COMS E6998-9: Algorithms for Massive Data (Fall'25)

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# Lecture 6: Compressed Sensing

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### 1 Review

**Definition 1.** A mapping  $\varphi : \mathbb{R}^n \to \mathbb{R}^m$  is called a  $(k,\delta)$  Oblivious Subspace Embedding (OSE) if for any k-dimensional subspace  $U \subseteq \mathbb{R}^n$ :

$$Pr_{\varphi}\left[\forall z \in U : \frac{\|\varphi(z)\|^2}{\|z\|^2} \in (1 \pm \epsilon)\right] \ge 1 - \delta$$

**Definition 2.** A matrix  $A \in \mathbb{R}^{m \times n}$  is said to satisfy the  $(k, \epsilon)$ Restricted Isometry Property (RIP) if for all k-sparse vectors  $z \in \mathbb{R}^n$ :

$$||Az||_2 \in (1 \pm \epsilon)||z||_2$$

### 2 Guarantees of OSE and RIP

We will start by showing that OSE implies RIP. So we want to prove the following theorem:

**Theorem 3.** if A is an  $(k, \frac{0.1}{\binom{n}{k}})$ -OSE, then A satisfies  $(k, \epsilon)$ -RIP.

*Proof.* We want to prove:

$$Pr_A\left[\forall z \in S_k : \frac{\|Az\|}{\|z\|} \in (1 \pm \epsilon)\right] \ge 0.9$$

where  $S_k$  is the set of all k-sparse vectors in  $\mathbb{R}^n$ .

We will write  $S_k$  as a union of subspaces:

$$S_k = \bigcup_{T \subseteq [n], |T| = k} U_T$$

where  $U_T = \{ z \in \mathbb{R}^n : z_i = 0 \text{ for } i \notin T \}$ 

Note that  $U_T$  is a k-dimensional subspace of  $\mathbb{R}^n$  and the number of such subspaces is  $\binom{n}{k}$ .

By the definition of OSE and using union bound, we have:

$$Pr_{A}\left[\exists z \in S_{k} : \frac{\|Az\|}{\|z\|} \notin (1 \pm \epsilon)\right]$$

$$\leq Pr_{A}\left[\exists z \in U_{T_{1}} : \frac{\|Az\|}{\|z\|} \notin (1 \pm \epsilon)\right] + Pr_{A}\left[\exists z \in U_{T_{2}} : \frac{\|Az\|}{\|z\|} \notin (1 \pm \epsilon)\right] + \dots$$

$$\leq \sum_{T \subseteq [n], |T| = k} Pr_{A}\left[\exists z \in U_{T} : \frac{\|Az\|}{\|z\|} \notin (1 \pm \epsilon)\right]$$

$$\leq \binom{n}{k} \cdot \frac{0.1}{\binom{n}{k}} = 0.1$$

Corollary 4. For fixed  $\epsilon > 0$ , if A is a scaled Gaussian  $m \times n$  matrix with  $m = O(k + \log \frac{1}{\delta})$  then there exists A that satisfies  $(k, \epsilon)$ -RIP  $m = O(\log \binom{n}{k}) \le O(\log (\frac{n}{k})^k) = O(k \cdot \log \frac{n}{k})$ 

## 3 proving the L1 theorem

Recall the Problem setup from last lecture:

We are given an input x and measurements y = Ax and we want to recover x from y using L1 minimization.

$$L_1(y) = \underset{Ax=y}{\operatorname{argmin}} ||x||_1 \triangleq x^*$$

Recall that we define the distance from k sparsity as:

$$\operatorname{Err}_{1}^{k}(x) = \min_{x' \text{ is k-sparse}} \|x - x'\|_{1}$$

We want to prove the following theorem:

**Theorem 5.** if A is  $(4k, \epsilon)$ -RIP for a small enough  $\epsilon$ , then  $||x^* - x|| \leq O(1) \cdot Err_1^k(x)$ 

The proof will rely on some intermediate definitions and lemmas.

**Definition 6.** We say that A satisfies  $(k,\epsilon)$  - nullspace property if  $\forall T \subseteq [n], |T| \leq k, \forall \eta \in \mathbb{R}^n$  s.t.  $A\eta = 0$ 

$$\|\eta_T\|_1 \le \epsilon \|\eta_{-T}\|_1 \iff \|\eta\|_1 \le (1+\epsilon)\|\eta_{-T}\|_1$$

We will also use the following two lemmas to prove the above theorem.

**Lemma 7.** Fix  $\forall r \in N$ . If A is  $((2+r)k, \epsilon)$ -RIP for a small enough  $\epsilon$ , then A satisfies  $(2k, \sqrt{\frac{2}{r}} \cdot \frac{1+\epsilon}{1-\epsilon})$  -nullspace property.

**Lemma 8.** if A satisfies  $(2k,\epsilon)$ -nullspace property for a small enough  $\epsilon$ , then for any x, the output of L1 minimization  $x^*$  satisfies:

$$||x^* - x||_1 \le 2 \cdot \frac{1 + \epsilon}{1 - \epsilon} \cdot Err_1^k(x)$$

We will first prove Lemma 8.

In the proof of the lemma, k-RIPness will be actually good enough. But it's easier to see why 2k is useful: consider k-sparse x and a k-sparse potential solution  $x^*$  such that  $Ax = Ax^* = y$ . Then, then the nullspace property of  $\eta = x - x^*$  (a 2k-sparse vectors) immediately implies that  $\eta = 0$ , and hence  $x = x^*$ .

#### **Proof of Lemma 8:**

let  $\eta \triangleq x^* - x \implies A\eta = Ax^* - Ax = y - y = 0$  Let T be the indices of the k largest entries of x in absolute value.

By the definition of L1 minimization, we have:

$$||x^*||_1 \le ||x||_1 \implies ||x_T^*||_1 + ||x_{-T}^*||_1 \le ||x_T||_1 + ||x_{-T}||_1$$

by triangle inequality, we have:

$$||x_T^*||_1 \ge ||x_T||_1 - ||x_T^* - x_T|| = ||x_T||_1 - ||\eta_T||_1$$
$$||x_{-T}^*||_1 \ge ||\eta_{-T}||_1 - ||x_{-T}||_1$$

So we have:

$$||x_T||_1 - ||\eta_T||_1 + ||\eta_{-T}||_1 - ||x_{-T}||_1 \le ||x_T||_1 + ||x_{-T}||_1$$

$$\implies ||\eta_{-T}||_1 \le ||\eta_T||_1 + 2||x_{-T}||_1$$

By the nullspace property, we have:

$$\|\eta_{-T}\|_{1}(1-\epsilon) \leq 2\|x_{-T}\|_{1}$$

$$\implies \|\eta\|_{1} \leq \frac{2 \cdot (1+\epsilon)}{1-\epsilon} \|x_{-T}\|_{1}$$

$$\implies \|x^{*} - x\|_{1} \leq 2 \cdot \frac{1+\epsilon}{1-\epsilon} \cdot \operatorname{Err}_{1}^{k}(x)$$

Where the last step follows from the definition of  $\operatorname{Err}_1^k(x)$ .

This completes the proof of Lemma 8.

We will now prove Lemma 7.

#### Proof of Lemma 7:

Let  $M \triangleq rk$ , define  $T_0 = T$ ,  $T_1 =$  indices of the largest M entries of  $\eta_{-T}$ ,  $T_2 =$  indices of the next largest M entries of  $\eta_{-T}$  and so on up to  $T_s$ .

define 
$$\eta_0 = \eta_{T_0}, \eta_1 = \eta_{T_1}, \dots, \eta_s = \eta_{T_s}.$$

We can write  $\eta = \eta_0 + \eta_1 + \ldots + \eta_s$ .

Since  $A\eta = 0$ , we have:

$$A(\eta_0 + \eta_1) = -A(\eta_2 + \ldots + \eta_s)$$

Taking norms on both sides, we have:

$$||A\eta_0 + \eta_1||_2 = ||A(\eta_2 + \ldots + \eta_s)||_2$$

By triangle inequality (and since  $\eta_0$  and  $\eta_1$  are non-zero on disjoint coordinates), we have:

$$||A\eta_0||_2 \le ||A\eta_0 + A\eta_1||_2 \le ||A\eta_2||_2 + \ldots + ||A\eta_s||_2$$

By the RIP property, we have:

$$(1 - \epsilon) \|\eta_0\|_2 \le \|A\eta_0\|_2 \le (1 + \epsilon) \|\eta_0\|_2$$

and hence

$$\|\eta_T\|_2 \le \frac{1}{1-\epsilon} (\|A\eta_2\|_2 + \dots + \|A\eta_s\|_2)$$
  
$$\le \frac{1+\epsilon}{1-\epsilon} (\|\eta_2\|_2 + \dots + \|\eta_s\|_2)$$

Let  $\eta^{(i)} = i^{th}$  coord of  $\eta$ . for all  $j \geq 2$  and for all  $i \in T_j, |\eta^{(i)}| \leq \frac{\|\eta_{T_{j-1}}\|_1}{M}$ This implies that:

$$\|\eta_j\|_2^2 \le \sum_{i \in T_j} \left(\frac{\|\eta_{T_{j-1}}\|_1}{M}\right)^2 = M \cdot \left(\frac{\|\eta_{T_{j-1}}\|_1}{M}\right)^2 = \frac{\|\eta_{T_{j-1}}\|_1^2}{M}$$

$$\implies \|\eta_j\|_2 \le \frac{\|\eta_{T_{j-1}}\|_1}{\sqrt{M}}$$

So we have:

$$\|\eta_T\|_2 \le \frac{1+\epsilon}{1-\epsilon} \cdot \sum_{j=2}^s \frac{\|\eta_{T_{j-1}}\|_1}{\sqrt{M}} \le \frac{1+\epsilon}{1-\epsilon} \cdot \frac{\|\eta_{-T}\|_1}{\sqrt{M}}$$

Now, by the standard inequality between  $\ell_1$  and  $\ell_2$  norms (for vectors of dimension 2k or sparsity 2k), we have:

$$\|\eta_T\|_1 \le \sqrt{2k} \|\eta_T\|_2 \le \sqrt{2k} \cdot \frac{1+\epsilon}{1-\epsilon} \cdot \frac{\|\eta_{-T}\|_1}{\sqrt{M}}$$

substituting M = rk, we have:

$$\|\eta_T\|_1 \le \sqrt{\frac{2}{r}} \cdot \frac{1+\epsilon}{1-\epsilon} \cdot \|\eta_{-T}\|_1.$$