

# ANTLR

# ANTLR Lexer Specifications

Look like

```
class MyLexer extends Lexer;
options {
    option = value
}
```

```
Token1 : 'char' 'char' ;
Token2 : 'char' 'char' ;
Token3 : 'char' ('char')? ;
```

Tries to match all non-protected tokens at once.

# The Esterel LRM

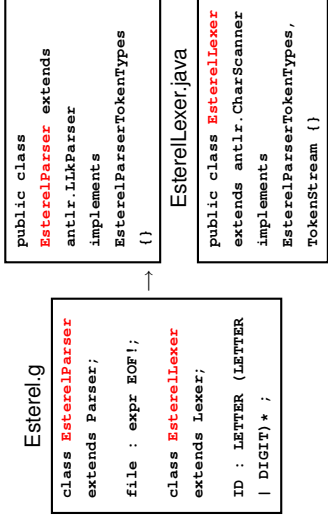
Lexical aspects are classical:

- Identifiers are sequences of letters, digits, and the underline character , starting with a letter.
- Integers are as in any language, e.g., 123, and floating-point numerical constants are as in C++ and Java; the values 12.3, .123E2, and 1.23E1 are constants of type double, while 12.3f, .123E2f, and 1.23E1f are constants of type float.
- Strings are written between double quotes, e.g., "a string", with doubled double quotes as in "a "" double quote".

# A Lexer for Esterel

Next, I wrote a rule for each punctuation character:

```
PERIOD : '.' ;
POUND : '#' ;
PLUS : '+' ;
DASH : '-' ;
SLASH : '/' ;
STAR : '*' ;
PARALLEL : "||" ;
```



# An ANTLR grammar for Esterel

Esterel: Language out of France. Programs look like

```
module ABRO:
input A, B, R;
output O;

loop
[ await A || await B ];
emit O
each R
end module
```

# A Lexer for Esterel

Operators from the language reference manual:

```
. # + - / * || < > , = ; := ( )
[ ] ? ?? <= >= <> =>
```

Main observation: none longer than two characters. Need  $k = 2$  to disambiguate, e.g., ? and ??.

```
class EsterelLexer extends Lexer;
options {
    k = 2;
}
```

# An ANTLR Grammar for Esterel

COMS W4115  
 Prof. Stephen A. Edwards  
 Fall 2006  
 Columbia University  
 Department of Computer Science

# ANTLR Parser Specifications

Look like

```
class MyParser extends Parser;
options {
    option = value
}
```

```
rule1 : Token1 Token2
      | Token3 rule2 ;
rule2 : (Token1 Token2)* ;
rule3 : rule1 ;
```

Looks at the next  $k$  tokens when deciding which option to consider next.

# The Esterel LRM

- Keywords are reserved and cannot be used as identifiers. Many constructs are bracketed, like "present ... end present". For such constructs, repeating the initial keyword is optional; one can also write "present ... end".
- Simple comments start with % and end at end-of-line. Multiple-line comments start with % { and end with }%.

## A Lexer for Esterel

Identifiers are standard:

```
ID
: ('a'..'z' | 'A'..'Z')
: ('a'..'z' | 'A'..'Z' | '_' | '0'..'9')*
```

## A Lexer for Esterel

String constants must be contained on a single line and may contain double quotes, e.g.,

"This is a constant with ""double quotes""

ANTLR makes this easy: annotating characters with ! discards them from the token text:

```
StringConstant
: '"'!
  ( ~('"' | '\n')
  | ('"'! '"'')
  )*
  '"'!
```

## A Lexer for Esterel

I got in trouble with the ~ operator, which inverts a character class. Invert with respect to what?

Needed to change options:

```
options {
  k = 2;
  charVocabulary = '\3'..'\'377';
  exportVocab = Esterel;
}
```

## A Lexer for Esterel

Another problem: ANTLR scanners check each recognized token's text against keywords by default.

A string such as "abort" would scan as a keyword!

```
options {
  k = 2;
  charVocabulary = '\3'..'\'377';
  exportVocab = Esterel;
  testLiterals = false;
}

ID options { testLiterals = true; }
: ('a'..'z' | 'A'..'Z') /* ... */;
```

## Numbers Defined

From the LRM:

Integers are as in any language, e.g., 123, and floating-point numerical constants are as in C++ and Java: the values 12.3, .123E2, and 1.23E1 are constants of type double, while 12.3f, .123E2f, and 1.23E1f are constants of type float.

## Numbers

With  $k = 2$ , for each rule ANTLR generates a set of characters that can appear first and a set that can appear second. But it doesn't consider the possible **combinations**.

I split numbers into Number and FractionalNumber to avoid this problem: if the two rules were combined, the lookahead set for Number would include a period (e.g., from ".1") followed by end-of-token e.g., from "1" by itself).

```
Example numbers:      First  Second
.1$                  .      EOT
.2                   1
1$                   2      1
```

## Number Rules

```
Number
: ('0'..'9')+
  ( '.' ('0'..'9')* (Exponent)?
  ( 'f'|'F') { setType(FloatConst); }
  | /* empty */ { setType(DoubleConst); }
  )
| /* empty */ { setType(Integer); }
```

## Number Rules Continued

```
FractionalNumber
: '.' ('0'..'9')+ (Exponent)?
  ( ('f'|'F') { setType(FloatConst); }
  | /* empty */ { setType(DoubleConst); }
  )
;

protected
Exponent
: ('e'|'E') ('+'|'-')? ('0'..'9')+
;
```

## Comments

From the LRM:

Simple comments start with % and end at end-of-line. Multiple-line comments start with %{ and end with }%.

## Comments

```
Comment
: '%'
( ('{' => '{'
  ( // Prevent .* from eating the whole file
    options {greedy=false};
  (
    ('\r' '\n') => '\r' '\n' { newline(); }
    | '\r'
    | '\n'
    | ~(''\n' | '\r' )
  )
)*
"}%"
| ((''\n')* '\n' { newline(); }
)
{ setType(Token.SKIP); }
```

## A Parser for Esterel

Esterel's syntax started out using ; as a separator and later allowed it to be a terminator.

The language reference manual doesn't agree with what the compiler accepts.

## Grammar from the LRM

*NonParallel:*  
*AtomicStatement*  
*Sequence*  
*Sequence:*  
*SequenceWithoutTerminator ; opt*  
*SequenceWithoutTerminator:*  
*AtomicStatement ; AtomicStatement*  
*SequenceWithoutTerminator ; AtomicStatement*  
*AtomicStatement:*  
*nothing*  
*pause*  
...

## Grammar from the LRM

But in fact, the compiler accepts

```
module TestSemicolon1:
nothing;
end module
module TestSemicolon2:
nothing; nothing;
end module
module TestSemicolon3:
nothing; nothing
end module
```

Rule seems to be "one or more statements separated by semicolons except for the last, which is optional."

## Grammar for Statement Sequences

Obvious solution:

```
sequence
: atomicStatement
( SEMICOLON atomicStatement ) *
( SEMICOLON ) ?
;
```

```
warning: nondeterminism upon
k==1:SEMICOLON
between alt 1 and exit branch of block
```

Which option do you take when there's a semicolon?

## Nondeterminism

```
sequence : atomicStatement
( SEMICOLON atomicStatement ) *
( SEMICOLON ) ? ;
Is equivalent to
sequence : atomicStatement seq1 seq2 ;
seq1 : SEMICOLON atomicStatement seq1
| /* nothing */ ;
seq2 : SEMICOLON
| /* nothing */ ;
```

## Nondeterminism

```
sequence : atomicStatement seq1 seq2 ;
seq1 : SEMICOLON atomicStatement seq1
| /* nothing */ ;
seq2 : SEMICOLON
| /* nothing */ ;
```

How does it choose an alternative in `seq1`?

First choice: next token is a semicolon.

Second choice: next token is one that may follow `seq1`.

But this may also be a semicolon!

## Nondeterminism

Solution: tell ANTLR to be greedy and prefer the iteration solution.

```
sequence
: atomicStatement
( options { greedy=true; }
: SEMICOLON! atomicStatement ) *
( SEMICOLON! ) ?
;
```

## Nondeterminism

Delays can be "A" "X A" "immediate A" or "[A and B]."

```
delay : expr bSigExpr
| bSigExpr
| "immediate" bSigExpr ;
```

```
bSigExpr : ID
| "[" signalExpression "]" ;
```

```
expr : ID | /* ... */ ;
```

Which choice when next token is an ID?

## Nondeterminism

```
delay : expr bSigExpr
      | bSigExpr
      | "immediate" bSigExpr ;
```

What do we really want here?

If the delay is of the form "expr bSigExpr," parse it that way.

Otherwise try the others.

## Nondeterminism

```
delay : ( (expr bSigExpr) => delayPair
         | bSigExpr
         | "immediate" bSigExpr
         ) ;
```

```
delayPair : expr bSigExpr ;
```

The => operator means "try to parse this first. If it works, choose this alternative."

## Greedy Rules

The author of ANTLR writes

I have yet to see a case when building a parser grammar where I did not want a subrule to match as much input as possible.

However, it is particularly useful in scanners:

```
COMMENT
: "/"* " ("*)* "*" /
;
```

This doesn't work like you'd expect...

## Turning Off Greedy Rules

The right way is to disable greedy:

```
COMMENT
: "/"*
  (options {greedy=false;} :)*
  "*" /
  ;
```

This only works if you have two characters of lookahead:

```
class L extends Lexer;
options {
  k=2;
}

CMT : "/"* (options {greedy=false;} :)* "*" /
;
```

## The Dangling Else Problem

```
class MyGram extends Parser;
```

```
stmt : "if" expr "then" stmt ("else" stmt)? ;
```

Gives

```
ANTLR Parser Generator Version 2.7.1
gram.g:3: warning: nondeterminism upon
gram.g:3: k=1: "else"
gram.g:3: between alts 1 and 2 of block
```

## Generated Code

```
stmt : "if" expr "then" stmt ("else" stmt)? ;
match(LITERAL_if);
expr();
match(LITERAL_then);
stmt();
if ((LA(1)==LITERAL_else)) {
  match(LITERAL_else); /* Close binding else */
  stmt();
} else if ((LA(1)==LITERAL_else)) {
  /* go on: else can follow a stmt */
} else {
  throw new SyntaxError(LT(1));
}
```

## Removing the Warning

```
class MyGram extends Parser;

stmt
: "if" expr "then" stmt
  (options {greedy=true;} : "else" stmt)?
;
```

## A Simpler Language

```
class MyGram
  extends Parser;

  match(LITERAL_if);
  expr();
  match(LITERAL_then);
  stmt();
  switch (LA(1)) {
  case LITERAL_else:
    match(LITERAL_else);
    stmt();
    break;
  case LITERAL_fi:
    break;
  default:
    throw new SyntaxError(LT(1));
  }
  match(LITERAL_fi);
```