

# Advanced Algorithms (4232)

Lecture 1. 1/18/2022.

1) Hashing:  $h: U \rightarrow [n] = \{1, 2, \dots, n\}$

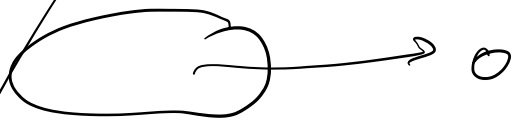
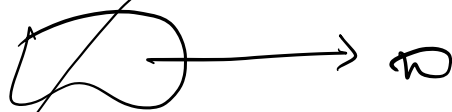
Dictionary problem: index set  $S \subseteq U$   
s.t. quickly  $x \in S$

$O(1)$  expected time.

perfect hashing:  $O(1)$  worst-case g.t.

2) Sketching / Streaming algo's.

↳ functional compression.



similar/diss.

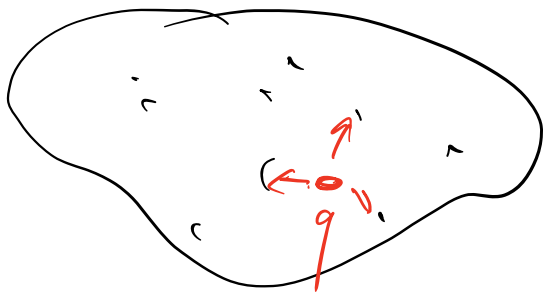
speedy up N.L.A.



### 3) Nearest Neighbor search (NNS).

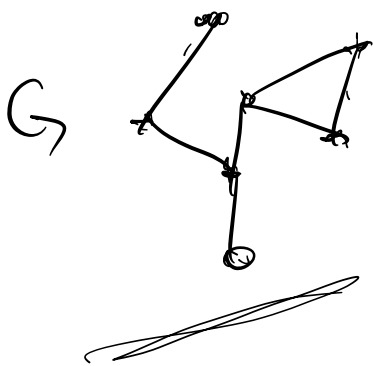
problem: given  $n$  points in high-d space

query: given  $q$ , report  
closest point from  $d_s$



### 4) Graph algorithms: max-flow.

### 5) Spectral graph theory.



$$A_G = \begin{matrix} & \begin{matrix} 1 & 2 & \dots & 7 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ \vdots \\ 7 \end{matrix} & \begin{matrix} & & & \\ & & & \\ & & \square & \\ & & & \end{matrix} \end{matrix}$$

$$\begin{matrix} \lambda_1 & \dots & \lambda_n \\ v_1 & \dots & v_n \end{matrix}$$

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## Spectral clustering

### 6) Optimization:

$$\min f(x)$$

$$x \in \mathbb{R}^n$$

$$x \in C$$

LP:  $f, C \rightarrow$  linear

IPM: interior-point method.

MWL: learning from experts.

### 7) Large-scale models.

- parallel / cluster computing.

Map Reduce.

- I/O model: cache.

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### 1) Randomized:

$\forall$  input  $x$ :  $\Pr_{\text{Alg}} [ \text{Alg correct} ] \geq 90\%$

2) Approximate:

$$a \in \hat{Q} \leq c \cdot a$$

↑                    ↑                    ↑  
correct            Alg.                    approx  $\geq 1$ .  
answer

1+2:

$$\text{Pr}_{\text{Alg}} [ a \in \hat{Q} \leq c \cdot a ] \geq 90\%$$

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Admin: Pre-reqs mathematical mat.

$\Theta(\cdot), \Theta(\cdot)$ .

Self-evaluation test

2022

Grading:

1) HW : 5 → 50%

2) Mid terms → 25%

3) Final project. team-based.  
2-4.

→ reading :

- implementations
- research focus.

Proposal: 1-4 pages.

Final: ~10 pages.

### Policy for HK3:

- collaboration

- late: .5 free late days.

◦ 10% penalty/day.

◦ max is 7 days. (exc. HK3).

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Problem: count up to  $n$ .

Space  $\lceil \lg_2 n \rceil$  bits to store number  
up to  $n$ .

Thm: unless allow random and approx,  
can't do better.

Morris's algo: randomized & approx.

use space  $O(\lg \lg n)$

$$C = O(1). \quad (1 + \epsilon, \quad \epsilon = 0.1)$$

$$n = 10^9$$

$$\lg_2 n \approx 30$$

$\lg \lg n =$  single digit #s.

Algo:

keep register  $X$ .

- init:  $X = 0$ .

- @ button press:  $X = \begin{cases} X+1, & w/Pr = 2^{-X} \\ X, & \text{others.} \end{cases}$

- Estimator:  $\hat{n} = 2^X - 1$ .

Intuition:

$$n = \overbrace{1100\dots 10}^{(\lg_2 n)}$$

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$$\overbrace{1100\dots 10}$$



$X$  stores the index of  
msb of  $n$ .

Def: expectation  $\mathbb{E}\{X\}$ ,  $X = \text{r.v.}$

Fact: linearity of  $\mathbb{E}$ :  $\mathbb{E}[X_1 + X_2]$   
 $= \mathbb{E}[X_1] + \mathbb{E}[X_2]$ .

Def: variance:  $\text{Var}[X] \stackrel{\text{def}}{=} \mathbb{E}[(X - \mathbb{E}[X])^2]$ .

Concentration bounds:

Markov's bound:  $\forall$  r.v.  $X \geq 0$ :

$$\Pr[X \geq \lambda] \leq \frac{\mathbb{E}[X]}{\lambda}.$$

Chebyshev's bound:  $\forall$  r.v.  $X$ :

$$\Pr[(X - \mathbb{E}[X])^2 \geq \lambda^2] \leq \frac{\text{Var}[X]}{\lambda^2}.$$