Adding a Flow-Oriented Paradigm to Commodity Operating Systems

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The Status Quo: Memory as Buffer



I/O Becoming Faster than CPUs



Our Idea: Flows Controlled by OS







Abstract Server Model



Signaling API

Need some sort of mechanism for creating, controlling, and tearing down flows, e.g.,

```
flow = flow_open();
flow_source(flow, "/usr/http/secretfile.html");
flow_xformer(flow, "/dev/crypto");
flow_xformer(flow, "/dev/http");
flow_sink(flow, "/dev/inet/192.168.1.3");
flow_start(flow);
flow_stop(flow);
```

Exception Handling

What happens if something goes wrong?

- OS may try to re-start device
- Redirect to a different device
- Error passed to application
- Switch to all-software flow
- Terminate

Resource Scheduler

What if two processes want access to /dev/crypto?

- Performance requirements
- Available resources
- Priorities

Flow-level scheduling costly, but infrequent

Detailed (e.g., bus access) scheduling more frequent

Programmable Peripherals



RadiSys ENP-2611 Network Card w/ IXP2400

Altera Stratix PCI

FPGAs: flow components or

wrappers around legacy peripherals (DMA absorbers)

Proof-of-Concept System



Experimental Results

Time to send 1 Million packets through the pipeline

Packet size	Memory- centric	Flow- centric*	Speed-up
64 bytes	15.5s	13.6	13%
1024 bytes	114	59.2	49%

*Plus main processor mostly idle for these packets

Conclusions

New flow-centric architecture for operating systems Have the OS manage inter-peripheral flows under application control Requires programmable peripherals: many already extant Proof-of-concept showed nearly a 2× speedup Minimal performance impact from additional computations (e.g., security)