Authentication

- password-based authentication
- address-based authentication
- cryptographic protocols
- passwords as keys
- eavesdropping
- trusted intermediaries
- session key establishment
- delegation

Slide 1

Password

- proof by knowledge, sharing
- eavesdropping
- needed for dumb end systems
- cellular phone cloning
- single password across multiple hosts

Password Guessing

on-line: limit tries, delay, alarm

off-line: dictionary attack \Longrightarrow capture f(p)

- 1. Your first, last, or kid's name
- 2. "secret"
- 3. stress-related words ("deadline", "work")
- 4. sports teams or terms ("bulls", "golfer")
- 5. "payday"
- 6. "bonkers"
- 7. The current season ("winter", "spring")
- 8. Your ethnic group

- 9. repeated characters ("aaaaa", "bbbbb")
- 10. obscenities, sexual terms

Storing Passwords

per-node: /etc/passwd

server: authentication storage server, retrieved by node (yp/NIS)

facilitator: server says yes/no

need to authenticate node asking

• store hash only

• store encrypted with good, protected key

• but: needs to be in non-volatile memory (ROM?)

Slide 5

Address-Based Authentication

- rcp, rsh: .rhosts node, user name
- per user
- reverse-lookup on IP address (in-addr.arpa)
- can use different login names
- /etc/hosts.equiv: trusted hosts

Address-Based Authentication: Threats

- break in one, break in all
- often: A trusts B, B trusts A
- address spoofing; not easy for connections, but "blind" sending
- easy to listen/send on broadcast network
- MAC address spoofing prevention: filter on port, scramble

Source routing to have T spoof A: $\langle A, T, D \rangle \implies \langle D, T, A \rangle$

Slide 7

Humans and Computers

humans: short, memorable key (8 characters, 48 bits) directly or as key for longer key (PGP, Netscape)

computers: hidden key, directly

Passwords as Keys

- directly as 56-bit key (e.g., use words)
- can't use for RSA p, q:
 - use as seed for rng
 - "simulation-style" rng, until primes found
 - do once, then give offset hints to user

Slide 9

Eavesdropping

- public key: need to secure Alice's private key
- use random challenge with signing
- difficult to protect against eavesdropping and disclosure * Lamport, S/Key

Trusted Intermediaries

- can't do pairwise authentication with secret keys: key explosion!
- Wey Distribution Center (KDC)
 - KDC knows all secrets
 - α asks KDC for secret (securely) to talk to any other node β
 - hand out session key $R_{\alpha\beta}$: ticket
 - single point of failure
 - bottleneck

Slide 11

Trusted Intermediaries: CA

CA: ensure validity of public keys

- small number, preconfigured
- CA: single PoF
- CA: typically off-line, protected
- certificates are not sensitive
- compromised CA cannot eavesdrop
- need revocation list (CRL) must be signed and recent

Multiple KDC Domains

Secret keys:

- KDCs share pairwise key
- topology of KDC: tree with shortcuts

Public keys:

- cross-certification of CAs
- example: Alice with CA_A, Boris CA_B
 - Alice gets CA_B's certificate signed by CA_A
 - Alice gets Boris' certificate signed by CA_B

Slide 13

Session Key Establishment

- use public keys to authenticate, generate private key
- trade-off: processing, exposure
- limit lifetime imit replay attacks
- only need to expose short-term key to semi-trusted software

Authorization

- authentication: *identity* (who)
- authorization: *capability* (what)
- may be implied (physical access)
- network: authentication access control list (ACL)
- groups: central server, signed certificate
- certificate: unwieldy, CRLs
- hierarchical groups
- typical: hiearchy (DH, director, ...) and organization

Slide 15

Solaris ACLs

- setfacl -r -m user:czen:r-- file
- default entries per directory
- getfacl:

```
# file: papers
# owner: hgs
# group: faculty
user::rwx
group::r-x #effective:r-x
group:irt:r-x #effective:r-x
other:---
```

Delegation

- short-term authorization for principals
- sign "letter of authority" (delegation)
- limit time, scope