# Fundamentals of Computer Systems <br> Finite State Machines 

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## Finite State Machine Components



## Moore and Mealy Machines



## The Moore Form:

Outputs are a function of only the current state.

## Moore and Mealy Machines



The Mealy Form:
Outputs may be a function of both the current state and the inputs.

A mnemonic: Moore machines often have more states.

## Mealy Machines are the Most General



Another, equivalent way of drawing Mealy Machines
This is exactly the synchronous digital logic paradigm

## Moore vs. Mealy FSMs

Alyssa P. Hacker has a snail that crawls down a paper tape with 1's and 0's on it. The snail smiles whenever the last four digits it has crawled over are 1101.
Design Moore and Mealy FSMs of the snail's brain.


## State Transition Diagrams: Looking for "1101"



Moore Machine: States indicate output

## State Transition Diagrams: Looking for "1101"



Moore Machine: States indicate output

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Moore Machine: States indicate output


Mealy Machine: Arcs indicate input/output

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## State Transition Diagrams: Looking for "1101"



Moore Machine: States indicate output


Mealy Machine: Arcs indicate input/output

## Moore Machine

| Next State |  |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
| S | A | S' | S | Y |
| S0 | 0 | S0 | S0 | 0 |
| S0 | 1 | S1 | S1 | 0 |
| S1 | 0 | S0 | S2 | 0 |
| S1 | 1 | S2 | S3 | 0 |
| S2 | 0 | S3 | S4 | 1 |
| S2 | 1 | S2 |  |  |
| S3 | 0 | S0 |  |  |
| S3 | 1 | S4 |  |  |
| S4 | 0 | S0 |  |  |
| S4 | 1 | S2 |  |  |

Moore Machine

| Next State |  |  | Output |  |
| :---: | :---: | :---: | :---: | :---: |
| S | A | S' | S | Y |
| 000 | 0 | 000 | 000 | 0 |
| 000 | 1 | 001 | 001 | 0 |
| 001 | 0 | 000 | 010 | 0 |
| 001 | 1 | 010 | 011 | 0 |
| 010 | 0 | 011 | 100 | 1 |
| 010 | 1 | 010 |  |  |
| 011 | 0 | 000 |  |  |
| 011 | 1 | 100 |  |  |
| 100 | 0 | 000 |  |  |
| 100 | 1 | 010 |  |  |



## Mealy Machine

| $\mathbf{S}$ | $\mathbf{A}$ | $\mathbf{S}$ | $\mathbf{Y}$ |
| :---: | :---: | :---: | :---: |
| S0 | 0 | S0 | 0 |
| S0 | 1 | S1 | 0 |
| S1 | 0 | S0 | 0 |
| S1 | 1 | S2 | 0 |
| S2 | 0 | S3 | 0 |
| S2 | 1 | S2 | 0 |
| S3 | 0 | S0 | 0 |
| S3 | 1 | S1 | 1 |

Mealy Machine

| $\mathbf{S}$ | $\mathbf{A}$ | $\mathbf{S}$ | $\mathbf{Y}$ |
| :---: | :---: | :---: | :---: |
| 00 | 0 | 00 | 0 |
| 00 | 1 | 00 | 0 |
| 01 | 0 | 00 | 0 |
| 01 | 1 | 10 | 0 |
| 10 | 0 | 11 | 0 |
| 10 | 1 | 10 | 0 |
| 11 | 0 | 00 | 0 |
| 11 | 1 | 01 | 1 |



## More Intuitive Solutions using Shift Registers



Mealy Form: Output Depends on Input Immediately


Moore Form: Output Depends Only on State

## FSM Example: A Traffic Light Controller



Source: Mead and Conway, Introduction to VLSI Systems, 1980, p. 85.

## State Transition Diagram for the TLC



Inputs:
C: Car sensor
S: Short Timeout
L: Long Timeout

Outputs:
T: Timer Reset
H: Highway color
F: Farm road color

## State Transition Diagram for the TLC

$$
\bar{C}+\bar{L} / \bar{T}
$$



Inputs:
C: Car sensor
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## State Transition Diagram for the TLC



Inputs:
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F: Farm road color

State Transition Diagram for the TLC


## State and Output Encoding

## $\mathbf{S} \mathbf{C} \mathbf{S} \mathbf{L} \quad \mathbf{S}$

HG $0 \times X \quad X \quad 0 \quad$ HG HG $\quad X \quad X \quad 0 \quad 0 \quad H G$ HG $1 \begin{array}{llllll} & X & 1 & 1 & H Y\end{array}$ MY X $0 \quad$ X $0 \quad$ MY WY $\quad$ X $1 \quad 1 \quad X \quad 1 \quad$ PG
FF $1 \begin{array}{lllll} & X & 0 & 0 & F G\end{array}$ PG $\quad 0 \quad X \quad X \quad X \quad 1 \quad F Y$ EG $\quad X \quad X \quad 1 \quad 1 \quad$ FY FY $\quad \mathrm{X} \quad 0 \quad \mathrm{X} \quad 0 \quad \mathrm{FY}$
FY $\quad$ X $\quad 1 \quad \mathrm{X} \quad 1 \mathrm{HG}$

## $\mathbf{S} \mathbf{H}$ F

HG GR
MY $Y$ R
EG $R \quad G$
FY R Y

State and Output Encoding

## $\mathbf{S} \mathbf{C} \mathbf{S} \mathbf{L}$ S'

| 00 | 0 | X | X | 0 | 00 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | X | X | 0 | 0 | 00 |
| 00 |  | x | 1 | 1 | 01 |
| 01 | X | 0 | X | 0 | 01 |
| 01 | x | 1 | X | 1 | 10 |
|  | 1 | $x$ | 0 | 0 | 10 |
| 10 | 0 | $x$ | $x$ | 1 | 11 |
| 10 | X | X | 1 | 1 | 11 |
| 11 | X | 0 | X | 0 |  |
|  |  |  |  |  |  |


| $\mathbf{S}$ | $\mathbf{H}$ | $\mathbf{F}$ |
| :---: | :---: | :---: |
| 00 | 00 | 10 |
| 01 | 01 | 10 |
| 10 | 10 | 00 |
| 11 | 10 | 01 |

State and Output Encoding

## $\mathbf{S} \mathbf{C} \mathbf{S} \mathbf{L} \quad \mathbf{S}$

| 00 | 0 | $X$ | $X$ | 0 | 00 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | $X$ | $X$ | 0 | 0 | 00 |
| 00 | 1 | $X$ | 1 | 1 | 01 |
| 01 | $X$ | 0 | $X$ | 0 | 01 |
| 01 | $X$ | 1 | $X$ | 1 | 10 |
| 10 | 1 | $X$ | 0 | 0 | 10 |
| 10 | 0 | $X$ | $X$ | 1 | 11 |
| 10 | $X$ | $X$ | 1 | 1 | 11 |
| 11 | $X$ | 0 | $X$ | 0 | 11 |
| 11 | $X$ | 1 | $X$ | 1 | 00 |


| $\mathbf{S}$ | $\mathbf{H}$ | $\mathbf{F}$ |
| :---: | :---: | :---: |
| 00 | 00 | 10 |
| 01 | 01 | 10 |
| 10 | 10 | 00 |
| 11 | 10 | 01 |

State and Output Encoding

## $\mathbf{S} \mathbf{C} \mathbf{S} \mathbf{L}$ S'

| 00 | 0 | $X$ | $X$ | 0 | 00 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 00 | $X$ | $X$ | 0 | 0 | 00 |
| 00 | 1 | $X$ | 1 | 1 | 01 |
| 01 | $X$ | 0 | $X$ | 0 | 01 |
| 01 | $X$ | 1 | $X$ | 1 | 10 |
| 10 | 1 | $X$ | 0 | 0 | 10 |
| 10 | 0 | $X$ | $X$ | 1 | 11 |
| 10 | $X$ | $X$ | 1 | 1 | 11 |
| 11 | $X$ | 0 | $X$ | 0 | 11 |
| 11 | $X$ | 1 | $X$ | 1 | 00 |
| $\mathbf{S}$ | $\mathbf{H}$ | $\mathbf{F}$ |  |  |  |
| 00 | 00 | 10 |  |  |  |
| 01 | 01 | 10 |  |  |  |
| 10 | 10 | 00 |  |  |  |
| 11 | 10 | 01 |  |  |  |

$$
\begin{aligned}
T= & \overline{S_{0}}\left(\overline{S_{1}} C L+S_{1}(\bar{C}+L)\right)+ \\
& S_{0} S \\
S_{1}^{\prime}= & S_{0} S+S_{1} \overline{S_{0}} \\
S_{0}^{\prime}= & \overline{S_{0}}\left(\overline{S_{1}} C L+S_{1}(\bar{C}+L)\right)+ \\
& S_{0} \bar{S} \\
H_{1}= & S_{1}
\end{aligned}
$$

$$
H_{0}=\overline{S_{1}} S_{0}
$$

$$
F_{1}=\overline{S_{1}}
$$

$$
F_{0}=S_{1} S_{0}
$$

## State and Output Encoding



$$
\begin{aligned}
& T= \overline{S_{0}}\left(\overline{S_{1}} C L+S_{1}(\bar{C}+L)\right)+ \\
& S_{0} S \\
& S_{1}^{\prime}= S_{0} S+S_{1} \overline{S_{0}} \\
& S_{0}^{\prime}= \overline{S_{0}}\left(\overline{S_{1}} C L+S_{1}(\bar{C}+L)\right)+ \\
& S_{0} \bar{S} \\
& H_{1}= S_{1} \\
& H_{0}= \overline{S_{1}} S_{0} \\
& F_{1}= \overline{S_{1}} \\
& F_{0}= S_{1} S_{0}
\end{aligned}
$$

## State and Output Encoding

## $\mathbf{S} \mathbf{C} \mathbf{S} \mathbf{L} \quad \mathbf{S}$

 HG $0 \times X \quad 0 \quad$ HG HG $\quad X \quad X \quad 0 \quad 0 \quad H G$ HG $1 \begin{array}{llllll} & X & 1 & 1 & H Y\end{array}$ WY X $0 \quad X \quad 0 \quad$ MY WY $\quad$ X $1 \quad 1 \quad X \quad 1 \quad$ PGFF $1 \begin{array}{lllll} & X & 0 & 0 & F G\end{array}$ EG $\quad 0 \quad X \quad X \quad 1 \quad \mathrm{FY}$ EG $\quad X \quad X \quad 1 \quad 1 \quad F Y$ FY $\quad \mathrm{X} \quad 0 \quad \mathrm{X} \quad 0 \quad \mathrm{FY}$
FY $\quad \mathrm{X} \quad 1 \quad \mathrm{X} \quad 1 \mathrm{HG}$

| $\mathbf{S}$ | $\mathbf{H}$ | $\mathbf{F}$ |
| :--- | :--- | :--- |


| HG | G | $R$ |
| :--- | :--- | :--- |
| FY | $Y$ | $R$ |
| PG | $R$ | $G$ |
| FY | $R$ | $Y$ |

## State and Output Encoding

$\mathbf{S} \quad \mathbf{C} \mathbf{S} \mathbf{T} \quad \mathbf{S}$

| 0001 | 0 | $X$ | $X$ | 0 | 0001 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0001 | $X$ | $X$ | 0 | 0 | 0001 |
| 0001 | 1 | $X$ | 1 | 1 | 0010 |
| 0010 | $X$ | 0 | $X$ | 0 | 0010 |
| 0010 | $X$ | 1 | $X$ | 1 | 0100 |
| 0100 | 1 | $X$ | 0 | 0 | 0100 |
| 0100 | 0 | $X$ | $X$ | 1 | 1000 |
| 0100 | $X$ | $X$ | 1 | 1 | 1000 |
| 1000 | $X$ | 0 | $X$ | 0 | 1000 |
| 1000 | $X$ | 1 | $X$ | 1 | 0001 |


| $\mathbf{S}$ | $\mathbf{H}$ | $\mathbf{F}$ |
| :---: | :---: | :---: |
| 0001 | 001 | 100 |
| 0010 | 010 | 100 |
| 0100 | 100 | 001 |
| 1000 | 100 | 010 |

$$
\begin{aligned}
T= & S_{0} C L+S_{1} S+ \\
& S_{2}(\bar{C}+L)+S_{3} S \\
S_{3}^{\prime}= & S_{2}(\bar{C}+L)+S_{3} \bar{S} \\
S_{2}^{\prime}= & S_{1} S+S_{2}(\overline{\bar{C}}+L) \\
S_{1}^{\prime}= & S_{0} C L+S_{1} \bar{S} \\
S_{0}^{\prime}= & S_{0} \overline{(C L)}+S_{3} S \\
H_{R}= & S_{2}+S_{3} \\
H_{Y}= & S_{1}
\end{aligned}
$$

$$
H_{G}=S_{0}
$$

$$
F_{R}=S_{0}+S_{1}
$$

$$
F_{Y}=S_{3}
$$

$$
F_{G}=S_{2}
$$

## State and Output Encoding



