

Lecture 7      2/2/21

Last time:  $f \in \mathbb{R}_+^n$

goal: estimate  $F_2 = \|f\|_2^2$ .

Toolt algos: for  $i=1..k$ :  $k = \Theta\left(\frac{1}{\epsilon^2}\right)$   
picked  $\sigma_{ij} \in \{\pm 1\}$ ,  $j \in [n]$ .

sketch:  $z_i = \sum_{j=1}^n \sigma_{ij} f_j$ .

est:  $\frac{1}{k} \sum z_i^2 \stackrel{\Delta}{=} Z^2$ .

Proved:

$$\mathbb{E}[Z^2] = F_2$$

$$\text{Var}[Z^2] \leq \frac{4F_2}{k}$$

} by Cheb.

$$\Rightarrow Z^2 = (1 \pm \epsilon) F_2$$

with prob  $\geq 90\%$

if  $k \geq \Omega\left(\frac{1}{\epsilon^2}\right)$ .

Also: what if Prob. fail  $\leq \delta$  param.

$$Z^2 = (1 \pm \epsilon) \cdot F_2 \text{ with prob. } \geq 1 - \delta,$$

as long as  $k \geq \Omega\left(\frac{1}{\delta} \cdot \frac{1}{\epsilon^2}\right)$ .  
(by Chebyshev).

A question today: can we get better dep. on  $\delta$ ?

Yes, in fact  $k = O\left(\frac{\log(1/\delta)}{\epsilon^2}\right)$  is enough.

$T_{0,k}$  is a linear sketch:

$$(z_1, z_2, \dots, z_k)^T = \frac{1}{\sqrt{k}} \begin{bmatrix} \sigma_{11} & \sigma_{12} & \dots & \sigma_{1n} \\ \sigma_{21} & \sigma_{22} & \dots & \sigma_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ \sigma_{k1} & \dots & \dots & \sigma_{kn} \end{bmatrix} \begin{bmatrix} f_1 \\ f_2 \\ \vdots \\ f_n \end{bmatrix}$$

each  $\sigma \in \{\pm 1\}$  i.i.d.

Est:  $\sum_{i=1}^k z_i^2$

Corollary: sketch of sum of  $f^1$  &  $f^2$   
 $=$  sum of sketches of  $f^1$  &  $f^2$

Can see sketch:  $\varphi: \mathbb{R}^n \rightarrow \mathbb{R}^k$   
 $\varphi(f) = \begin{bmatrix} \sigma's \end{bmatrix} \cdot \begin{bmatrix} f \end{bmatrix}$

Est:  $\sum_i (\varphi(f)_i)^2 \approx F_2$

$$\downarrow$$

$$\| \varphi(f) \|_2^2 = (1 \pm \epsilon) \cdot \|f\|_2^2 = (1 \pm \epsilon) \cdot \|f\|_2^2$$

$\uparrow$   $l_2$  norm of a vector of dim.  $k \approx O(\frac{k^{1/2}}{\epsilon^2})$ .

$\uparrow$   $l_2$  norm of a vector of dim.  $n$

An example of Dimension Reduction

Lemma [Johnson-Lindenstrauss '84]:  
 $\forall \epsilon > 0$ , there is a random  $\varphi: \mathbb{R}^n \rightarrow \mathbb{R}^k$   
 s.t.  $\forall x, y \in \mathbb{R}^n$ :

$$\Pr_{\varphi} [ \| \varphi(x) - \varphi(y) \| = (1 \pm \epsilon) \cdot \|x - y\| ] \geq 1 - e^{-\epsilon^2 k / 9}$$

Proof: If  $\varphi$  is linear:

$\varphi(x) - \varphi(y)$  approximates  $\|x - y\|$

equivalent:

$$\varphi(x) - \varphi(y) = \varphi(x - y) \text{ approximates } \|z\|, z \stackrel{\Delta}{=} x - y.$$

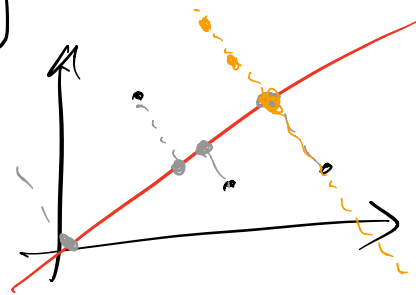
Enough to show  $\|\varphi(z)\| \approx \|z\|$ ,  
 $\forall z \in \mathbb{R}^n$ .

Which  $\varphi$ ?

1) Total  $\varphi$  actually works.

$$\varphi(z) = \begin{matrix} n \\ \pm 1 \end{matrix} \cdot z.$$

2)  $\varphi$ : pick a random  $k$ -dim.  
 linear subspace of  $\mathbb{R}^n$ , project  
 argument onto subspace.



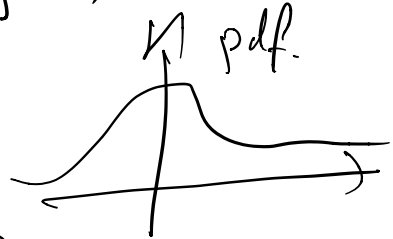
$$3) \varphi(z) = \begin{matrix} g_{ij} \end{matrix} \cdot z.$$

a random iid  
 Gaussian / Normal.

Obs: if no randomness, then  $\mathcal{F} \times \mathcal{G}$  will collapse.  $\forall \varphi: \mathbb{R}^n \rightarrow \mathbb{R}^{n+1}$

$$\varphi(z) = \frac{1}{\sqrt{k}} \left( \sum_{j=1}^n g_{1j} z_j, \sum g_{2j} z_j, \dots, \sum g_{kj} z_j \right).$$

$g_{ij} \sim \text{Gaussian } N(0,1)$ .



$g$  has pdf  $f(g) = \frac{1}{\sqrt{2\pi}} e^{-g^2/2}$ .

Property: spherically symmetric.

$(g_1, \dots, g_n) \leftarrow$  pdf is i.i.d.

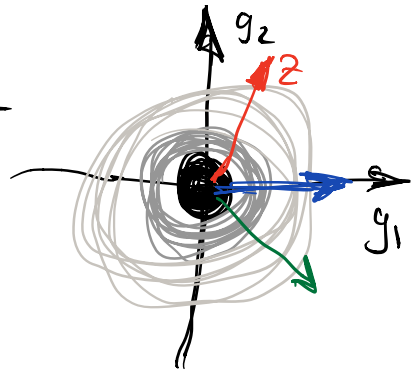
$\uparrow \quad \uparrow$   
i.i.d.  $\in N(0,1)$

pdf  $(g_1, \dots, g_n)$  depends only

on  $\|g\|$ .

$$\begin{aligned} \text{pdf}(g_1, g_2, \dots, g_n) &= \prod_{i=1}^n \left( \frac{1}{\sqrt{2\pi}} \right) \cdot e^{-g_i^2/2} \\ &= \frac{1}{(2\pi)^{n/2}} \cdot e^{-\sum_{i=1}^n g_i^2/2} \\ &= \frac{1}{(2\pi)^{n/2}} \cdot e^{-\|g\|_2^2/2}. \end{aligned}$$

Corollary:  $z \cdot g$  has same  
distribution as if  
 $z = (\|z\|, \vartheta, 0, \dots, 0)$ .



$\Rightarrow z \cdot g$  has distribution  $\|z\| \cdot g'$ ,  
 $g' \in N(0, 1)$

Why Gaussian distrib. can replace  $\pm 1$   
in  $T_0 \mathcal{W}$ ?

1)  $\mathbb{E}[\sigma_{ij}] = 0$

2)  $\mathbb{E}[\sigma_{ij}^2] = 1$ .

3)  $\mathbb{E}[\sigma_{ij}^4] = 1$ . ok if  $O(1)$ .

To get guarantees of  $T_0 \mathcal{W}$ , enough  
to replace  $\{\sigma_{ij} \in \pm 1\}$  with  $\#$  r.v.  
from dist. sat  $1 \div 3$ .

E.g.  $N(0, 1)$  satisfies them.

So far:  $\varphi(z) = \frac{1}{\sqrt{k}} \left( \underbrace{\sum g_{1j} z_j}_{\text{distributed as } g^{(1)} \cdot \|z\|}, \underbrace{\sum g_{2j} z_j}_{\dots}, \dots \right)$

$\varphi(z)$  is distributed as

$$\frac{1}{\sqrt{k}} \left( \|z\| \cdot g^{(1)}, \|z\| \cdot g^{(2)}, \dots, \|z\| \cdot g^{(k)} \right)$$

↑      ↑      ↗  
iid  $\in N(0,1)$ .

$$= \|z\| \cdot \frac{1}{\sqrt{k}} \cdot \underbrace{(g^{(1)}, g^{(2)}, \dots, g^{(k)})}_{k\text{-dim } N(0,1)}$$

$$\|\varphi(z)\|_2^2 = \|z\|_2^2 \cdot \frac{1}{k} \cdot \sum_{i=1}^k (g^{(i)})^2$$

$\chi_k^2$       chi-squared  
dist w/ k  
degrees of fr.

Fact:  $\Pr [ \chi_k^2 \notin (1 \pm \epsilon)k ] \leq e^{-\epsilon^2 k / 8}$  for  $\epsilon < 1/2$ .

$$\Rightarrow \Pr_{\varphi} [\|\varphi(z)\|_2^2 = (1 \pm \epsilon) \|z\|^2] \geq 1 - e^{-\epsilon^2 k/g}. \quad \square \quad \square$$

$$\sqrt{1-\epsilon} \approx 1 - \epsilon/2.$$

$$\alpha \in (1 \pm \epsilon)^{1/2} \cdot \alpha' \Rightarrow \alpha \in (1 \pm \epsilon) \alpha'.$$

Corollary [of JL]: fix  $N$  vectors

$x_1, \dots, x_N \in \mathbb{R}^n$ . Pick  $\varphi$  as in JL,  
with  $k = \Theta\left(\frac{\lg N}{\epsilon^2}\right)$ . Then

$$\Pr_{\varphi} [\forall i, j \in [N], \|\varphi(x_i) - \varphi(x_j)\| = (1 \pm \epsilon) \|x_i - x_j\|] \geq 1 - \frac{1}{N}.$$

proof: fix  $k = \frac{3 \cdot g \cdot \ln N}{\epsilon^2}$ .

Then, by JL,  $\forall i, j \in [N]$ :

$$\Pr [\|\varphi(x_i) - \varphi(x_j)\| \notin (1 \pm \epsilon) \|x_i - x_j\|] \leq e^{-\epsilon^2 k/g} = \frac{1}{N^3}.$$



By union bound over all pairs  $i, j \in [N]$ :

$$\Rightarrow \Pr[\exists i, j \in [N] \text{ where } \dots \neq \dots] \\ \leq N^2 \cdot \frac{1}{N^3} = \frac{1}{N}. \quad \square$$

Rephases for  $N$  vectors in  $\mathbb{R}^n$ ,  
can map them into  $\mathbb{R}^k$ ,  
 $k = O\left(\frac{\log N}{\epsilon^2}\right)$ ,  
while preserving dist  $1 \pm \epsilon$ .

Remark: is  $k$  tight? yes,  
even for non-linear  
maps  $\psi$ .  
(even if  $\psi$  depends on  
 $x_1, \dots, x_N$ ).

Remark: what about dim-reduction in  
other spaces?

Can we do the same in  $\ell_1$ ?

$$f: \ell_1^n \rightarrow \ell_1^k \text{ s.t.}$$

distances preserved with "good" prob.

or map  $N$  vectors in  $\ell_1$  into lower-dim.  $\ell_1$ ?

Basically no: min dim is

[Brinkman  
Charikar ~2003]  $k = N^{2(1/d)}$  for

approx  $d \geq 1$ .

[Naor, Johnson 2009] good dim-red.  $\Rightarrow$  must be  $\ell_2$ .