

# **Data Structures and Algorithms**

**Session 10. February 23, 2009**

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# Announcements

- \* Homework 2 due now.
- \* Homework 3 to be posted after class. Due 3/9
- \* Midterm review March 9<sup>th</sup>
- \* Midterm Exam March 11<sup>th</sup>
  - \* closed book, closed notes

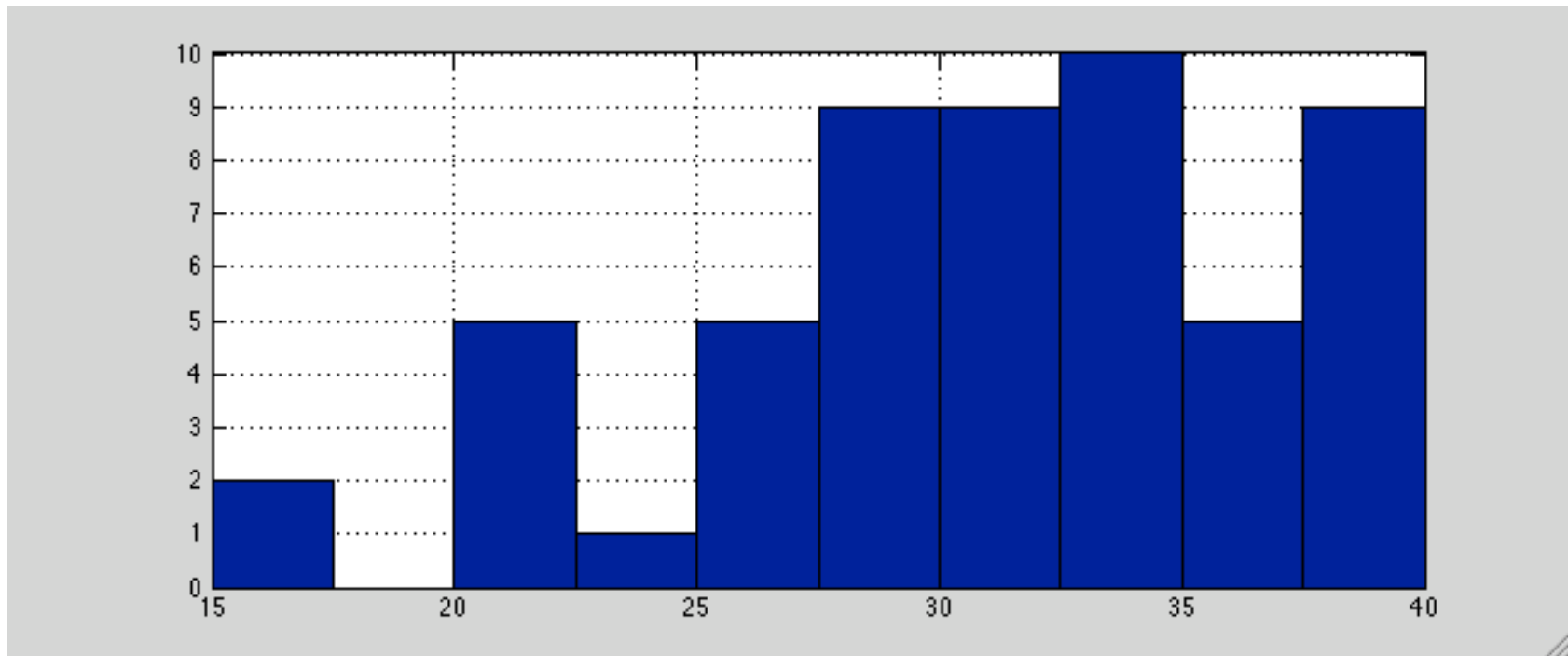
# Review

- \* Brief look at tradeoffs
- \* Balanced (AVL) Binary Search Trees
  - \* AVL Tree property
  - \* Tree Rotations
  - \* Worst case depth analysis

# Today's Plan

- \* HW1 solutions (long overdue)
- \* Splay Trees
- \* Prefix Trees (tries)

# HW1 Histogram



\* Average was 31.25, std-deviation 6

# Amortized Running Time

- \* Don't guarantee *each* operation is  **$O(\log N)$**
- \* Instead, prove that  **$M$**  operations take  **$O(M \log N)$**
- \* Then each operation has an **amortized** running time of  **$O(\log N)$**

# Splay Trees

- \* Like AVL trees, use the standard binary search tree property
- \* After any operation on a node, make that node the new root of the tree
- \* Make the node the root by repeating one of two moves that make the tree more spread out

# Informal Justification

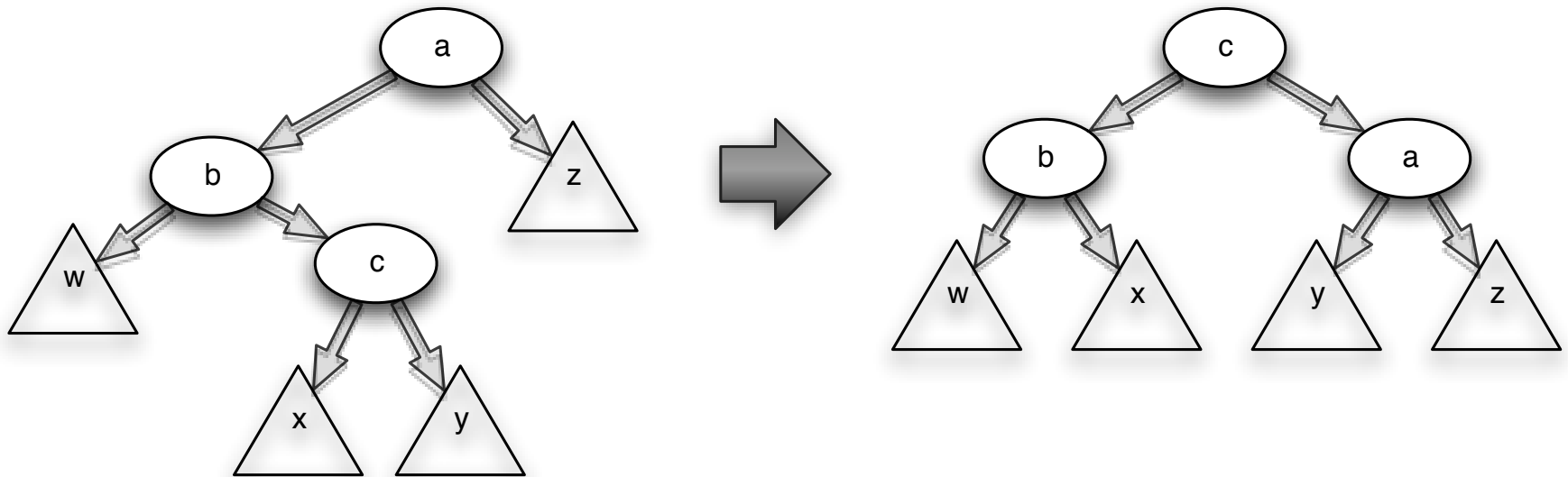
- \* Similar to *caching*.
- \* Heuristically, data that is accessed tends to be accessed often.
- \* Easier to implement than AVL trees
  - \* No height info



# Easy cases

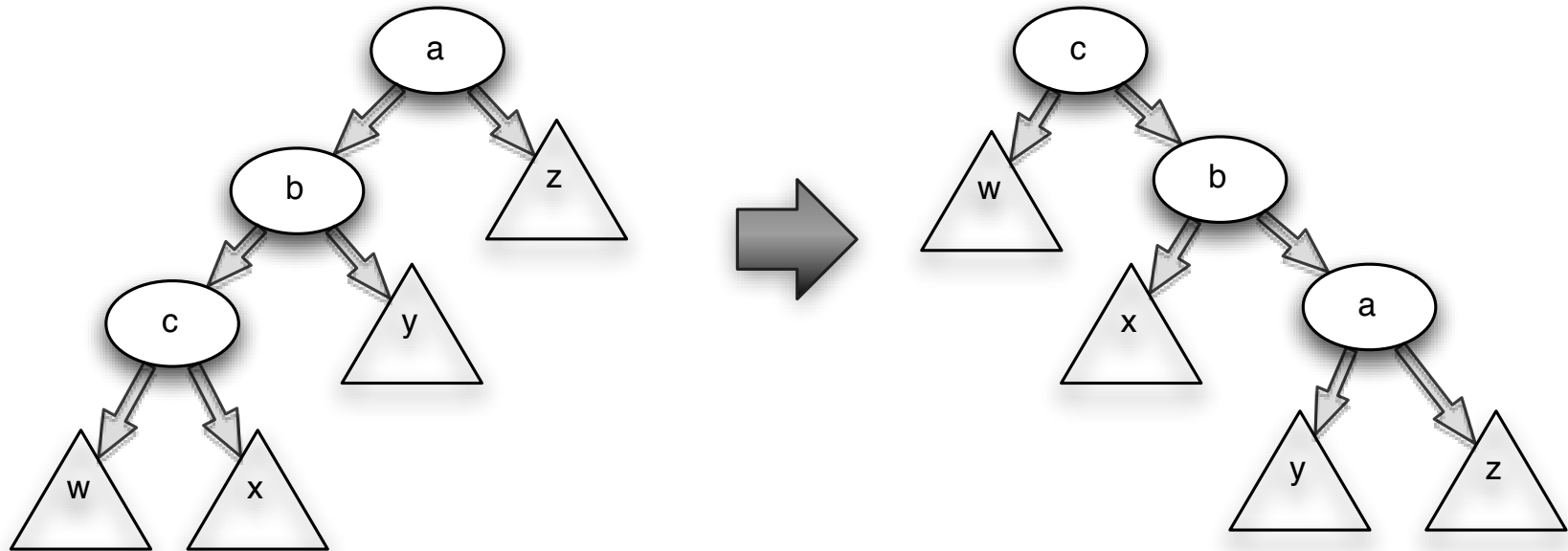
- \* If node is root, do nothing
- \* If node is child of root, do single AVL rotation
- \* Otherwise, node has a grandparent, and there are two cases

# Case 1: zig-zag



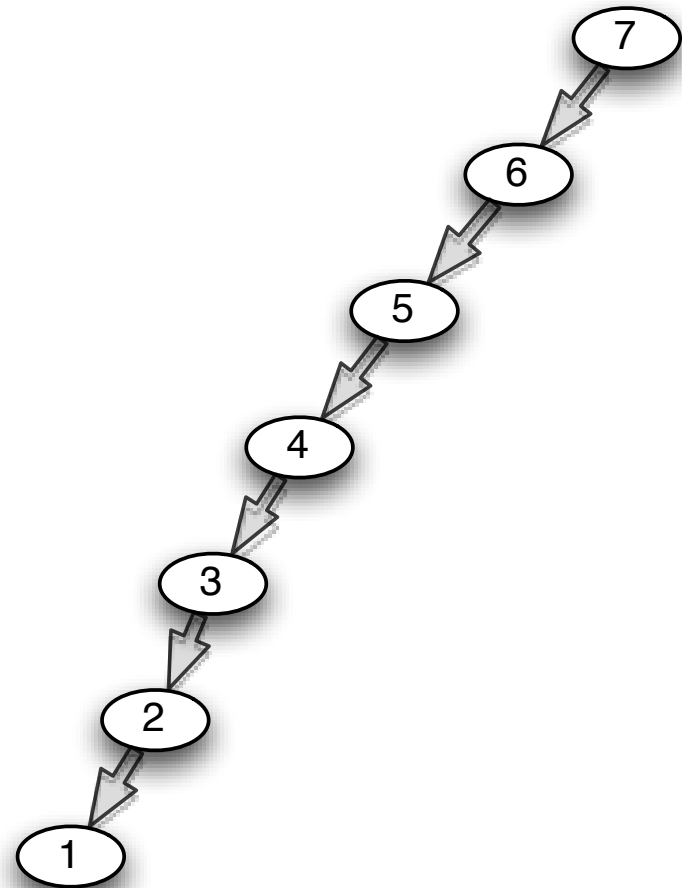
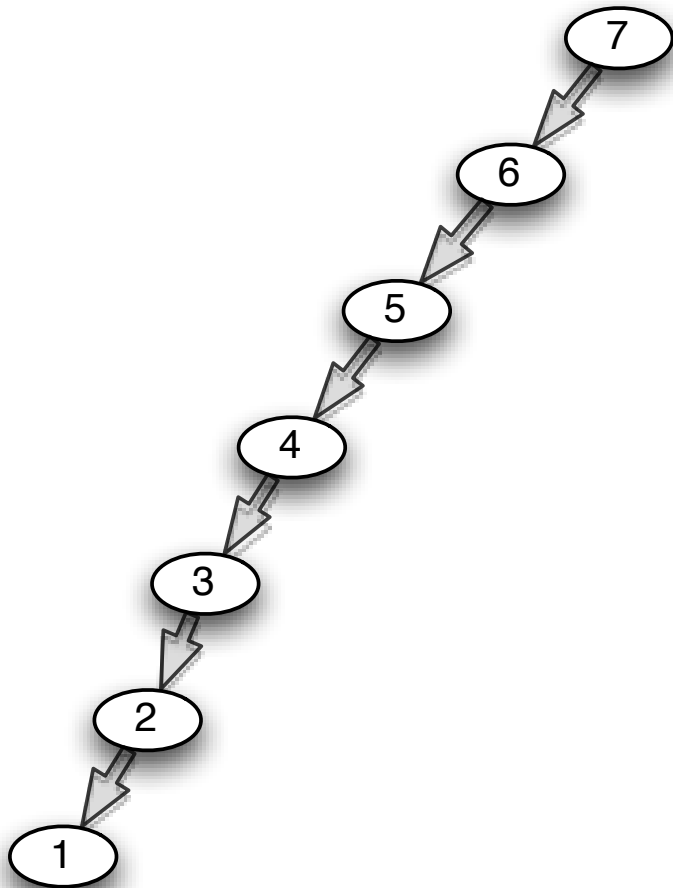
- \* Use when the node is the right child of a left child (or left-right)
- \* Double rotate, just like AVL tree

# Case 2: zig-zig

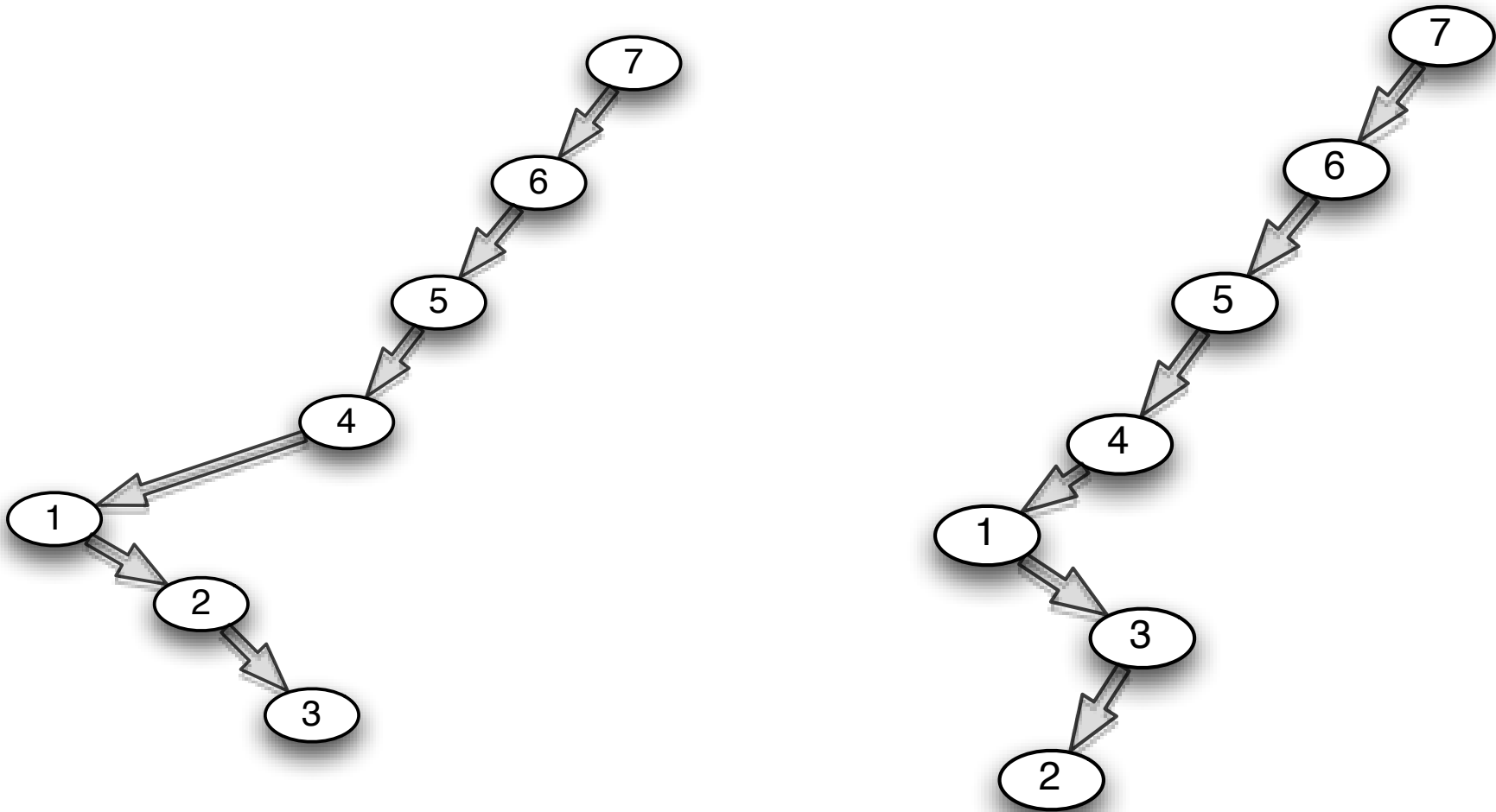


- \* Use when node is the right-right child (or left-left)
- \* Reverse the order of grandparent->parent->node
- \* Make it node->parent->grandparent

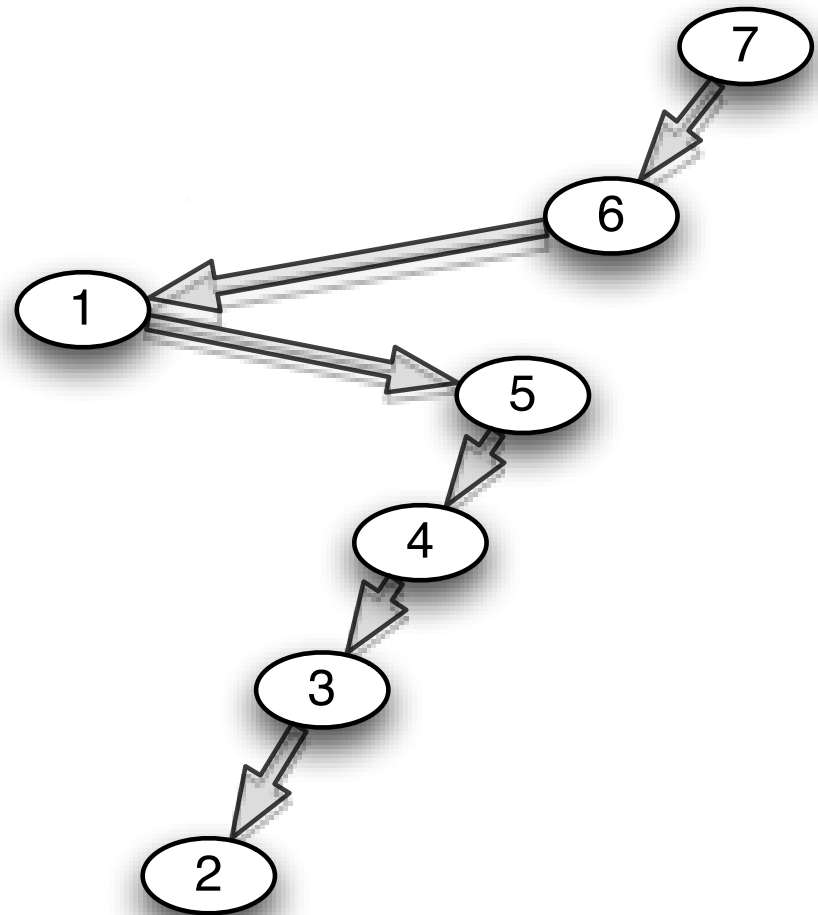
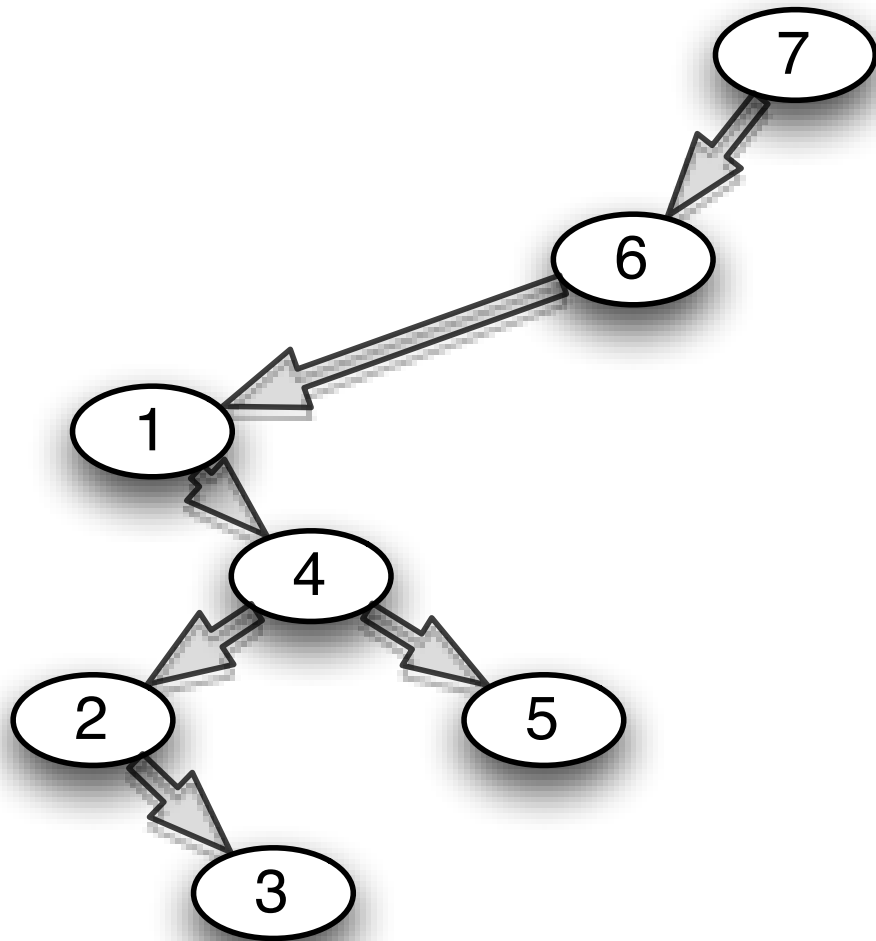
# Case 2 versus Single Rotations 1



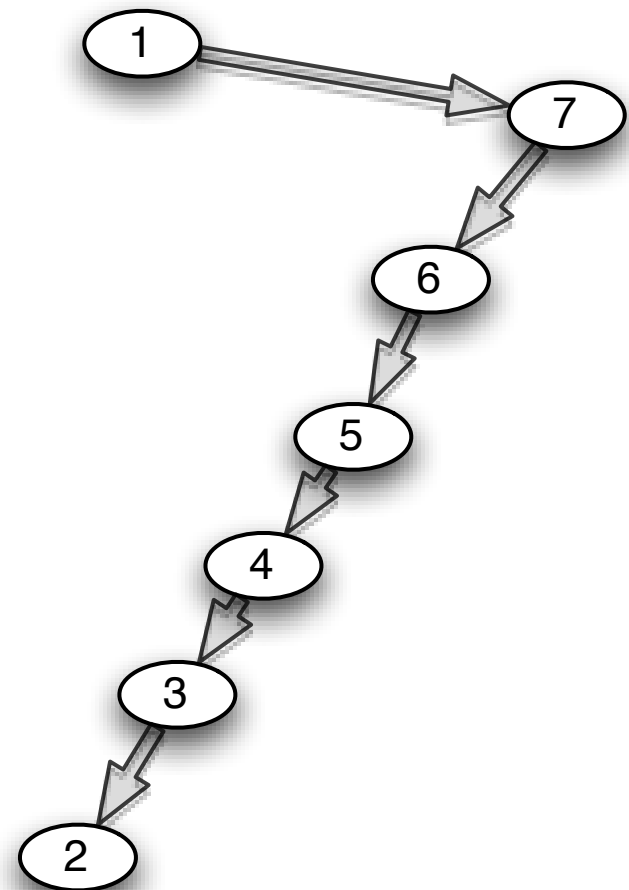
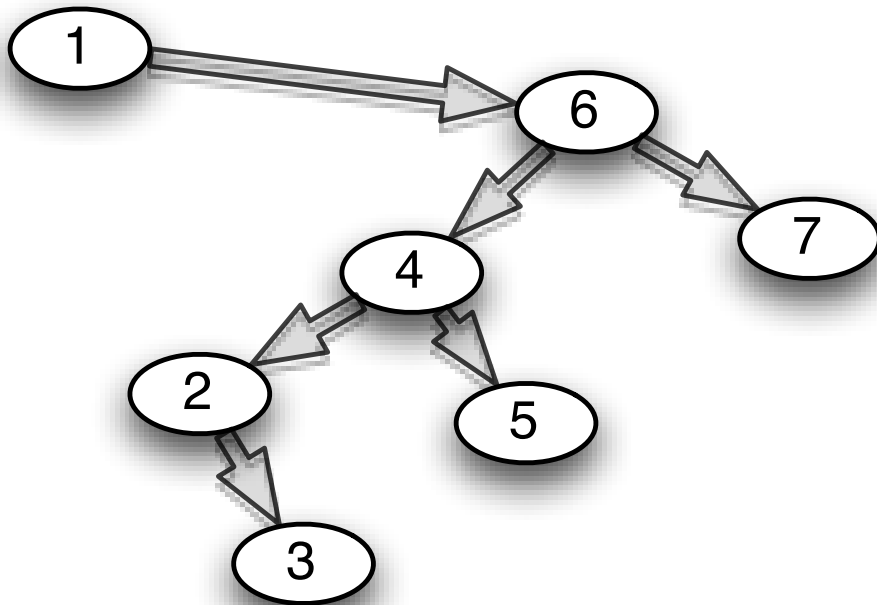
# Case 2 versus Single Rotations 2



# Case 2 versus Single Rotations 3



# Case 2 versus Single Rotations 4



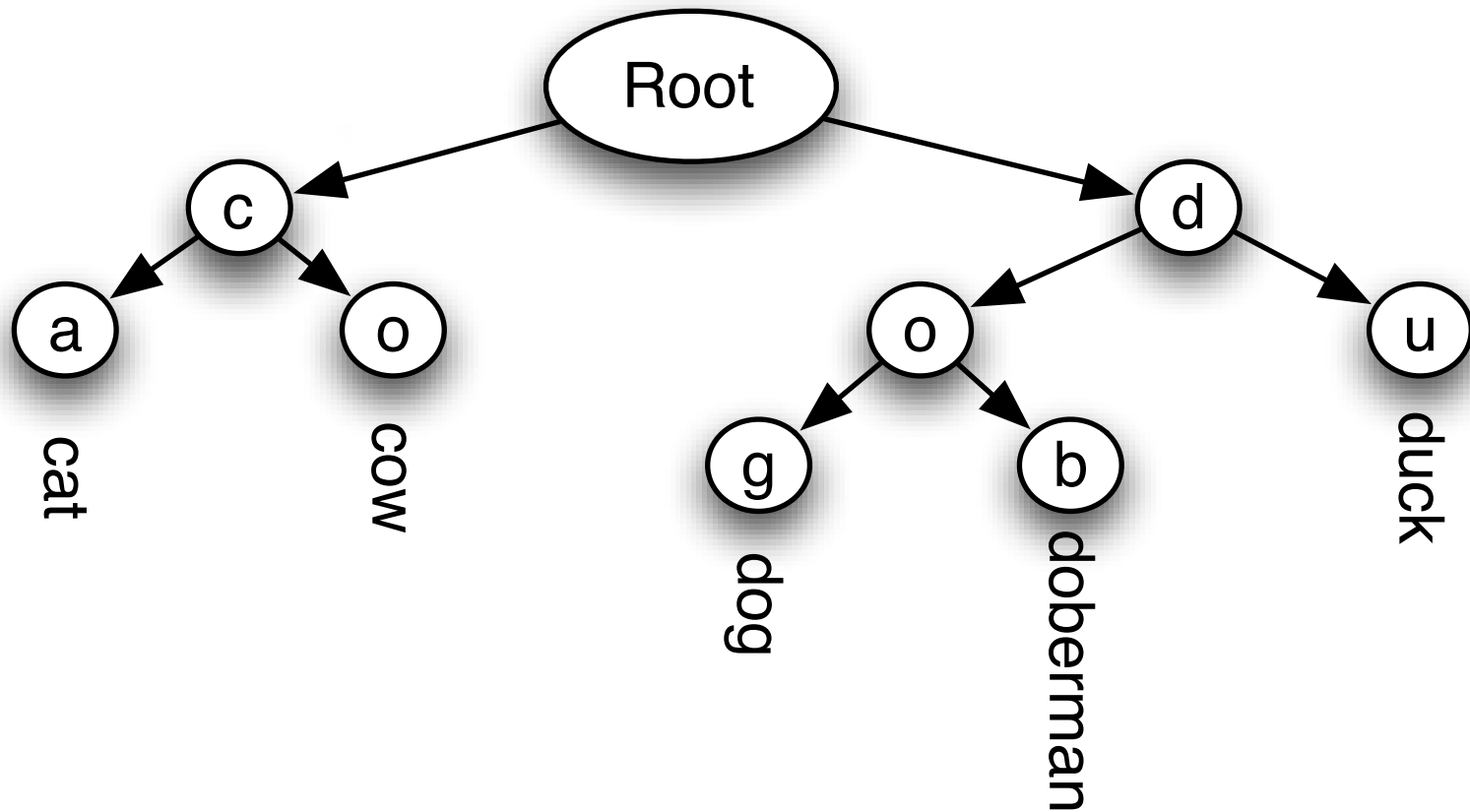
# Prefix Trees (Tries)

- \* Nicknamed “Trie”, short for **retrieval**
- \* Efficiently store objects for fast retrieval via keys
  - \* Usually key is a String
- \* Basic strategy:
  - \* split into sub-tries based on current letter



# Trie Example

- \* “cat”, “cow”, “dog”, “doberman”, “duck”



# Trie Details

- \* Not all words are at leaves
  - \* cat, cataclysm, cataclysmic
- \* Initially, one letter is enough to uniquely identify
- \* When a new word is inserted that conflicts, need to branch
  - \* Originally-unique word must be moved to lower level

# Trie Analysis

- \* In the worst case, inserting a key of length **k** or (looking up) is  **$O(k)$**
- \* This is not dependent on **N!** (surprise, not factorial)
- \* Much better than  **$\log(N)$**  for huge data like dictionaries
- \* Sometimes we can access words even faster.
  - \* E.g., we can find qwerty uniquely with just “qw”