

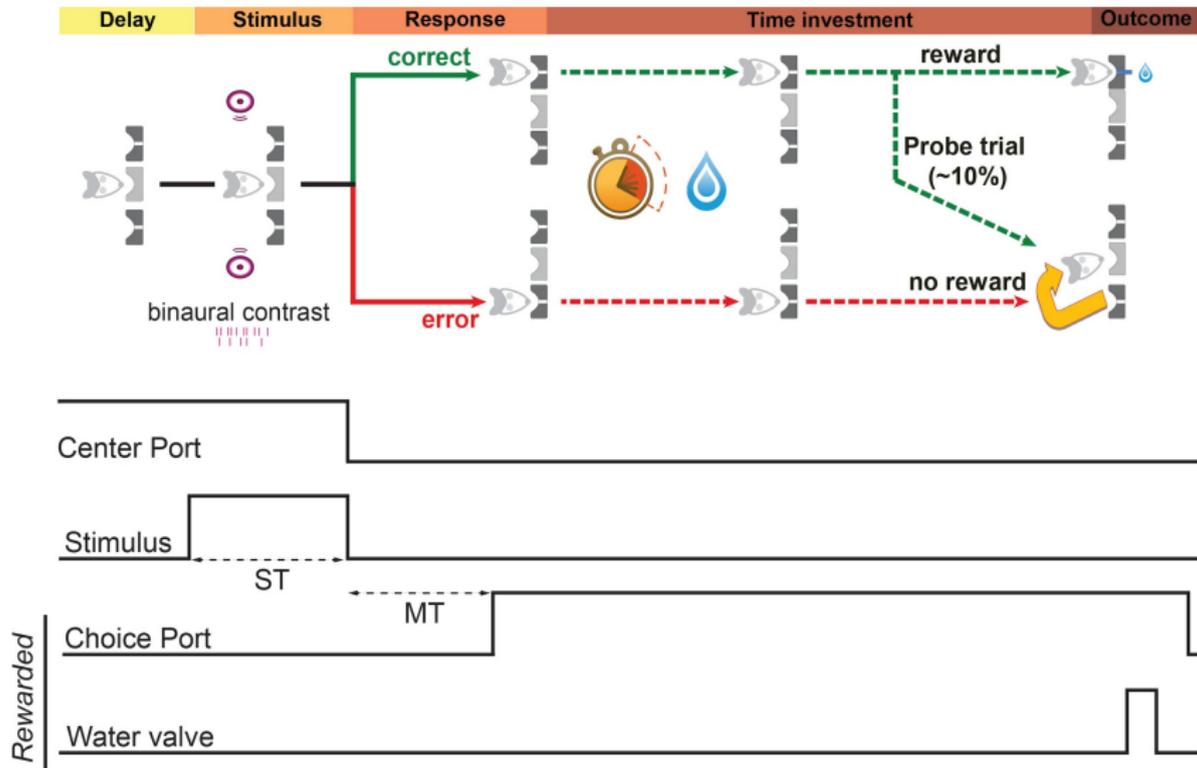
# BEADL: A New Real-Time Language for Behavioral Experiments

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Adam Kepecs, Cold Spring Harbor Laboratory

[Lak et al., Neuron 84(1), 2014]

# Bpod: An Open Hardware Platform for Behavioral Monitoring and Control

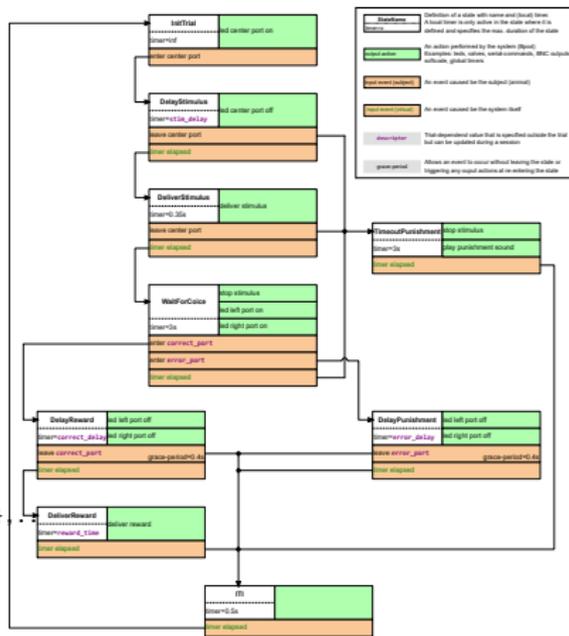


Sanworks.io, spun out of Kepecs' lab.  
Teensy 3.6: ARM Cortex M4, 180 MHz

```

sma = NewStateMatrix();
sma = AddState(sma,'Name', 'ITI',...
'Timer',S.ITI,...
'StateChangeConditions', {'Tup', 'PreState'},...
'OutputActions',{ });
%Pre task states
sma = AddState(sma, 'Name', 'PreState',...
'Timer', S.GUI.PreCue,...
'StateChangeConditions', {'Tup', 'CueDelivery'},...
'OutputActions', {'BNCState', 1});
%Cue
sma=AddState(sma,'Name', 'CueDelivery',...
'Timer',S.GUI.CueDuration,...
'StateChangeConditions',{'Tup', 'Delay'},...
'OutputActions', {'SoftCode',S.Cue});
%Delay
sma=AddState(sma,'Name', 'Delay',...
'Timer', S.Delay,...
'StateChangeConditions',{'Tup', 'ExtraCueDelivery'},...
'OutputActions',{ });
%Extra Cue for L3-SecondaryCue
sma=AddState(sma,'Name', 'ExtraCueDelivery',...
'Timer',S.ExtraCueDuration,...
'StateChangeConditions', {'Tup', 'ExtraDelay'},...
'OutputActions', {'SoftCode',S.ExtraCue});
%Extra Delay for L3-SecondaryCue
sma=AddState(sma,'Name', 'ExtraDelay',...
'Timer',S.ExtraDelay,...
'StateChangeConditions',{'Tup', 'Outcome'},...
'OutputActions',{ });
%Reward
sma=AddState(sma,'Name', 'Outcome',...
'Timer',S.Outcome,...
'StateChangeConditions', {'Tup', 'PostOutcome'},...
'OutputActions', {'ValveState', S.Valve});
%Post task states
sma=AddState(sma,'Name', 'PostOutcome',...
Timer',S.GUI.PostCue);

```



Describe as FSM

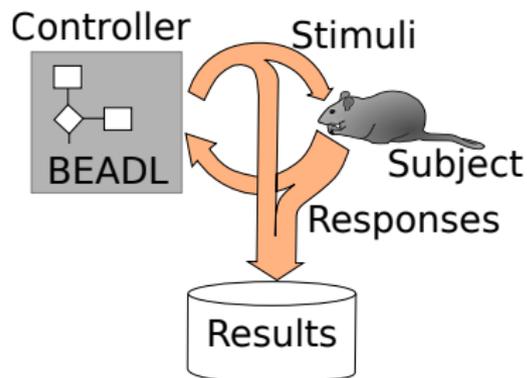
Build FSM with Matlab API

Download FSM to firmware

Firmware: FSM interpreter

w/100  $\mu$ s heartbeat

# BEADL: The Idea



## outputs

*valve dispense* # Channel w/ 1 event  
*led on off* # Two possible events

## inputs

*gate enter exit*

## task simple\_example:

```
"Subject Attraction" # State label
valve dispense # Generate event
await
  10 s: # Timeout
    "Failure"
  goto "Subject Attraction"
gate enter: # Event arrived

"Light Stimulus"
led on # Generate event
await 100 ms # Simple delay
led off
```

# BEADL: Philosophy

Deterministic formal semantics

Explicit model-time delays only; platform-independent timing above some minimum delay (synchronous logic)

“Bare metal” microcontroller implementations: hardware counter/timer drives timing, timer interrupts for scheduling

Schedulability/static timing analysis done at compile time

# BEADL: Possible Single-Threaded Implementation

## outputs

*valve dispense*  
*led on off*

## inputs

*gate enter exit*

"Attract"

*valve dispense*

await

10 s:

"Fail"

goto "Attract"

*gate enter:*

"Stimulus"

*led on*

await 100 ms

*led off*

```
void interrupt1() {
    now = get_platform_time();
    switch (state) {
    case S1:
        Attract: report("Attract");
        valve_dispense();
        state = S2, schedule(now + SEC(10)); return;
    case S2:
        switch (get_interrupting_event()) {
        case TIMEOUT:
            Fail: report("Fail");
            goto Attract;
        case GATE_ENTER: break;
        default: return; }
        Stimulus: report("Stimulus");
        led_on();
        state = S3, schedule(now + MS(100)); return;
    case S3:
        led_off();
        state = STOPPED; return;
    case STOPPED:
        return;
    }
}
```

# BEADL: Parallel Composition

*Language Design is Library Design* —Stroustrup

A desired BEADL library: input debounce

Nervous rats often jitter before making a decision; want a library that discards “on” events shorter than  $x$  ms

⇒ Parallel composition?

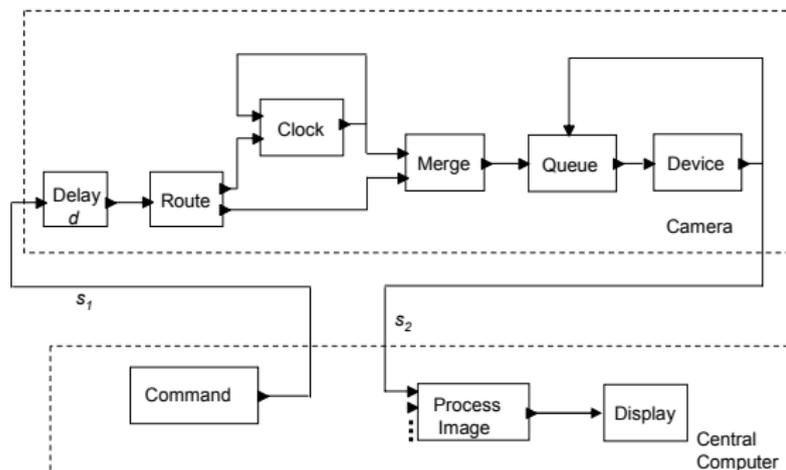


Feedback loops?

Simultaneous events?

Contradictions?

# PTIDES/PtidyOS



[Zhao et al., RTAS 2007]

“Real-time  
discrete-event  
simulation”

Parallel actors  
with multiple  
inputs, outputs,  
channels

Event tags:  $\langle \text{model time, microstep, topological level} \rangle$

Feedback loops must have delay actors (microstep or longer)

Scheduler w/ event queue; timer interrupt delays keep model time “close” to platform time.

[Edward Lee et al.]

# The Lustre Synchronous Dataflow Language

```
node Watchdog (set,reset,u_tps: bool; delay: int) returns (alarm: bool);  
var is_set : bool;  
    remain : int;  
let  
    alarm = is_set and (remain = 0) and pre(remain) > 0;  
    is_set = false -> if set then true  
                        else if reset then false  
                        else pre(is_set);  
    remain = 0 -> if set then delay  
                  else if u_tps and pre(remain) > 0  
                  then pre(remain) - 1  
tel
```

Declarative dataflow style; expressing control is awkward

Every loop must have a unit delay (“pre”)

Implicit clock not tied to wallclock time

[Halbwachs, Caspi, et al.]

# The Esterel Synchronous Programming Language

```
module ABRO:  
input A, B, R;  
output O;  
  
loop  
  [ await A || await B ];  
  emit O  
each R  
  
end module
```

Imperative style with sequencing, concurrency, conditionals, and exceptions

More subtle causality constraints; “constructive” causality requires a per-state analysis

[Berry et al.]

# BEADL: A Work in Progress

- ▶ Semantics
  - Event-driven with explicit model time advances
  - Synchronous: Esterel-like with Lustre-style causality?
  - No reactions to absent events (timeouts only)
- ▶ Implementation
  - PTIDES-like: Interrupt driven
  - Events: timeouts, input arrivals
  - Model time “matched” to platform time
- ▶ Schedulability
  - Run-time deadline checking
  - Compile-time WCET analysis?