# An Efficient Algorithm for the Analysis of Cyclic Circuits

Osama Neiroukh Stephen A. Edwards Xiaoyu Song Intel Columbia University Portland State University

An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 1/

#### What is a Combinational Circuit?

Malik [ICCAD 93]: "A circuit is combinational *for an input pattern* if three-valued simulation starting from Xs converges to 0s and 1s."

Shiple [96]: "Equivalent to stability in Brzozowski and Seger's [95] model."





# Goal

Given a cyclic circuit that is combinational for some inputs, create an acyclic circuit that computes the same combinational function.



# Applications

Fixing cyclic circuits from high-level synthesis Stok [ICCAD 92]: cycles from resource sharing Berry [92]: cycles from Esterel programs Acyclic circuits easier to simulate

#### **Related Work**

Malik [ICCAD 93]: basic definitions, unrolling Edwards [DAC 03]: basis of our work Gupta and Selvidge [ICCAD 05]: fix single loops Riedel [DAC 03]: a technique for creating them

An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 6/

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 6/



An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 6/



An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 6/1

# First Observation [Edwards 2003]



For an input pattern to be combinational, at least one input coming from outside each strongly-connected component must have a controlling value.

If all external inputs were non-controlling, the gates in the SCC would stay at X.



An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 8/1



An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 8/1



Frontier gate: some inputs defined, output remains X

An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 8/1



*Frontier gate*: some inputs defined, output remains X

Input is combinational  $\Leftrightarrow$  frontier is empty

An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 8/

# Our Algorithm Step 1: Apply a Controlling Value to Each Input



# Our Algorithm Step 1: Apply a Controlling Value to Each Input



a = 0 acyclic

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 9/



a = 0 acyclic

b = 0

An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 9/

# Our Algorithm Step 1: Apply a Controlling Value to Each Input



a = 0 acyclic

An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 9/

# Our Algorithm Step 1: Apply a Controlling Value to Each Input



a = 0 acyclic



An Efficient Algorithm for the Analysis of Cyclic Circuits – p. 9/

# Our Algorithm Step 1: Apply a Controlling Value to Each Input



a = 0 acyclic d = 1 0

b = 0

 $\mathbf{c} = \mathbf{0}$ 

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 9

# Our Algorithm Step 1: Apply a Controlling Value to Each Input



 $d = 1_0 - V$ a = 0 acyclic e = 0b = 0Ζ f = 1 $\mathbf{C} = \mathbf{0}$ Ζ

 $g = 0 \begin{array}{c} 1 \end{array} = \begin{bmatrix} z \\ z \end{bmatrix}$  $g = 1 \begin{array}{c} 1 \end{array} = \begin{bmatrix} z \\ z \end{bmatrix}$ 

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 9/

#### Our Algorithm Step 2: Attack Frontier Gates with Combinations



 $b = 0 \quad 0 \quad v \quad + \quad c = 0 \quad 0 \quad v \quad + \quad c = 0 \quad 0 \quad v \quad + \quad c = 0 \quad 0 \quad v \quad - \quad v \quad$ 

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 10/

#### Our Algorithm Step 2: Attack Frontier Gates with Combinations



- + c = 0 0 - + d = 1 0 b = 0b = 0

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 10/

#### Our Algorithm Step 2: Attack Frontier Gates with Combinations



 $b = 0 \quad 0 \quad V \quad + \quad c = 0 \quad 0 \quad V \quad + \quad d = 1 \quad 0 \quad V \quad + \quad d = 1 \quad 0 \quad V \quad + \quad d = 1 \quad 0 \quad V \quad + \quad g = 0 \quad 1 \quad z \quad + \quad f = 1 \quad 1 \quad z \quad + \quad g = 0 \quad 1 \quad z \quad + \quad z \quad = \quad z \quad$ 

An Efficient Algorithm for the Analysis of Cyclic Circuits - p. 10/2

# **Experimental Results**

| Netlist | SCC  | DAC 03   |  | Ours   |   | Acyclic   |
|---------|--|--|--|--|---|---|
| Gates   | Gates  | PAs  | time   | PAs  | time  | PAs   |
| 213     | 25   | 257  | 1.3  | 25   | 0.1   | 14  |
| 248     | 30   | 745  | 8  | 29   | 0.1   | 16  |
| 283     | 35   | 2205   | 69   | 33   | 0.2   | 18  |
| 318     | 40   | 6581   | 656  | 37   | 0.3   | 20  |
| 124     | 69   | 54517  | 2868   | 23260  | 2.0   | 338   |
| 150     | 47   | 43777  | 2341   | 232  | 1.0   | 10  |
| 177     | 32   |  | $\infty$   | 290  | 0.6   | 11  |
| 253     | 51   |  | $\infty$   | 1489   | 0.3   | 22  |
| 272     | 61   |  | $\infty$   | 588  | 0.2   | 89  |
| 311     | 49   |  | $\infty$   | 3604   | 1.0   | 38  |
|         | Netlist<br>Gates<br>213<br>248<br>283<br>318<br>124<br>150<br>177<br>253<br>272<br>311 | NetlistSCCGatesGates21325248302833531840124691504717732253512726131149 | NetlistSCCDACGatesGatesPAs213252572483074528335220531840658112469545171504743777177321272611311491 | NetlistSCCDAC $\cup$ 3GatesGatesPAstime213252571.3248307458283352205693184065816561246954517286815047743777234117732 $\infty$ 25351 $\infty$ 27261 $\infty$ 31149 $\infty$ | NetlistSCCDAC $\cup$ 3OuGatesGatesPAstimePAs213252571.3252483074582928335220569333184065816563712469545172868232601504743777234123217732 $\infty$ 29025351 $\infty$ 148927261 $\infty$ 58831149 $\infty$ 3604 | NetlistSCC $DAC \cdot 03$ $Ou \cdot s$ GatesGatesPAstimePAstime213252571.3250.1248307458290.128335220569330.2318406581656370.312469545172868232602.0150474377723412321.017732 $\infty$ $\infty$ 14890.327261 $\infty$ 5880.231149 $\infty$ $3604$ 1.0 |

### Conclusions

- More focused exploration of search space
- Idea: combine partial assignments to attack frontier gates
- Exponential improvement compared to Edwards [DAC 03]
- Future work
  - Even better pruning
  - Symbolic approach?