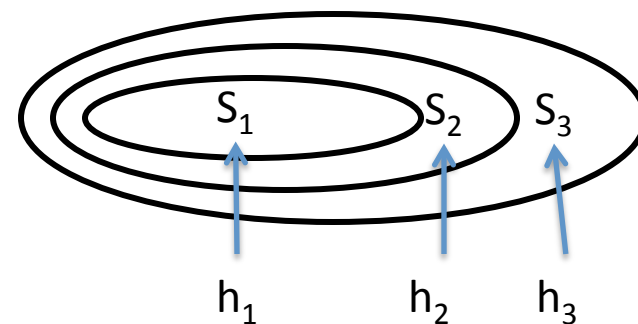
SRM Principle (idea)

- Instead of minimizing empirical risk at any cost, search for the optimal relationship between:
 1. Amount of empirical data
 2. Quality of approximation by the function chosen from a given set of functions
 3. Value that characterizes the capacity of a set of functions
- Lets impose a *structure* (S^*) on the set of loss functions
- We assume that any element S_k of the structure S^* has a finite VC dimension h_k
- The sequence $\{h_k\}$ for elements $\{S_k\}$ of S^* is non-decreasing (as k is increased)

$$S_1 \subset S_2 \subset \dots \subset S_n \subset \dots$$

$$S^* = \bigcup_k S_k, \quad S_k = \{L(z, \alpha) : \alpha \in \Lambda_k\}$$

$$h_1 \leq h_2 \leq \dots \leq h_n \leq \dots$$



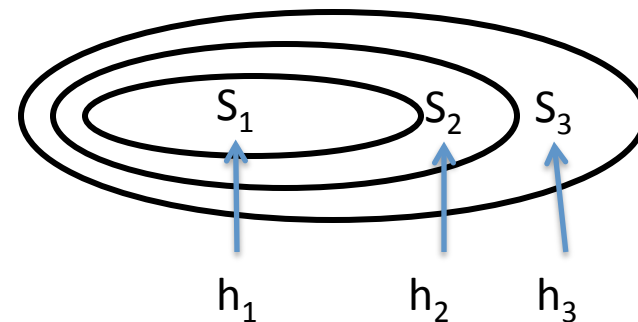
SRM Principle (idea)

- For a given sample, the SRM principle chooses the element S_k of the structure for which the smallest bound on the risk (the smallest guaranteed risk) is achieved
- Within the element S_k , we choose the function that minimizes empirical risk
- General model of capacity control
- We need to provide an *admissible structure* (which satisfies conditions) and then choose the function that yields the best guaranteed risk
- Support Vector Machine (SVM) does just that

$$S_1 \subset S_2 \subset \dots \subset S_n \subset \dots$$

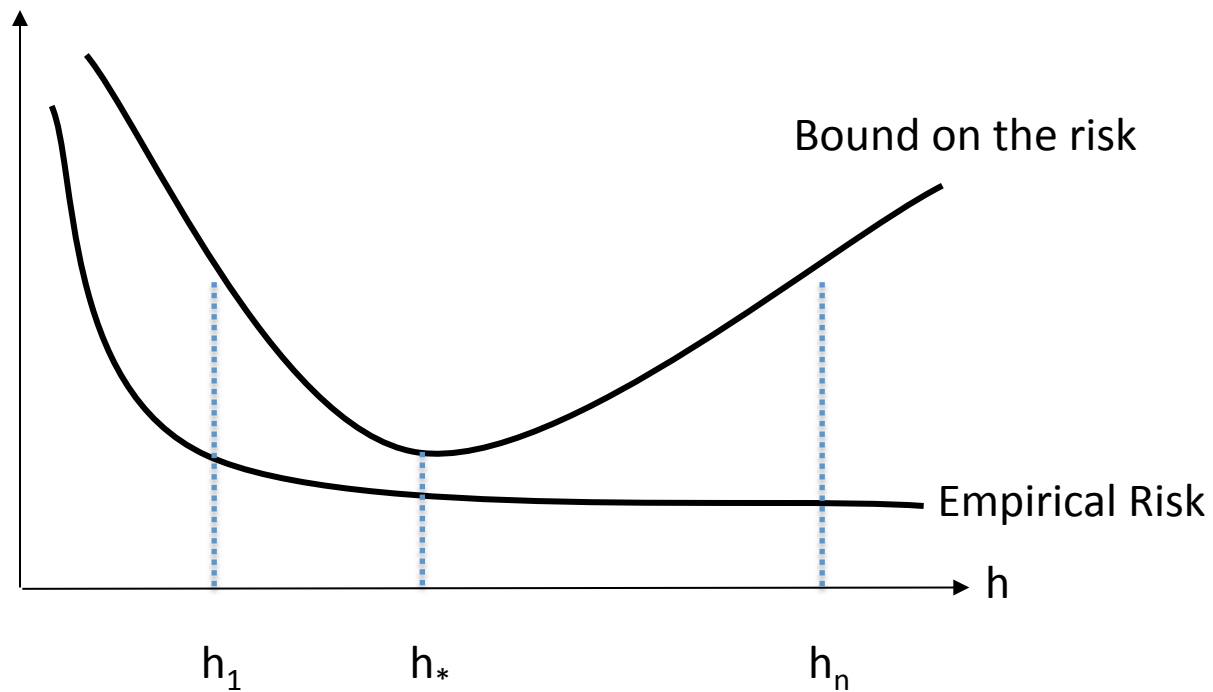
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SRM Principle (idea)

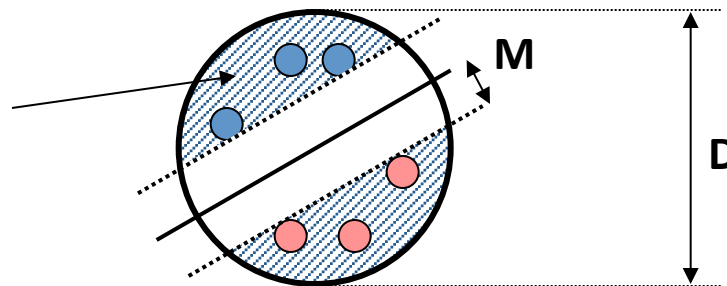
- How do we justify SRM?
- Result: SRM is always consistent and defines a bound on the rate of convergence



Gap Tolerant Classifiers (definition)

- Recall: for N-D linear classifiers, $h = N+1$
- Not quite satisfactory in practice!
- What if I have lots of redundant features (dimensions)? h should be less than $N+1$
- But VC estimate does not distinguish between such cases and cases where features are valuable!
- Solution: constrain linear classifiers to data inside a sphere
- *Gap Tolerant Classifier*: linear classifier whose activity is constrained to a sphere & outside a margin

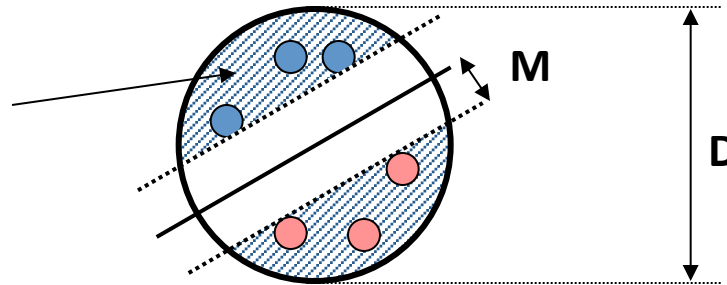
Only count errors
in shaded region
Elsewhere have
 $L(x,y)=0$



M=margin
D=diameter
d=dimensionality

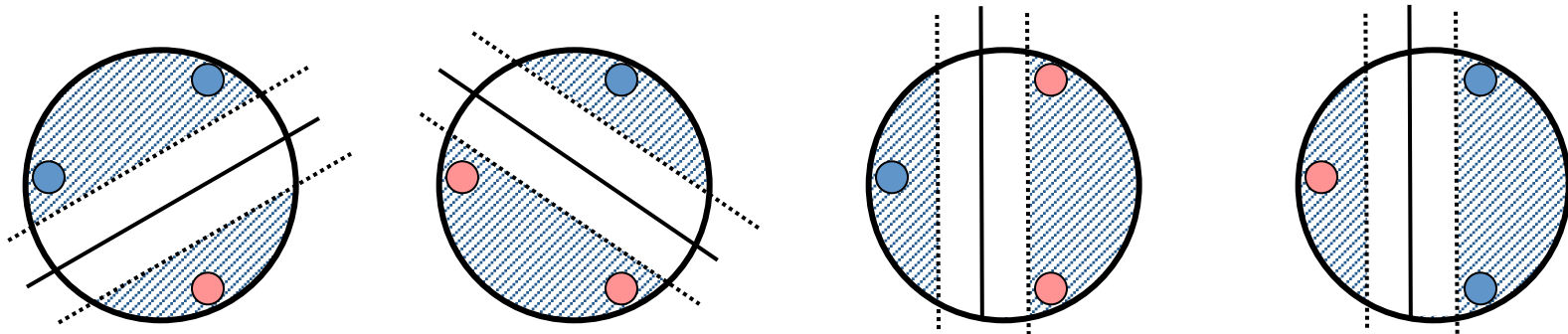
Gap Tolerant Classifiers (idea)

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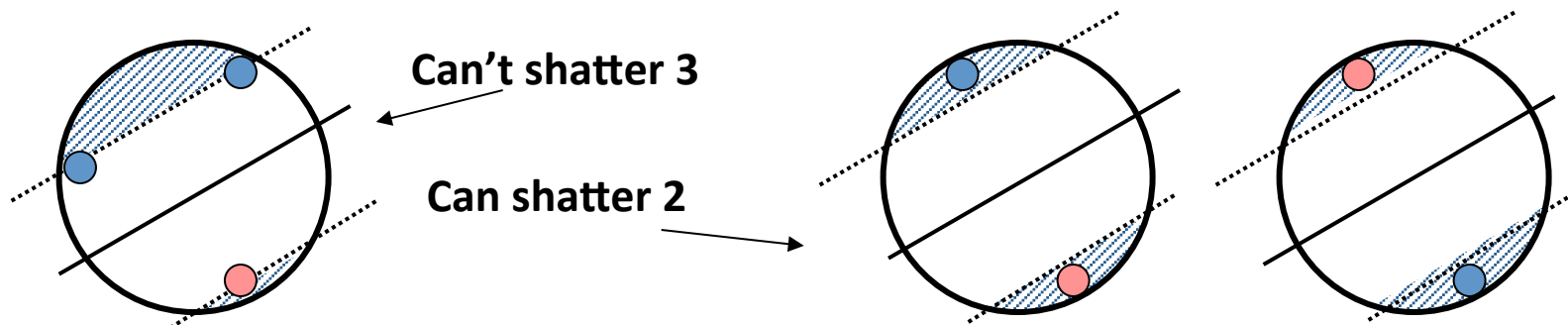


M =margin
 D =diameter
 d =dimensionality

- If M is small relative to D , can still shatter 3 points:



- But as M grows relative to D , can only shatter 2 points!



Large Margin

- We have observed that: as the margin grows relative to data sphere, we can shatter fewer points
- In other words, the larger the margin, the smaller the VC dimension
- The general relation between h & M is expressed as:

$$h \leq \min \left\{ \left\lceil \frac{r^2}{M^2} \right\rceil, N \right\} + 1, \quad r = \max_i \|x_i\|$$

- Previously we just had $h = N+1$.
- Now we have a bound on h in terms of M and radius (r) of the data sphere
- This reflects a fairly typical case where the real data is bounded (if its not, then by default $h = N+1$)
- Note: sometimes bound is expressed in terms of diameter (margin is taken to be the width between the hyperplanes)
- General rule: maximizing margin reduces the VC dimension (inverse relation)

Relation to Perceptron

• **Theorem:** assuming conditions {1,2} below are satisfied, the sequence of weight vectors determined by the online perceptron algorithm will converge to a solution vector in finite number of steps

1. Assume all data lies inside a sphere of radius r : $r = \max_i \|x_i\|$
2. Assume that the data is linearly separable: $\forall i: y_i((w^*)^T x_i) \geq \gamma > 0$

• The bound on the number of steps (k) is expressed in terms of the margin:

$$k \leq \frac{r^2}{\gamma^2} \|\mathbf{w}^*\|^2$$

Optimal Hyperplane (idea)

- Consider a linearly separable 2-class problem:

- Data set: $\{(x_1, y_1), \dots, (x_\ell, y_\ell)\}, \quad x_i \in \mathcal{R}^n, y_i \in \{-1, 1\}$

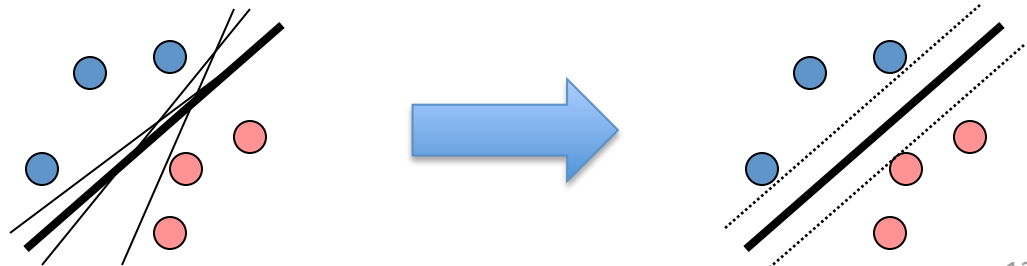
- Decision boundary:

$$f(x; w) = w^T x + b = 0$$

- Symmetry:

$$\frac{w^T x_i + b > 0: \text{ assign } 1}{w^T x_i + b < 0: \text{ assign } -1} \Rightarrow y_i (w^T x_i + b) > 0$$

- There are many solutions (solution region). Perceptron chooses some solution vector
- Can we require that the hyperplane with maximum margin is selected?
- Can we guarantee it is unique?



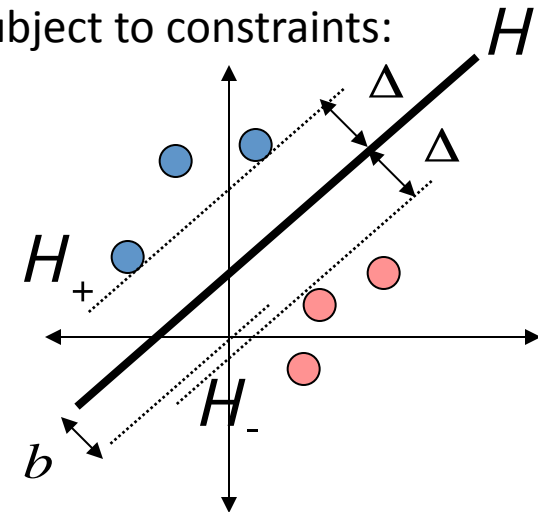
Optimal Hyperplane (definition)

- Define two quantities: $h_1(w) = \min_{i: y_i=1} (w^T x_i)$, $h_2(w) = \max_{i: y_i=-1} (w^T x_i)$

- Consider the unit vector w_0 which maximizes margin subject to constraints:

$$\max_w \Delta(w) = \frac{h_1(w) - h_2(w)}{2}$$

$$s.t \quad \|w\| = 1, \quad \forall i: y_i(w^T x_i + b) > 0$$



- The vector w^* and the constant b^* determine the *maximal margin hyperplane* or the *optimal hyperplane* H

$$b^* = -(h_1(w^*) + h_2(w^*)) / 2$$

- Note: the optimal hyperplane is unique (not proved here)

Better Formulation

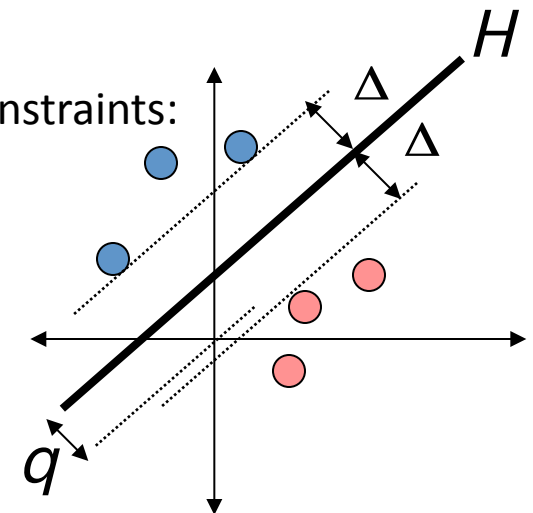
- Goal: find effective methods for constructing the optimal hyperplane
- Consider equivalent problem: instead of restricting the norm of the weight vector (hyperplane), let's scale the value of $f(x)$ for the closest points to the hyperplane

$$\forall i: y_i(w^T x_i + b) \geq 1$$

- Now we are trying to minimize the norm subject to these constraints:

$$\min_w \frac{1}{2} \|w\|^2$$

$$s.t \quad \forall i: y_i(w^T x_i + b) \geq 1$$



- Not hard to show: if we normalize the vector which minimizes the above we obtain the unit vector solution w^* on the previous slide
- Note: the distance to the origin is not just the value of b anymore (denoted q above)

Quadratic Program

- Recall geometry of linear surface: discriminant function $f(\mathbf{x})$ is proportional to the distance from \mathbf{x} to H

$$dist = \frac{f(x)}{\|w\|} = \frac{(w^T x + b)}{\|w\|}, \quad dist2origin = q = \frac{f(0)}{\|w\|} = \frac{|b|}{\|w\|}$$

$$margin = \Delta = \frac{|f(x) = \pm 1|}{\|w\|} = \frac{1}{\|w\|}, \quad width = 2\Delta = \frac{2}{\|w\|}$$

- We have a quadratic program (QP), just plug into a solver (matlab: quadprog), done!

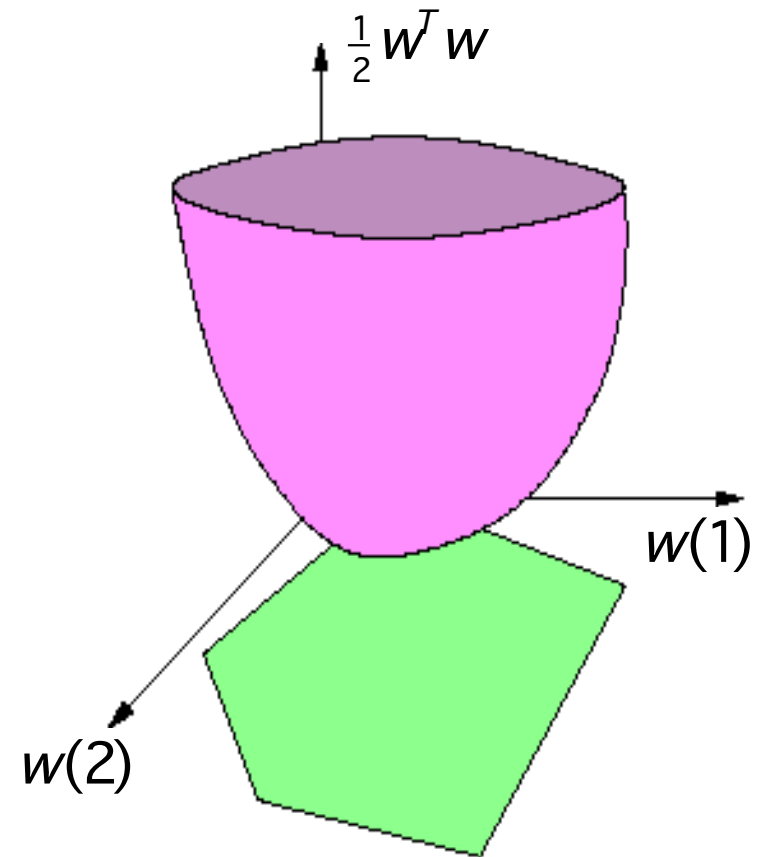
$$\min_w \frac{1}{2} \|w\|^2$$

$$s.t \quad \forall i: y_i (w^T x_i + b) \geq 1$$

- We would solve the problem in *primal space*, but can also solve it in dual space

QP Visualization

- Each data point adds a linear inequality to QP
- Each point cuts a half plane of allowable planes and reduces green region
- The optimal hyperplane is the closest point to the origin that is still in the green region
- The perceptron algorithm just puts us randomly in the green region



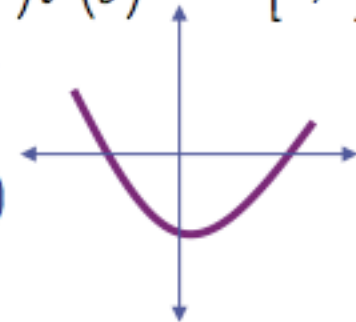
Convexity

- **Convex functions:** $f(tx + (1-t)y) \leq tf(x) + (1-t)f(y) \quad t \in [0,1]$

$$f(x) = \exp(x), \quad f(\vec{x}) = \vec{x}^T b + \frac{1}{2} \vec{x}^T H \vec{x}, \quad f(\vec{x}) = \vec{x}$$

Have non-negative second derivatives (bowls)

$$\frac{\partial^2 f(x)}{\partial x^2} = \exp(x), \quad \frac{\partial^2 f(\vec{x})}{\partial \vec{x} \partial \vec{x}} = H, \quad \frac{\partial^2 f(\vec{x})}{\partial \vec{x} \partial \vec{x}} = 0$$

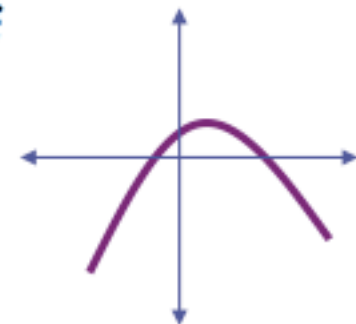


- **Concave functions:** $f(tx + (1-t)y) \geq tf(x) + (1-t)f(y) \quad t \in [0,1]$

$$f(x) = \log(x), \quad f(\vec{x}) = \vec{x}^T b - \frac{1}{2} \vec{x}^T H \vec{x}, \quad f(\vec{x}) = \vec{x}$$

Have non-positive second derivatives (caves)

$$\frac{\partial^2 f(x)}{\partial x^2} = -\frac{1}{x^2}, \quad \frac{\partial^2 f(\vec{x})}{\partial \vec{x} \partial \vec{x}} = -H, \quad \frac{\partial^2 f(\vec{x})}{\partial \vec{x} \partial \vec{x}} = 0$$



Duality

- Every convex function f has a dual f^* :
All tangent lines below it form an epigraph
The f^* gives the intercept for each slope.

$$f(x) = \max_{\lambda} (x^T \lambda - f^*(\lambda))$$

- Every concave function f has a dual f^* :
All tangent lines above it form an epigraph
The f^* gives the intercept for each slope.

$$f(x) = \min_{\lambda} (x^T \lambda - f^*(\lambda))$$

- This $*$ is called the **Legendre Transform** or **Fenchel Dual**

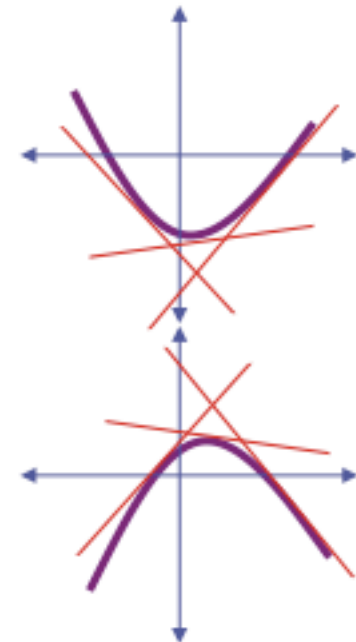
- The dual of the dual f^{**} is f

- Example: $f(x) = \frac{1}{2}cx^2 \rightarrow f^*(\lambda) = \frac{1}{2c}\lambda^2$

- We can replace a minimization over x like this

$$\min_x f(x) = \min_x \max_{\lambda} (\lambda x - f^*(\lambda))$$

...and can work with a maximization of its dual instead



Optimization: Inequality Constraints

- Problem: given a function of several variables, find its stationary point subject to one inequality constraint

- Formally (general case):
$$\begin{aligned} \max_{\mathbf{x}} \quad & f(\mathbf{x}) \\ \text{s.t.} \quad & g(\mathbf{x}) \geq 0 \end{aligned}$$

- Consider the geometry of the problem, there are now two solutions possible:

1. On the boundary (constraint is *active*, $g(x) = 0$)
2. Inside the region (constraint is *inactive*, $g(x) > 0$)

- For case 2, the constraint has no effect. Case 1 is analogous to equality constraint discussed previously, but the sign of the multiplier is crucial (gradient should be oriented away from the region $g(x) > 0$)

Optimization: Inequality Constraints

- For case 2 (region), the constraint has no effect.
- Case 1 (boundary) is analogous to equality constraint discussed previously, but the sign of the multiplier is crucial (gradient should be oriented away from the region defined by the constraint $g(x) > 0$)

1. Boundary: $\nabla f(x) = -\lambda \nabla g(x), \lambda > 0$

2. Region: $\nabla f(x) = 0 \quad \equiv \quad \nabla L(x, \lambda = 0)$

- We can combine both cases into one: $\lambda g(x) = 0$

KKT Conditions

- A. Define a function: $L(x, \lambda) = f(x) + \lambda g(x)$
- B. Find the stationary point of L with respect to $\{x, \lambda\}$ and subject to:

$$g(x) \geq 0, \quad \lambda \geq 0, \quad \lambda g(x) = 0$$

- These are known as the *Karush-Kuhn-Tucker* (KKT) conditions
- If we wish to minimize the function $f(x)$ we need to define the Lagrangian as:

$$L(x, \lambda) = f(x) - \lambda g(x)$$

Multiple Constraints

- Problem: given a function of several variables, find its stationary point subject to one or more equality and inequality constraints

- Formally (general case): $\max_{\mathbf{x}} f(\mathbf{x})$

$$s.t. \quad g_j(\mathbf{x}) = 0, \quad j = 1, \dots, J$$

$$h_k(\mathbf{x}) \geq 0, \quad k = 1, \dots, K$$

- Define the Lagrangian:

$$L(\mathbf{x}, \{\lambda\}, \{\mu\}) = f(\mathbf{x}) + \sum_{j=1}^J \lambda_j g_j(\mathbf{x}) + \sum_{k=1}^K \mu_k h_k(\mathbf{x}),$$

$$s.t.: \quad \forall k: \mu_k \geq 0, \quad \mu_k h_k(\mathbf{x}) = 0$$

Dual Form Derivation

- Recall optimal hyperplane problem in primal space:

$$\min_w \frac{1}{2} \|w\|^2 \quad s.t \quad \forall i: y_i(w^T x_i + b) \geq 1$$

- This is a convex program, define the Lagrangian and find stationary point:

$$L(w, b, \alpha) = \frac{1}{2} \|w\|^2 - \sum_{i=1}^{\ell} \alpha_i [y_i(w^T x_i + b) - 1], \quad \alpha_i \geq 0$$

- Minimize L over {w,b}, maximize over {alphas}:

$$\frac{\partial L(w, b, \alpha)}{\partial w} = w - \sum_{i=1}^{\ell} \alpha_i y_i x_i = 0 \Rightarrow w = \sum_{i=1}^{\ell} y_i \alpha_i x_i$$

$$\frac{\partial L(w, b, \alpha)}{\partial b} = - \sum_{i=1}^{\ell} \alpha_i y_i = 0 \Rightarrow \sum_{i=1}^{\ell} y_i \alpha_i = 0$$

Dual Form

This is a convex program, define the Lagrangian and find stationary point:

$$L(w, b, \alpha) = \frac{1}{2} \|w\|^2 - \sum_{i=1}^{\ell} \alpha_i [y_i (w^T x_i + b) - 1], \quad \alpha_i \geq 0$$

- Minimize L over {w,b}, maximize over {alphas}:

$$\frac{\partial L(w, b, \alpha)}{\partial w} \Rightarrow w = \sum_{i=1}^{\ell} y_i \alpha_i x_i, \quad \frac{\partial L(w, b, \alpha)}{\partial b} \Rightarrow \sum_{i=1}^{\ell} y_i \alpha_i = 0, \quad \alpha_i \geq 0$$

- Plug back into the Lagrangian and get the dual form:

$$\max_{\alpha} D(\alpha) = \sum_{i=1}^{\ell} \alpha_i - \frac{1}{2} \sum_{i,j=1}^{\ell} y_i y_j \alpha_i \alpha_j (x_i \cdot x_j)$$

$$s.t: \sum_{i=1}^{\ell} y_i \alpha_i = 0, \quad \alpha_i \geq 0$$

Why Solve in Dual Space?

- QP runs in cubic polynomial time (in terms of # of variables)
- QP in primal space has complexity $O(d^3)$, where d is the dimensionality of the input vectors (weight vector)
- QP in dual space has complexity $O(\ell^3)$, where ℓ is the number of examples
- More importantly: dual space yields “deeper results”

$$\min_w \frac{1}{2} \|w\|^2$$

$$s.t \quad \forall i: y_i(w^T x_i + b) \geq 1$$

$$\max_{\alpha} D(\alpha) = \sum_{i=1}^{\ell} \alpha_i - \frac{1}{2} \sum_{i,j=1}^{\ell} y_i y_j \alpha_i \alpha_j (x_i \cdot x_j)$$

$$s.t: \sum_{i=1}^{\ell} y_i \alpha_i = 0, \alpha_i \geq 0$$

$$\max_{\alpha} D(\alpha) \Rightarrow \alpha^* \Rightarrow w^* = \sum_{i=1}^{\ell} y_i \alpha_i^* x_i$$