

Clamping Variables and Approximate Inference

Adrian Weller



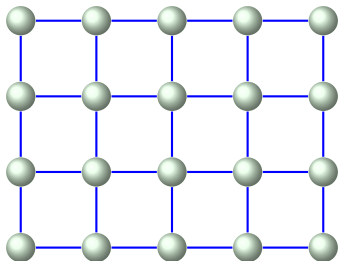
Slides and full paper at
www.cs.columbia.edu/~adrian

Work with Tony Jebara, Columbia University

UCL CSML Seminar Part 1, March 27 2015

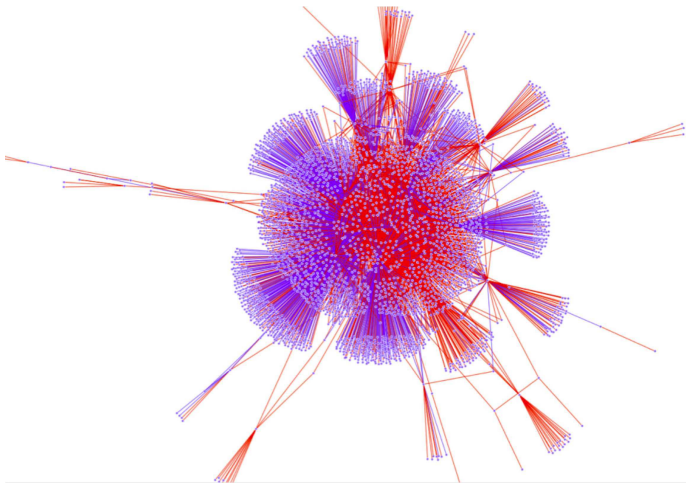
Motivation: *undirected graphical models*

- Powerful way to represent relationships across variables
- Many applications including: computer vision, social network analysis, deep belief networks, protein folding...
- In this talk, focus on binary pairwise (Ising) models



Example: Grid for computer vision (attractive)

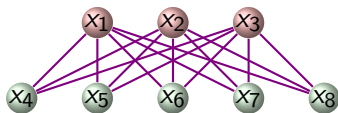
Motivation: *undirected graphical models*



Example: Part of opinions social network (general)

Figure courtesy of N. Ruozi

Motivation: *undirected graphical models*



Example: Restricted Boltzmann machine (general)

A fundamental problem is *marginal inference*

- Estimate marginal probability distribution of one variable

$$p(x_1) = \sum_{x_2, \dots, x_n} p(x_1, x_2, \dots, x_n)$$

- Closely related to computing the *partition function*
- Computationally intractable, focus on approximate methods
- Will show that combining approximate methods with *clamping* can be very fruitful for marginal inference

Outline: *Clamping can be very helpful*

1. Motivation
2. Background on inference and clamping



Combining clamping variables with variational inference, we obtain

3. Strong theoretical results
4. Promising empirical results

Background: *Binary pairwise models*

- Binary variables $X_1, \dots, X_n \in \{0, 1\}$
- Pairwise potentials θ
- Write x for one complete configuration of all variables,
 $\theta \cdot x$ for its score
- Probability distribution given by

$$p(x) = \frac{1}{Z} \exp(\theta \cdot x)$$

- To ensure probabilities sum to 1, need normalizing constant

$$Z = \sum_x \exp(\theta \cdot x)$$

- Z is called the *partition function*, a fundamental quantity we'd like to compute or approximate



Background: *A variational approximation*

$$\text{Recall } p(x) = \frac{1}{Z} \exp(\theta \cdot x)$$

- Exact inference may be viewed as *optimization*,

$$\log Z = \max_{\mu \in \mathbb{M}} [\theta \cdot \mu + S(\mu)]$$

\mathbb{M} is the space of marginals that are *globally consistent*, S is the (Shannon) entropy

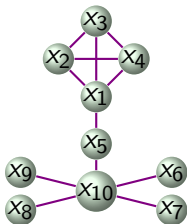
- Bethe makes two pairwise approximations,

$$\log Z_B = \max_{q \in \mathbb{L}} [\theta \cdot q + S_B(q)]$$

\mathbb{L} is the space of marginals that are *pairwise consistent*, S_B is the Bethe entropy approximation

- Loopy Belief Propagation finds stationary points of Bethe
- On acyclic models, Bethe is exact $Z_B = Z$

Background: *What is clamping?*

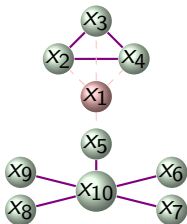


Example 'lamp' graph

To compute the partition function Z , can enumerate all states and sum

$x_1 x_2 \dots x_{10}$	score	$\exp(\text{score})$
0 0 ... 0	1	2.7
0 0 ... 1	2	7.4
...
0 1 ... 1	1.3	3.7
1 0 ... 0	-1	0.4
1 0 ... 1	0.2	1.2
...
1 1 ... 1	1.8	6.0
Total $Z =$		47.1

Background: *What is clamping?*



Can split Z in two: **clamp** variable X_1 to each of $\{0, 1\}$, then add the two sub-partition functions:

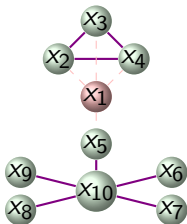
$$Z = Z|_{X_1=0} + Z|_{X_1=1}$$

When clamp a variable, remove it from the graph

$x_1 x_2 \dots x_{10}$	score	$\exp(\text{score})$	
0 0 ... 0	1	2.7	
0 0 ... 1	2	7.4	
...	
0 1 ... 1	1.3	3.7	27.5
1 0 ... 0	-1	0.4	
1 0 ... 1	0.2	1.2	
...	
1 1 ... 1	1.8	6.0	19.6
Total $Z =$		47.1	

$$p(X_1 = 1) = \frac{Z|_{X_1=1}}{Z}$$

Background: *What is clamping?*



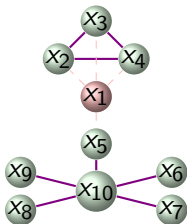
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When clamp a variable, remove it from the graph

- After removing the clamped variable, if the remaining sub-models are **acyclic** then can find sub-partition functions efficiently (BP, Bethe approximation is exact on trees)

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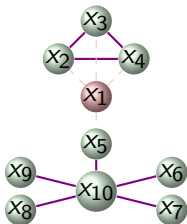
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When clamp a variable, remove it from the graph

- After removing the clamped variable, if the remaining sub-models are **acyclic** then can find sub-partition functions efficiently (BP, Bethe approximation is exact on trees)
- If not,
 - Can repeat: clamp and remove variables until acyclic, *or*
 - Settle for **approximate inference** on sub-models

$$Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1}$$

Background: *What is clamping?*



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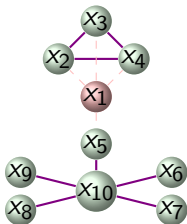
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Will this always lead to a better estimate than approximate inference on the original model?

Background: *What is clamping?*



Can split Z in two: **clamp** variable X_1 to each of $\{0, 1\}$, then add the two sub-partition functions:

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$$Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1}$$

Will this always lead to a better estimate than approximate inference on the original model? *Often but not always*

A variational perspective on clamping

- Bethe approximation

$$\log Z_B = \max_{q \in \mathbb{L}} [\theta \cdot q + S_B(q)]$$

- Observe that when X_i is clamped, we optimize over a subset

$$\log Z_{B|X_i=0} = \max_{q \in \mathbb{L}: q_i=0} [\theta \cdot q + S_B(q)]$$

$$\Rightarrow Z_{B|X_i=0} \leq Z_B, \text{ similarly } Z_{B|X_i=1} \leq Z_B$$

Recap of Notation

Z

true partition function

Z_B

Bethe optimum partition function

$$Z_B^{(i)} := Z_{B|X_i=0} + Z_{B|X_i=1}$$

approximation obtained when
clamp and sum approximate
sub-partition functions

Clamping variables: *an upper bound on Z*

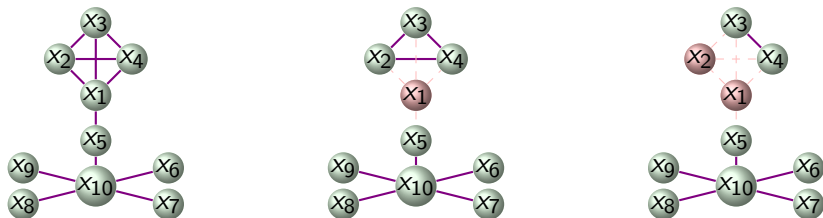
- From before,

$$Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1} \leq 2Z_B$$

- Repeat: **clamp and remove variables**, until remaining model is **acyclic**, where **Bethe is exact**
- For example, if must delete 2 variables X_i, X_j , obtain

$$Z_B^{(ij)} := \sum_{a,b \in \{0,1\}} Z_B|_{X_i=a, X_j=b} \leq 2^2 Z_B$$

But sub-partition functions are *exact*, hence $\text{LHS} = Z$



Clamping variables: *an upper bound on Z*



$$Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1} \leq 2Z_B$$

- Repeat: clamp and remove variables, until remaining model is acyclic, where Bethe is exact
- Let $\nu(G)$ be the minimum size of a **feedback vertex set**

Theorem (result is tight)

$$Z \leq 2^\nu Z_B$$

Clamping variables: *an upper bound on Z*

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Attractive models: *a lower bound on Z*

- An *attractive* model is one with all edges attractive
- Recall definition,

$$Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1}$$

Theorem

For an attractive binary pairwise model and any X_i , $Z_B \leq Z_B^{(i)}$

Repeat as before: $Z_B \leq Z_B^{(i)} \leq Z_B^{(ij)} \leq \dots \leq Z$

Corollary (similar proof to earlier result; first proved Ruozi, 2012)

For an attractive binary pairwise model, $Z_B \leq Z$

Attractive models: *a lower bound on Z*

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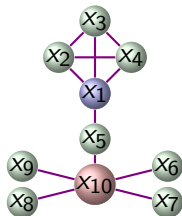
\Rightarrow each clamp and sum can only *improve* Z_B

Experiments: *Which variable to clamp?*

Compare error $|\log Z - \log Z_B^{(i)}|$ to original error $|\log Z - \log Z_B|$ for various ways to choose which variable X_i to clamp:

- **best Clamp** best improvement in error of Z in hindsight
- **worst Clamp** worst improvement in error of Z in hindsight
- **avg Clamp** average performance

- **maxW** max sum of incident edge weights $\sum_{j \in N(i)} |W_{ij}|$
- **Mpower** more sophisticated, based on powers of related matrix



Experiments: *attractive random graph* $n = 10, \rho = 0.5$

unary $\theta_i \sim U[-2, 2]$,
edge $W_{ij} \sim U[0, W_{max}]$

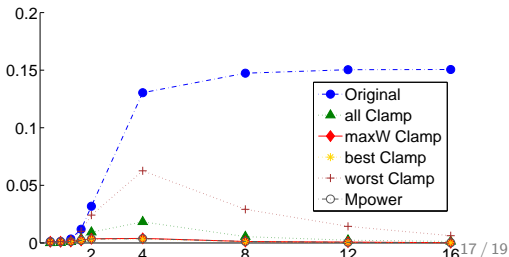
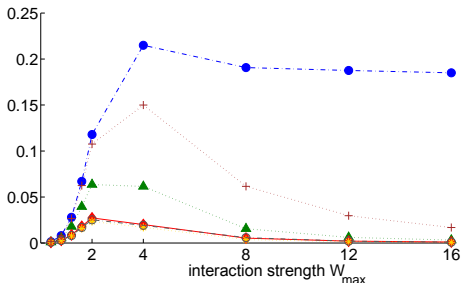
Error of estimate of $\log Z$

Observe

- Clamping any variable helps significantly
- Our selection methods perform well

Avg ℓ_1 error of singleton
marginals

Using Frank-Wolfe to optimize
Bethe free energy



Experiments: *general random graph* $n = 10, p = 0.5$

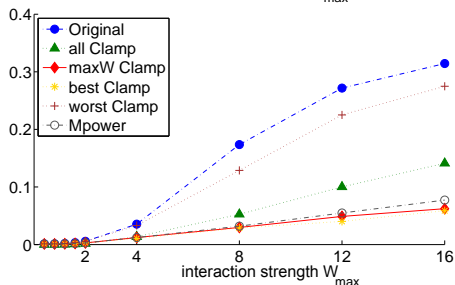
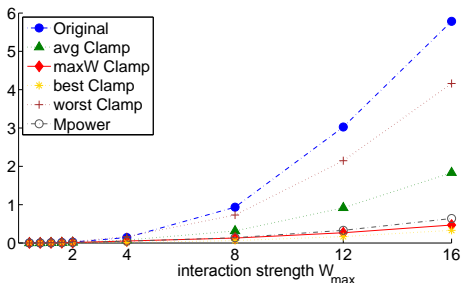
unary $\theta_i \sim U[-2, 2]$,
edge $W_{ij} \sim U[-W_{max}, W_{max}]$

Error of estimate of $\log Z$

Results remain promising
for higher n

Avg ℓ_1 error of singleton
marginals

Using Frank-Wolfe to optimize
Bethe free energy



Recap of theoretical results

- Simple observation on variational view of clamping variables gives $Z_B^{(i)} \leq 2Z_B$
- Repeat until graph is acyclic, where Bethe is exact
- Yields effective upper bound on Z

For *attractive* models,

- Theorem: $Z_B \leq Z_B^{(i)}$ for any X_i
- Then argue as above to yield simple new proof of $Z_B \leq Z$
- Clamping any variable and summing can only improve Z_B
- To prove Theorem above, derive stronger result on convexity of function combining conditioned Bethe optimum with singleton entropy, ask if interested

Thank you!

Slides and full paper at www.cs.columbia.edu/~adrian

Extra slides for questions or further explanation

Clamping variables: *strongest result for attractive models*

$$\log Z_B = \max_{q \in \mathbb{L}} [\theta \cdot q + S_B(q)]$$

- For any variable X_i and $x \in [0, 1]$, let $q_i = q(X_i = 1)$ and

$$\log Z_{B_i}(x) = \max_{q \in \mathbb{L}: q_i = x} [\theta \cdot q + S_B(q)]$$

- $Z_{B_i}(x)$ is '*Bethe partition function constrained to $q_i = x$* '

Note: $Z_{B_i}(0) = Z_B|_{X_i=0}$, $Z_{B_i}(x^*) = Z_B$, $Z_{B_i}(1) = Z_B|_{X_i=1}$

Clamping variables: *strongest result for attractive models*

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Note: $Z_{B_i}(0) = Z_B|_{X_i=0}$, $Z_{B_i}(x^*) = Z_B$, $Z_{B_i}(1) = Z_B|_{X_i=1}$

- Define new function,

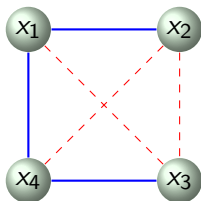
$$A_i(q_i) := \log Z_{B_i}(q_i) - S_i(q_i)$$

Theorem (implies all other results for attractive models)

For an attractive binary pairwise model, $A_i(q_i)$ is convex

- Builds on derivatives of Bethe free energy from [WJ13]

Example: here clamping any variable *worsens* Z_B estimate



Blue edges are attractive with edge weight $+2$

Red edges are repulsive with edge weight -2

No unary potentials

(performance is only slightly worse with clamping)

Experiments: *attractive complete graph* $n = 10$, TRW

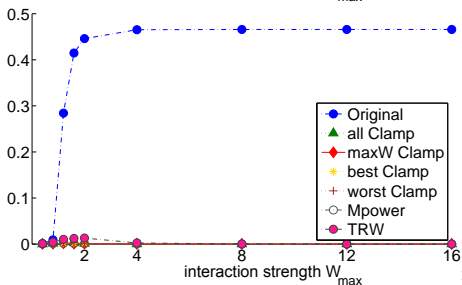
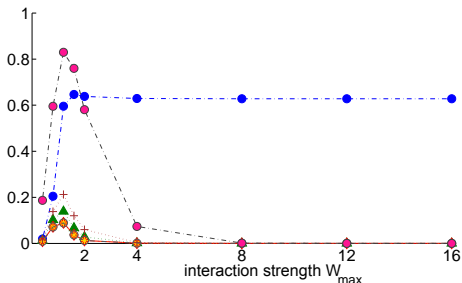
unary $\theta_i \sim U[-0.1, 0.1]$,
edge $W_{ij} \sim U[-W_{max}, W_{max}]$

Error of estimate of $\log Z$

Note low unary potentials

Avg ℓ_1 error of singleton
marginals

Clamping a variable 'breaks
symmetry' and overcomes
TRW advantage



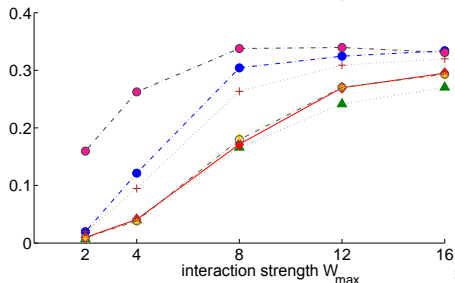
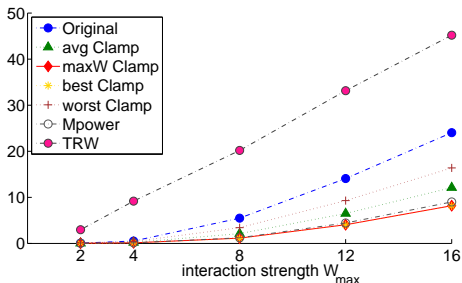
Experiments: *general complete graph* $n = 10$, *TRW*

unary $\theta_i \sim U[-2, 2]$,
edge $W_{ij} \sim U[0, W_{max}]$

Error of estimate of $\log Z$

Note regular singleton
potentials

Avg ℓ_1 error of singleton
marginals



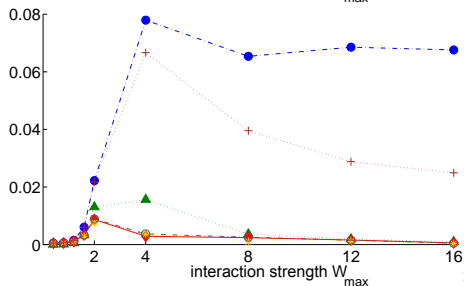
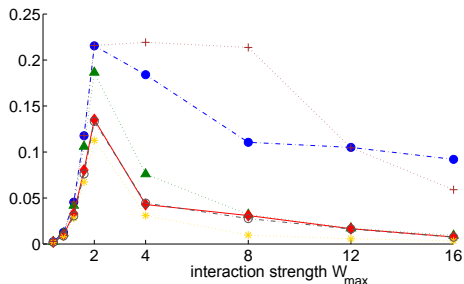
Experiments: *attractive random graph* $n = 50, \rho = 0.1$

unary $\theta_i \sim U[-2, 2]$,
edge $W_{ij} \sim U[0, W_{max}]$

Error of estimate of $\log Z$

'worst Clamp' performs worse
here due to suboptimal
solutions found by Frank-Wolfe

Avg ℓ_1 error of singleton
marginals



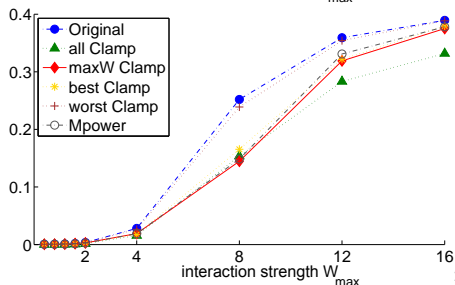
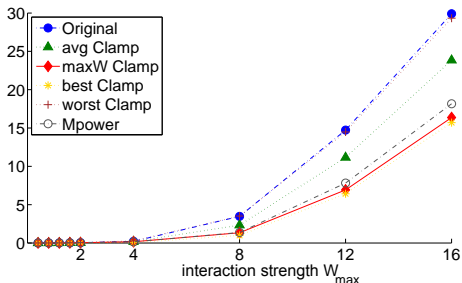
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Error of estimate of $\log Z$

Performance still good for
clamping just one variable

Avg ℓ_1 error of singleton
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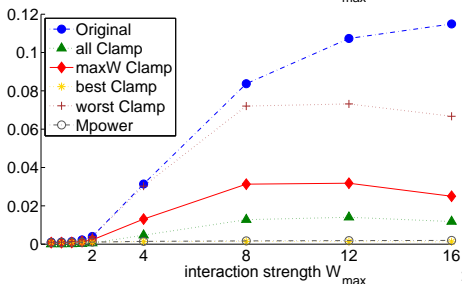
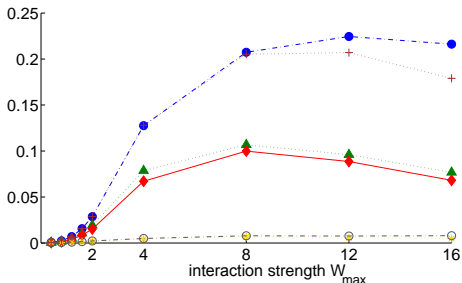
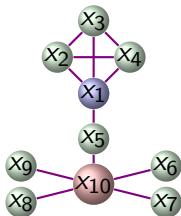
Experiments: *attractive 'lamp' graph*

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Error of estimate of $\log Z$

Mpower performs well,
significantly better than maxW

Avg ℓ_1 error of singleton
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Experiments: *general 'lamp' graph*

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